

Algorithm foo(n)

```

C1 { i ← 1
      K ← 1
      • while i < n do {
          { if i = K then {
              A[i] ← K
              K ← K + 1
              i ← 1
            }
          }
          i ← i + 1
        }
      }
  
```

Let us count the number of iterations performed by the while loop:

	# iterations
K = 1, i = <u>1</u>	1
K = 2, i = 1 → <u>2</u>	1
K = 3, i = 1 → 2 → <u>3</u>	2
K = 4, i = 1 → 2 → <u>3</u> → <u>4</u>	3
K = 5, i = 1 → <u>2</u> → <u>3</u> → <u>4</u> → <u>5</u>	4

⋮	
K = n-1, i = 1 → <u>2</u> → <u>3</u> → <u>4</u> → ⋯ → <u>n-1</u>	n-2
K = n, i = 1 → <u>2</u> → <u>3</u> → <u>4</u> → ⋯ → <u>n-1</u> → n	<u>n-2</u>

Each underlined number means one iteration of the while loop where the value of K is compared with the underlined value

$$\begin{aligned}
 \text{total} &= \underline{1} + 1 + 2 + 3 + \dots + \\
 &\quad n-2 + \underline{n-2} \\
 &= 1 + 2 + 3 + \dots + n-2 \\
 &\quad + n-1
 \end{aligned}$$

The total number of iterations is $1 + 2 + 3 + \dots + n-2 + n-1 = \underbrace{\sum_{j=1}^{n-1} j}_{\text{Arithmetic sum}} = \frac{n(n-1)}{2} = \frac{n^2}{2} - \frac{n}{2}$

The time complexity of the algorithm is then $C_1 + C_2 \times \# \text{ iterations} = C_1 - \frac{C_2}{2}n + \frac{C_2}{2}n^2$ is $O(n^2)$

$Kf(n)$ is $O(f(n))$ for any constant $K > 0$:

Find constants $c > 0$, $n_0 \geq 1$ such that

$$Kf(n) \leq \underline{c} f(n) \quad \forall n \geq \underline{n_0}$$

$$0 \leq cf(n) - Kf(n) = (c-K)f(n), \quad \forall n \geq \underline{n_0}$$

$$\boxed{c=K}: \quad 0 \leq 0f(n)=0, \quad \forall n \geq \boxed{n_0=1}$$

$f(n) + g(n)$ is $O(\max\{f(n), g(n)\})$:
Find constants $c > 0$, $n_0 \geq 1$ such that

$$\rightarrow f(n) + g(n) \leq \underline{\underline{c}} \max\{f(n), g(n)\}, \forall n \geq \underline{\underline{n_0}}$$

$$+ f(n) \leq \max\{f(n), g(n)\}, \forall n \geq 1$$

$$g(n) \leq \max\{f(n), g(n)\}, \forall n \geq 1$$

$$f(n) + g(n) \leq \underline{\underline{2}} \max\{f(n), g(n)\}, \forall n \geq \underline{\underline{1}}$$

$$c = 2, n_0 = 1$$