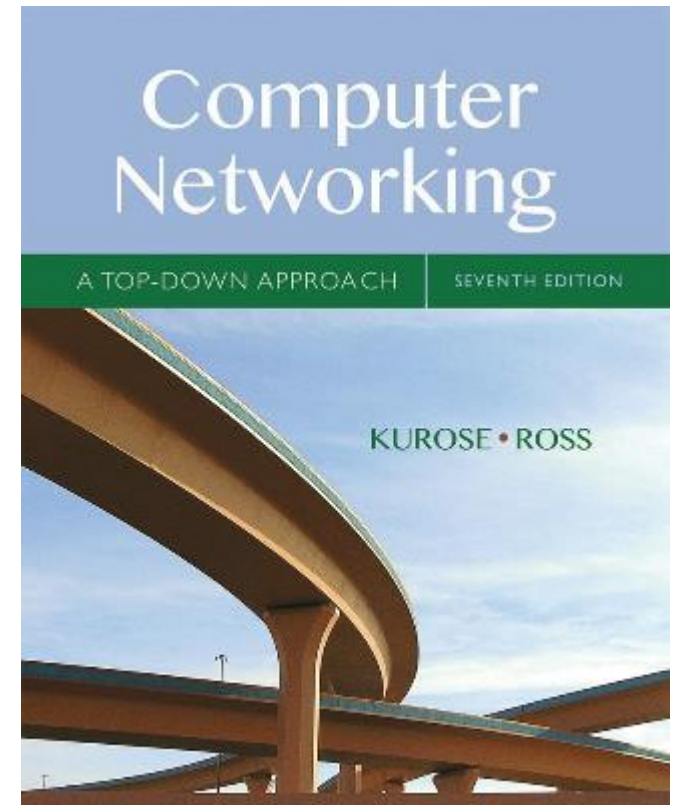


# Chapter 8

## Security

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## *Computer Networking: A Top Down Approach*

7<sup>th</sup> edition

Jim Kurose, Keith Ross

Pearson/Addison Wesley

April 2016

# Chapter 8: Network Security

## *Chapter goals:*

- understand principles of network security:
  - cryptography and its *many* uses beyond “confidentiality”
  - authentication
  - message integrity
- security in practice:
  - firewalls and intrusion detection systems
  - security in application, transport, network, link layers

# Chapter 8 roadmap

*8.1 What is network security?*

8.2 Principles of cryptography

8.3 Message integrity, authentication

8.4 Securing e-mail

8.5 Securing TCP connections: SSL

8.6 Network layer security: IPsec

8.7 Securing wireless LANs

8.8 Operational security: firewalls and IDS

# What is network security?

**confidentiality:** only sender, intended receiver should “understand” message contents

- sender encrypts message
- receiver decrypts message

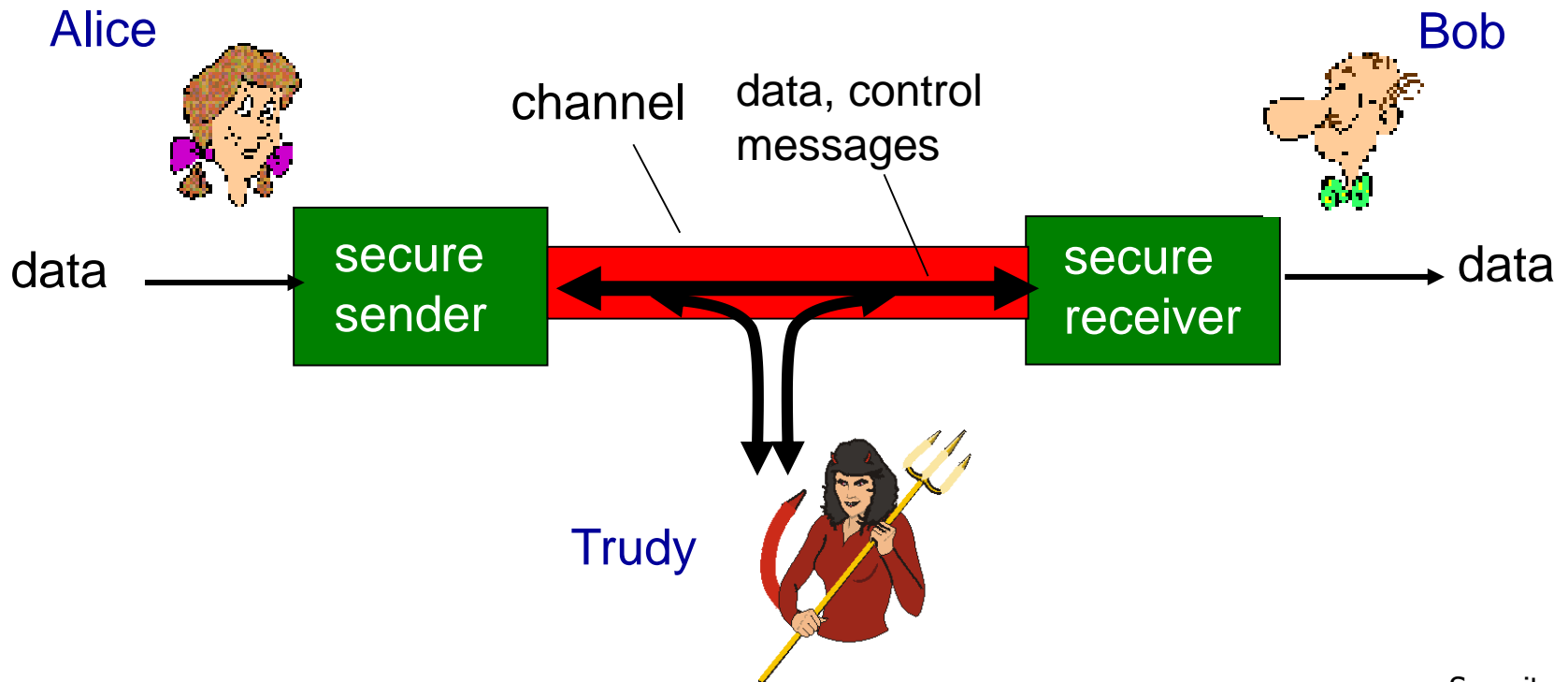
**authentication:** sender, receiver want to confirm identity of each other

**message integrity:** sender, receiver want to ensure message not altered (in transit, or afterwards) without detection

**access and availability:** services must be accessible and available to users

# Friends and enemies: Alice, Bob, Trudy

- well-known in network security world
- Bob, Alice (very close friends) want to communicate “securely”
- Trudy (intruder) may intercept, delete, add messages



# Who might Bob, Alice be?

- ... well, *real-life* Bobs and Alices!
- Web browser/server for electronic transactions (e.g., on-line purchases)
- on-line banking client/server
- DNS servers
- routers exchanging routing table updates
- other examples?

# There are bad guys (and girls) out there!

Q: What can a “bad guy” do?

A: A lot! See section 1.6

- *eavesdrop*: intercept messages
- actively *insert* messages into connection
- *impersonation*: can fake (spoof) source address in packet (or any field in packet)
- *hijacking*: “take over” ongoing connection by removing sender or receiver, inserting himself in place
- *denial of service*: prevent service from being used by others (e.g., by overloading resources)

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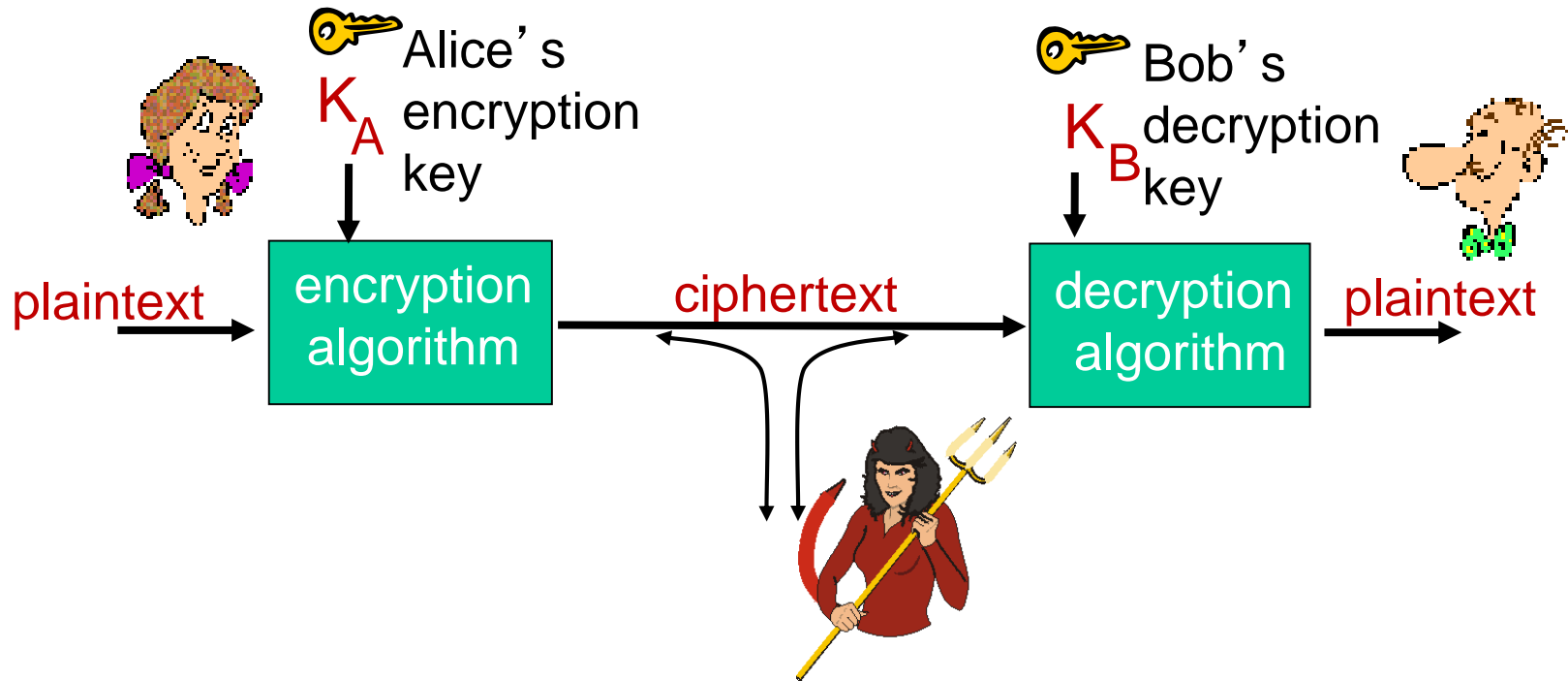
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# The language of cryptography



$m$  plaintext message

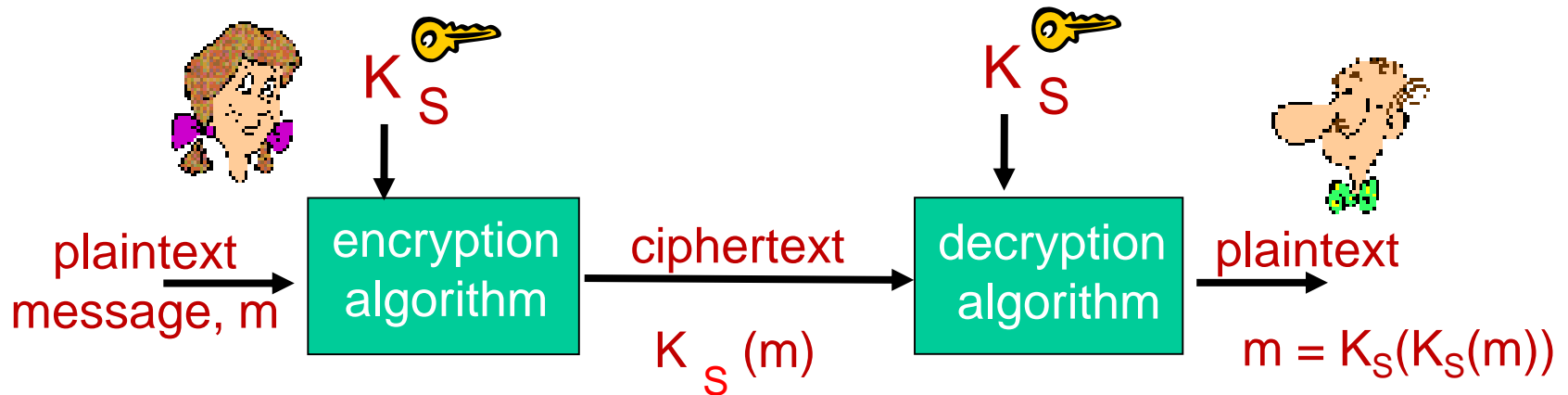
$K_A(m)$  ciphertext, encrypted with key  $K_A$

$m = K_B(K_A(m))$

# Breaking an encryption scheme

- **cipher-text only attack:**  
Trudy has ciphertext she can analyze
- **two approaches:**
  - brute force: search through all keys
  - statistical analysis
- **known-plaintext attack:**  
Trudy has plaintext corresponding to ciphertext
  - e.g., in monoalphabetic cipher, Trudy determines pairings for a,l,i,c,e,b,o,
- **chosen-plaintext attack:**  
Trudy can get ciphertext for chosen plaintext

# Symmetric key cryptography



**symmetric key crypto:** Bob and Alice share same (symmetric) key:  $K_S$

- e.g., key is knowing substitution pattern in mono alphabetic substitution cipher

Q: how do Bob and Alice agree on key value?

# Simple encryption scheme

*substitution cipher*: substituting one thing for another

- monoalphabetic cipher: substitute one letter for another

plaintext:	a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v	w	x	y	z
ciphertext:	m	n	b	v	c	x	z	a	s	d	f	g	h	j	k	l	p	o	i	u	y	t	r	e	w	q

e.g.: Plaintext: bob. i love you. alice  
ciphertext: nkn. s gktc wky. mgsbc

🔑 *Encryption key*: mapping from set of 26 letters  
to set of 26 letters

# A more sophisticated encryption approach

- n substitution ciphers,  $M_1, M_2, \dots, M_n$
- cycling pattern:
  - e.g.,  $n=4$ :  $M_1, M_3, M_4, M_3, M_2$ ;  $M_1, M_3, M_4, M_3, M_2$ ; ..
- for each new plaintext symbol, use subsequent substitution pattern in cyclic pattern
  - dog: d from  $M_1$ , o from  $M_3$ , g from  $M_4$

*Encryption key:* n substitution ciphers, and cyclic pattern



- key need not be just n-bit pattern

# Symmetric key crypto: DES

## DES: Data Encryption Standard

- US encryption standard [NIST 1993]
- 56-bit symmetric key, 64-bit plaintext input
- block cipher with cipher block chaining
- how secure is DES?
  - DES Challenge: 56-bit-key-encrypted phrase decrypted (brute force) in less than a day
  - no known good analytic attack
- making DES more secure:
  - 3DES: encrypt 3 times with 3 different keys

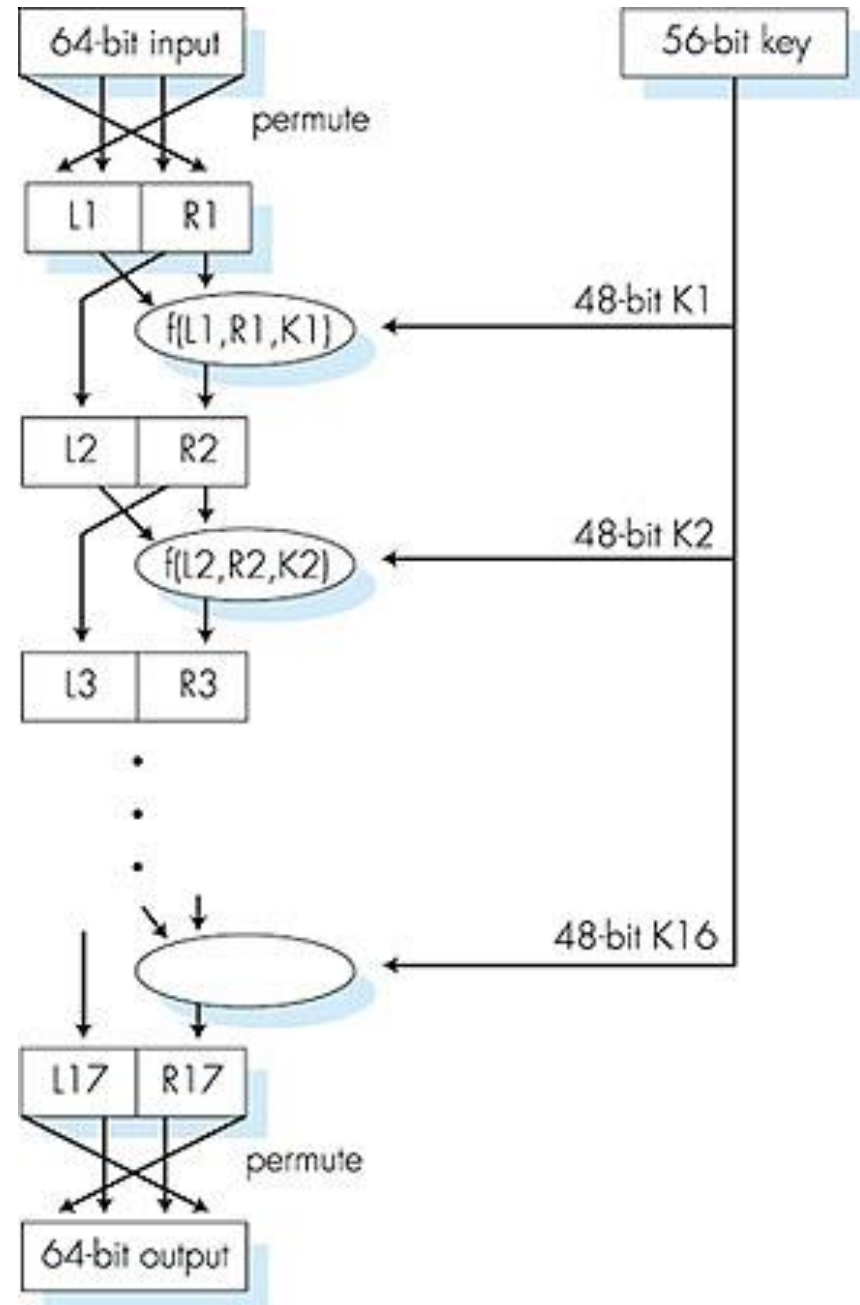
# Symmetric key crypto: DES

## *DES operation*

initial permutation

16 identical “rounds” of  
function application,  
each using different 48  
bits of key

final permutation



# AES: Advanced Encryption Standard

- symmetric-key NIST standard, replaced DES (Nov 2001)
- processes data in 128 bit blocks
- 128, 192, or 256 bit keys
- brute force decryption (try each key) taking 1 sec on DES, takes 149 trillion years for AES



# Public Key Cryptography



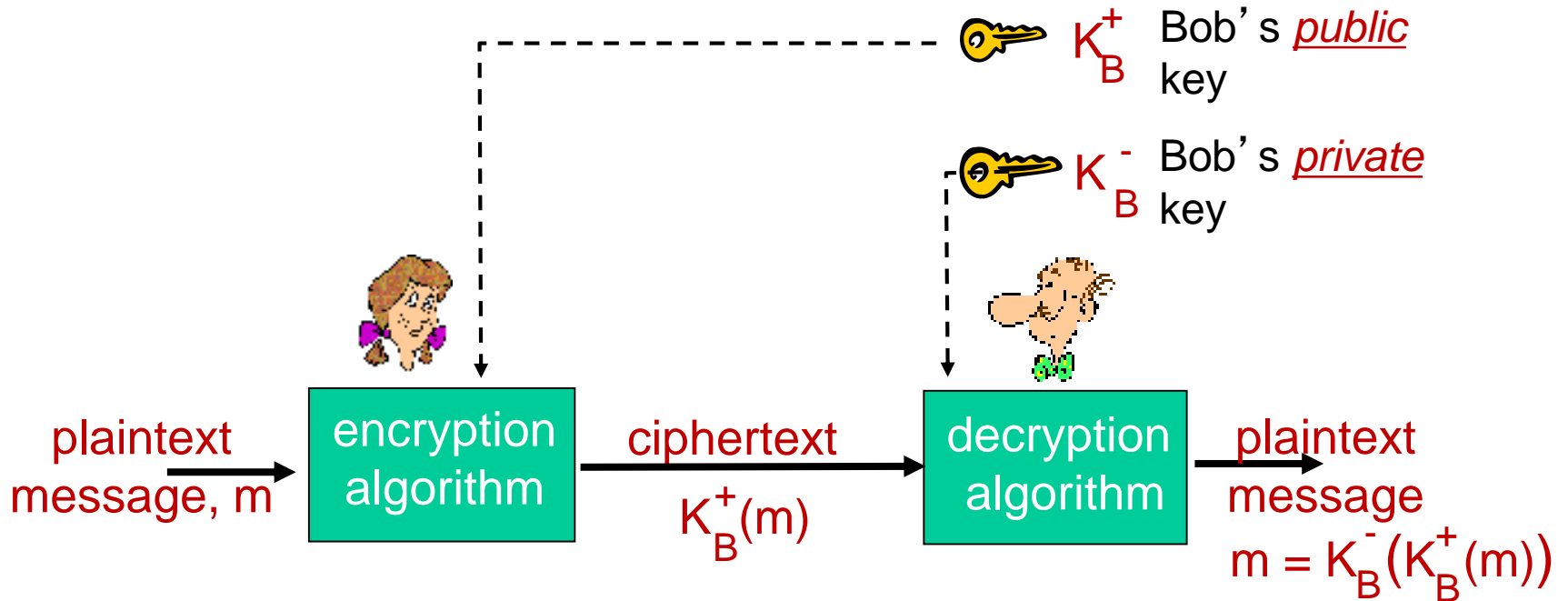
## *symmetric key crypto*

- requires sender, receiver know shared secret key
- Q: how to agree on key in first place (particularly if never “met”)?

## *public key crypto*

- radically different approach [Diffie-Hellman76, RSA78]
- sender, receiver do *not* share secret key
- *public* encryption key known to *all*
- *private* decryption key known only to receiver

# Public key cryptography



# Public key encryption algorithms

requirements:

- ① need  $K_B^+(\cdot)$  and  $K_B^-(\cdot)$  such that

$$K_B^-(K_B^+(m)) = m$$

- ② given public key  $K_B^+$ , it should be impossible to compute private key  $K_B^-$

**RSA:** Rivest, Shamir, Adelson algorithm

# RSA:

- message: just a bit pattern
- bit pattern can be uniquely represented by an integer number
- thus, encrypting a message is equivalent to encrypting a number

## *example:*

- $m = 10010001$ . This message is uniquely represented by the decimal number 145.
- to encrypt  $m$ , we encrypt the corresponding number, which gives a new number (the ciphertext).

# RSA: important property!!

The following property will be *very* useful later:

$$\underbrace{K_B^-(K_B^+(m))}_{\text{use public key first, followed by private key}} = m = \underbrace{K_B^+(K_B^-(m))}_{\text{use private key first, followed by public key}}$$

use public key first,  
followed by  
private key

use private key  
first, followed by  
public key

*result is the same!*

# Why is RSA secure?

- suppose you know Bob's public key  $(n,e)$ . How hard is it to determine  $d$ ?
- essentially need to find factors of  $n$  without knowing the two factors  $p$  and  $q$ 
  - fact: factoring a big number is hard

# RSA in practice: session keys

- exponentiation in RSA is computationally intensive
- DES & AES is at least 100 times faster than RSA
- use public key crypto to establish secure connection, then establish second key – symmetric session key – for encrypting data

## *session key, $K_S$*

- Bob and Alice use RSA to exchange a symmetric key  $K_S$
- once both have  $K_S$ , they use symmetric key cryptography

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# Authentication

*Goal:* Bob wants Alice to “prove” her identity to him

*Protocol ap1.0:* Alice says “I am Alice”



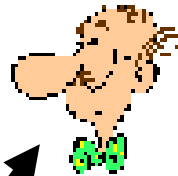
Failure scenario??



# Authentication

*Goal:* Bob wants Alice to “prove” her identity to him

*Protocol ap1.0:* Alice says “I am Alice”

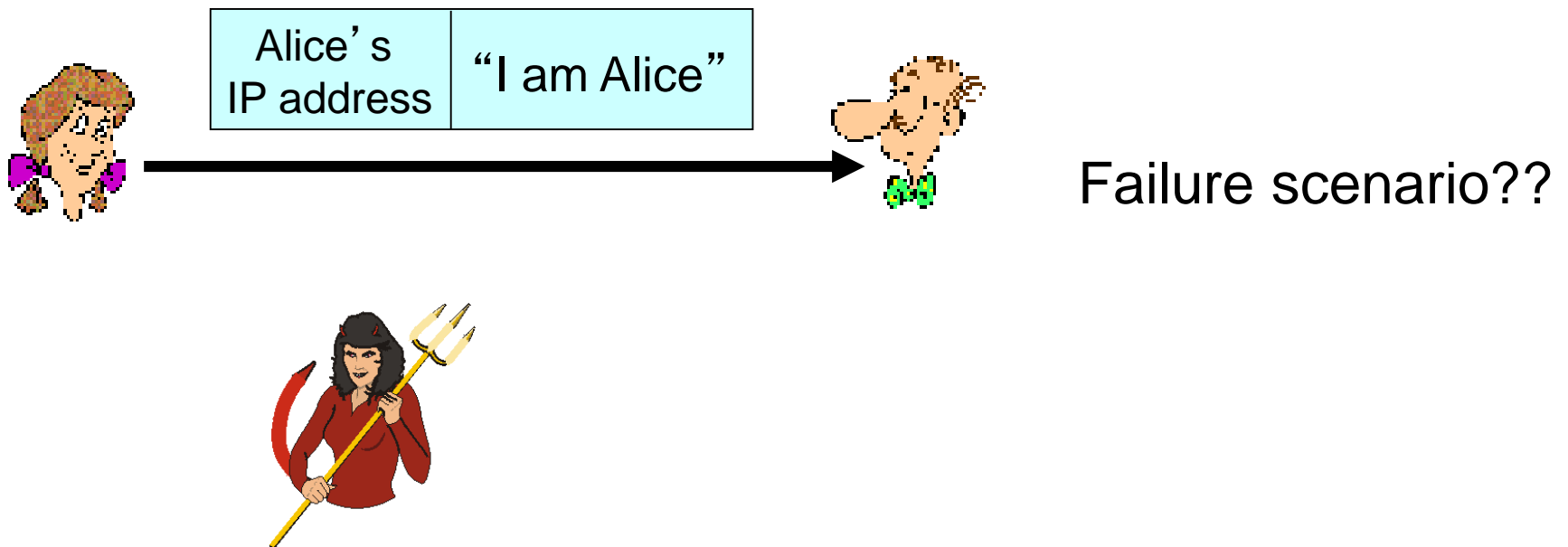


“I am Alice”

in a network,  
Bob can not “see” Alice,  
so Trudy simply declares  
herself to be Alice

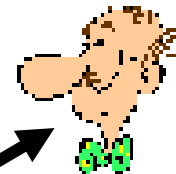
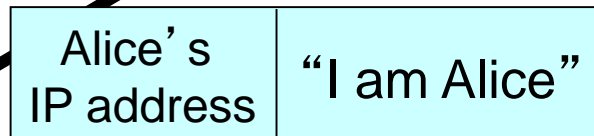
# Authentication: another try

*Protocol ap2.0:* Alice says “I am Alice” in an IP packet containing her source IP address



# Authentication: another try

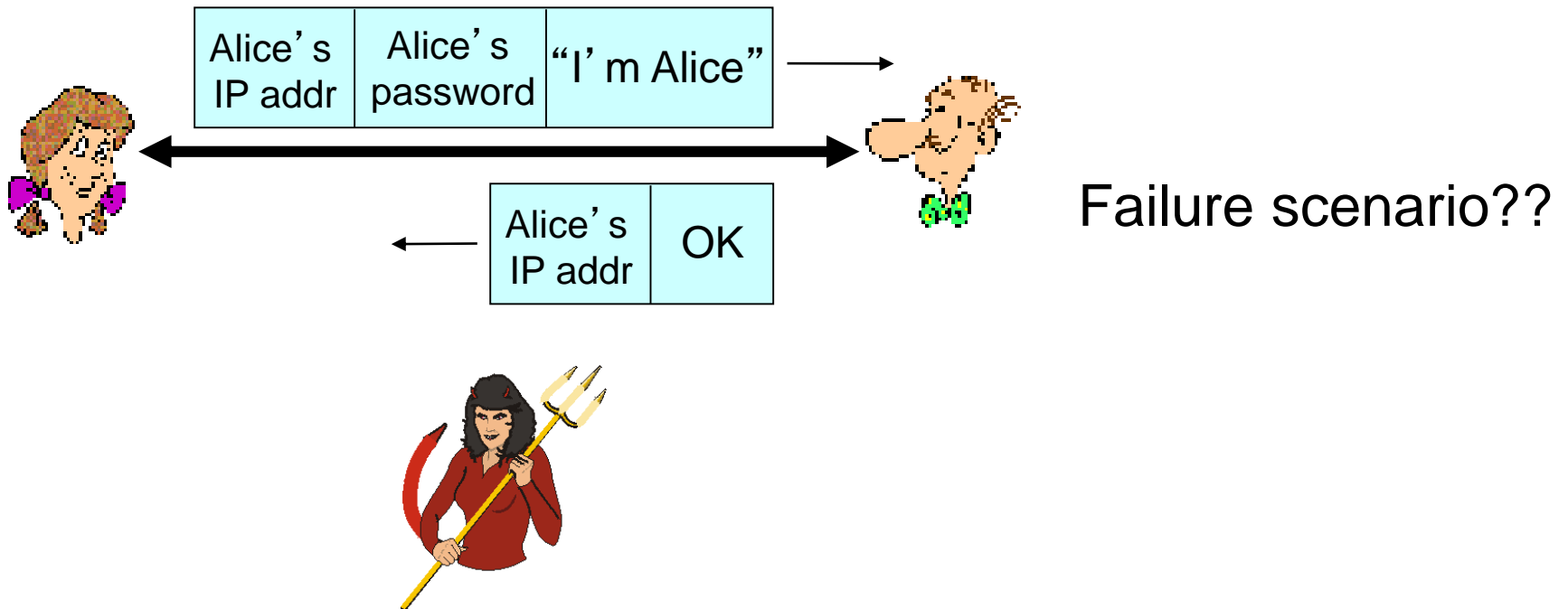
*Protocol ap2.0:* Alice says “I am Alice” in an IP packet containing her source IP address



Trudy can create  
a packet  
“spoofing”  
Alice’s address

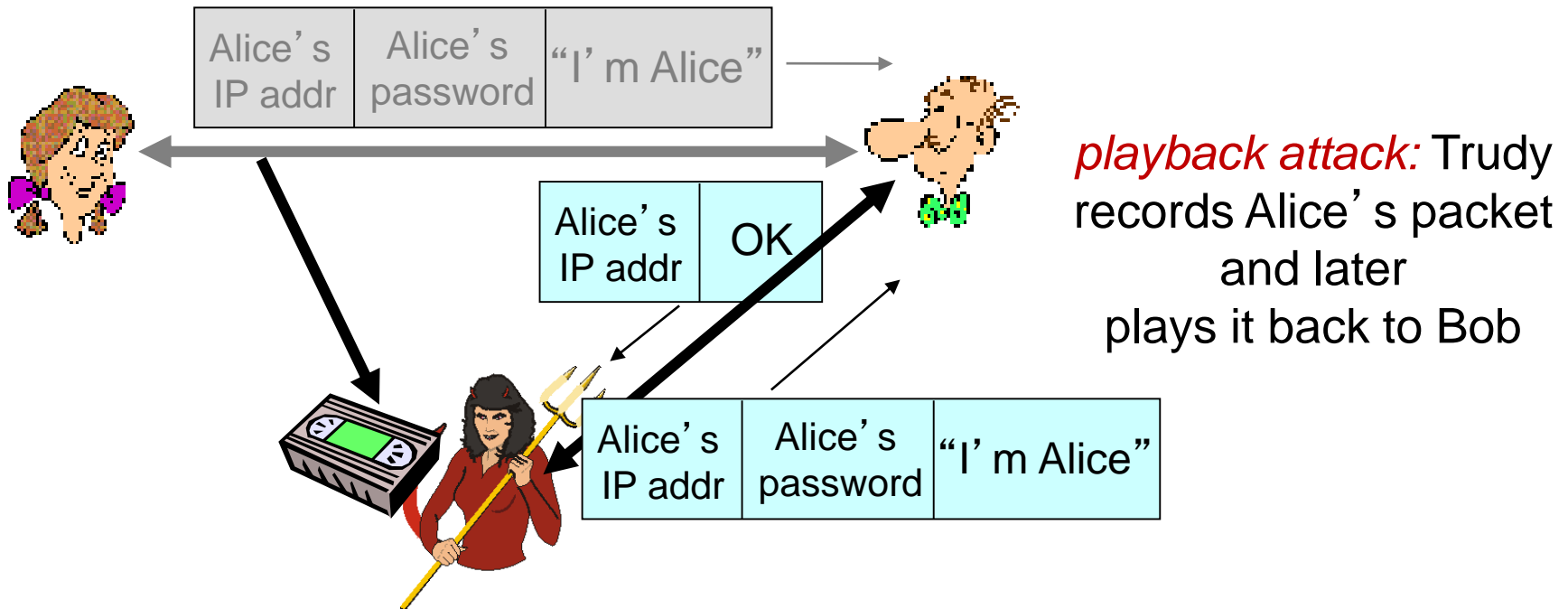
# Authentication: another try

*Protocol ap3.0:* Alice says “I am Alice” and sends her secret password to “prove” it.



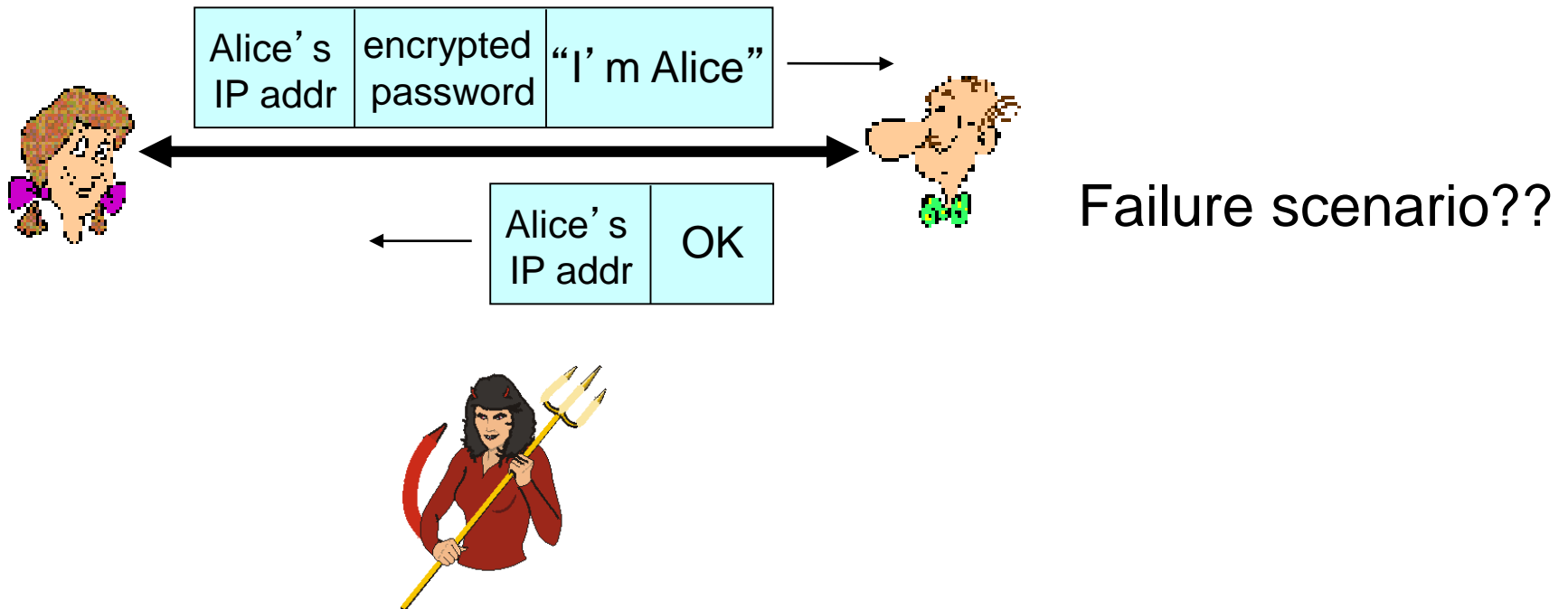
# Authentication: another try

*Protocol ap3.0:* Alice says “I am Alice” and sends her secret password to “prove” it.



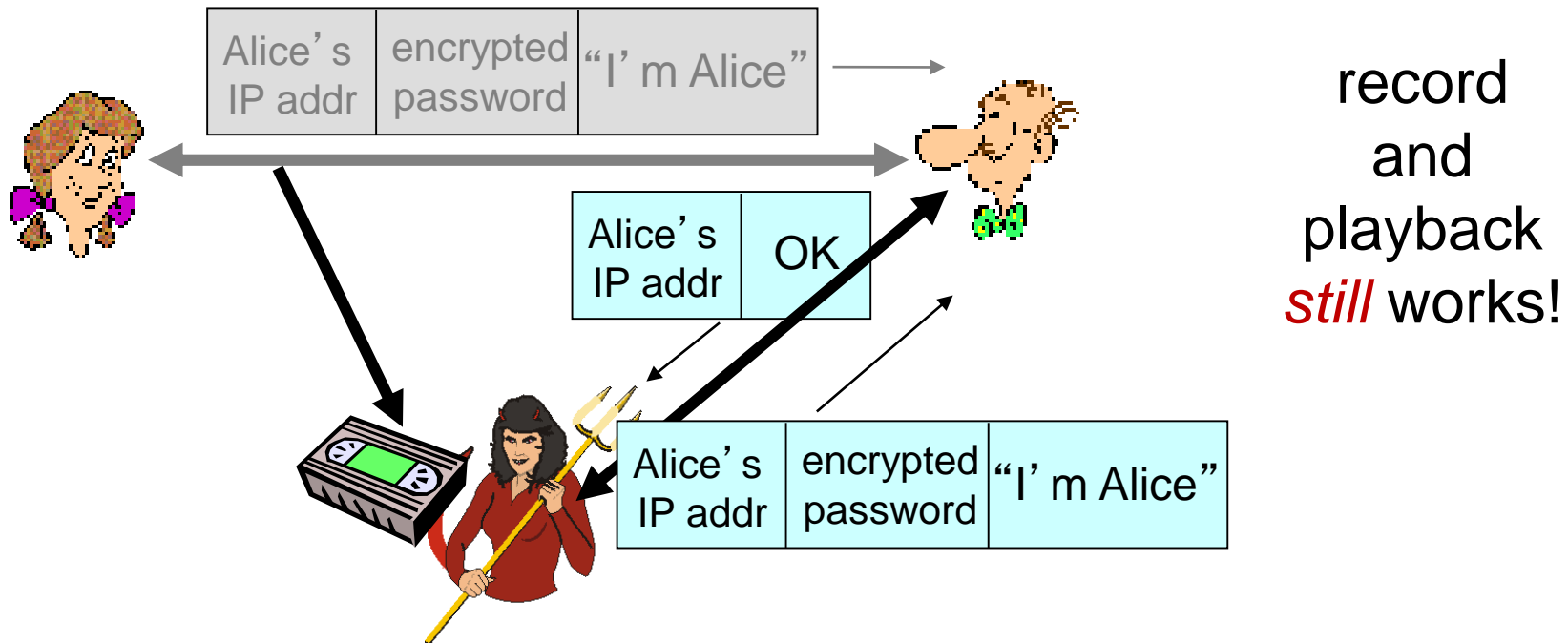
# Authentication: yet another try

*Protocol ap3.1:* Alice says “I am Alice” and sends her *encrypted* secret password to “prove” it.



# Authentication: yet another try

*Protocol ap3.1:* Alice says “I am Alice” and sends her *encrypted* secret password to “prove” it.



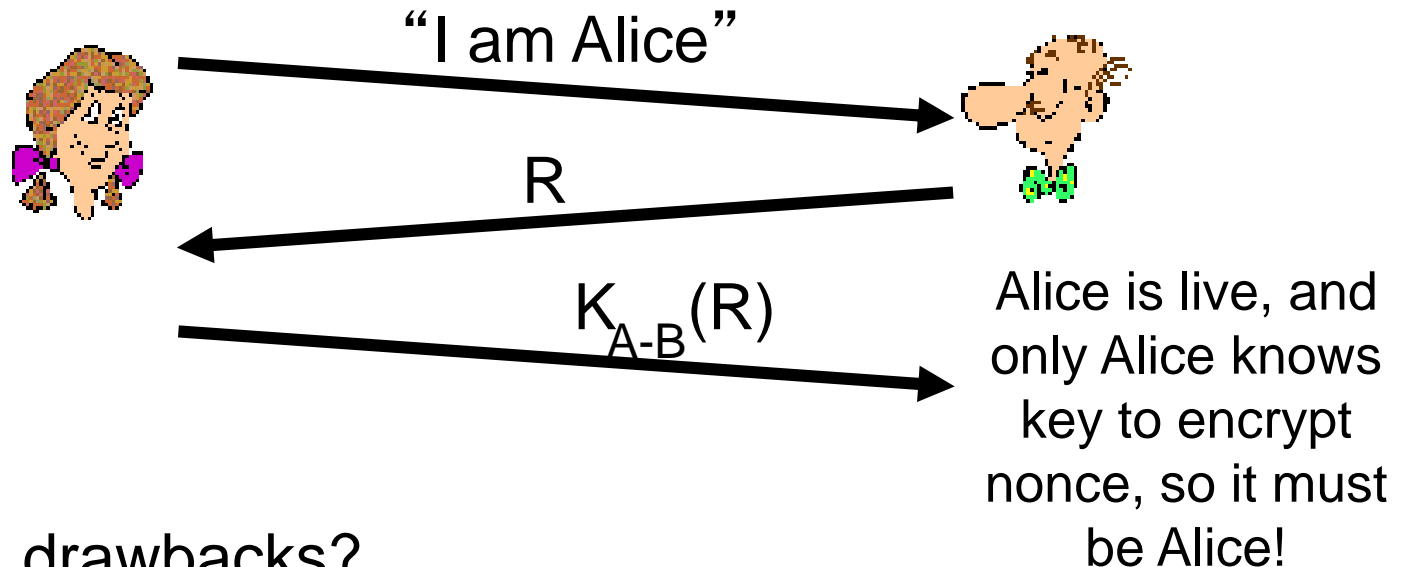


# Authentication: yet another try

**Goal:** avoid playback attack

**nonce:** number (R) used only *once-in-a-lifetime*

**ap4.0:** to prove Alice “live”, Bob sends Alice **nonce**, R. Alice must return R, encrypted with shared secret key



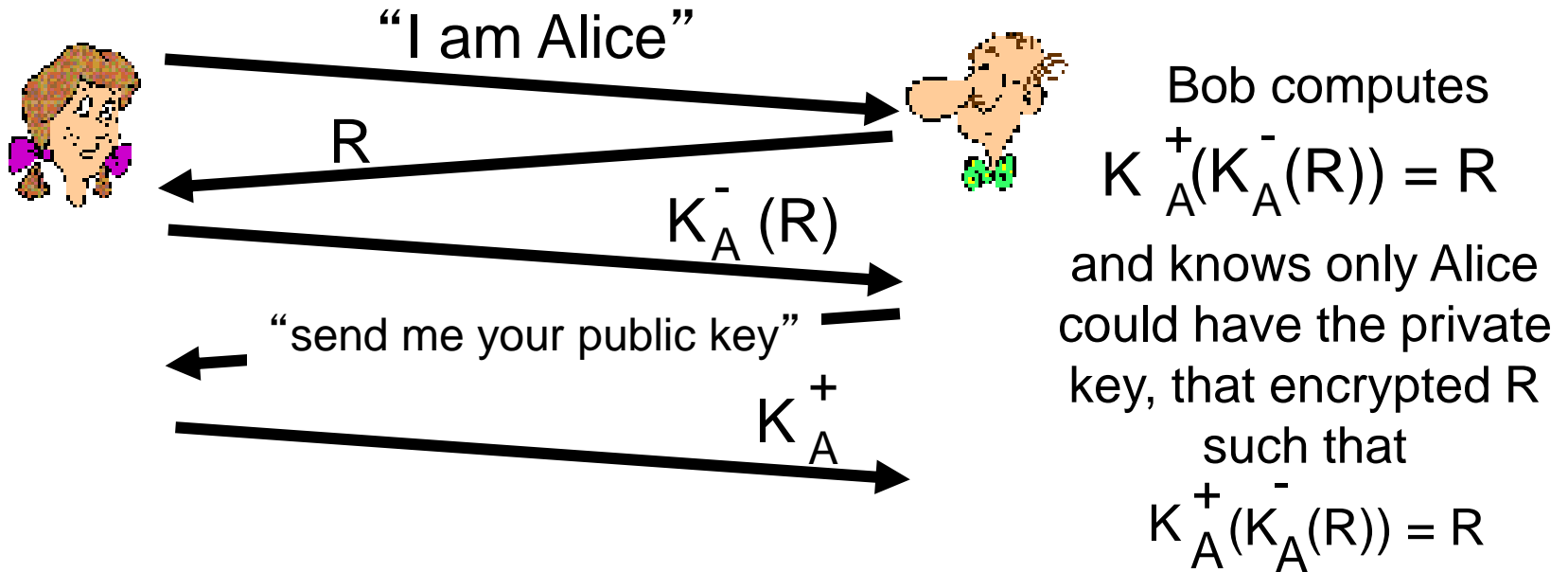
Failures, drawbacks?

# Authentication: ap5.0

ap4.0 requires shared symmetric key

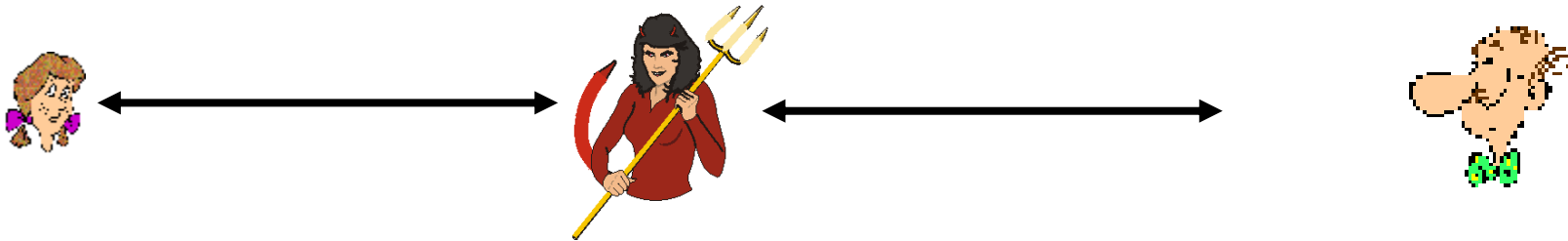
- can we authenticate using public key techniques?

*ap5.0*: use nonce, public key cryptography



# ap5.0: security hole

*man (or woman) in the middle attack:* Trudy poses as Alice (to Bob) and as Bob (to Alice)



difficult to detect:

- Bob receives everything that Alice sends, and vice versa. (e.g., so Bob, Alice can meet one week later and recall conversation!)
- problem is that Trudy receives all messages as well!

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# Digital signatures

cryptographic technique analogous to hand-written signatures:

- sender (Bob) digitally signs document, establishing he is document owner/creator.
- *verifiable, nonforgeable*: recipient (Alice) can prove to someone that Bob, and no one else (including Alice), must have signed document


# Digital signatures

simple digital signature for message  $m$ :

- Bob signs  $m$  by encrypting with his private key  $K_B^-$ , creating “signed” message,  $K_B^-(m)$

Bob's message,  $m$

Dear Alice  
Oh, how I have missed  
you. I think of you all the  
time! ... (blah blah blah)  
Bob

  $K_B^-$  Bob's private  
key

Public key  
encryption  
algorithm

$m, K_B^-(m)$

Bob's message,  
 $m$ , signed  
(encrypted) with  
his private key

# Digital signatures

- suppose Alice receives msg  $m$ , with signature:  $m, K_B^-(m)$
- Alice verifies  $m$  signed by Bob by applying Bob's public key  $K_B$  to  $K_B^-(m)$  then checks  $K_B(K_B^-(m)) = m$ .
- If  $K_B(K_B^-(m)) = m$ , whoever signed  $m$  must have used Bob's private key.

Alice thus verifies that:

- Bob signed  $m$
- no one else signed  $m$
- Bob signed  $m$  and not  $m'$

non-repudiation:

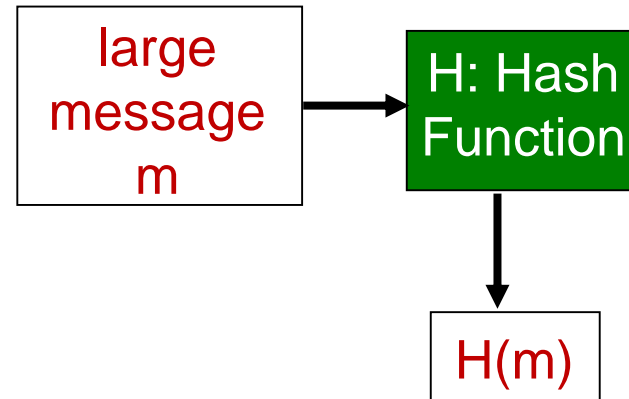
- ✓ Alice can take  $m$ , and signature  $K_B^-(m)$  to court and prove that Bob signed  $m$

# Message digests

computationally expensive to public-key-encrypt long messages

**goal:** fixed-length, easy-to-compute digital “fingerprint”

- apply hash function  $H$  to  $m$ , get fixed size message digest,  $H(m)$ .



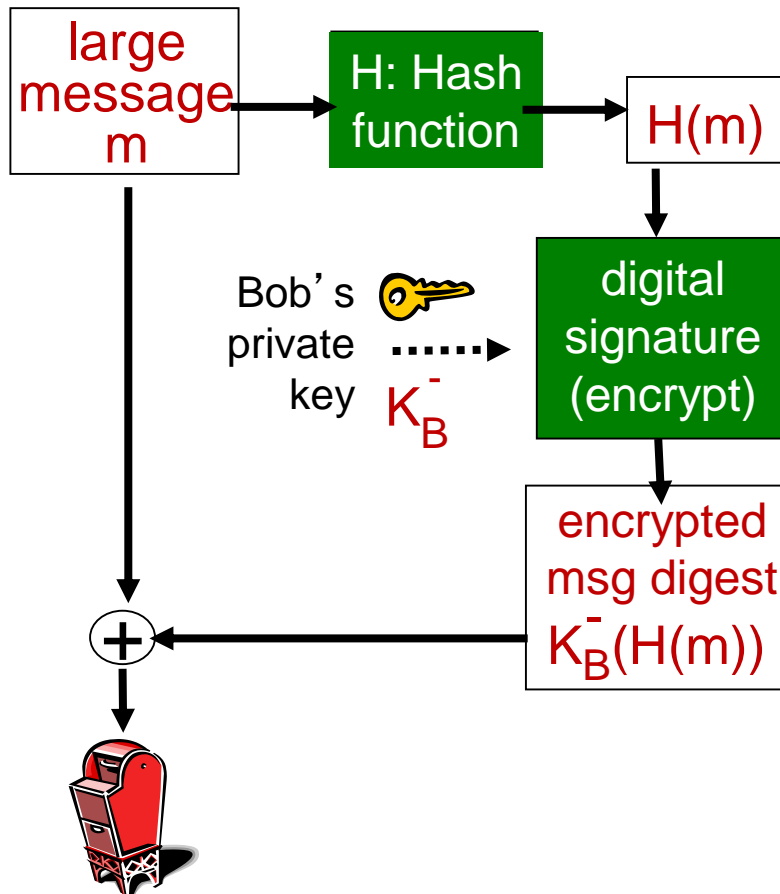
**Hash function properties:**

- many-to-1
- produces fixed-size msg digest (fingerprint)
- given message digest  $x$ , computationally infeasible to find  $m$  such that  $x = H(m)$

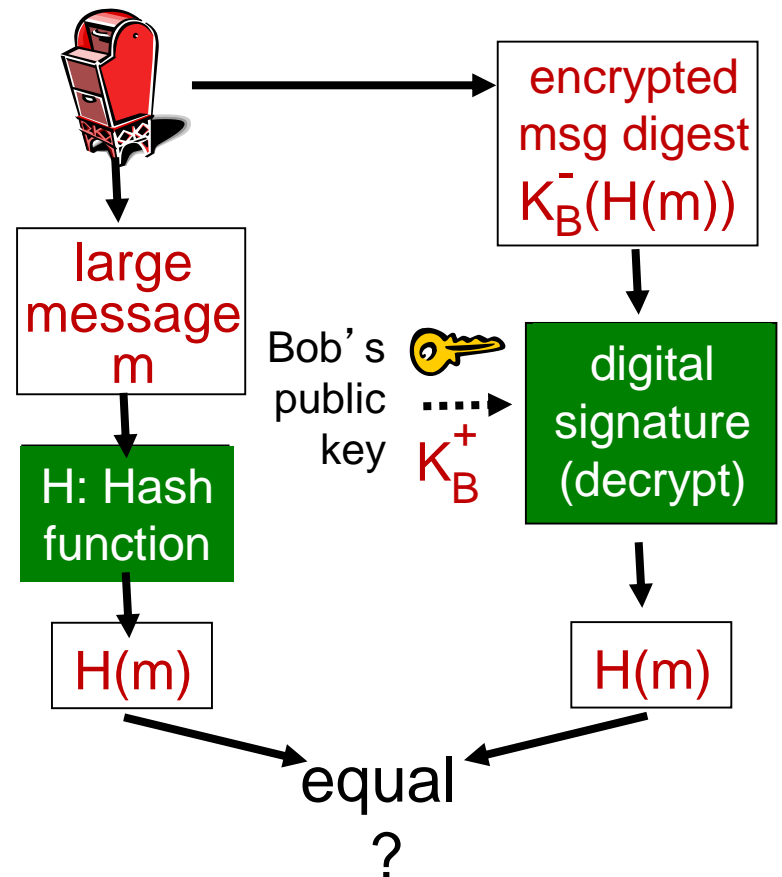


# Digital signature = signed message digest

Bob sends digitally signed message:



Alice verifies signature, integrity of digitally signed message:

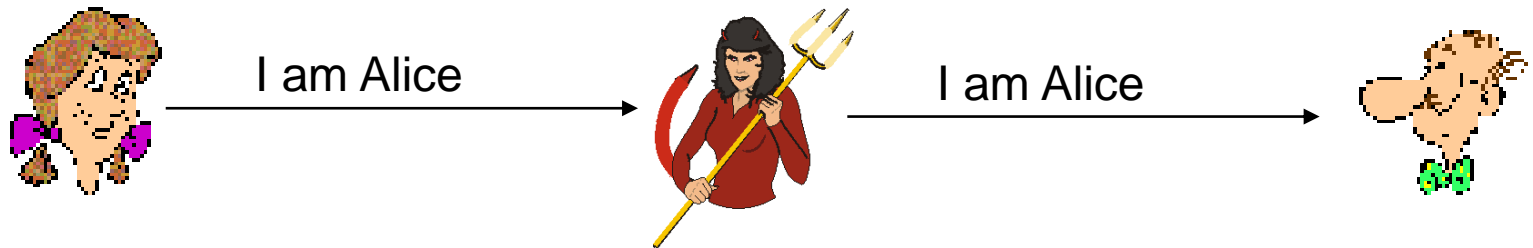


# Hash function algorithms

- **MD5 hash function widely used (RFC 1321)**
  - computes 128-bit message digest in 4-step process.
  - arbitrary 128-bit string  $x$ , appears difficult to construct msg  $m$  whose MD5 hash is equal to  $x$
- **SHA-1 is also used**
  - US standard [NIST, FIPS PUB 180-1]
  - 160-bit message digest

# Recall: ap5.0 security hole

*man (or woman) in the middle attack:* Trudy poses as Alice (to Bob) and as Bob (to Alice)

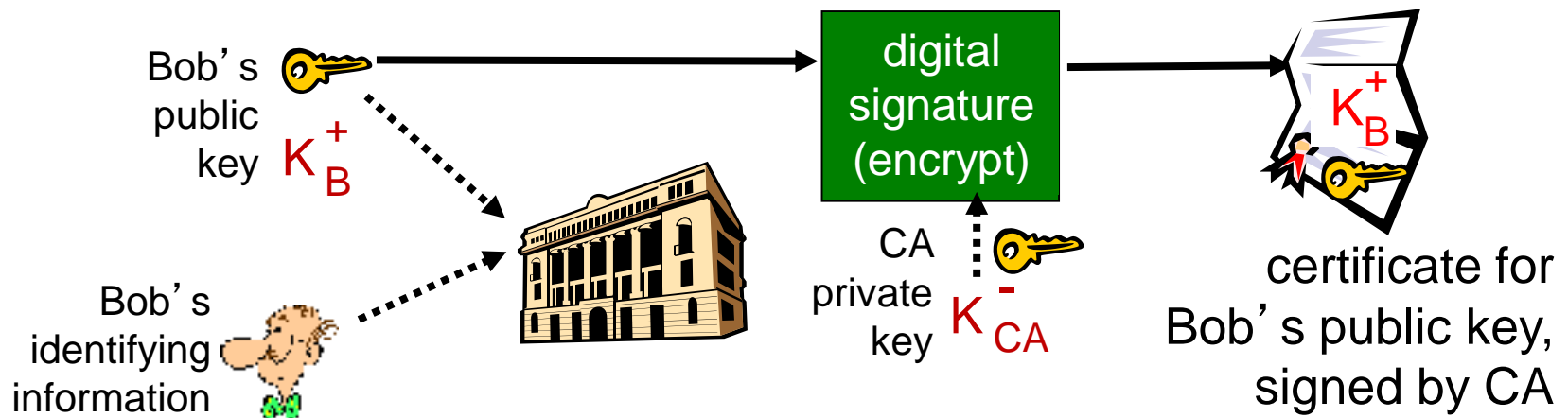


# Public-key certification

- motivation: Trudy plays pizza prank on Bob
  - Trudy creates e-mail order:  
*Dear Pizza Store, Please deliver to me four pepperoni pizzas. Thank you, Bob*
  - Trudy signs order with her private key
  - Trudy sends order to Pizza Store
  - Trudy sends to Pizza Store her public key, but says it's Bob's public key
  - Pizza Store verifies signature; then delivers four pepperoni pizzas to Bob
  - Bob has to pay for the pizzas, and he doesn't even like pizza.
  - We need some way to verify public keys!!!!

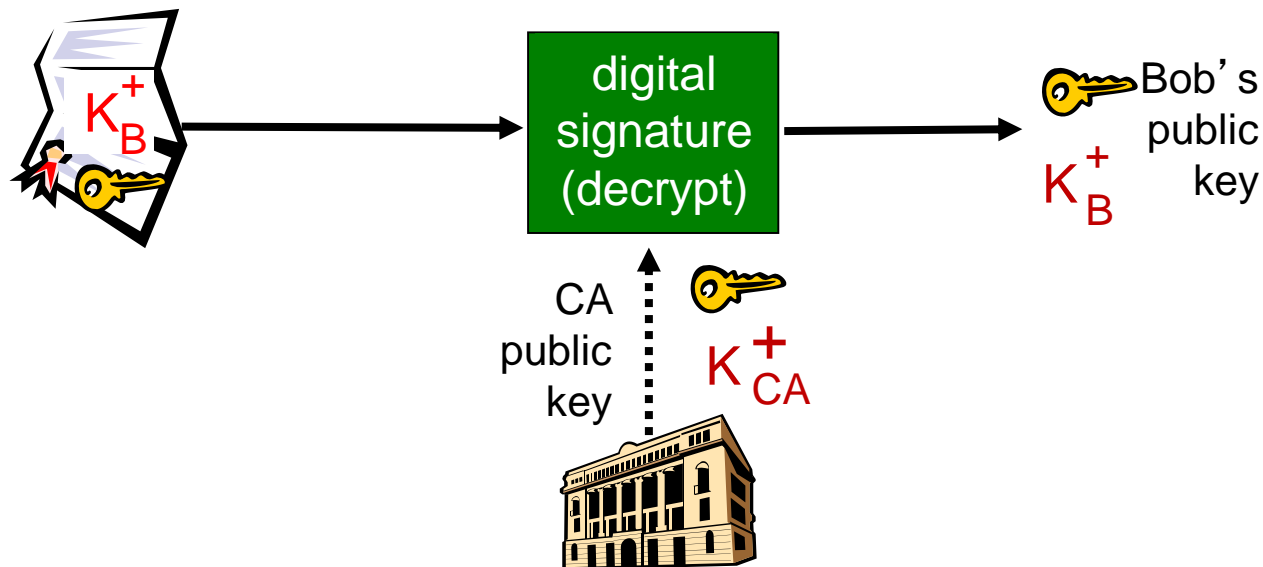
# Certification authorities

- **certification authority (CA):** binds public key to particular entity, E.
- E (person, router) registers its public key with CA.
  - E provides “proof of identity” to CA.
  - CA creates certificate binding E to its public key.
  - certificate containing E’s public key digitally signed by CA – CA says “this is E’s public key”



# Certification authorities

- when Alice wants Bob's public key:
  - gets Bob's certificate (Bob or elsewhere).
  - apply CA's public key to Bob's certificate, get Bob's public key



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8.3 Message integrity, authentication

*8.4 Securing e-mail*

8.5 Securing TCP connections: SSL

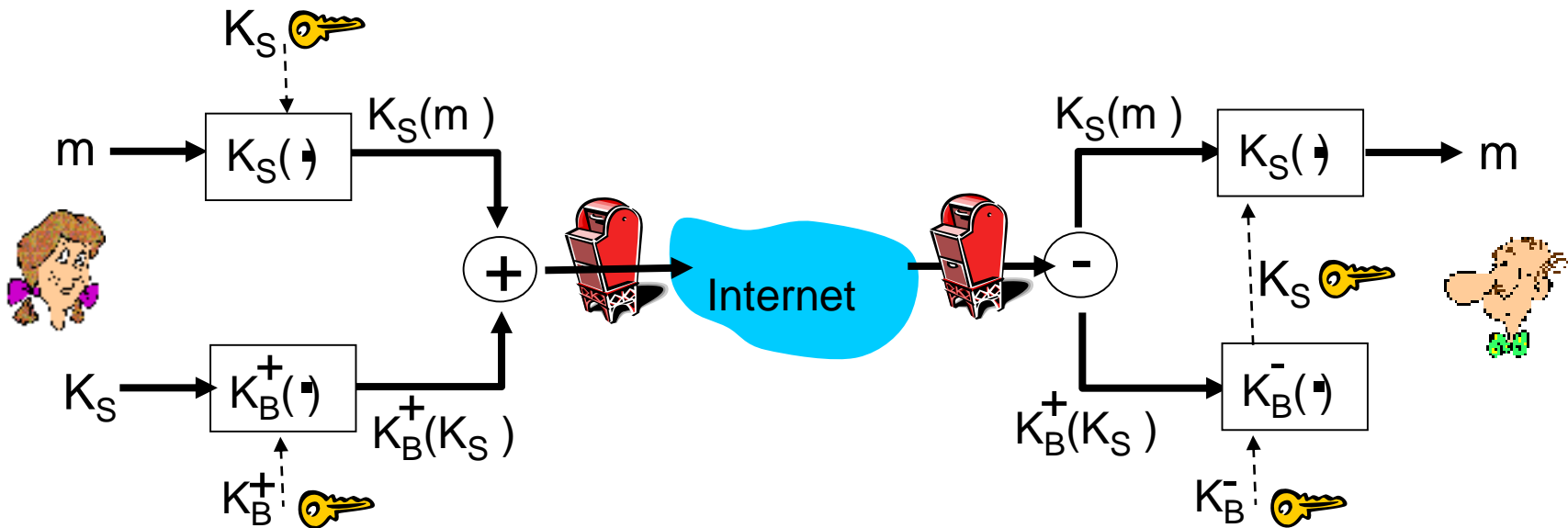
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# Secure e-mail

Alice wants to send confidential e-mail,  $m$ , to Bob.



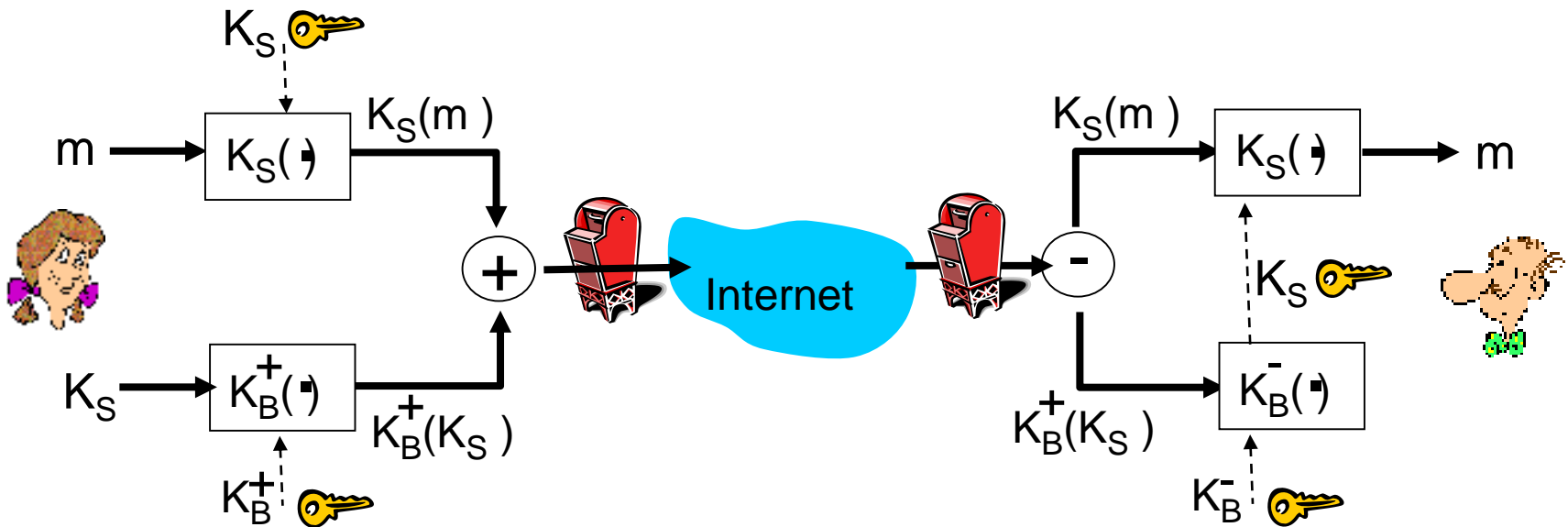
**Alice:**

- generates random *symmetric* private key,  $K_S$
- encrypts message with  $K_S$  (for efficiency)
- also encrypts  $K_S$  with Bob's public key
- sends both  $K_S(m)$  and  $K_B(K_S)$  to Bob



# Secure e-mail

Alice wants to send confidential e-mail,  $m$ , to Bob.

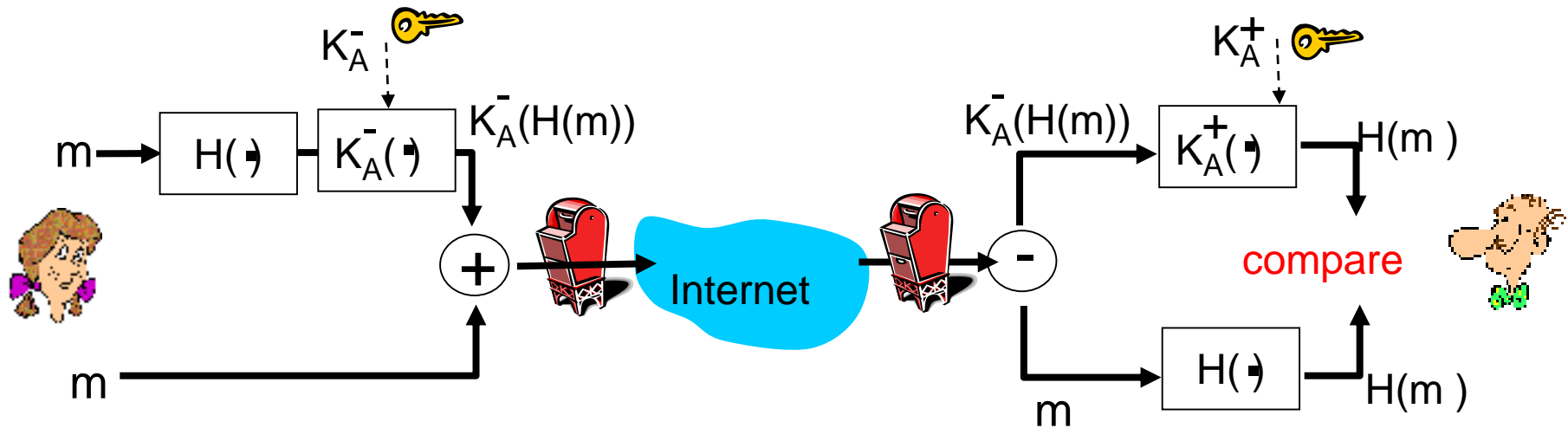


**Bob:**

- uses his private key to decrypt and recover  $K_S$
- uses  $K_S$  to decrypt  $K_S(m)$  to recover  $m$

# Secure e-mail (continued)

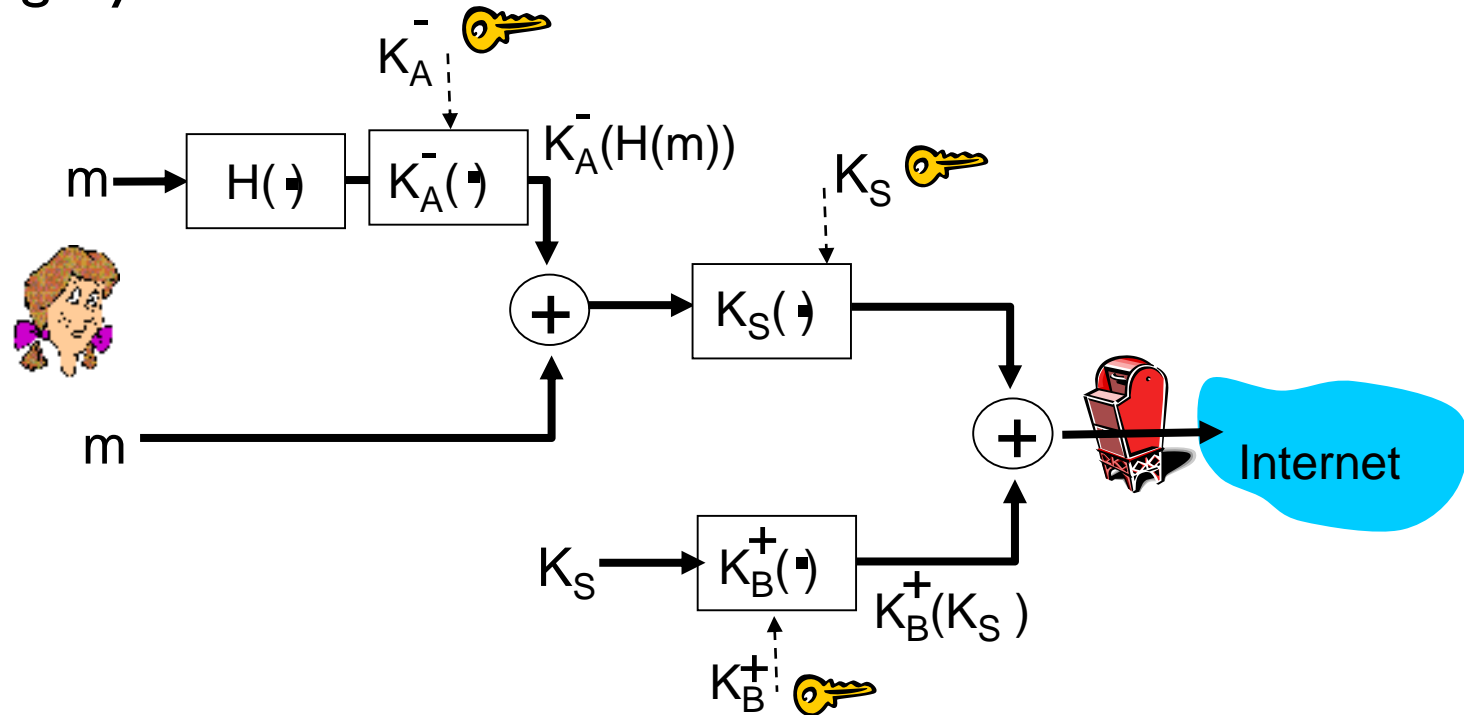
Alice wants to provide sender authentication message integrity



- Alice digitally signs message
- sends both message (in the clear) and digital signature

# Secure e-mail (continued)

Alice wants to provide secrecy, sender authentication, message integrity.



*Alice uses three keys:* her private key, Bob's public key, newly created symmetric key

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*8.5 Securing TCP connections: SSL*

8.6 Network layer security: IPsec

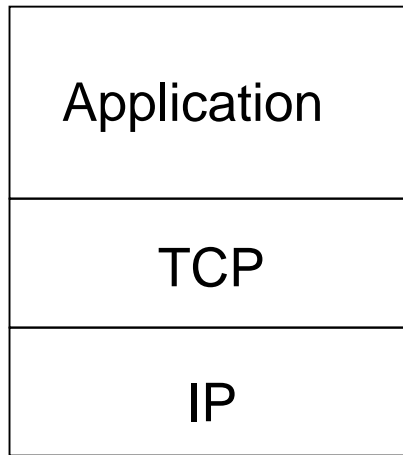
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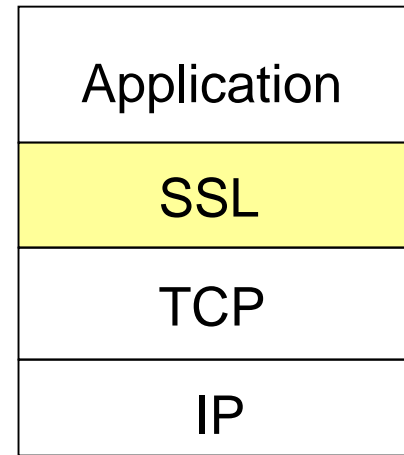
# SSL: Secure Sockets Layer

- widely deployed security protocol
  - supported by almost all browsers, web servers
  - https
  - billions \$/year over SSL
- mechanisms: [Woo 1994], implementation: Netscape
- variation -TLS: transport layer security, RFC 2246
- provides
  - *confidentiality*
  - *integrity*
  - *authentication*
- original goals:
  - Web e-commerce transactions
  - encryption (especially credit-card numbers)
  - Web-server authentication
  - optional client authentication
  - minimum hassle in doing business with new merchant
- available to all TCP applications
  - secure socket interface

# SSL and TCP/IP



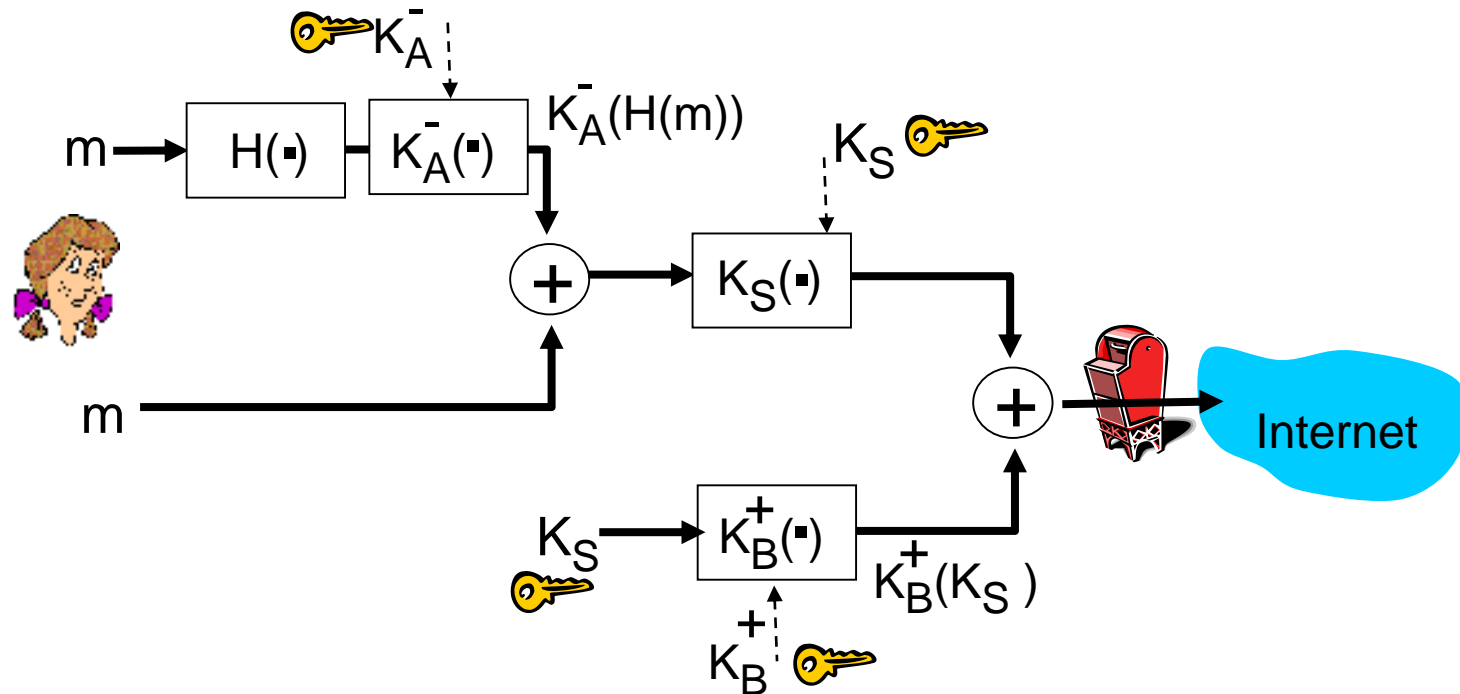
*normal application*



*application with SSL*

- SSL provides application programming interface (API) to applications
- C and Java SSL libraries/classes readily available

# Could do something like PGP:



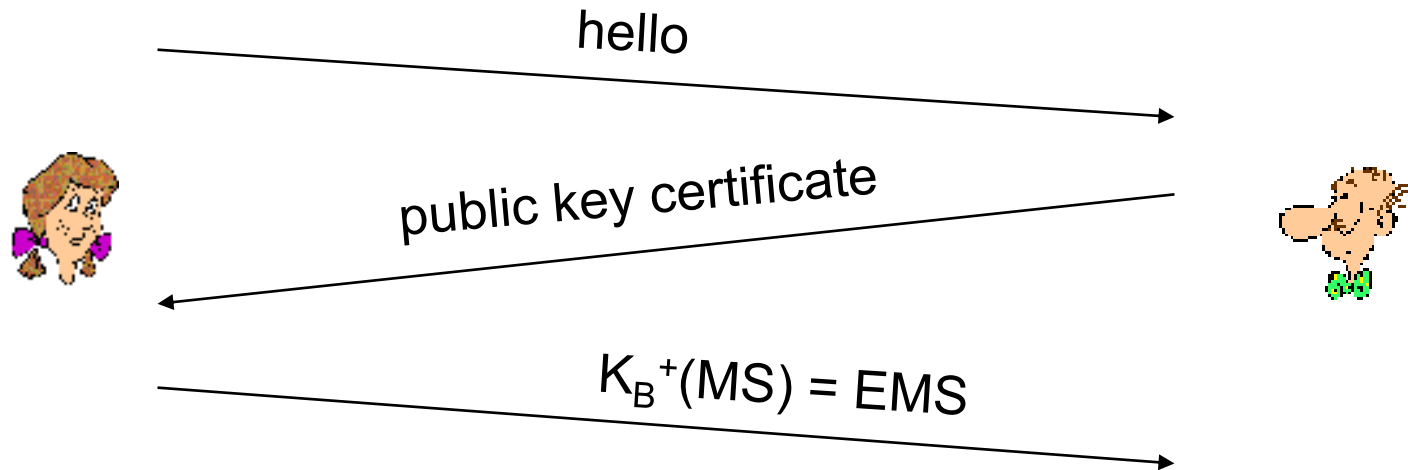
- but want to send byte streams & interactive data
- want set of secret keys for entire connection
- want certificate exchange as part of protocol: handshake phase

# Toy SSL: simple secure channel

- *handshake*: Alice and Bob use their certificates, private keys to authenticate each other and exchange shared secret
- *key derivation*: Alice and Bob use shared secret to derive set of keys
- *data transfer*: data to be transferred is broken up into series of records
- *connection closure*: special messages to securely close connection



# Toy: a simple handshake



**MS:** master secret

**EMS:** encrypted master secret

# Toy: key derivation

- considered bad to use same key for more than one cryptographic operation
  - use different keys for message authentication code (MAC) and encryption
  - $\text{MAC} = H(m + s)$
- four keys:
  - $K_c$  = encryption key for data sent from client to server
  - $M_c$  = MAC key for data sent from client to server
  - $K_s$  = encryption key for data sent from server to client
  - $M_s$  = MAC key for data sent from server to client
- keys derived from key derivation function (KDF)
  - takes master secret and (possibly) some additional random data and creates the keys

# Toy: data records

- why not encrypt data in constant stream as we write it to TCP?
  - where would we put the MAC? If at end, no message integrity until all data processed.
  - e.g., with instant messaging, how can we do integrity check over all bytes sent before displaying?
- instead, break stream in series of records
  - each record carries a MAC
  - receiver can act on each record as it arrives
- issue: in record, receiver needs to distinguish MAC from data
  - want to use variable-length records, need to define the length of the data field



# Toy: sequence numbers

- *problem:* attacker can capture and replay record or re-order records
- *solution:* put sequence number into MAC:
  - $MAC = MAC(M_x, \text{sequence}||\text{data})$
  - note: no sequence number field
- *problem:* attacker could replay all records
- *solution:* use nonce to establish initial connection, keys should always be different.

# Toy: control information

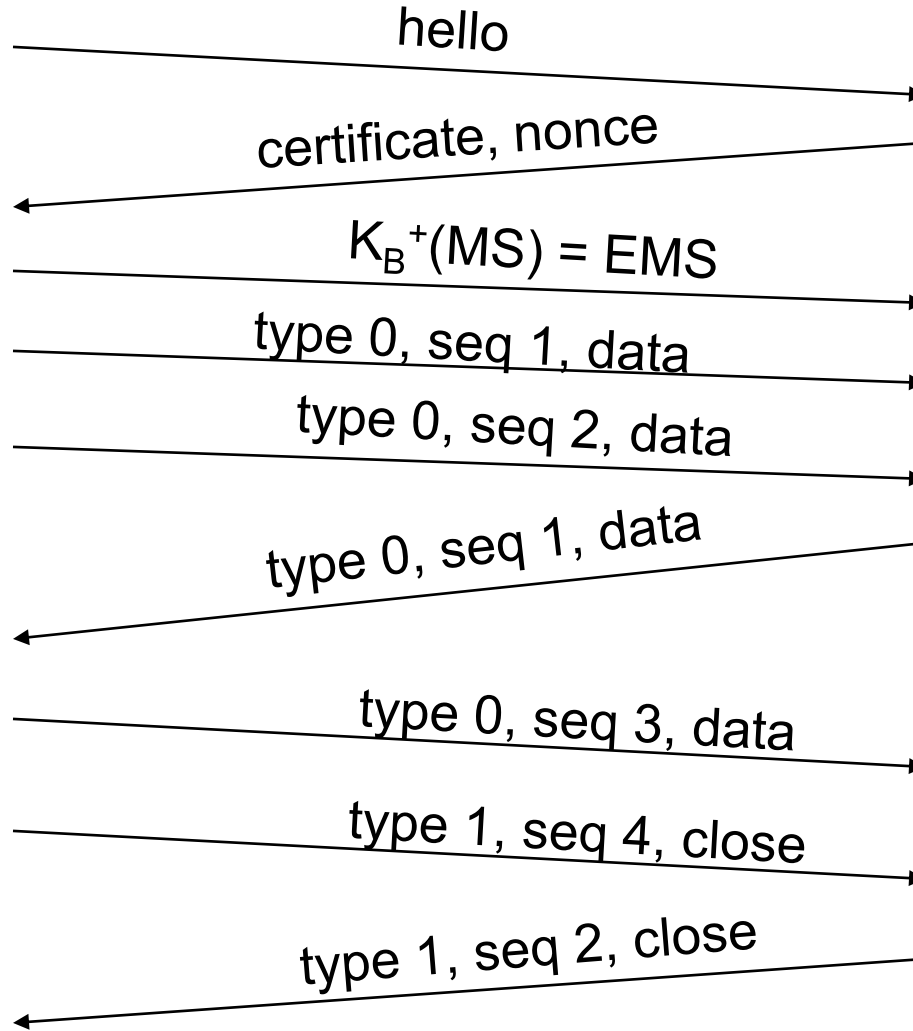
- *problem:* truncation attack:
  - attacker forges TCP connection close segment
  - one or both sides thinks there is less data than there actually is.
- *solution:* record types, with one type for closure
  - type 0 for data; type 1 for closure
- $MAC = MAC(M_x, \text{sequence} || \text{type} || \text{data})$



# Toy SSL: summary



*encrypted*



bob.com

# Toy SSL isn't complete

- how long are fields?
- which encryption protocols?
- want negotiation?
  - allow client and server to support different encryption algorithms
  - allow client and server to choose together specific algorithm before data transfer

# SSL cipher suite

- cipher suite
  - public-key algorithm
  - symmetric encryption algorithm
  - MAC algorithm
- SSL supports several cipher suites
- negotiation: client, server agree on cipher suite
  - client offers choice
  - server picks one

## common SSL symmetric ciphers

- DES – Data Encryption  
Standard: block
- 3DES – Triple strength: block
- RC2 – Rivest Cipher 2: block
- RC4 – Rivest Cipher 4: stream

## SSL Public key encryption

- RSA



# Real SSL: handshake (Part I)

## *Purpose*

1. server authentication
2. negotiation: agree on crypto algorithms
3. establish keys
4. client authentication (optional)

# Real SSL: handshake (Part 2)

1. client sends list of algorithms it supports, along with client nonce
2. server chooses algorithms from list; sends back: choice + certificate + server nonce
3. client verifies certificate, extracts server's public key, generates pre\_master\_secret, encrypts with server's public key, sends to server
4. client and server independently compute encryption and MAC keys from pre\_master\_secret and nonces
5. client sends a MAC of all the handshake messages
6. server sends a MAC of all the handshake messages

# Real SSL: handshaking (Part 3)

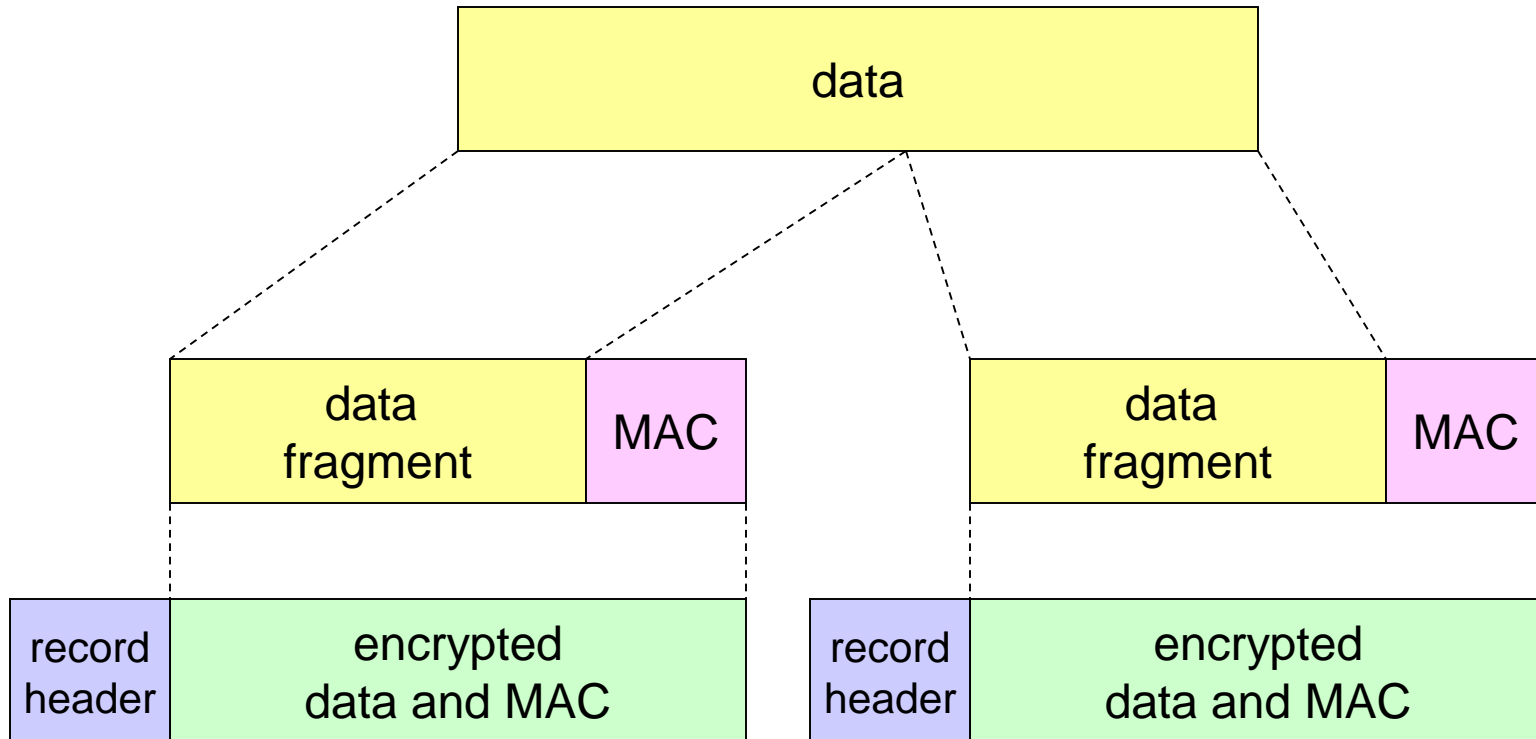
last 2 steps protect handshake from tampering

- client typically offers range of algorithms, some strong, some weak
- man-in-the middle could delete stronger algorithms from list
- last 2 steps prevent this
  - last two messages are encrypted

# Real SSL: handshaking (Part 4)

- why two random nonces?
- suppose Trudy sniffs all messages between Alice & Bob
- next day, Trudy sets up TCP connection with Bob, sends exact same sequence of records
  - Bob (Amazon) thinks Alice made two separate orders for the same thing
  - solution: Bob sends different random nonce for each connection. This causes encryption keys to be different on the two days
  - Trudy's messages will fail Bob's integrity check

# SSL record protocol

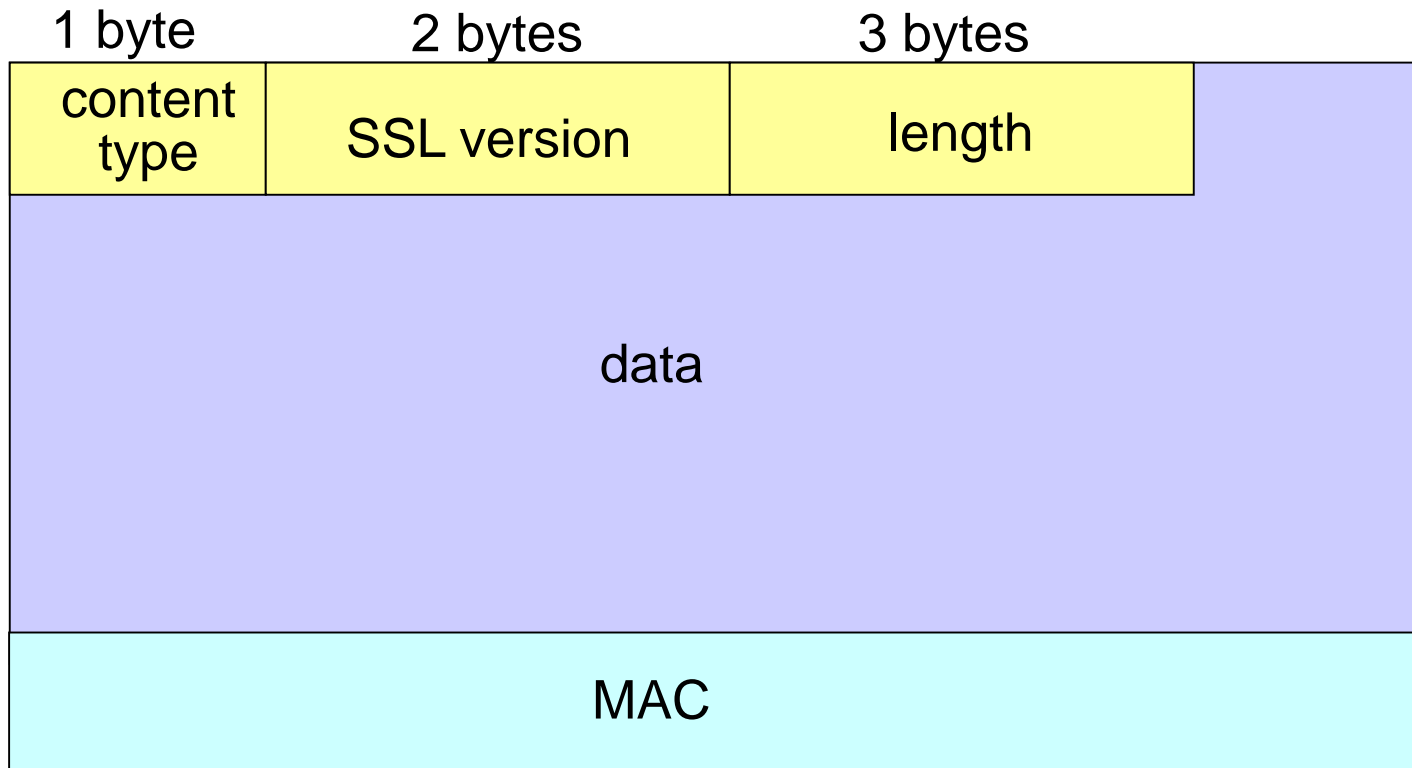


*record header*: content type; version; length

*MAC*: includes seq number, MAC key, Type,  $M_x$

*fragment*: each SSL fragment  $2^{14}$  bytes (~16 Kbytes)

# SSL record format



data and MAC encrypted (symmetric algorithm)

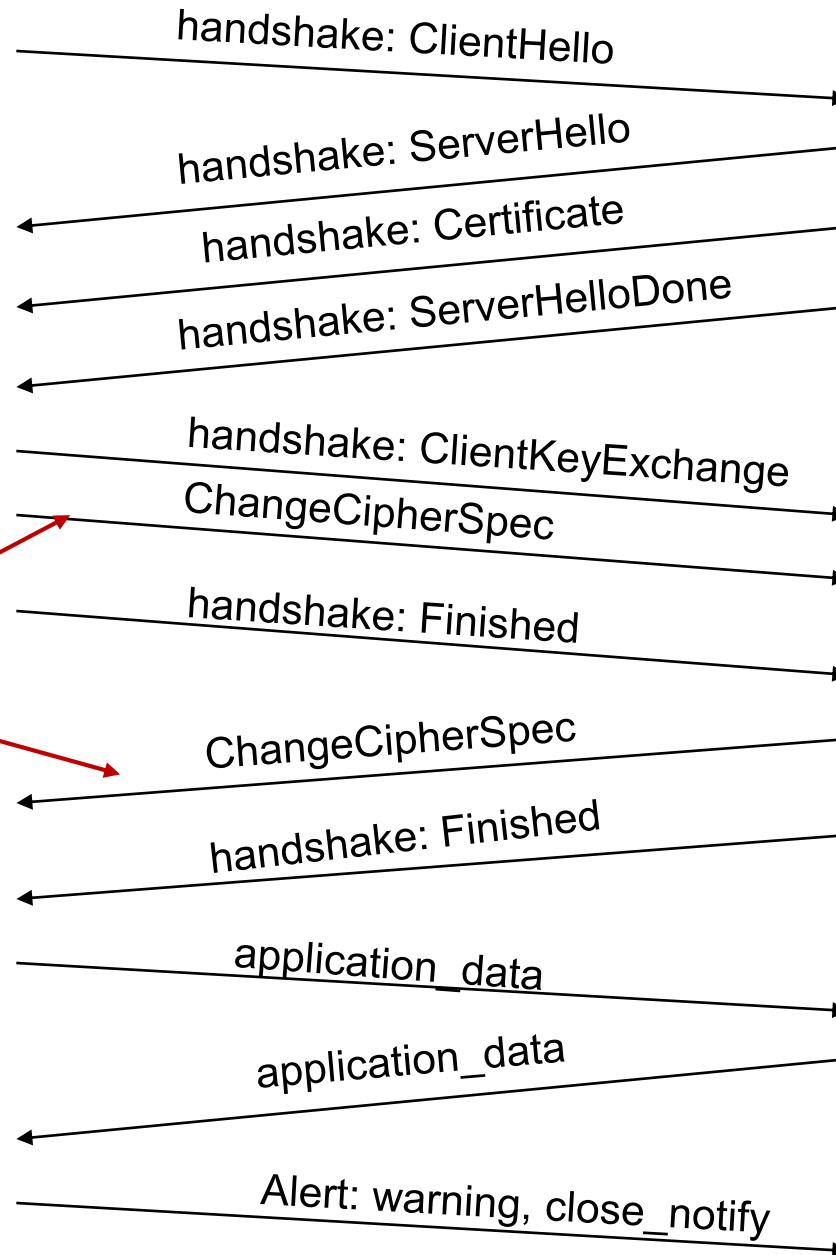
# Real SSL connection

Client has server nonce, server certificate, and encryption scheme. Creates pre master secret, sends to server. Can build keys.



*everything  
henceforth  
is encrypted*

TCP FIN follows



Server now has client nonce & pre master secret. Can build keys.

# Key derivation

- client nonce, server nonce, and pre-master secret input into pseudo random-number generator.
  - produces master secret
- master secret and new nonces input into another random-number generator: “key block”
  - because of resumption: TBD
- key block sliced and diced:
  - client MAC key
  - server MAC key
  - client encryption key
  - server encryption key
  - client initialization vector (IV)
  - server initialization vector (IV)



# Chapter 8 roadmap

8.1 What is network security?

8.2 Principles of cryptography

8.3 Message integrity

8.4 Securing e-mail

8.5 Securing TCP connections: SSL

*8.6 Network layer security: IPsec*

8.7 Securing wireless LANs

8.8 Operational security: firewalls and IDS

# What is network-layer confidentiality ?

*between two network entities:*

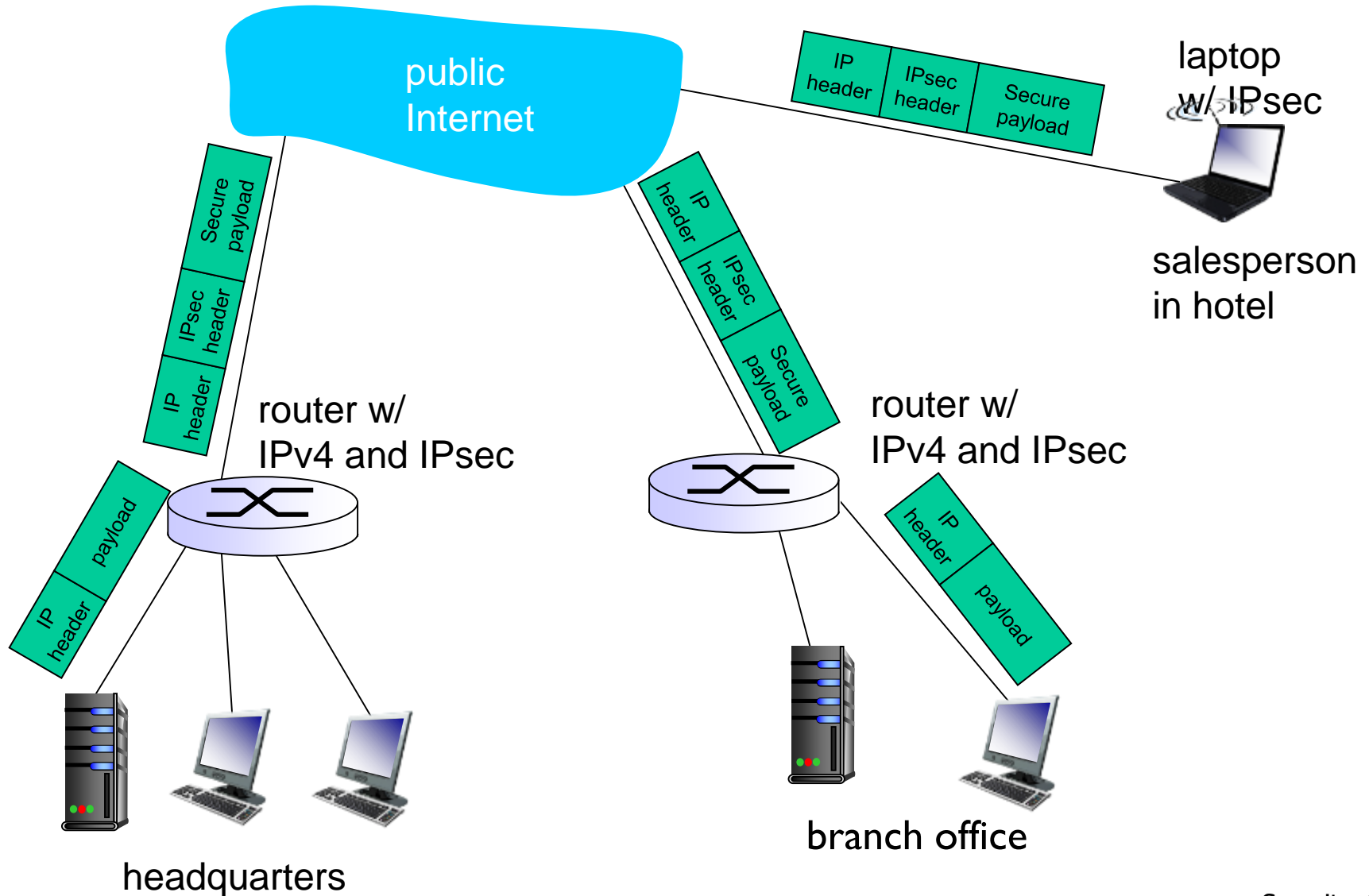
- sending entity encrypts datagram payload, payload could be:
  - TCP or UDP segment, ICMP message, OSPF message ....
- all data sent from one entity to other would be hidden:
  - web pages, e-mail, P2P file transfers, TCP SYN packets
  - ...
- “blanket coverage”

# Virtual Private Networks (VPNs)

## *motivation:*

- institutions often want private networks for security.
  - costly: separate routers, links, DNS infrastructure.
- VPN: institution's inter-office traffic is sent over public Internet instead
  - encrypted before entering public Internet
  - logically separate from other traffic

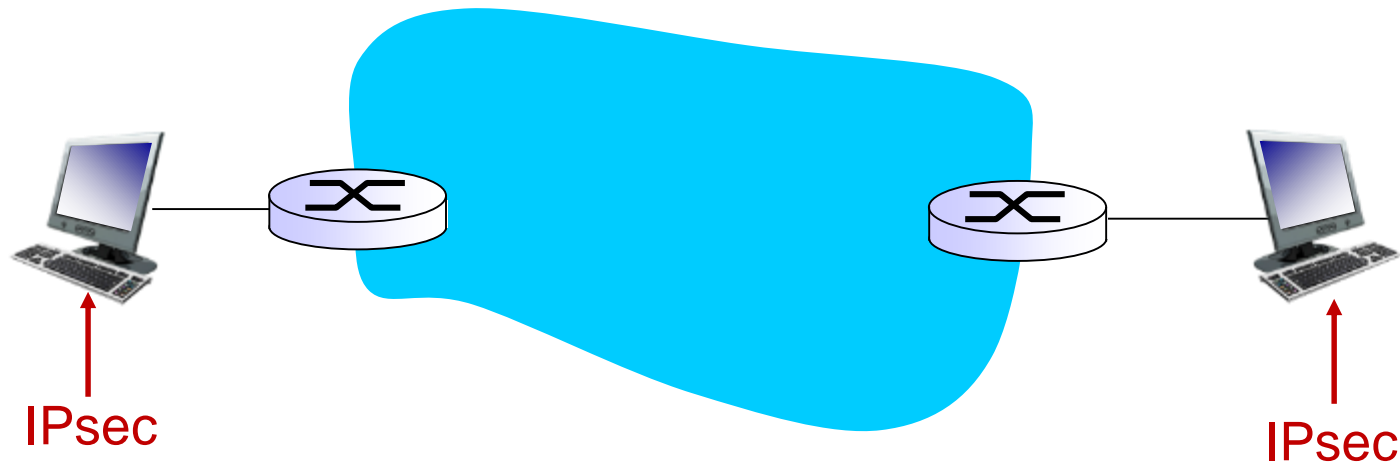
# Virtual Private Networks (VPNs)



# IPsec services

- data integrity
- origin authentication
- replay attack prevention
- confidentiality
  
- two protocols providing different service models:
  - AH
  - ESP

# IPsec transport mode



- IPsec datagram emitted and received by end-system
- larger MTUs possible than with tunneling
- typically used within another tunnel
- IP header is not encrypted, only the payload
- RDP, Encrypted Telnet are examples

# IPsec – tunneling mode



- edge routers IPsec-aware
- hosts IPsec-aware
- IP datagram is completely encrypted
- Easier to traverse NATs (TCP checksums intact)
- Increased overhead, smaller MTUs
- Very secure! (IKEv2/IPSec)

# Two IPsec protocols

- Authentication Header (AH) protocol
  - provides source authentication & data integrity but *not* confidentiality
- Encapsulation Security Protocol (ESP)
  - provides source authentication, data integrity, *and* confidentiality
  - more widely used than AH



# Four combinations are possible!

Host mode with AH	Host mode with ESP
Tunnel mode with AH	Tunnel mode with ESP



most common and  
most important

# IPsec summary

- Used for 'blanket coverage' of IP traffic
- Typically used to provide
  - Data integrity
  - Origin Authentication
  - Confidentiality
- IPsec peers can be two end systems, two routers/firewalls, or a router/firewall and an end system

# TOR & Darkweb – very briefly

- TOR uses relay servers and encryption in an attempt to provide anonymous network access
- TOR provides access to Darkweb sites which are:
  - Not accessible via a conventional web browser
  - Not indexed by conventional search engines (ie Google)
- TOR can be used for good (free speech) or bad (illegal activities)

# Chapter 8 roadmap

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- 8.5 Securing TCP connections: SSL
- 8.6 Network layer security: IPsec
- 8.7 Securing wireless LANs*
- 8.8 Operational security: firewalls and IDS

# WEP design goals



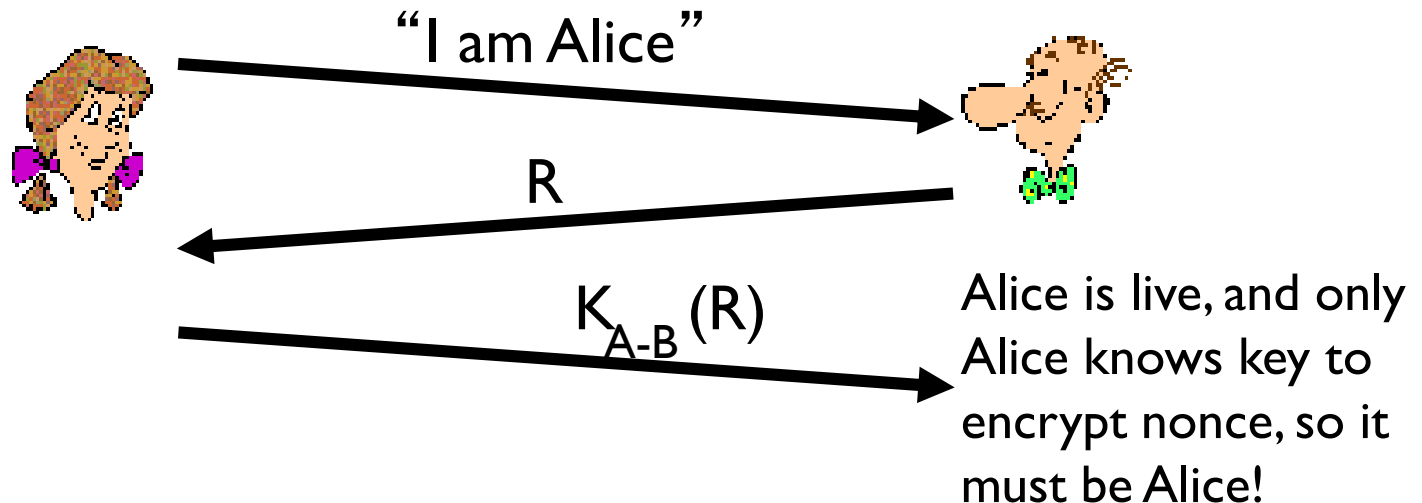
- symmetric key crypto
  - confidentiality
  - end host authorization
  - data integrity
- self-synchronizing: each packet separately encrypted
  - given encrypted packet and key, can decrypt; can continue to decrypt packets when preceding packet was lost (unlike Cipher Block Chaining (CBC) in block ciphers)
- Efficient
  - implementable in hardware or software



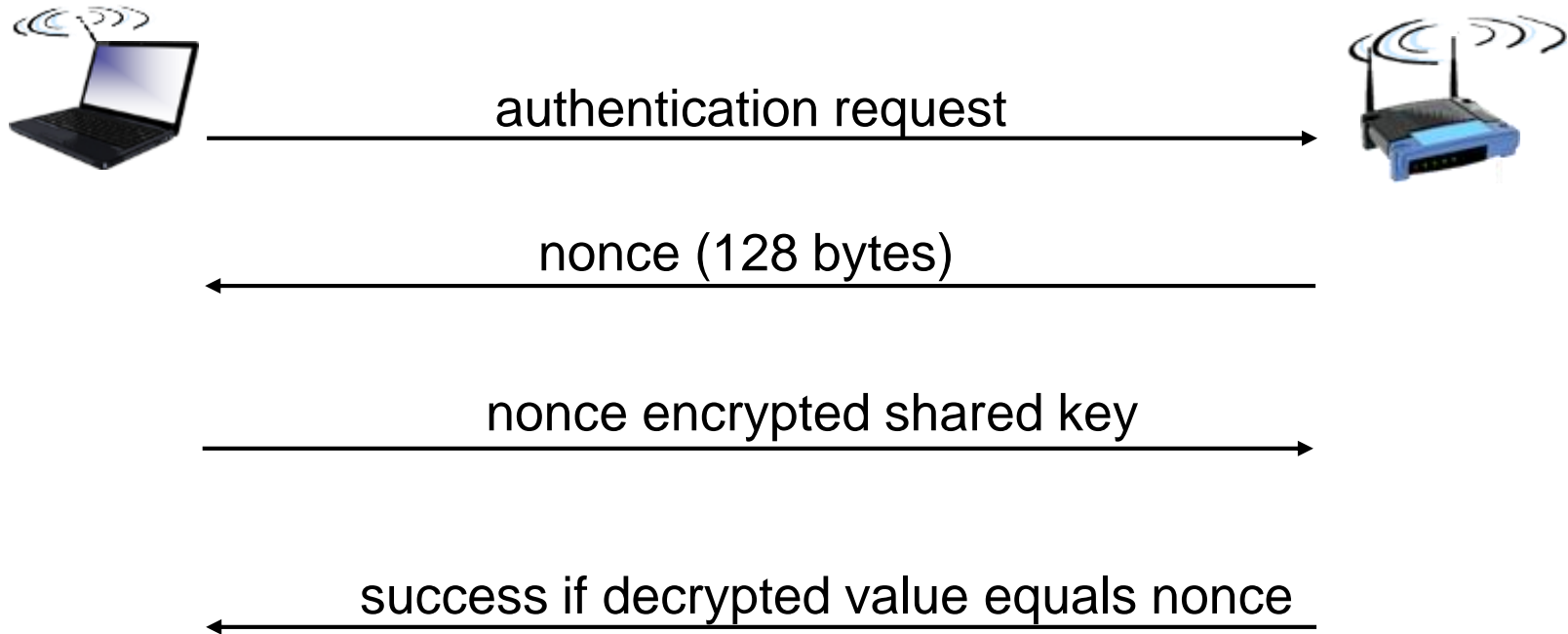
# End-point authentication w/ nonce

**Nonce:** number (R) used only *once* –*in-a-lifetime*

**How to prove Alice “live”:** Bob sends Alice **nonce**, R. Alice must return R, encrypted with shared secret key



# WEP authentication



## *Notes:*

- not all APs do it, even if WEP is being used
- AP indicates if authentication is necessary in beacon frame
- done before association

# Breaking 802.11 WEP encryption

## *security hole:*

- 24-bit initialization vector (IV), one IV per frame, -> IV's eventually reused
- IV transmitted in plaintext -> IV reuse detected

## *attack:*

- Trudy causes Alice to encrypt known plaintext  $d_1 d_2 d_3 d_4 \dots$
- Trudy sees:  $c_i = d_i \text{ XOR } k_i^{\text{IV}}$
- Trudy knows  $c_i d_i$ , so can compute  $k_i^{\text{IV}}$
- Trudy knows encrypting key sequence  $k_1^{\text{IV}} k_2^{\text{IV}} k_3^{\text{IV}} \dots$
- Next time IV is used, Trudy can decrypt!



# 802.11i: improved security

- numerous (stronger) forms of encryption possible
- provides mechanism for key distribution (EAP)
- Can use authentication server separate from access point (Enterprise mode)
- WEP < WPA < WPA2
- WPA still uses RC4 but has larger IVs and uses a 256 bit key. TKIP makes sure each client gets a new key. Fixes major holes in WEP, but only meant as a stop-gap before WPA2
- WPA2 – 256 bit keys, TKIP & RC4 replaced with CCMP & AES. More secure, but still not 100%

# 802.11i: continued

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- Security issues mainly related to enterprise attacks, home networks can be considered secure with WPA2
- WPS on home routers – don't do it.

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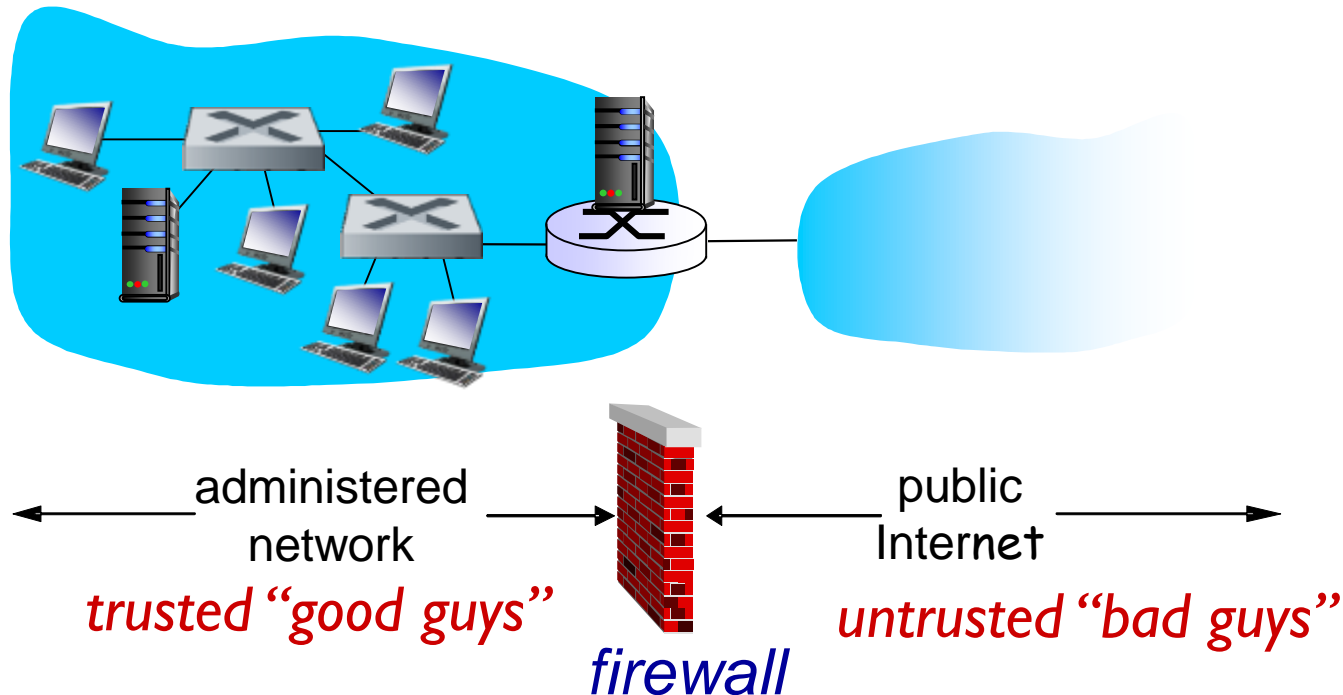
8.7 Securing wireless LANs

*8.8 Operational security: firewalls and IDS*

# Firewalls

## *firewall*

isolates organization's internal net from larger Internet, allowing some packets to pass, blocking others



# Firewalls: why

prevent denial of service attacks:

- SYN flooding: attacker establishes many bogus TCP connections, no resources left for “real” connections

prevent illegal modification/access of internal data

- e.g., attacker replaces CIA’s homepage with something else

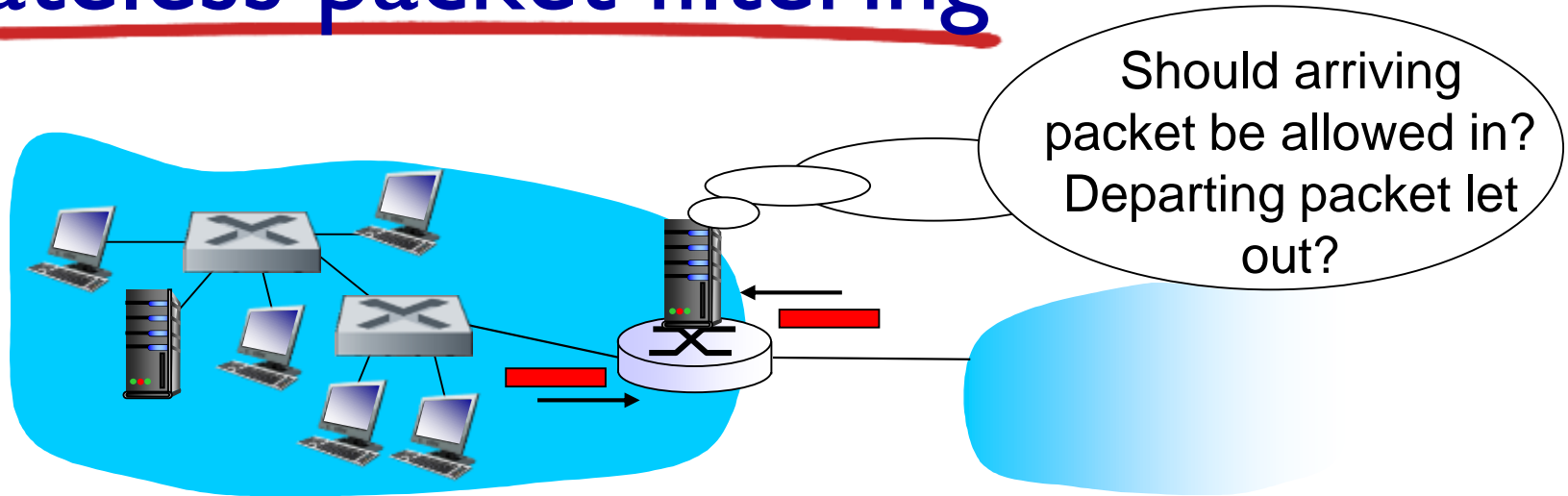
allow only authorized access to inside network

- set of authenticated users/hosts

three types of firewalls:

- stateless packet filters
- stateful packet filters
- application gateways

# Stateless packet filtering



- internal network connected to Internet via *router firewall*
- router *filters packet-by-packet*, decision to forward/drop packet based on:
  - source IP address, destination IP address
  - TCP/UDP source and destination port numbers
  - ICMP message type
  - TCP SYN and ACK bits

# Stateless packet filtering: example

- *example 1:* block incoming and outgoing datagrams with IP protocol field = 17 and with either source or dest port = 23
  - *result:* all incoming, outgoing UDP flows and telnet connections are blocked
- *example 2:* block inbound TCP segments with ACK=0.
  - *result:* prevents external clients from making TCP connections with internal clients, but allows internal clients to connect to outside.

# Stateless packet filtering: more examples

<i>Policy</i>	<i>Firewall Setting</i>
No outside Web access.	Drop all outgoing packets to any IP address, port 80
No incoming TCP connections, except those for institution's public Web server only.	Drop all incoming TCP SYN packets to any IP except 130.207.244.203, port 80
Prevent Web-radios from eating up the available bandwidth.	Drop all incoming UDP packets - except DNS and router broadcasts.
Prevent your network from being used for a smurf DoS attack.	Drop all ICMP packets going to a "broadcast" address (e.g. 130.207.255.255).
Prevent your network from being tracerouted	Drop all outgoing ICMP TTL expired traffic



# Access Control Lists

**ACL:** table of rules, applied top to bottom to incoming packets:  
(action, condition) pairs: looks like OpenFlow forwarding (Ch. 4)!

action	source address	dest address	protocol	source port	dest port	flag bit
allow	222.22/16	outside of 222.22/16	TCP	> 1023	80	any
allow	outside of 222.22/16	222.22/16	TCP	80	> 1023	ACK
allow	222.22/16	outside of 222.22/16	UDP	> 1023	53	---
allow	outside of 222.22/16	222.22/16	UDP	53	> 1023	----
deny	all	all	all	all	all	all

# Stateful packet filtering

- *stateless packet filter*: heavy handed tool
  - admits packets that “make no sense,” e.g., dest port = 80, ACK bit set, even though no TCP connection established:

action	source address	dest address	protocol	source port	dest port	flag bit
allow	outside of 222.22/16	222.22/16	TCP	80	> 1023	ACK

- *stateful packet filter*: track status of every TCP connection
  - track connection setup (SYN), teardown (FIN): determine whether incoming, outgoing packets “makes sense”
  - timeout inactive connections at firewall: no longer admit packets

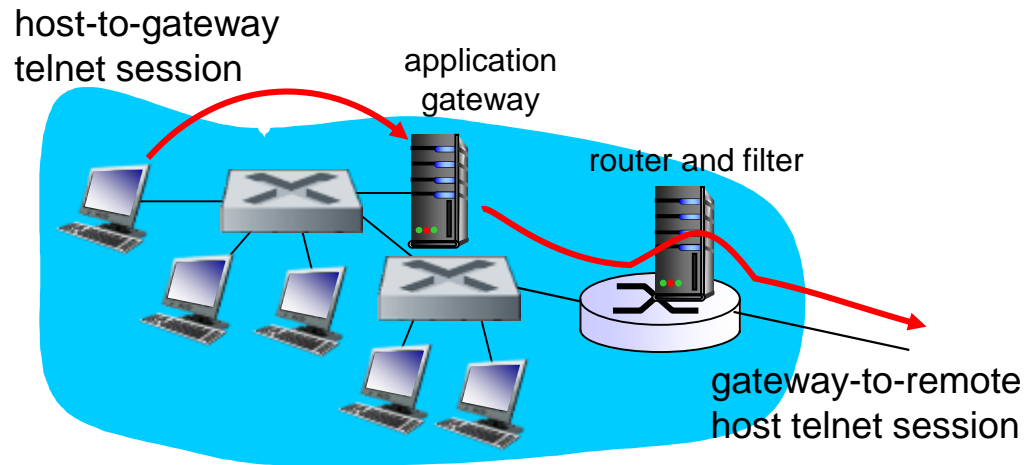
# Stateful packet filtering

ACL augmented to indicate need to check connection state table before admitting packet

action	source address	dest address	proto	source port	dest port	flag bit	check conxion
allow	222.22/16	outside of 222.22/16	TCP	> 1023	80	any	
allow	outside of 222.22/16	222.22/16	TCP	80	> 1023	ACK	X
allow	222.22/16	outside of 222.22/16	UDP	> 1023	53	---	
allow	outside of 222.22/16	222.22/16	UDP	53	> 1023	----	X
deny	all	all	all	all	all	all	

# Application gateways

- filter packets on application data as well as on IP/TCP/UDP fields.
- *example:* allow select internal users to telnet outside



1. require all telnet users to telnet through gateway.
2. for authorized users, gateway sets up telnet connection to dest host. Gateway relays data between 2 connections
3. router filter blocks all telnet connections not originating from gateway.

# Limitations of firewalls, gateways

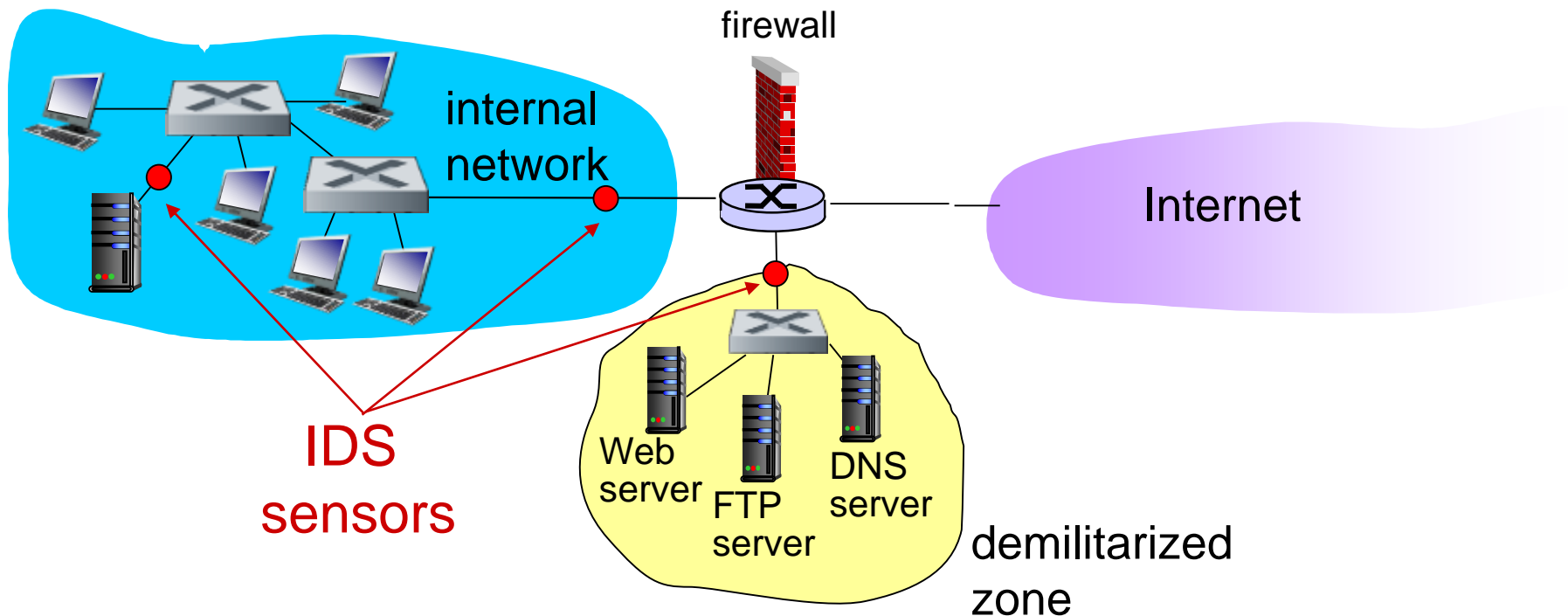
- *IP spoofing*: router can't know if data “really” comes from claimed source
- if multiple app's. need special treatment, each has own app. gateway
- client software must know how to contact gateway.
  - e.g., must set IP address of proxy in Web browser
- filters often use all or nothing policy for UDP
- *tradeoff*: degree of communication with outside world, level of security
- many highly protected sites still suffer from attacks

# Intrusion detection systems

- packet filtering:
  - operates on TCP/IP headers only
  - no correlation check among sessions
- *IDS: intrusion detection system*
  - *deep packet inspection*: look at packet contents (e.g., check character strings in packet against database of known virus, attack strings)
  - *examine correlation* among multiple packets
    - port scanning
    - network mapping
    - DoS attack

# Intrusion detection systems

multiple IDSs: different types of checking at different locations



# Network Security (summary)

## basic techniques.....

- cryptography (symmetric and public)
- message integrity
- end-point authentication

## .... used in many different security scenarios

- secure email
- secure transport (SSL)
- IP sec
- 802.11

## operational security: firewalls and IDS