# COMP1917: Computing 1

## 15. Stacks and Queues

Reading: Moffat, Section 10.1-10.2

#### **Overview**

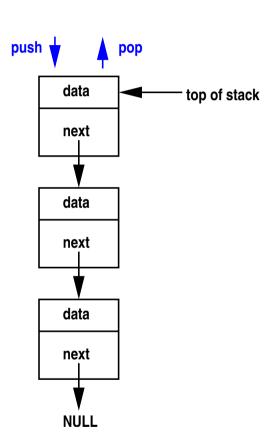
- Stacks
- Queues
- Adding to the Tail of a List
- Efficiency Issues
- Queue Structure
- Stack Application: Postfix Calculator

#### **Stacks and Queues**

- Stacks and Queues are examples of Abstract Data Types
- Stacks and Queues are used in many computing applications, as well as forming auxiliary data structures for common algorithms, and appearing as components of larger structures.

#### **Stacks**

- a stack is a collection of items such that the last item to enter is the first one to exit, i.e. "last in, first out" (LIFO)
- based on the idea of a stack of books, or plates



#### **Stack Functions**

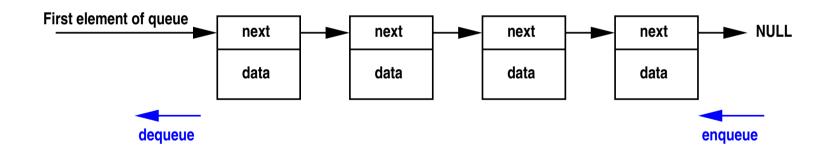
- **Essential Stack functions:** 
  - push() // add new item to stack
  - pop() // remove top item from stack
- Additional Stack functions:
  - ▶ top() // fetch top item (but don't remove it)
  - ▶ size() // number of items
  - ▶ isEmpty()

## **Stack Applications**

- page-visited history in a Web browser
- undo sequence in a text editor
- checking for balanced brackets
- HTML tag matching
- postfix calculator
- chain of function calls in a program

#### Queues

- a queue is a collection of items such that the first item to enter is the first one to exit, i.e. "first in, first out" (FIFO)
- based on the idea of queueing at a bank, shop, etc.



#### **Queue Functions**

- Essential Queue functions:
  - enqueue() // add new item to queue
  - ▶ dequeue() // remove front item from queue
- Additional Queue functions:
  - front() // fetch front item (but don't remove it)
  - ▶ size() // number of items
  - ▶ isEmpty()

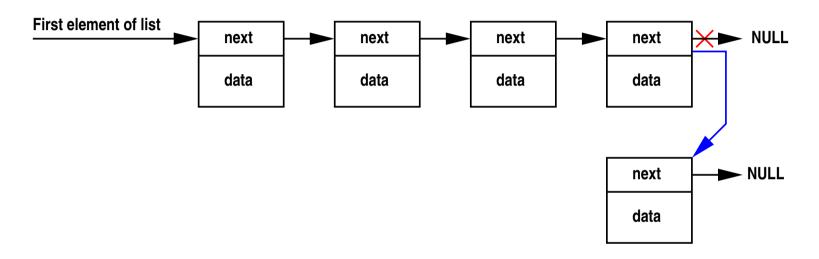
## **Queue Applications**

- waiting lists, bureaucracy
- access to shared resources (printers, etc.)
- phone call centres
- multiple processes in a computer

## **Implementing Stacks and Queues**

- a stack can be implemented using a linked list, by adding and removing at the head [push() and pop()]
- for a queue, we need to either add or remove at the tail
  - can either of these be done efficiently?

## Adding to the Tail of a List



- adding an item at the tail is achieved by making the last node of the list point to the new node
- we first need to scan along the list to find the last item

## Adding to the Tail of a List

```
Lnode * add_to_tail( Lnode *new_node, Lnode *head )
  head = new_node;
  else {
                         // list not empty
     Lnode *node = head;
     while( node->next != NULL ) {
         node = node->next; // scan to end
     node->next = new_node;
  return( head );
```

### **Efficiency Issues**

Unfortunately, this implementation is very slow. Every time a new item is inserted, we need to traverse the entire list (which could be very large).

We can do the job much more efficiently if we retain a direct link to the last item or "tail" of the list:

Note: there is no way to efficiently remove items from the tail. (Why?)

#### **Queue Structure**

We can use this structure to implement a queue efficiently:

```
typedef struct queue Queue;
struct queue {
    Lnode *head;
    Lnode *tail;
    int size;
};
```

### Making a new Queue

```
Queue * makeQueue()
   Queue *q = (Queue *)malloc( sizeof( Queue ));
   if( q == NULL ) {
      fprintf(stderr, "Error, failed to allocate Queue.\n");
      exit( 1 );
   q->head = NULL;
   q->tail = NULL;
   q \rightarrow size = 0;
   return( q );
```

### Adding a new Item to a Queue

```
void enqueue( Lnode *new_node, Queue *q )
   if( q->tail == NULL ) { // queue is empty
       q->head = new_node;
  else {
                          // queue not empty
       q->tail->next = new_node;
  q->tail = new_node;
  q->size++;
```

### Removing an Item from a Queue

```
Lnode * dequeue( Queue *q )
   Lnode *node = q->head;
   if( q->head != NULL ) {
       if( q->head == q->tail ) { // only one item
           q->tail = NULL;
       q->head = q->head->next;
       q->size--;
   return( node );
```

### **Example:** queue.c

```
int main( void )
  Queue *q = makeQueue();
  Lnode *node;
   int ch;
  while(( ch = getchar()) != EOF ) {
      if( ch == '-' ) {
          node = dequeue( q );
          if( node != NULL ) {
              printf("Dequeueing %c\n", node->data );
              free( node );
```

## Example: queue.c

```
. . .
   else if( ch == '\n' ) {
       printList( q->head );
   else {
        enqueue( makeNode(ch), q );
freeList( q->head );
return 0;
```

#### **Reverse Polish Notation**

Some early calculators and programming languages used a convention known as Reverse Polish Notation (RPN) where the operator comes after the two operands rather than between them:

```
1 2 +
result = 3
3 2 *
result = 6
4 3 + 6 *
result = 42
1 2 3 4 + * +
result = 15
```

#### **Postfix Calculator**

A calculator using RPN is called a Postfix Calculator; it can be implemented using a stack:

- when a number is entered: push it onto the stack
- when an operator is entered: pop the top two items from the stack, apply the operator to them, and push the result back onto the stack.

#### postfix.c

```
int main( void )
  Lnode *list = NULL;
  int num;
  int a,b, num;
  while(( ch = getc(stdin)) != EOF ) {
     if( ch == '\n' ) {
         printf("Result: %d\n", list->data );
     else if( isdigit(ch)) {
         ungetc( ch, stdin ); // put first digit back to file
         scanf( "%d", &num ); // now scan entire number
         list = push( makeNode(num), list );
```

#### postfix.c

```
else if( ch == '+' || ch == '-' || ch == '*') {
    if( list != NULL ) {
       a = list->data;  // fetch top item
       list = pop( list );
       if( list != NULL ) {
           b = list->data; // fetch 2nd item
           list = pop( list );
           switch( ch ) {
             case '+': num = b + a; break;
             case '-': num = b - a; break;
             case '*': num = b * a; break;
           list = push( makeNode(num), list );
```