

Genomes Comparision via de Bruijn graphs

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Synteny Blocks: Algorithmic challenge

- ▶ Suppose that we are given two genomes
- ▶ The question is: how are they evolutionary related to each other?
- ▶ In order to do rearrangements analysis we must decompose genomes into synteny blocks
- ▶ Synteny blocks are evolutionary conserved segments of the genome
- ▶ These blocks cover most of the genome
- ▶ Occur in both genomes with possible variations

Academic Project

Project: Identify synteny blocks for duplicated genomes represented as sequences of **nucleotides**.

- ▶ **None** of the previous synteny blocks reconstruction software (DRIMM-Synteny (Pham And Pevzner 2010) included) can efficiently solve this problem.
- ▶ DRIMM-Synteny can find the synteny blocks for complicated genomes. But:

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- ▶ It requires the genome to be represented as sequence of genes.

General Idea: de Bruijn Graph

- ▶ We are given an alphabet Σ and a string S over it, $|\Sigma| = m$
- ▶ A substring T , $|T| = k$ is called *k-mer*
- ▶ De Bruijn graph is a multigraph $G_k = (V, E)$, where
$$V = \Sigma^{k-1} = \{\text{all possible } (k-1)\text{-mers}\}$$
- ▶ If *k-mer* T is presented in S , then we add an oriented edge $(T[1, k-1], T[2, k])$ to the graph
- ▶ Create de Bruijn graph from the nucleotide sequence
- ▶ Conserved regions will yield non-branching paths

Challenges

- ▶ Variations in synteny blocks generate cycles, so we need to simplify the graph
- ▶ Double strandness: conserved regions may occur on both strands. Example:
5' AACCGGTT 3'
3' TTGGCAA 5'
Such blocks are reverse complementary to each other \Rightarrow no non-branching paths
- ▶ Spurious similarity
- ▶ Memory efficiency

Colored graph

- ▶ We use colored de Bruijn graphs [Iqbal et al., 2012] to handle double-strandness
- ▶ Suppose that S^+ and S^- are positive and negative strands of the chromosome
- ▶ Colored de Bruijn graph is a multigraph $G_k = (V, E)$ where $V = \Sigma^{k-1}$
- ▶ For each k -mer T^+ in S^+ add edge $(T^+[1, k-1], T^+[2, k])$ to G_k and mark it *blue*
- ▶ For each k -mer T^- in S^- add edge $(T^-[1, k-1], T^-[2, k])$ to G_k and mark it *red*

Edge labeling

- ▶ Note that our graph is built from a string, not set of reads
- ▶ Each walk in the graph represents a string
- ▶ We are interested only in walks that represent substrings of the source string
- ▶ Assign to each edge e label $L(e) =$ position of the corresponding k -mer on the positive strand
- ▶ Walk $W = (v_1 e_1 v_2 e_2 \dots)$ is considered valid iff:
 1. e_i and e_{i+1} are of the same color
 2. $|L(e_i) - L(e_{i+1})| = 1$

Example

5' ACCTGTCAGT 3'
3' TGGACAGTCA 5'

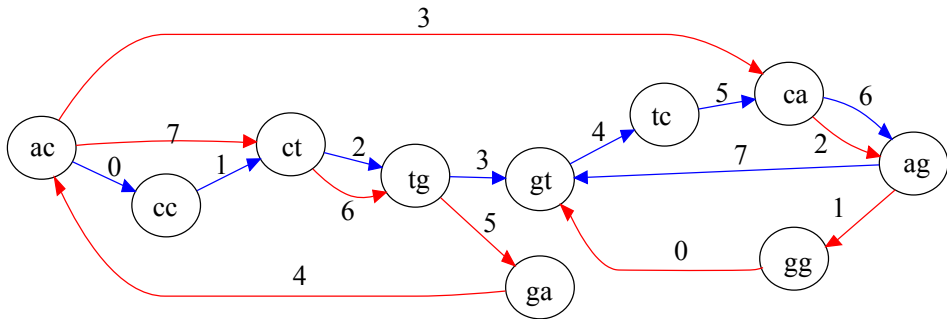


Figure 1: Colored de Bruijn graph built from two strands

Graph simplification

- ▶ Bulges spoil long non-branching paths and indicate indels/mismatches
- ▶ A pair of walks (W_1, W_2) is a bulge iff:
 - 1) Start and end vertices of W_1 and W_2 coincide
 - 2) W_1 and W_2 have exactly 2 common vertices
 - 3) There are no edges $u \in W_1$ and $v \in W_2$ such that $L(u) = L(v)$
 - 4) $|W_1| \leq \delta$ and $|W_2| \leq \delta$

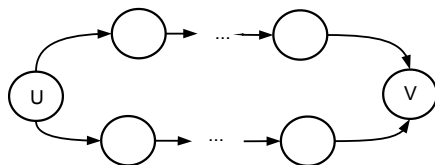


Figure 2: A bulge

General pipeline

- ▶ Build de Bruijn graph from the genome
- ▶ Remove bulges (BFS-like algorithm)
- ▶ Bulges are removed by replacing long branches with shorter ones
- ▶ Output non-branching paths

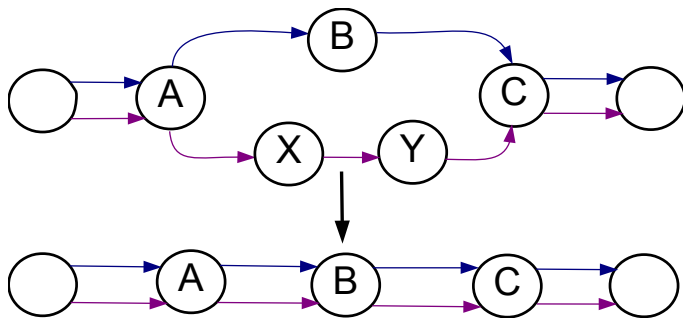


Figure 3: Bulge removal illustration

Parameters selection

- ▶ How should we choose K and δ ?
- ▶ Duplicated genes can have no long ($K > 50$) shared K - mers
- ▶ Big $K \sim 50$ – we find only few synteny blocks
- ▶ Small $K \sim 10$ and small $\delta \sim 15$ – we find very short synteny blocks
- ▶ Small $K \sim 10$ and big $\delta \sim 200$ – the genome will be disrupted completely

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- ▶ Solution – do simplification in multiple stages

New pipeline

- ▶ General idea – "align" similar regions first, then glue them together into syntenic blocks
- ▶ Start with small K and small δ to smooth duplicated regions and obtain long K -mers
- ▶ Rebuild and simplify the graph with higher K and δ
- ▶ Continue this process several times
- ▶ Final step can be done with $K \sim$ several hundreds

Current Progress

Now:

- ▶ Program for finding syntenic blocks using above approach
- ▶ Covers 20 – 25% of Arabidopsis genome

Future work:

- ▶ Improve coverage somehow?
- ▶ Propose new alignment tool?

References

- ▶ 1. Pevzner P and Tesler G, (2003) Human and mouse genomic sequences reveal extensive breakpoint reuse in mammalian evolution.
- ▶ 2. Pham S and Pevzner P, (2010) DRIMM-Synteny: Decomposing Genomes into Evolutionary Conserved Segments
- ▶ 3. Iqbal Z, Caccamo M, Turner I, Flicek P, McVean G, (2012) De novo assembly and genotyping of variants using colored de Bruijn graphs

Thank you!