Verification Assignment

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System Verification & Testing

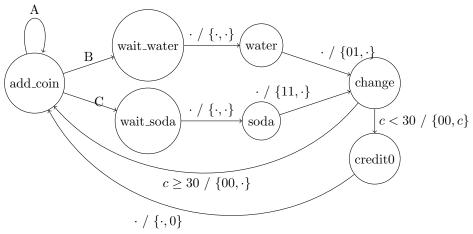
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1 RTL Design

1.1 Finite State Machine

The output is represented as the set: {beverage_out,change_out}, where the symbol · represents an unchanged signal.

In order to keep the graphical representation as clean as possible, I've decided to label with letters the transition that have long descriptions and explain them in a bulleted list.



Where:

- A: $coin_i \in \{10c, 20c, 50c, 1 \in 20\} / \{\cdot, \cdot\}$
- **B**: $c \ge 30 \&\&$ button_in == 01 / $\{\cdot, \cdot\}$
- C: $c \ge 50 \&\& button_in == 11 / \{\cdot, \cdot\}$

1.2 Implementation choices

- wait_water and wait_soda implement the delay between the selection and the delivery of a beverage. Without them, the machine would be forced to stay in one of the other states: for example, if the delay was computed in add_coin, the FSM would accept coins while the delay is passing; if it was computed in water or soda, the signal beverage_out would be high for too long
- credit0 is necessary because change_out <= c and c <= 0 can't be computed in the same clock cycle: the change would be set to zero every time
- \bullet If there is no change after the delivery of a beverage, the state machine accepts an input (a new coin or a pressed button) without the delay M

2 Verilog-style Test Bench

After initializing the machine and adding some coins, the test bench inputs an illegal button code but it gets ignored. Then, it starts acting like a normal user until the final part, in which it inputs a coin while the vending machine is returning the change. Just like the other occurrence, the machine manages to ignore the wrong timing of the input.

3 Constrained Random Test Bench

3.1 Transaction constraints

- coin_in has to be in the set of admitted coins: $\{10c, 20c, 50c, 1 \in 20\}$
- button_in = 01 or 11 implies coin_in = 0: this allows the FSM's state to evolve correctly
- button_in is set to be no choice (00), water (01) or soda (11)

3.2 Asserted properties

```
    G((state = water)->(nexttime[1] beverage_out = 01))
    G((state = soda)->(nexttime[1] beverage_out = 11))
    G((state = add_coin && c < 30)->(nexttime[1] state = add_coin))
    G((state = credit0)->(change_out > 0 && change_out < 30))</li>
    G((c ≥ 30 && state = change)->(nexttime[1] change_out = 0))
```

3.3 Contingency tables

The contingency tables were computed performing the negation of the antecedent and the consequent using De Morgan's laws. To implement them in the code, there are four instances of each property. For example, regarding the first one:

- p0 : ATCT
- pon_: AFCT (N indicates the antecedent is negated, _ as the second character indicates the consequent is untouched)
- p0_N : ATCFp0NN : AFCF

The results are divided in two cases: the first one has the standard constraints for the input values. The second one uses a subset of allowed coins in which 1 \in and $2 \in$ coins were removed. This choice was taken to allow a higher number of cases in which some change has to be given. With them included in the set, the credit in the machine would be always higher than 30 cents and the ATCT of p3 is zero.

Additionally, only property p3 can be asserted 1000 times: the other properties get covered 999 times because they contain a reference to the following time instant, which doesn't exist in the last clock period of the simulation. Finally, property p4 has a very high number of AFCTs because when there is no output change, its signal is set to zero.

Contingency tables with the whole set of coins

$\frac{\mathbf{AT}}{\mathbf{AF}}$	CT 59	0 94	_	$\frac{\overline{\mathbf{AT}}}{\mathbf{AF}}$	60 0	CF 0 939		$\frac{\mathbf{AT}}{\mathbf{AF}}$	3 520	C		
Table 1: p0				Та	Table 2: p1			Table 3: p2				
		\mathbf{CT}	\mathbf{CF}					CT	CI	7		
$\overline{\mathbf{A}}$	\mathbf{T}	0	0				$\overline{\mathbf{AT}}$	119	0			
$\overline{\mathbf{A}}$	F	0	1000				$\overline{ m AF}$	880	0	_		

Table 4: p3 (Note: ATCT = 0) Table 5: p4

Contingency tables without 1 and $2 \in \text{coins}$

$\frac{\mathbf{AT}}{\mathbf{AF}}$	CT C1 68 0 0 93	_	$\frac{\overline{\mathbf{AT}}}{\mathbf{AF}}$	68 0	0 931	_	AT AF	CT 22 415	CF 0 562
Tab	le 6: p0		Table 7: p1			Table 8: p2			
$\frac{\overline{\mathbf{A}}}{\mathbf{A}}$	_	CF 0 992				$\frac{\overline{\mathbf{AT}}}{\mathbf{AF}}$	130 861	0 8	_
				Table 10: p4					

4 HARM

The tool extracted the properties p0 and p1, in their exact form. This was obtained by removing three templates and leaving only two that were similar to the manually defined assertions. A filter was applied in order to obtain a manageable set of results: it selects the assertions that have at least 90% of causality.

As expected, the credit quickly rises in the .vcd trace that includes every coin and the interesting assertions will be placed at rows 6 and 7, due to their low frequency.



Figure 1: Assertions with the full set of coins

However, the trace with the subset of coins allows the same properties to be asserted more often, therefore they're placed on top of the table.

Figure 2: Assertions without 1 and $2 \in \text{coins}$

These results, as well as the contingency tables of the previous section, are contained in Part_3/GoodAssertions.txt.