

# From Fine Analysis to Instance Optimality

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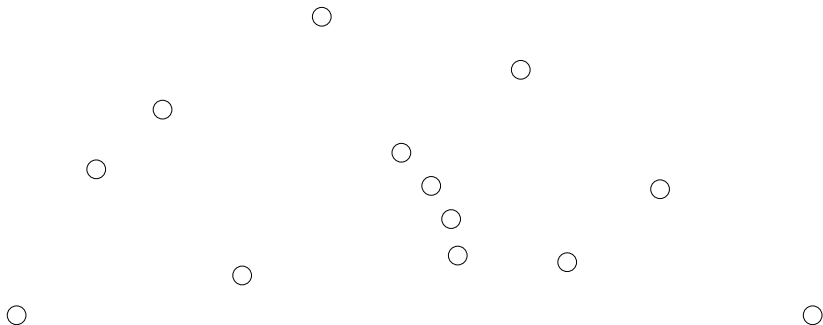
# Outline

- 1 The Convex Hull Paradox
  - $O(n \lg n)$
  - $O(nh)$
  - Worst Case Complexity?
- 2 Fine analysis of the convex hull
  - $O(n \lg h)$  in 2D
  - $O(n \lg h)$  in 3D
  - $O(nH(C))$ , instance optimal
- 3 Similar Paradoxes

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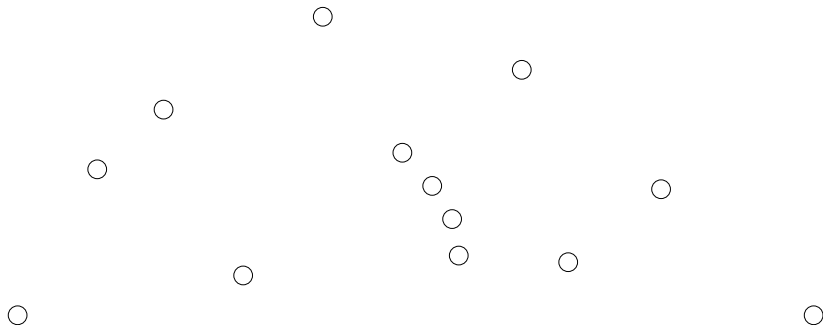
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# The Planar Convex Hull



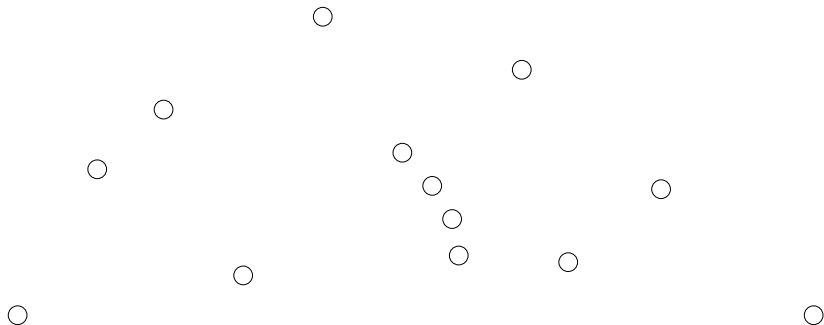
- Can one of you define it?
- What is the best complexity known for it?

## 2d Convex Hull in $O(n \lg n)$



- 1 Sort the points by  $x$ -coordinates;
- 2 Scan them, backtracking if necessary.

## 2d Convex Hull in $O(nh)$



- 1 Find the left-most point
- 2 Compute the  $n - 1$  slopes with the other points
- 3 Choose the highest slope
- 4 Iterate

# Worst case Complexity of 2d Convex Hull

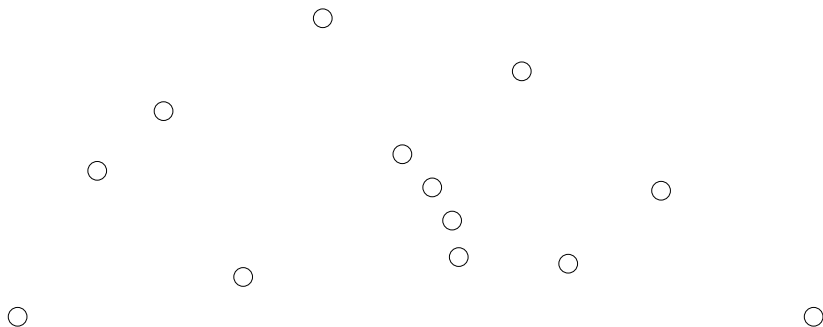
- Question ill-defined: the worst case over what?
  - all instances of fixed size  $n$ ?
  - all instances of fixed input size  $n$  and output size  $h$ ?
- For each we have distinct lower bounds:
  - $\Omega(n \lg n)$ , which is tight; and
  - $\Omega(n \lg h)$ , which is **not** tight!
- So what is the complexity of 2d convex hull?

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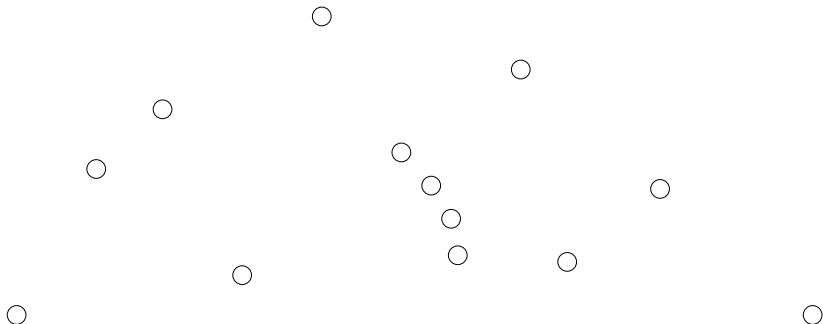


## 2d Convex Hull in $O(n \lg h)$ in 2D



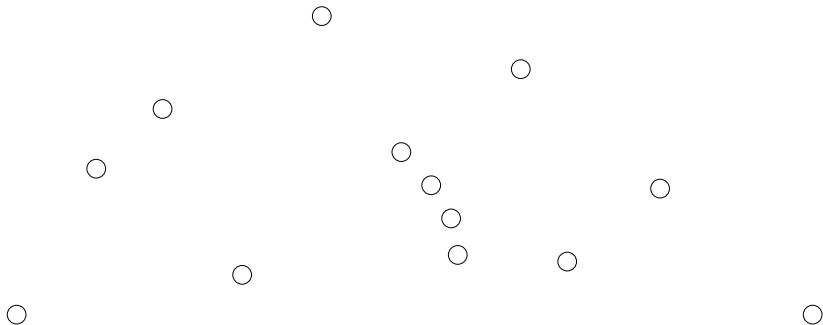
- 1 Compute the point  $m$  of median  $x$ -coordinate;
- 2 Partition the points by  $m.x$ ;
- 3 Compute the highest edge  $(a, b)$  intersecting the line  $x = m.x$ ;
- 4 Recurse on each side;

## 2d Convex Hull in $O(n \lg h)$ in 3D



- 1 Start with a small guess for  $h$ ;
- 2 Group the instances in  $n/h$   $x$ -sorted groups of size  $h$ ;
- 3 Simulate the  $O(nh)$  algorithms on the groups;
- 4 If it did not suffice, merge the group two by two and iterate.

# Convex Hull in $O(n(1 + H(C)))$



- Algorithm: a variant of [Kirkpatrick, Seidel]
  - 1 Compute the points leftmost  $l$  and rightmost  $r$ ;
  - 2 Compute the point  $m$  of median  $x$ -coordinate;
  - 3 Compute the highest edge  $(a, b)$  intersecting the line  $x = m$ ;
  - 4 Remove all points contained in the polygon  $(l, a, b, r)$ ;
  - 5 Recurse on each side;

## Instance Optimality: definitions

### Definition (Instance Optimality)

An algorithm is **instance-optimal** if its cost is at most a constant factor from the cost of any other algorithm  $A'$  running on the same input, for every input instance.

Unfortunately, for many problems, this requirement is too stringent.

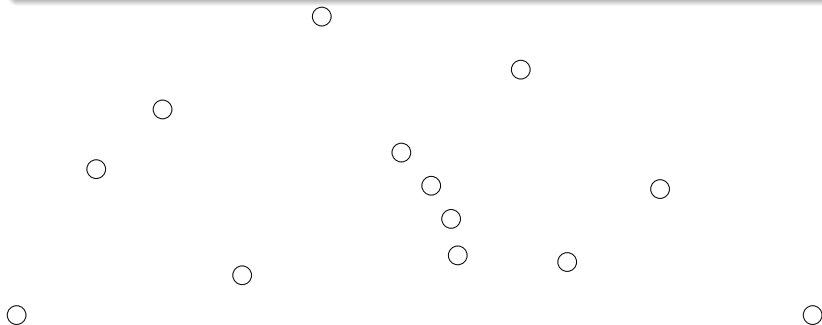
### Definition (Input Order Oblivious Instance Optimality)

For a set  $S$  of  $n$  elements in  $\mathcal{D}$ , let  $T_A(S)$  denote the maximum running time of  $A$  on input  $\sigma$  over all  $n!$  possible permutations  $\sigma$  of  $S$ . Let  $\text{OPT}(S)$  denote the minimum of  $T_{A'}(S)$  over all correct algorithms  $A' \in \mathcal{A}$ . If  $A \in \mathcal{A}$  is a correct algorithm such that  $T_A(S) \leq O(1) \cdot \text{OPT}(S)$  for every set  $S$ , then we say  $A$  is **instance-optimal in the order-oblivious setting**.

# Certificate and Instance Optimal Proof

## Definition (Certificate)

A *Certificate* for an instance  $I$  and a solution  $S$  is the description of a sequence of steps to **check** the validity of  $S$  for  $I$ .



## Example

For the convex hull, a list of triangles and the points they cover.

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## Optimal Prefix Free Codes [Closed]

- $O(n \lg n)$  classical algorithm;
- $O(n)$  algorithm when frequencies are sorted;
- $O(nk)$  for  $k$  distinct code lengths.
- A lower bound of  $\Omega(n \lg k)$  in the worst case over instances resulting in  $k$  distinct code lengths;
- $O(nH(n_1, \dots, n_k))$  instance optimal algorithm and
- $O(n)$  algorithm in word-RAM [Barbay]!

## Optimal Minimax Trees [Semi Closed]

- Tree minimizing the max weight+height of a leaf.
- $O(n \lg n)$  classical algorithm [Golumbic];
- $O(n)$  algorithm when weights partially sorted by fractional part [Drmota, Szpankowski];
- $O(nd \lg \lg n)$  where  $d$  is the number of distinct values  $\lceil w_i \rceil$  [Kirkpatrick and Klawe]
- $O(n)$  algorithm in word-RAM [Gawrichowski, Gagie]!



# Optimal Alphabetic Binary Search Tree [Open]

- $O(n \lg n)$  classical Hu-Tucker algorithm;
- $o(n \lg n)$  algorithms in many particular cases;
- $O(n)$  algorithm when frequencies “can be sorted in linear time”;
- A lower bound of  $\Omega(n \lg k)$  in the worst case over instances resulting in  $k$  distinct code lengths.

# Summary

- $O(nk)$  and  $O(n \lg n)$  suggests  $O(n \lg k)$
- and (Input Order Oblivious) **Instance Optimality**.
- Outlook
  - Input Order Adaptive Instance Optimality.
  - Full Instance Optimality (Kolmogorov's complexity?)
  - 1-instance optimality!