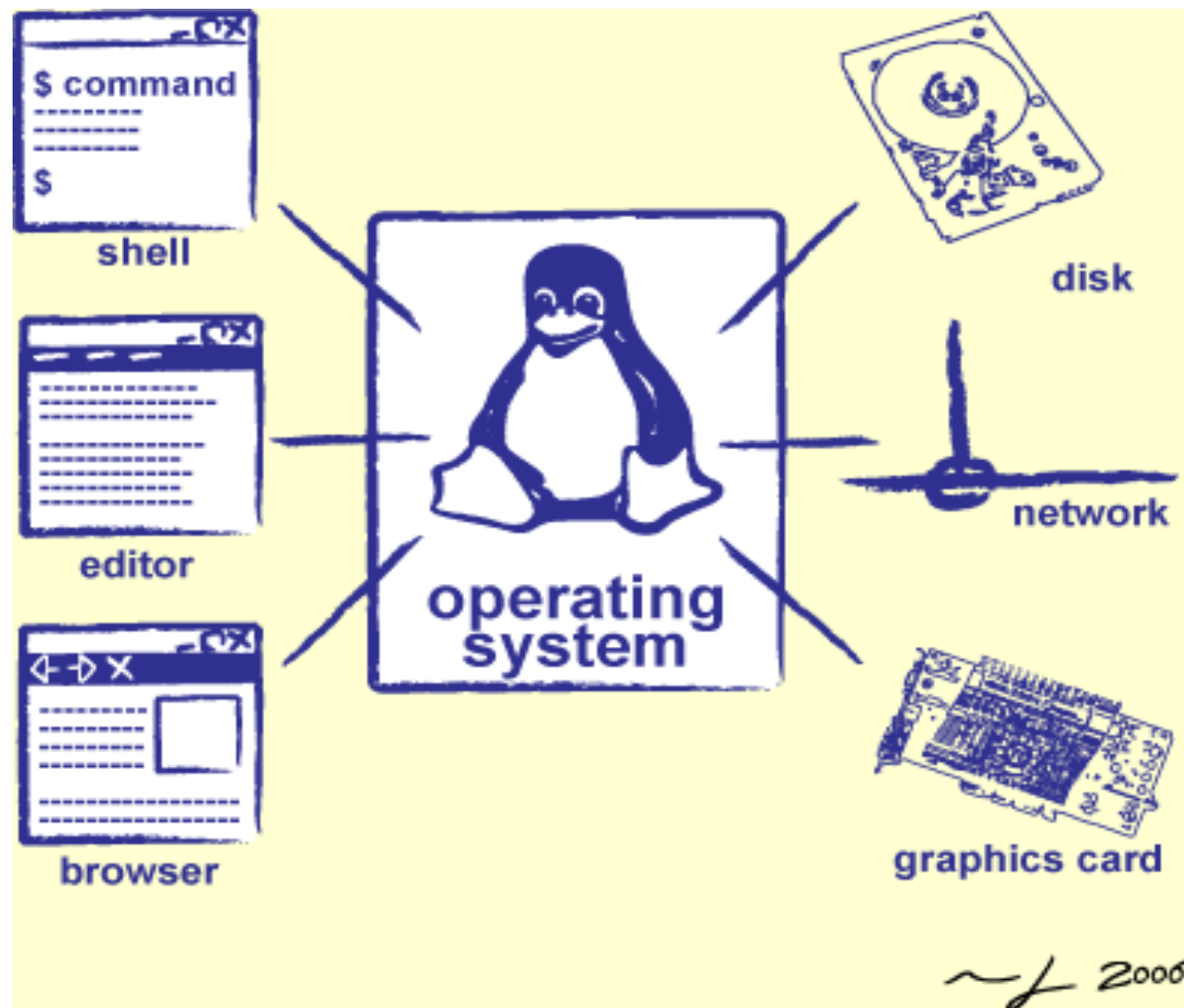


# Concurrent and Distributed Systems

## CSCU9V5

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4B86



# PART 1: processes

# Introduction

- About Processes
  - Execution of a program
  - Processes – an Operating System abstraction
  - Life cycle of a process
- Process Implementation
  - Parts of a process
  - Process Creation and Termination
- Cooperating Processes
  - Producer - Consumer Problem
  - Race Conditions



# About Processes

- Informally: a program in execution
- Computer execute predefined actions
  - Actions are specified by a program
  - Program is a *self-contained* entity
- Execution of a program requires resources
  - E.g., CPU
  - Running programs compete for CPU
  - Programs do not know when they get the CPU
- Actions of CPU and program logic are unrelated



# Program – Process Relationship

- A program is a static text that specifies a range of instructions which have to be executed by one or more processors (passive entity).
- A Process is an entity executing the range of instructions, which are specified by the program.

Its thread of execution is somewhere in the middle between the first and the last of the actions specified by the program (active entity).

Image Name	User Name	CPU	Mem Usage
Skype.exe	evan	00	20,940 K
ctfmon.exe	evan	00	1,096 K
zlclient.exe	evan	00	2,804 K
jusched.exe	evan	00	556 K
iTunesHelper.exe	evan	00	2,240 K
QTTask.exe	evan	00	1,272 K
apdproxy.exe	evan	00	1,768 K
LogiTray.exe	evan	00	2,024 K
LVComS.exe	evan	00	960 K
ashDisp.exe	evan	00	1,188 K
Directcd.exe	evan	00	1,612 K
hkcmd.exe	evan	00	1,060 K
AcroRd32.exe	evan	00	2,360 K
ieexplore.exe	evan	00	92,740 K
ieexplore.exe	evan	00	88,760 K
EXPLORER.EXE	evan	02	23,596 K
skypePM.exe	evan	00	5,488 K
NclUSBSrv.exe	SYSTEM	00	944 K
ServiceLayer.exe	SYSTEM	00	1,188 K
ALG.EXE	LOCAL SERVICE	00	924 K
ashServ.exe	SYSTEM	00	17,844 K
iPodService.exe	SYSTEM	00	1,360 K
aswUpdSv.exe	SYSTEM	00	192 K
ashWebSv.exe	SYSTEM	00	14,868 K
SVCHOST.EXE	LOCAL SERVICE	00	1,164 K
ashMaiSv.exe	SYSTEM	00	516 K
SVCHOST.EXE	NETWORK SERVICE	00	2,124 K
SVCHOST.EXE	SYSTEM	00	9,888 K
POWERPNT.EXE	evan	00	24,008 K
VSMON.EXE	SYSTEM	02	23,792 K
TextPad.exe	evan	00	936 K
DvzIncMsgr.exe	evan	00	1,376 K
AcroTray.exe	evan	00	468 K
SVCHOST.EXE	NETWORK SERVICE	00	2,084 K
WZQKPICK.EXE	evan	00	868 K
WDFMGR.EXE	LOCAL SERVICE	00	428 K
SVCHOST.EXE	SYSTEM	00	1,360 K
LSASS.EXE	SYSTEM	00	1,336 K
SERVICES.EXE	SYSTEM	00	1,824 K
WINLOGON.EXE	SYSTEM	00	452 K
CSRSS.EXE	SYSTEM	00	2,424 K
LowLight.exe	evan	00	584 K
SVCHOST.EXE	SYSTEM	00	2,296 K

☐ Show processes from all users

End Process

# Processes Lifecycle

- Number of processes  $\gg$  number of CPUs
- Some processes cannot run
- Processes waiting for resources
  - CPU ('ready processes')
  - Other resources ('blocked processes')
- Operating System provides for liveness
  - State transitions of a process

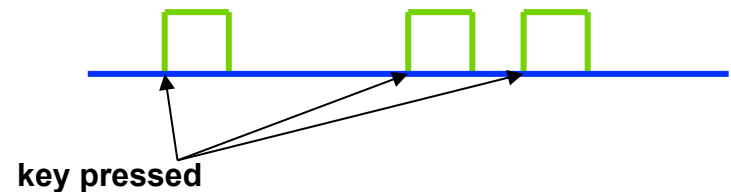




# Example: Multitasking

- The OS virtual machine can (in effect) run several programs at the "same" time.
- Efficient usage of CPU:

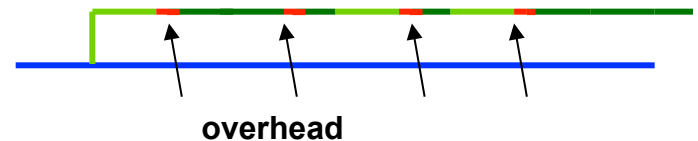
I/O bound process (editor):



CPU bound process (sci calc.):



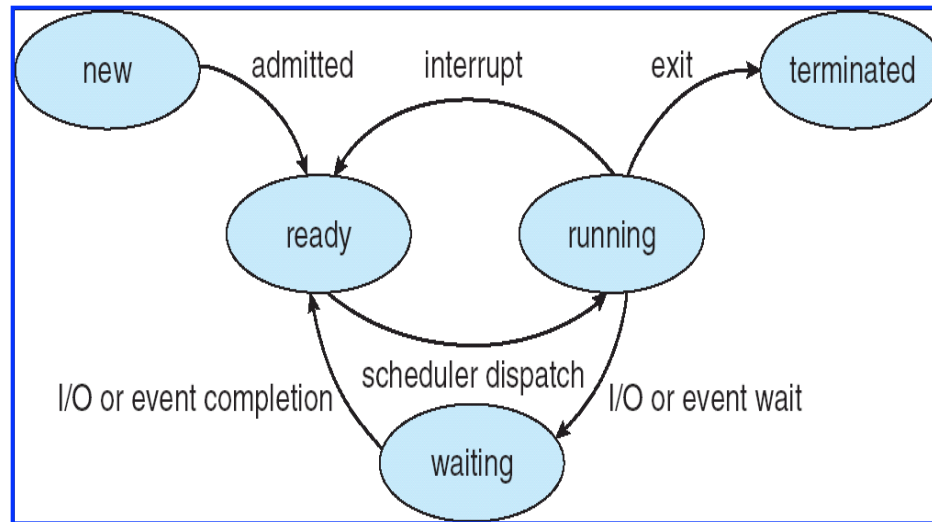
Efficient sharing of CPU:



# Process States

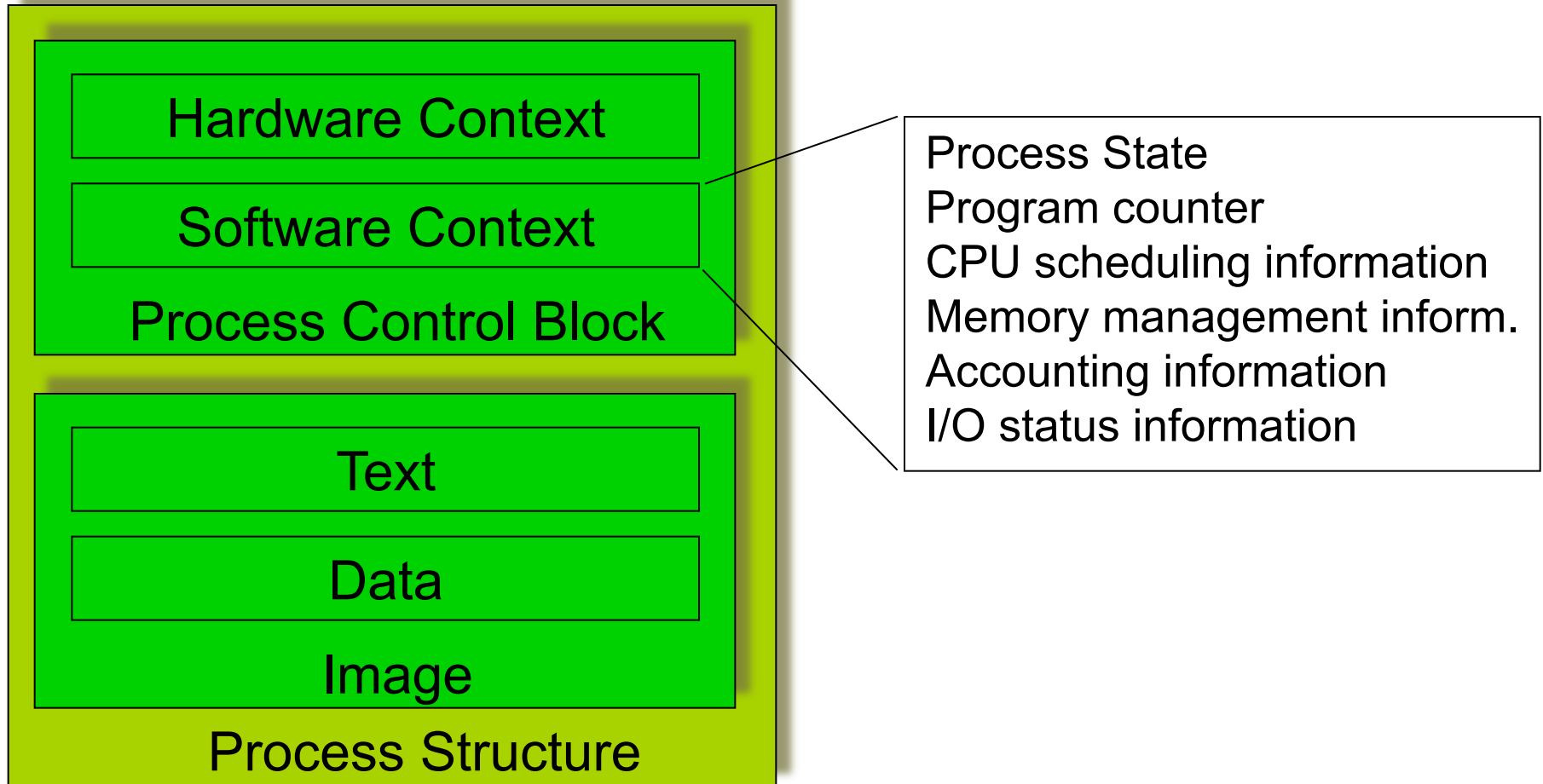
- As a process executes, it changes **state**
  - **initiated**: The process is being created.
  - **running**: Instructions are being executed.
  - **waiting**: The process is waiting for an event to occur.
  - **ready**: The process is waiting to be assigned.
  - **terminated**: The process has finished execution.

# Process Lifecycle

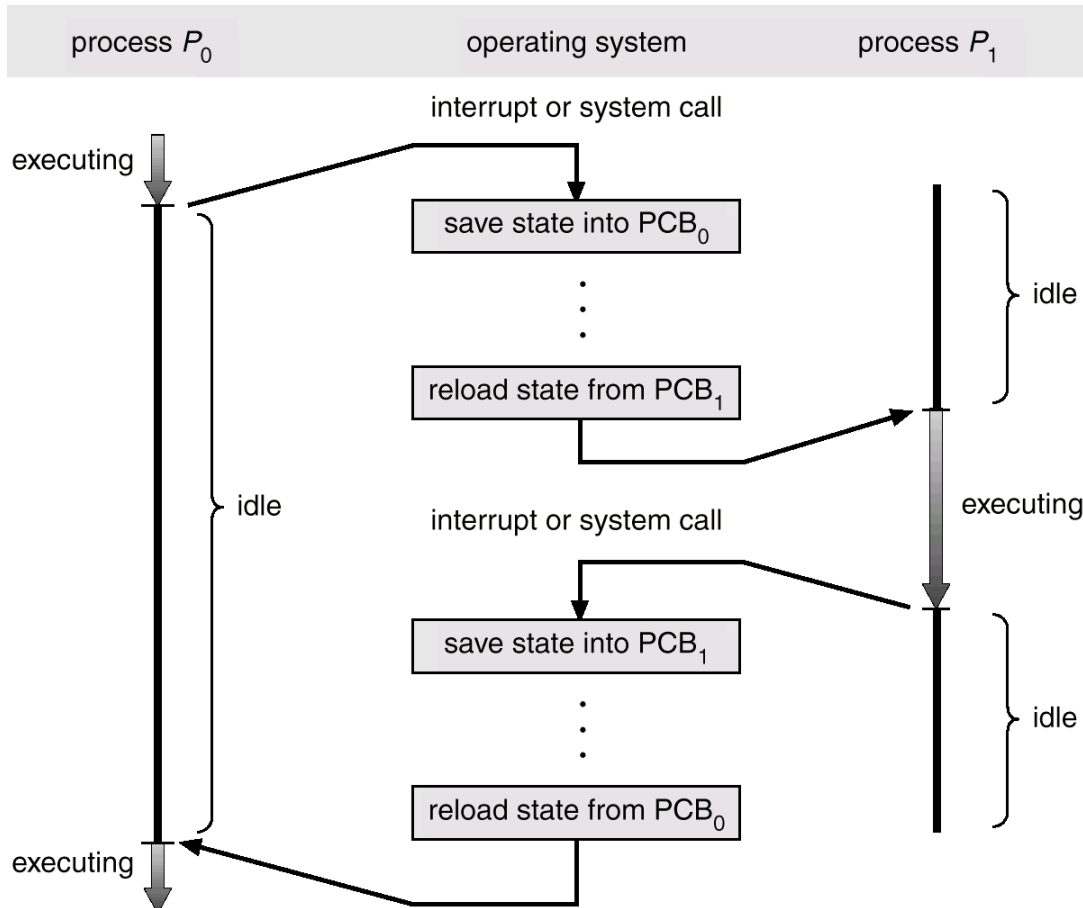


- **new** process are created (a program is launched) and put in the **ready** queue (list)
- then the use of the CPU is assigned to a ready process, which becomes **running**
- a running process can either **terminate**, suspend while **waiting** for some event to occur or be forced to **release** the CPU (interrupted), for instance because it is being running for too long or because the OS has to load a new page in memory, and hence be put back to the ready queue
- a waiting process can eventually become **ready again**, when the event it was waiting for occurs.

# Parts of a Process



# Context Switch



- When CPU switches to another process, the system must save the state of the old process and load the saved state for the new process.
- Context-switch time is overhead; the system does no useful work while switching.
- Time dependent on hardware support.



# Process Creation

- Parent process creates children processes, which, in turn create other processes, forming a tree of processes.
- Resource sharing
  - Parent and children share all resources.
  - Children share subset of parent's resources.
  - Parent and child share no resources.
- Execution
  - Parent and children execute concurrently.
  - Parent waits until children terminate.
- Address space
  - Child duplicate of parent.
  - Child has a program loaded into it.



# Example in c

```
#include <stdio.h>
```

```
void main(int argc, char *argv[])
```

```
{  
    int pid;
```

```
    /* fork another process */  
    pid = fork();
```

```
    if (pid < 0) { /* error occurred */  
        fprintf(stderr, "Fork  
                        Failed\n");
```

```
        exit(-1);
```

```
    }
```

```
    else if (pid == 0) { /* child process */  
        execlp("/bin/ls", "ls", NULL);
```

```
    }
```

```
    else { /* parent process */  
        /* parent waits for the child to  
                                                complete */
```

```
        wait(NULL);  
        printf("Child Complete\n");  
        exit(0);
```

```
    }
```

```
}
```



# Process Termination

- Process executes last statement and asks the operating system to delete it.
  - Output data from child to parent.
  - Process' resources are deallocated by operating system (open files, physical and virtual memory, I/O buffers).
- Parent may terminate execution of children processes.
  - Child has exceeded allocated resources.
  - Task assigned to child is no longer required.
  - Parent is exiting (e.g., OS may not allow child to continue if its parent terminates).





# Cooperating Processes

- **Independent** process cannot affect or be affected by the execution of another process.
- **Cooperating** process can affect or be affected by the execution of another process
- Advantages of process cooperation
  - Information` sharing (shared files)
  - Computation speed-up (split-up of a task into subtasks and run them in parallel; Note: number of processors!)
  - Modularity (divide a system into separate processes)
  - Convenience (a user has many tasks; e.g. printing, editing, compiling)



# Producer-Consumer Problem

- Paradigm for cooperating processes, **producer** produces information that is consumed by a **consumer**.
  - Examples: printer queue, keyboard buffer
  - Consumer and Producer processes need to be synchronised
  - Buffer may be provided by OS (IPC mechanism) or be explicitly coded by the programmer
  - *unbounded-buffer* places no practical limit on the buffer size
    - Producer can always produce elements
    - Consumer gets blocked when there are no elements
  - *bounded-buffer* assumes that there is a fixed buffer size.
    - Producer is blocked when the buffer is full
    - Consumer is blocked when the buffer is empty



# Race Conditions

- Concurrent access to shared data may result in data inconsistency (the bank example!).
- Maintaining data consistency requires mechanisms to ensure the orderly execution of cooperating processes.
- Example: Print Spooler

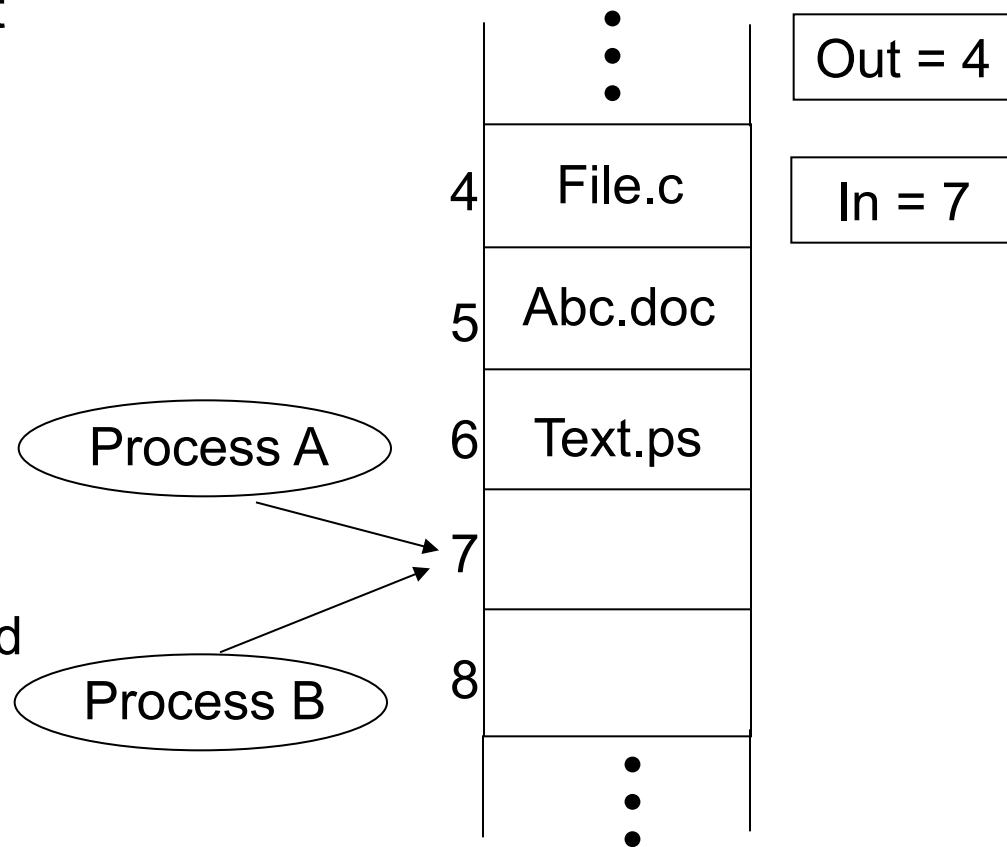
# Example of Race Condition

- Print Spooler
  - User processes put file names into a spooler directory
  - Printer daemon periodically checks for any new files in the spooler directory
    - If there are any – print the file and remove file name from spooler directory
  - Directory has infinite number of slots
  - Two globally available variables:
    - *out* points at the next file to be printed
    - *in* points at the next available slot
  - At a time: slots 0 – 3 are empty (files printed)  
slots 4 – 6 are full (files to be printed)



# Example cont.

- Almost simultaneously processes A and B want to print
- Process A reads *in* (value 7) and assigns it to local variable
- Process context switch occurs to process B
- Process B reads *in* (value 7) and stores a file in slot 7, updates *in* to 8
- Process A runs again, continuing from where it stopped
- Process A reads local variable and stores a file at slot 7
- Process B's file is erased!
- **RACE CONDITION!**



# Summary

- Process is an Operating System abstraction (program in execution)
- One program may result in many processes
- Processes may be interrupted while processing (process states)
- Processes cooperate (communicate)
- Race Conditions