

Algorithm BinarySearch (L,x, first, last)

if first > last **then return** -1 } C_1

else {

mid \leftarrow (first+last)/2

if x = L[mid] **then return** mid \Leftarrow

else if x < L[mid] **then**

return BinarySearch (L,x,first,mid -1)

else return BinarySearch (L,x,mid +1,last

)

$f(n)$ = number of operations performed by
the algorithm in the worst case ($x \notin L$)
when the input has size n

$$f(0) = C_1$$

$$f(n) = C_2 + f\left(\frac{n-1}{2}\right)$$

Recurrence equation

Repeated Substitutions

$$f(0) = C_1$$

$$f(n) = C_2 + f\left(\frac{n-2^0}{2}\right), n > 0$$

$$f\left(\frac{n-2^0}{2}\right) = C_2 + f\left(\frac{\frac{n-2^0}{2}-1}{2}\right) = C_2 + f\left(\frac{n-2^0-2^1}{2^2}\right)$$

$$f\left(\frac{n-2^0-2^1}{2^2}\right) = C_2 + f\left(\frac{\frac{n-2^0-2^1}{2^2}-1}{2}\right) = C_2 + f\left(\frac{n-2^0-2^1-2^2}{2^3}\right)$$

$$f\left(\frac{n-2^0-2^1-2^2}{2^3}\right) = C_2 + f\left(\frac{n-2^0-2^1-2^2-2^3}{2^4}\right)$$

$$f\left(\frac{n-2^0-2^1-2^2-2^3}{2^4}\right) = C_2 + f\left(\frac{n-2^0-2^1-2^2-2^3-2^4}{2^5}\right)$$

⋮

$$f\left(\frac{n-2^0-2^1-\dots-2^k}{2^{k+1}}\right) = C_2 + f\left(\frac{n-2^0-2^1-\dots-2^{k+1}}{2^{k+2}}\right)$$

$= 0$

$$f(n) = \underbrace{C_2 + C_2 + C_2 + C_2 + C_2 + \dots + C_2}_{k+2} + \underbrace{f(0)}_{C_1}$$

$$= C_2(k+2) + C_1$$

$$\boxed{f(n) = C_2 + \log_2(n+1) + C_1} \text{ is } O(\log n)$$

$$f\left(\frac{n-1}{2}\right) = C_2 + f\left(\frac{\frac{n-1}{2}-1}{2}\right)$$

$$n - 2^0 - 2^1 - \dots - 2^{k+1} = 0$$

$$n = 2^0 + 2^1 + \dots + 2^{k+1}$$

$$= \sum_{i=0}^{k+1} 2^i = 2^{k+2} - 1$$

Geometric Sum

$$\log_2(n+1) = \log_2 2^{k+2}$$

$$= k+2$$

$$k = \log_2(n+1) - 2$$