Quiz (Week 1)

Typing

Assuming for the sake of simplicity that all numeric literals are of type Int, What is the type of the following Haskell expressions?

Question 1

```
"hello world"

1. X string
2. ✓ [Char]
3. X char*
4. X [String]
```

In Haskell, strings are actually just lists of characters, and the names of types (like Char) are always written with an initial upper-case letter.

Question 2

In Haskell, a tuple (x,y) is typed according to the following rule:

$$\frac{x:\tau_1 \quad y:\tau_2}{(x,y):(\tau_1,\tau_2)}$$

This can be read as (x,y) is of type (τ_1, τ_2) if x is of type τ_1 and y is of type τ_2 . A similar rule exists for triples like (1, 'x', [True]), and as 1 is of type [Int], [True] is of type [Int], we have answer 4 as the

only correct answer.

Question 3

4. X String

Keeping in mind that String is a synonym for [Char], we have the type of cons (the (:) operator) as:

```
(:) :: a -> [a] -> [a]
```

If we apply the first argument, "x":

```
("x":) :: [[Char]] -> [[Char]]
```

Then apply the second argument, [], and we have:

```
("x":[]) :: [[Char]]
```

Lastly, this list of list of characters is in turn put in a list, as it is surrounded by square brackets. So the answer is number 2, or a list of lists of characters.

Question 4

```
map (\x → x + 1)

1. X [a] → [b]
2. X Int → Int
3. ✓ [Int] → [Int]
4. X (Int → Int) → [Int] → [Int]
5. X Invalid, as not enough arguments are given to map.
```

It's worth noting that *all functions in Haskell accept one argument and return one result*. Multi-argument functions are emulated by writing a function that, given its first argument, returns a *function* that awaits further arguments. This technique is called *currying*.

For example, the function map has the following type:

```
map :: (a -> b) -> [a] -> [b]
```

This can be more explicitly expressed with the right-associated parentheses, as follows:

```
map :: (a -> b) -> ([a] -> [b])
```

Given the argument function $(\x -> x + 1)$, a lambda expression of type Int -> Int, map shall return a function of type [Int] -> [Int], or option 3.

Evaluation

Choose all expressions that are equivalent to the following expressions¹:

Question 5

```
3: [40] ++ [50] ++ 5: [60]

1.  3: [40] ++ ([50] ++ 5: [60])

2.  3: ([40] ++ [50] ++ 5): [60]

3.  (3: [40] ++ [50]) ++ (5: [60])

4.  3: ([40] ++ [50] ++ (5: [60]))

5.  3: [40, 50] ++ [5, 60]
```

It's important to note that the (++) operator is associative, that is:

```
(xs ++ ys) ++ zs == xs ++ (ys ++ zs)
```

This can be proven by induction on \times s, with the aid of some helper lemmas. Because of this associativity, the placement of parentheses around ++-terms is not important, which makes options 1,3 and 4 correct. In addition, option 5 is also correct as we know that [40] ++ [50] = [40,50] and that [5:[60] = [5,60].

```
map ($ 5) [(-),(+),(*)]
```

- 1. \times map (\f x -> f x 5) [(-),(+),(*)]
- 2. \checkmark map (\f x -> f 5 x) [(-),(+),(*)]
- $3. \times [(-5), (+5), (*5)]$
- 4. \checkmark [(5 -),(5 +),(5 *)]
- 5. X The expression is invalid.

The (\$) operator is defined as follows:

```
($): (a -> b) -> a -> b
f $ x = f x
```

```
[\x -> (-) 5 x, \x -> (+) 5 x, \x -> (*) 5 x]
```

Which is equivalent to the operator sections used in answer 4. Answers 1 and 3 are incorrect as they flip the order of arguments used for the function. Answer 3 is even more incorrect as (-5) will be interpreted as a negative number, not an operator section, and thus produce a type error.

Question 7

Note: The functions ord and chr are from Data. Char They convert Char values to/from their ASCII (or unicode) numbers, respectively. For these questions, the answers may have a more general type than the original expression. So long as a given answer has equivalent behaviour *for the type of the original expression*, we consider the answer to be equivalent.

```
let increment x = 1 + x
in \xs -> map chr (map increment (map ord xs))
```

```
1. \checkmark map chr \cdot map (1+) \cdot map ord
```

```
2.  map (chr . (1+) . ord)
3.  map succ
4.  map chr $ map (1+) $ map ord
5.  \( \cdot \xx -> map chr . map (1+) $ map ord xs
```

The following bit of equational reasoning hits every answer in this question, except 4, which is not type correct.

```
let increment x = 1 + x
in \xs -> map chr (map increment (map ord xs))
= -- Shift argument into lambda
let increment = \xspace x -> 1 + x
in \xs -> map chr (map increment (map ord xs))
= -- Simplify lambda to operator section
let increment = (1 +)
in \xs -> map chr (map increment (map ord xs))
= -- Reduce let expression by substitution
\xs \rightarrow map chr (map (1 +) (map ord xs))
= -- Introduce composition, f(g x) = (f \cdot g) x
\xs \rightarrow (map chr \ map (1 +)) (map ord xs))
= -- Remove parentheses with ($)
\xs -> map chr \ map (1 +) \ map ord xs -- Answer 5
\xs \rightarrow (map chr map (1 +) map ord) xs
= -- \eta-reduction
map chr \cdot map (1 +) \cdot map ord -- Answer 1
= -- Map (functor) law, map f . map g = map (f . g)
map (chr.(1+).ord) -- Answer 2
= -- succ is defined for Char values as chr . (1 +) . ord
map succ -- Answer 3
```

Question 8

```
foldr (&&) True . map (>= 0)
1. ✓ and . map (>= 0)
2. ✓ all (>= 0)
3. ✗ any (>= 0)
4. ✓ foldr (\a b -> a >= 0 && b) True
5. ✗ foldl (\a b -> a && b > 0) True
```

The following derivation shows the equivalence to answers 1 and 2.

```
foldr (&&) True . map (>=0)

= -- and = foldr (&&) True

and . map (>=0) -- Answer 1

= -- all f = and . map f

all (>=0) -- Answer 2
```

Furthermore, Answer 4 is also equivalent, as the following derivation shows:

```
foldr (&&) True . map (>=0)
= -- \eta - expansion \ on \ the \ (\&\&)
foldr (\a b -> a && b) True . map (>=0)
= -- \ We \ have \ a \ fold/map \ rule
-- foldr (\x y -> z) \ . map \ f = foldr \ (\x y -> z[x := f x])
-- (where \ z[x := f x] \ is \ a \ substitution).
foldr (\a b -> (>= 0) a && b) True
= -- \ Nicer \ syntax
foldr (\a b -> a >= 0 && b) True -- Answer 4
```

Footnotes:

¹: By "equivalent", we mean will evaluate to equal results. We consider two functions equal if, for any input, they produce equal outputs (functional extensionality).

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