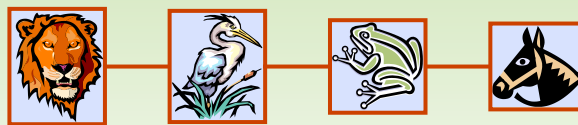


## Linked Lists

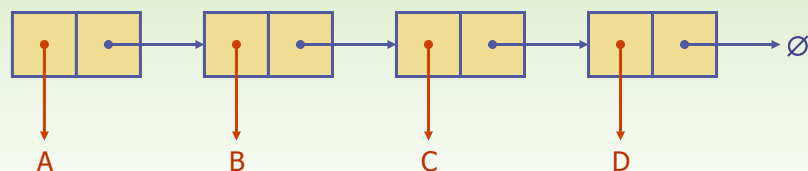
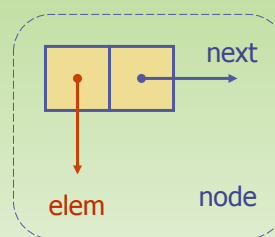


Linked Lists

1

## Singly Linked List

- A singly linked list is a concrete data structure consisting of a sequence of nodes
- Each node stores:
  - an element
  - a link to the next node



Linked Lists

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## The Node Class for List Nodes

```
public class StringNode
{
    // Attributes
    private String element; // The data to be stored in this node
    private StringNode next; // A link to the next node in the chain

    /** Constructor
     * Creates a node with the given element and next node.
     */
    public StringNode(String e, StringNode n)
    {
        element = e;
        next = n;
    }

    /** Constructor
     * Creates a node with null references to its element and next node.
     */
    public StringNode()
    {
        this(null, null);
    }
}
```

Linked Lists

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## The Node Class for List Nodes

```
// Accessor methods

public String getElement()
{
    return element;
}

public StringNode getNext()
{
    return next;
}

// Modifier methods:
public void setElement(String newElem)
{
    element = newElem;
}

public void setNext(StringNode newNext)
{
    next = newNext;
}
}
```

Linked Lists

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## For any data structure/storage:

### Operations

- How to add
- How to retrieve
- How to remove/delete.

### Location of Operation

- Beginning
- End
- Middle

Linked Lists

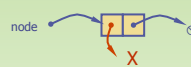
5

## Inserting at the Head of a List

1. The current list...

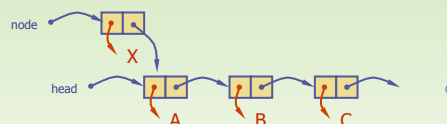


2. Create a new node with content X

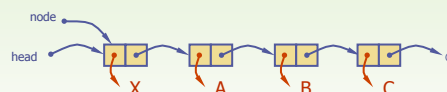


3. Insert the node

- A. Have new node point to old head



- B. Update head to point to new node



Linked Lists

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## Linked List - addFirst

```
public class StringLinkedList {
    StringNode head = null;           // The start of the list
    StringNode tail = null;           // The last element in the list

    public void addFirst(StringNode n)
    {
        // Check we have a node to add...
        if (n == null) return;

        // Set our new node to point to the head
        // of the current list.
        n.setNext(head);

        // Our new node 'n' will now become the head of the list
        head = n;

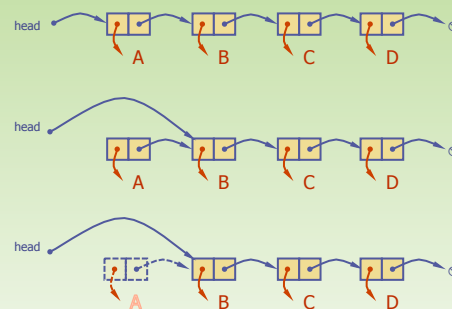
        // If the list was empty, make the tail point to
        // the new node as well
        if (tail == null) tail = n;
    }
}
```

Linked Lists

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## Removing from the Head

1. Update head to point to next node in the list
2. Allow garbage collector to reclaim the former first node



Linked Lists

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## Linked List - removeFirst

```
public void removeFirst()
{
    // If list is empty, we can't remove anything so leave
    if (head == null) return;

    // Move head to the next item in the list
    head = head.getNext();

    // If the list is empty, set the tail reference to null
    if (head == null) tail = null;

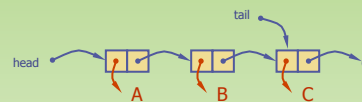
    // The original item that was at the head of the list
    // no longer has anything referencing it and will be
    // garbage collected in Java. In other programming languages
    // e.g. C++, you would have to delete it other wise you
    // would get a memory leak.
}
```

Linked Lists

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## Inserting at the Tail

1. The current list...

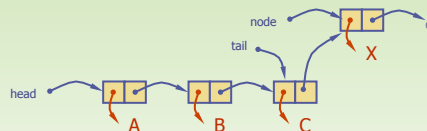


2. Create a new node pointing to null

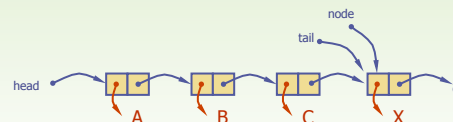


3. Insert new element

- A. Have old last node point to new node



- B. Update tail to point to new node



Linked Lists

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## Linked List - addLast

```
public void addLast(StringNode node)
{
    // If we were not given a node, leave
    if (node == null) return;

    // If list is empty, our new node will
    // be the head and tail of list
    if (head == null)
    {
        head = node;
        tail = node;
        return;
    }

    // Make the current last node point to our new node
    tail.setNext(node);

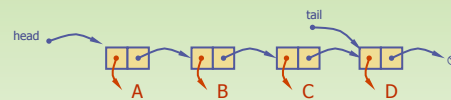
    // Now update the tail to be our new node
    tail = node;
}
```

Linked Lists

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## Removing at the Tail

- Removing at the tail of a singly linked list is not efficient!
- There is no constant-time way to update the tail to point to the previous node



Linked Lists

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## Linked List - removeLast

```
public void removeLast() {
    if (head == null) return; // If list is empty, leave

    // If head is also the tail, the list
    // will be empty
    if (head == tail) {
        head = null;
        tail = null;
        return;
    }

    // Start at the head of the list
    StringNode n = head;

    // Now look for the last item
    while (n.getNext() != tail)
        n = n.getNext();

    // n should now be pointing to the last but one
    // node in the list. This will be the new tail
    // We are going to drop the last element in the list
    // so make the current node's next pointer null
    n.setNext(null);

    // The old tail node is now replaced with 'n'. The
    // old tail node has no reference and will be garbage
    tail = n;
}
```

Linked Lists

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## Node List ADT

- The **Node List** ADT models a sequence of positions storing arbitrary objects
- It establishes a before/after relation between positions
- Generic methods:
  - ♦ **size()**, **isEmpty()**

Lists

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## Node List ADT

Accessor methods:

- `first()`, `last()`
- `prev(p)`, `next(p)`

Update methods:

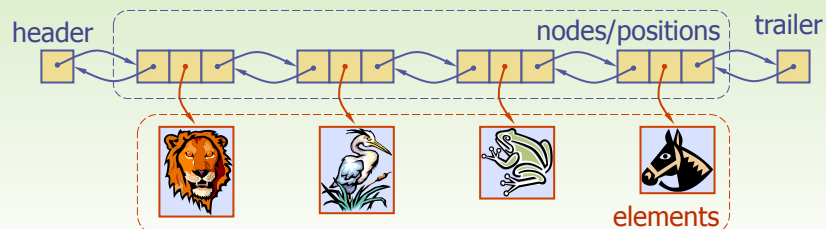
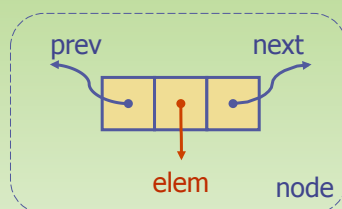
- `set(p, e)`
- `addBefore(p, e)`, `addAfter(p, e)`,
- `addFirst(e)`, `addLast(e)`
- `remove(p)`

Lists

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## Doubly Linked List

- A doubly linked list provides a natural implementation of the Node List ADT
- Nodes implement Position and store:
  - element
  - link to the previous node
  - link to the next node
- Special trailer and header nodes



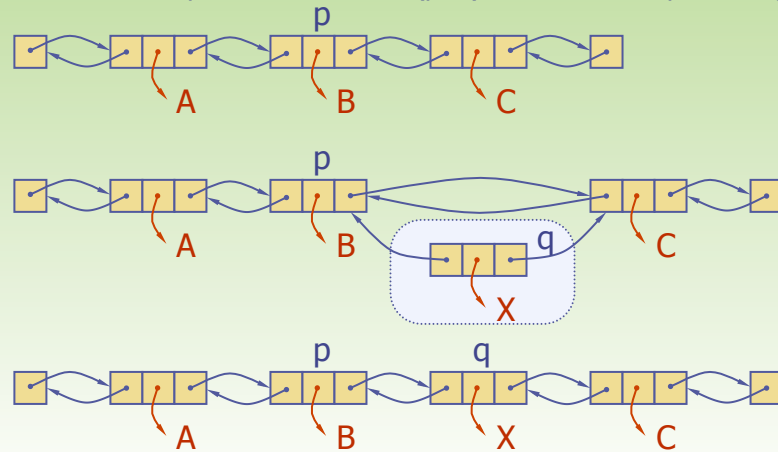
Lists

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## Insertion

- We visualize operation `insertAfter(p, X)`, which returns position `q`



Lists

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## Insertion Algorithm

**Algorithm** `addAfter(p,e)`:

Create a new node `v`

`v.setElement(e)`

`v.setPrev(p)`      {link `v` to its predecessor}

`v.setNext(p.getNext())`    {link `v` to its successor}

`(p.getNext()).setPrev(v)` {link `p`'s old successor to `v`}

`p.setNext(v)`      {link `p` to its new successor, `v`}

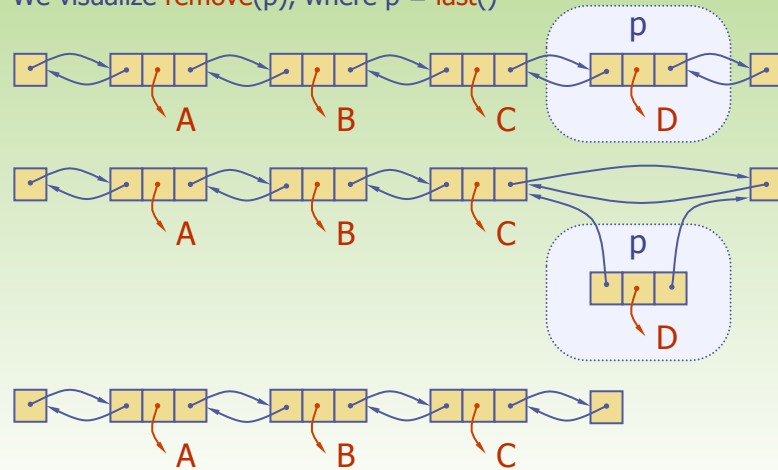
**return** `v`    {the position for the element `e`}

Lists

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## Deletion

- We visualize `remove(p)`, where `p = last()`



Lists

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## Deletion Algorithm

### Algorithm `remove(p)`:

```

t = p.element    {a temporary variable to hold the
                  return value}
(p.getPrev()).setNext(p.getNext())    {linking out p}
(p.getNext()).setPrev(p.getPrev())
p.setPrev(null)    {invalidating the position p}
p.setNext(null)
return t
  
```

Lists

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## Performance

- In the linked list implementation of the List ADT
  - The space used by a list with  $n$  elements is  $O(n)$
  - The space used by each position of the list is  $O(1)$
  - Operations of the List ADT run in up to  $O(n)$  time