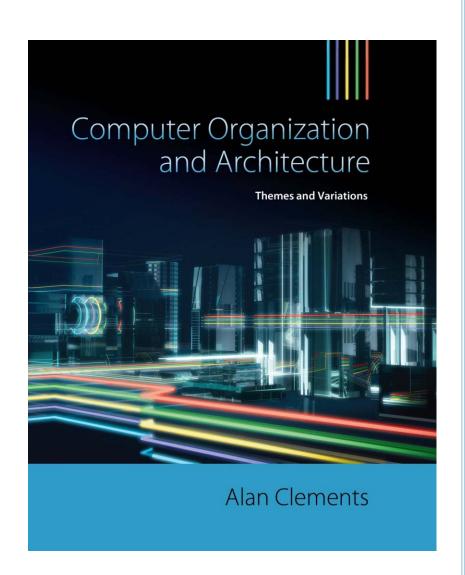
# CHAPTER 4

Computer
Organization
and
Architecture





## ISAs Breadth and Depth

- ☐ This chapter extends the overview of ISAs in both breadth and depth.
  - Yet, we will only cover the depth part in lectures this term
- ☐ In particular, we will look at the role of the stack and architectural support for subroutines and parameter passing.

- ☐ Let's begin by looking at some background issues concerning *data storage*, *procedures*, and *parameter passing*.
- □ Computer programs consist of
  - data elements and
  - procedures which operate on these data elements
- ☐ High-level language programmers use variables to represent *data elements*
- □ Variables are declared by:
  - *assigning* names to them and by *reserving* storage for them.
- □ *Reserving* memory storage for variables can be performed
  - at compilation time (static memory allocation)
  - at runtime (dynamic memory allocation)
- □ Statically allocated variables will have static addresses which will not be changed during execution
- □ Dynamically allocated variables will have dynamic addresses which will be changed during execution, as they will be allocated at runtime

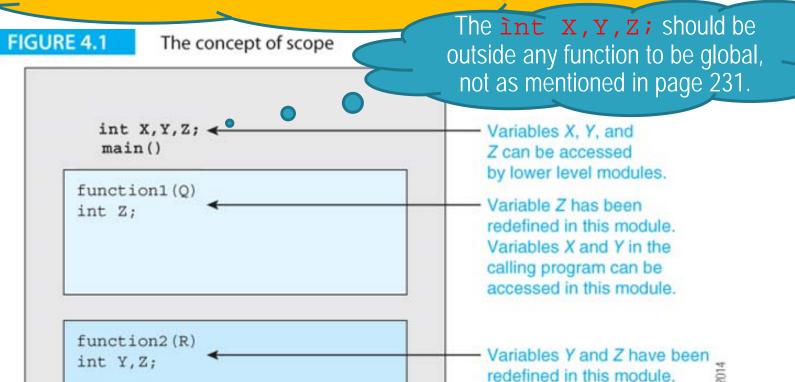
- ☐ Procedures often require *local workspace* for their temporary variables.
- ☐ The term *local* means that the workspace is private to the procedure and is never accessed by the calling program or by other subroutines.
- ☐ If a procedure is to be made re-entrant or to be used recursively, its local variables must be bounded up not only with the procedure itself, but with the occasion of its use.
  - Each time the procedure is called, a new workspace must be assigned to it.

- $\square$  A variable has a *scope* associated with it.
  - The scope of a variable defines the range of its *visibility* or *accessibility* within a program.
    - o *Global* variables are *visible* (accessible) from the moment they are loaded into memory to the moment when the program stops running *(static memory allocation)*
    - Variables declared within a procedure, as well as parameters are visible (accessible) within that procedure but invisible (inaccessible) outside the procedure (dynamic memory allocation)
- ☐ Here, we are interested to learn more about *dynamic memory allocation*

☐ Figure 4.1 illustrates the scope of variables

The duration of local variables and parameters are "automatically"

- allocated when the enclosing function is called and
- deallocated when the function returns



Global variable X can be accessed in this module.

#### Storage and the Stack

- ☐ When a language invokes a procedure, it is said to *activate* the procedure.
- Associated with each invocation (activation) of a procedure, there is an *activation record* containing all the information necessary to execute the procedure, including
  - parameters,
  - local variables, and
  - return address,