

# Genomes Comparision via de Bruijn graphs

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# Synteny Blocks: Algorithmic challenge

- ▶ Suppose that we are given two genomes
- ▶ The question is: how are they evolutionary related to each other?
- ▶ In order to do rearrangements analysis we must decompose genomes into synteny blocks
- ▶ Synteny blocks are evolutionary conserved segments of the genome
- ▶ These blocks cover most of the genome
- ▶ Occur in both genomes with possible variations

# Academic Project

Project: Identify synteny blocks for duplicated genomes represented as sequences of **nucleotides**.

- ▶ **None** of the previous synteny blocks reconstruction software (DRIMM-Synteny (Pham And Pevzner 2010) included) can efficiently solve this problem.
- ▶ DRIMM-Synteny can find the synteny blocks for complicated genomes. But:

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- ▶ DRIMM-Synteny can find the synteny blocks for complicated genomes. But:
- ▶ It requires the genome to be represented as sequence of genes.

# General Idea: de Bruijn Graph

- ▶ We are given an alphabet  $\Sigma$  and a string  $S$  over it,  $|\Sigma| = m$
- ▶ A substring  $T$ ,  $|T| = k$  is called *k-mer*
- ▶ De Bruijn graph is a multigraph  $G_k = (V, E)$ , where
$$V = \Sigma^{k-1} = \{\text{all possible } (k-1)\text{-mers}\}$$
- ▶ If *k-mer*  $T$  is presented in  $S$ , then we add an oriented edge  $(T[1, k-1], T[2, k])$  to the graph
- ▶ Create de Bruijn graph from the nucleotide sequence
- ▶ Conserved regions will yield non-branching paths

# Challenges

- ▶ Variations in synteny blocks generate cycles, so we need to simplify the graph
- ▶ Double strandness: conserved regions may occur on both strands. Example:  
5' AACCGGTT 3'  
3' TTGGCAA 5'  
Such blocks are reverse complementary to each other  $\Rightarrow$  no non-branching paths
- ▶ Spurious similarity
- ▶ Memory efficiency

# Colored graph

- ▶ We use colored de Bruijn graphs [Iqbal et al., 2012] to handle double-strandness
- ▶ Suppose that  $S^+$  and  $S^-$  are positive and negative strands of the chromosome
- ▶ Colored de Bruijn graph is a multigraph  $G_k = (V, E)$  where  $V = \Sigma^{k-1}$
- ▶ For each  $k$ -mer  $T^+$  in  $S^+$  add edge  $(T^+[1, k-1], T^+[2, k])$  to  $G_k$  and mark it *blue*
- ▶ For each  $k$ -mer  $T^-$  in  $S^-$  add edge  $(T^-[1, k-1], T^-[2, k])$  to  $G_k$  and mark it *red*

# Edge labeling

- ▶ Note that our graph is built from a string, not set of reads
- ▶ Each walk in the graph represents a string
- ▶ We are interested only in walks that represent substrings of the source string
- ▶ Assign to each edge  $e$  label  $L(e) =$  position of the corresponding  $k$ -mer on the positive strand
- ▶ Walk  $W = (v_1 e_1 v_2 e_2 \dots)$  is considered valid iff:
  1.  $e_i$  and  $e_{i+1}$  are of the same color
  2.  $|L(e_i) - L(e_{i+1})| = 1$



# Example

5' ACCTGTCAGT 3'  
3' TGGACA GTCA 5'

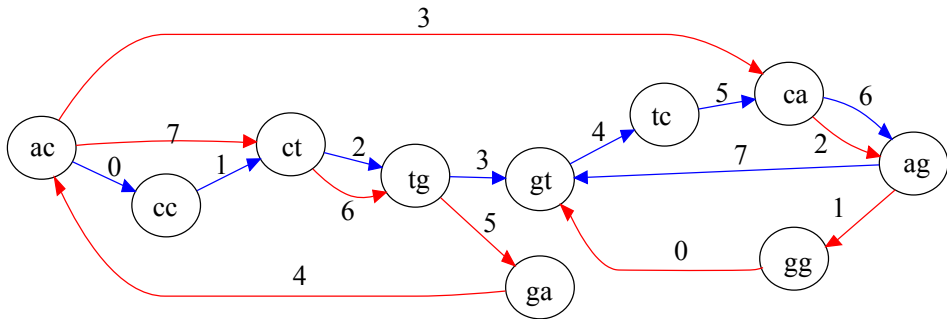


Figure 1: Colored de Bruijn graph built from two strands

# Graph simplification

- ▶ Bulges spoil long non-branching paths and indicate indels/mismatches
- ▶ A pair of walks ( $W_1, W_2$ ) is a bulge iff:
  - 1) Start and end vertices of  $W_1$  and  $W_2$  coincide
  - 2)  $W_1$  and  $W_2$  have exactly 2 common vertices
  - 3) There are no edges  $u \in W_1$  and  $v \in W_2$  such that  $L(u) = L(v)$
  - 4)  $|W_1| \leq \delta$  and  $|W_2| \leq \delta$

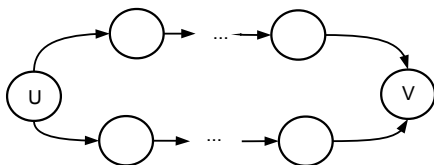


Figure 2: A bulge

# General pipeline

- ▶ Build de Bruijn graph from the genome
- ▶ Remove bulges (BFS-like algorithm)
- ▶ Bulges are removed by replacing long branches with shorter ones
- ▶ Output non-branching paths

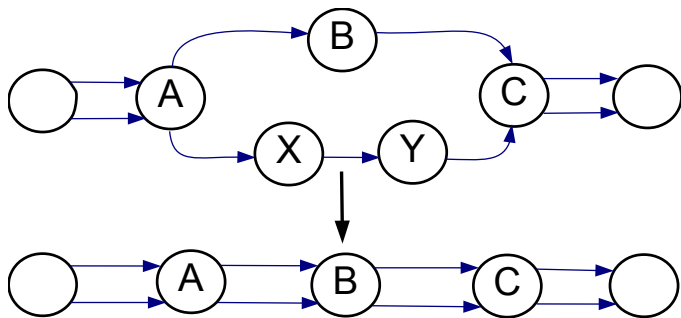


Figure 3: Bulge removal illustration

# Parameters selection

- ▶ How should we choose  $K$  and  $\delta$ ?
- ▶ Duplicated genes can have no long ( $K > 50$ ) shared  $K$  - mers
- ▶ Big  $K \sim 50$  – we find only few synteny blocks
- ▶ Small  $K \sim 10$  and small  $\delta \sim 15$  – we find very short synteny blocks
- ▶ Small  $K \sim 10$  and big  $\delta \sim 200$  – the genome will be disrupted completely

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- ▶ Solution – do simplification in multiple stages

# New pipeline

- ▶ General idea – "align" similar regions first, then glue them together into syntenic blocks
- ▶ Start with small  $K$  and small  $\delta$  to smooth duplicated regions and obtain long  $K$ -mers
- ▶ Rebuild and simplify the graph with higher  $K$  and  $\delta$
- ▶ Continue this process several times
- ▶ Final step can be done with  $K \sim$  several hundreds

# Experiment

- ▶ We have attempted to identify duplications in *Arabidopsis thaliana*
- ▶ Arabidopsis is known to be highly duplicated genome [Arabidopsis Genome Initiative]
- ▶ Size of the genome is  $\sim 120 \text{ Mbp}$
- ▶ We used 4 stages and following parameters:

Stage number	$K$	$\delta$
1	15	150
2	50	500
3	100	1000
4	500	5000

# Computation results

- ▶ We have found 4722 syntenic blocks in Arabidopsis
- ▶ These blocks cover 28 % of the genome
- ▶ Minimum length of the block is 1000 *bp*
- ▶ Largest block found has length  $\sim 95\,000$  *bp*
- ▶ We tried to verify blocks by aligning instances of the same block
- ▶ At least 74 % of blocks have 70 % of exact matches



# Computation results

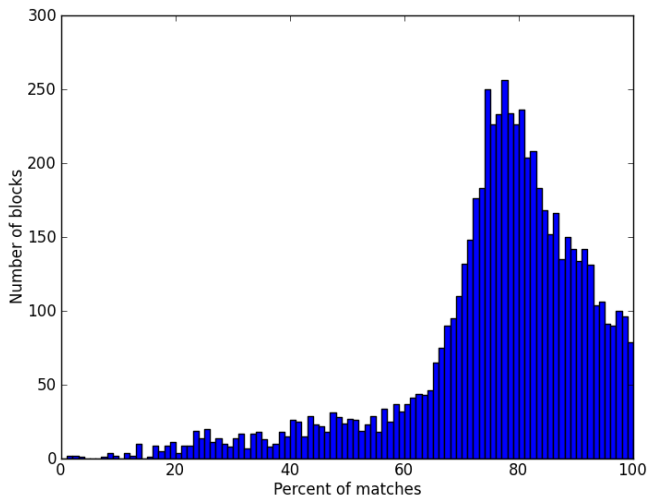


Figure 4: Matches percent vs. number of blocks plot

# Computation results

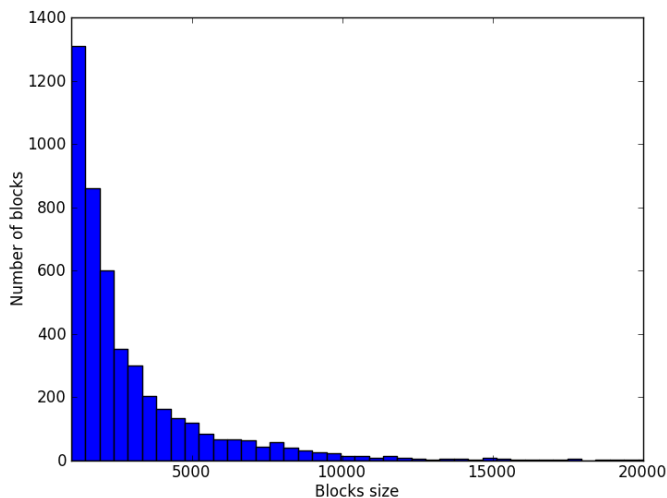


Figure 5: Synteny blocks length distribution

# Future plans and summary

## Summary

- ▶ We have covered 28 % of Arabidopsis genome with synteny blocks
- ▶ But we have missed some duplicated regions, described in [Arabidopsis Genome Initiative]
- ▶ Most of the blocks are short ( $< 5000$  bp)

## Future plans

- ▶ Improve coverage
- ▶ Examine other genomes
- ▶ Optimize algorithms to handle larger genomes

# References

- ▶ 1. Pevzner P and Tesler G, (2003) Human and mouse genomic sequences reveal extensive breakpoint reuse in mammalian evolution.
- ▶ 2. Pham S and Pevzner P, (2010) DRIMM-Synteny: Decomposing Genomes into Evolutionary Conserved Segments
- ▶ 3. Iqbal Z, Caccamo M, Turner I, Flicek P, McVean G, (2012) De novo assembly and genotyping of variants using colored de Bruijn graphs
- ▶ 4. Arabidopsis Genome Initiative, (2000) Analysis of the genome sequence of the flowering plant *Arabidopsis thaliana*

Thank you!