



CP Learning

Websites for learning

👉 To find any sequence use this website:
<https://oeis.org/>

Short Tricks

▼ Count Odds Numbers Between range L-R

```
int countOdd(int L, int R){  
  
    int N = (R - L) / 2;  
  
    // if either R or L is odd  
    if (R % 2 != 0 || L % 2 != 0)  
        N += 1;  
  
    return N;  
}
```

▼ Find the Least Significant Bit of a Number OR the minimum value of power of 2 divide the N

```
int a = 16;  
cout << (a & (-a)) << endl;  
// output << 16  
  
int a = 60;  
cout << (a & -a) << endl;  
// 4  
  
// So, 4 is the minimum value Power of 2 divide the 60  
  
// We can use log2(); to find the bit number;
```

▼ Most greatest Number divisible by 'a' and less than 'n'

Identify the highest number divisible by a that is less than or equal to n

$$Result = \lfloor a/n \rfloor \times a$$

```
int n, a;
cout << (n/a) * a;

// n = 11, a = 5;
// 11 / 5 = 2
// 2 * 5 = 10
```

▼ log2 function

```
int s(int a) {
    int ans=0;
    while (a>1) {
        ans++;
        a=(a)/2;
    }
    return ans;
}

// ### Short Cut

int s(int a){
    return __lg(a);
}

// log2() in ceil

int s(int a) {
    int ans=0;
    while (a>1) {
        ans++;
        a=(a+1)/2;
    }
    return ans;
}

// ## Shortcut

int s(int a){
    return __lg(2 * a - 1);
}
```

▼ Bitmask Equation

👉 if $a \wedge b = x$ then $b = x \wedge a$
 if $a|b = x$ then, $b = x \& \sim a$
 if $a \& b = x$, then b can't be unique value,
 but, the maximum value of b can be,
 $b = x | \sim a$

▼ How to check segments that intersect

To determine when two line segments $[a,b]$ and $[c,d]$ intersect (on a 1D number line), you need to check whether the intervals **overlap**.

The condition for intersection is:

$$\max(a, c) \leq \min(b, d)$$

if there are many segments use multiset.

▼ How calculate nC2 in real time using Map

If we have N values, and we want to find $nC2$ using map,

```
map<int,int> mp;
int ans = 0;

for(int i = 0; i < n; i++){
    int x; cin >> x;
    ans += mp[x]++;
}
```

▼ Lambda Expression for DFS

```
function<void(int)> dfs = [&](int x){
    if(x == 0) return;
    cout << x << endl;
    dfs(x - 1);
};

dfs(10);

// function<return_type(arguments type)>
```

▼ Priority Queue Comparator

```
auto cmp = [](pair<int, int>& a, pair<int, int>& b) {
    return a.second > b.second;
};
```

```
priority_queue<
    pair<int, int>,
    vector<pair<int, int>>,
    decltype(cmp)
> pq(cmp);
```

▼ Count how many odd multiples of k are in [L, R]

```
// Count how many odd multiples of k are in [L, R]
ll countOddMultiples(ll L, ll R, ll k) {
    // First multiple of k ≥ L
    ll start = (L + k - 1) / k * k;
    // First odd multiple ≥ L
    if (start % 2 == 0) start += k;
    if (start > R) return 0;

    // Last multiple of k ≤ R
    ll end = R / k * k;
    if (end % 2 == 0) end -= k;
    if (end < L) return 0;

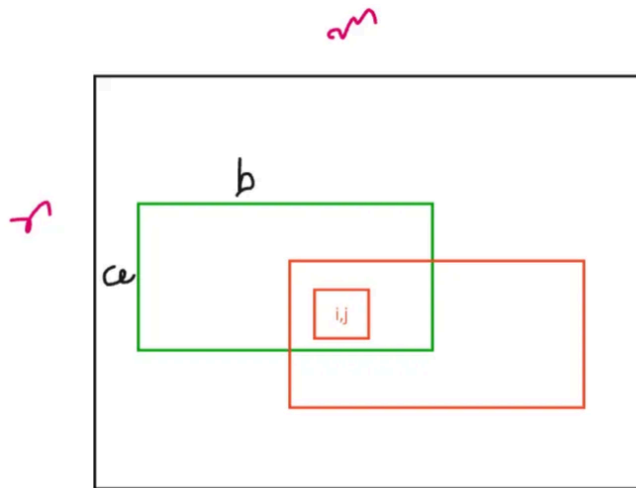
    return ((end - start) / (2 * k)) + 1;
}
```

▼ Unordered_map Pair declaration

```
struct pair_hash{
    size_t operator()(const pair<int, int>& p) const {
        return hash<int>()(p.first) ^ (hash<int>()(p.second) << 1);
    }
};
//unordered_map<pair<int,int>, int, pair_hash> m;
```

2D Array

▼ If we have $n \times m$ matrix and we have $a \times b$ rectangle, and if we put this rectangle to the matrix, so that it does not overflow, we have to find, how many rectangles will cover the every cells.



```

int n = 3;
int m = 3;
int a = 2, b = 2;
for(int i = 0; i < 3; i++){
    for(int j = 0; j < 3; j++){

        int lbx = max(0, i - b + 1);
        int ubx = min(i, n - b);
        int lby = max(0, j - a + 1);
        int uby = min(j, m - a);
        cout << (ubx - lbx + 1) * (uby - lby + 1) << " ";
    }
    cout << endl;
}

```

▼ If we have a rectangle of $N * M$, and we want know how many sub-rectangle are there of area of K

```

int countSubrectangles(int n, int m, int k) {
    int count = 0;
    for (int a = 1; a * a <= k; ++a) {
        if (k % a == 0) {
            int b = k / a;
            // (a, b) subrectangle
            if (a <= n && b <= m) {
                count += (n - a + 1) * (m - b + 1);
            }
            // (b, a) subrectangle, only if a != b
            if (a != b && b <= n && a <= m) {
                count += (n - b + 1) * (m - a + 1);
            }
        }
    }
}

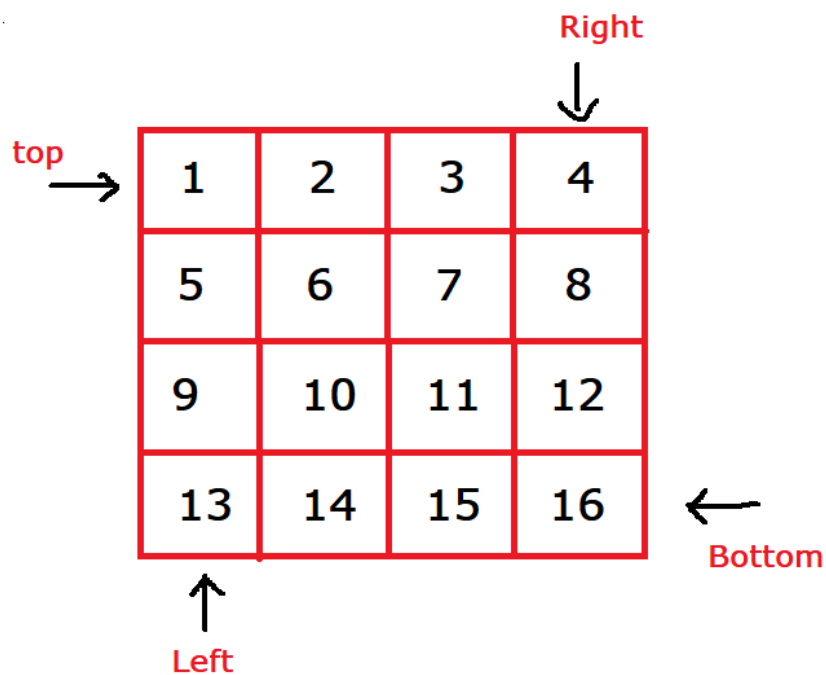
```

```
    return count;
}
```

PDF

attachment:b50936c0-b8fa-4f89-826d-f6cd642ddb5e:subrectangle_count.pdf

▼ Spiral Traversal of Matrix



```
#include <bits/stdc++.h>

using namespace std;

vector<int> printSpiral(vector<vector<int>> mat) {

    // Define ans array to store the result.
    vector<int> ans;

    int n = mat.size(); // no. of rows
    int m = mat[0].size(); // no. of columns

    // Initialize the pointers reqd for traversal.
    int top = 0, left = 0, bottom = n - 1, right = m - 1;
```

```

// Loop until all elements are not traversed.
while (top <= bottom && left <= right) {

    // For moving left to right
    for (int i = left; i <= right; i++)
        ans.push_back(mat[top][i]);

    top++;

    // For moving top to bottom.
    for (int i = top; i <= bottom; i++)
        ans.push_back(mat[i][right]);

    right--;

    // For moving right to left.
    if (top <= bottom) {
        for (int i = right; i >= left; i--)
            ans.push_back(mat[bottom][i]);

        bottom--;
    }

    // For moving bottom to top.
    if (left <= right) {
        for (int i = bottom; i >= top; i--)
            ans.push_back(mat[i][left]);

        left++;
    }
}
return ans;
}

int main() {

    //Matrix initialization.
    vector<vector<int>>> mat  {{1, 2, 3, 4},
                             {5, 6, 7, 8},
                             {9, 10, 11, 12},
                             {13, 14, 15, 16}};

    vector<int> ans = printSpiral(mat);

    for(int i = 0;i<ans.size();i++){

        cout<<ans[i]<<" ";
    }
}

```

```
cout<<endl;

return 0;
}
```

Math

▼ Sequence

Arithmetic Progression

▼ $a + ad + a2d + a3d + \dots$

$$a + a \cdot d + a \cdot 2d + a \cdot 3d + \dots$$

Nth Term (General Term):

$$a_n = a + (n - 1) \times d$$

where:

- a is the first term,
- d is the common difference, and
- n is the term number.

Sum of the First N Terms:

$$S_n = 2n \cdot (2a + (n - 1) \cdot d)$$

$$S_n = \frac{n}{2} \cdot (a + a_n)$$

- S_n is the sum of the first n terms, n
- a is the first term,
- a_n is the n th term, and
- d is the common difference.

Geometric Progression

▼ $a + rd + r^2d + \dots$

Nth Term (General Term):

$$a_n = a \cdot r^{n-1}$$

where:

- a is the first term,

- r is the common ratio, and
- n is the term number.

Sum of the First N Terms:

$$S_n = a \cdot \frac{r^n - 1}{r - 1} \quad ; r > 1$$

$$S_n = a \cdot \frac{1 - r^n}{1 - r} \quad ; r < 1$$

▼ Pythagorean Triplets

Any Pythagorean Triplets can be generated by these equation, where $m > n$.

$$a = m^2 - n^2, b = 2mn, c = m^2 + n^2, (m > n)$$

Number Theory

▼ Bezout's Identity



For any two integers a and b , **not both zero**, there exist integers x and y such that:
 $\gcd(a,b) = ax + by$

▼ Modular Arithmetic



You can only change a number X into Y if and only if either:

1. $x \equiv y \pmod{k}$ or
2. $x + y \equiv 0 \pmod{k}$

See solution C: <https://codeforces.com/blog/entry/145439>

▼ McNugget Theorem



If you have two **relatively prime** ($\gcd(m,n) = 1$) positive integers m and n , the largest number you **cannot** express as a nonnegative combination of m and n is:

$$mn - m - n$$

For **two** positive coprime pack sizes m and n , the number of positive integers

not representable as

$am + bn$ with $a, b \geq 0$ is

$$(m-1)(n-1) / 2$$

▼ Catalan Number

<https://cp-algorithms.com/combinatorics/catalan-numbers.html>

<http://www.geometer.org/mathcircles/catalan.pdf>

```
const int MOD = (int)1e9 + 7;
const int MAX = (int)1e5;
int catalan[MAX];
void init() {
    catalan[0] = catalan[1] = 1;
    for (int i=2; i<=n; i++) {
        catalan[i] = 0;
        for (int j=0; j < i; j++) {
            catalan[i] += (catalan[j] * catalan[i-j-1]) % MOD;
            if (catalan[i] >= MOD) {
                catalan[i] -= MOD;
            }
        }
    }
}
```

▼ GCD Sum function

link: <https://forthright48.com/gcd-sum-function/#more-502>

$$g(n) = gcd(1, n) + gcd(2, n) + gcd(3, n) + \dots + gcd(n, n) = \sum_{i=1}^n gcd(i, n)$$

$$g(6) = gcd(1, 6) + gcd(2, 6) + gcd(3, 6) + gcd(4, 6) + gcd(5, 6) + gcd(6, 6) \quad (1)$$

$$= 1 + 2 + 3 + 2 + 1 + 6 \quad (2)$$

$$= 15 \quad (3)$$

In short, there is a direct formula for calculating the value of **g(n)**.

If the prime factorization of n is $p_1^{a_1} \times p_2^{a_2} \times \dots \times p_k^{a_k}$

$$g(n) = \prod_{i=1}^k (a_i + 1) p_i^{a_i - 1}$$

```
int const MAXN = (int)1e7 + 5;

vector<int> spf(MAXN + 1, 1);

//from
//https://www.geeksforgeeks.org/dsa/prime-factorization-using-sieve-olog-n-multiple-queries/
```

```

// Time Complexity : O(nloglogn)

void sieve()
{

    spf[0] = 0;
    for (int i = 2; i <= MAXN; i++) {
        if (spf[i] == 1) {
            for (int j = i; j <= MAXN; j += i) {
                if (spf[j] == 1)
                    spf[j] = i;
            }
        }
    }
}

//log(n)
map<int,int> getFactorization(int x)
{
    map<int,int> mp;
    while (x != 1) {
        mp[spf[x]]++;
        x = x / spf[x];
    }
    return mp;
}

int _pow(int base, int power){
    int res = 1;
    while(power){
        if(power&1){
            res = res * base;
            power--;
        }else{
            base = base * base;
            power = power / 2;
        }
    }
    return res;
}

int gcdSumFunction(int n){
    map<int,int> pf = getFactorization(n);

    int ans = 1;

    for(auto i: pf){
        ans *= ((i.second + 1) * _pow(i.first, i.second)) - ((i.second) * _pow(i.first, i.second - 1));
    }
}

```

```
    return ans;
}
```

String

- ▼ Find Lexicographically largest subsequence

```
vector<int> ans;
for (ll i = 0; i < n; i++) {
    while (!ans.empty() && (str[i] > str[ans.back()])) ans.pop_back();
    ans.push_back(i);
}
```

Miscellaneous

- ▼ MEX (minimal excluded) of a sequence

<https://cp-algorithms.com/sequences/mex.html>

- ▼ Generate All Palindromic Numbers upto N

URL: <https://www.geeksforgeeks.org/generate-palindromic-numbers-less-n/>

```
vector<int> pnums;

int createPalindrome(int input, int b, bool isOdd)
{
    int n = input;
    int palin = input;

    if (isOdd)
        n /= b;

    while (n > 0)
    {
        palin = palin * b + (n % b);
        n /= b;
    }
    return palin;
}

void generatePalindromes(int n)
{
    int number;
    for (int j = 0; j < 2; j++)
    {
        int i = 1;
        while ((number = createPalindrome(i, 10, j % 2)) < n)
```

```

    {
        pnums.push_back(number);
        i++;
    }
}

// Have to sort, bcz, generates numbers ar not sorted;

sort(all(pnums));
}

```

Bitmask

▼ Sum of XOR of all possible subsets of a given array

if the length of the array is N, and the OR of each element of array is X,
the sum = $X * (2^N - 1)$

Example problem: <https://codeforces.com/problemset/problem/1614/C>

```

int subsetXORSum(vector<int> &arr) {
    int n = arr.size();
    int bits = 0;

    // Finding bitwise OR of all elements
    for (int i=0; i < n; ++i)
        bits |= arr[i];

    int ans = bits * (1<<N-1);

    return ans;
}

```

▼ CF Blog

<https://codeforces.com/blog/entry/72437>

▼ BitMask Equations



Some properties of bitwise operations:

- $a|b = a \oplus b + a \& b$
- $a \oplus (a \& b) = (a|b) \oplus b$
- $b \oplus (a \& b) = (a|b) \oplus a$
- $(a \& b) \oplus (a|b) = a \oplus b$

Addition:

- $a + b = a|b + a \& b$
- $a + b = a \oplus b + 2(a \& b)$

Subtraction:

- $a - b = (a \oplus (a \& b)) - ((a|b) \oplus a)$
- $a - b = ((a|b) \oplus b) - ((a|b) \oplus a)$
- $a - b = (a \oplus (a \& b)) - (b \oplus (a \& b))$
- $a - b = ((a|b) \oplus b) - (b \oplus (a \& b))$

Combinatorics

▼ Pigeonhole Principle



if i have, N elements, and k box. and if i put, element to boxex. them, there will be minimum, $\text{ceil}(N/k)$ element in a single box.

i.e $N = mk + 1$

m = how many object per box (minimum)



Theorem:

Every sequence of $n^2 + 1$ distinct real numbers contains a subsequence of length $n + 1$ that is either strictly increasing or strictly decreasing.

$n = 3; 3^2 + 1 = 10;$

1, 10, 100, 3, 4, 9, 11, 20, 23, 33, 200

there is always, subsequence of length $3 + 1$ which is strictly increasing or strictly decreasing.

▼ Intersection points of regular N' gons

Intersection points of regular N' gons: (N must be odd)

$$(nC4) = n(n-1)(n-2)(n-3)/24$$

Number of diagonals:

$$nC_2 - n = n(n-3)/2$$

Tricks

▼ Logarithmic Subarray Aggregator (LSA Trick)

https://youkn0wwho.academy/topic-list/logarithmic_bruteforce

Problem: <https://codeforces.com/gym/106057/problem/C>

Sol: [https://github.com/MHKhanCou/CoU-IUPC-2025-Regional/blob/main/problemC_Prime_Dominion_\(lca\).cpp](https://github.com/MHKhanCou/CoU-IUPC-2025-Regional/blob/main/problemC_Prime_Dominion_(lca).cpp)

editorial:

https://codeforces.com/gym/106057/attachments/download/33229/CoU_IUPC_2025_Editorial.pdf

Graph

▼ Topological sorting

Topological sorting:

DFS: <https://www.youtube.com/watch?v=5lZ0iJMrUMk>

BFS: <https://www.youtube.com/watch?v=73sneFXuTEg>

<https://codeforces.com/gym/105981/problem/E>

<https://codeforces.com/contest/2143/problem/C>

```
// Including necessary header file
#include <bits/stdc++.h>
using namespace std;

// We mainly take input graph as a set of edges. This function is
// mainly a utility function to convert the edges to an adjacency
// list
vector<vector<int>> > constructadj(int V,vector<vector<int>> &edges){

    // Graph represented as an adjacency list
    vector<vector<int> > adj(V);

    // Constructing adjacency list
    for (auto i : edges) {
        adj[i[0]].push_back(i[1]);
    }

    return adj;
}

// Function to return list containing vertices in
```

```

// Topological order.
vector<int> topologicalSort(int V, vector<vector<int> >& edges)
{
    vector<vector<int>> adj = constructadj(V,edges);

    // Vector to store indegree of each vertex
    vector<int> indegree(V);
    for (int i = 0; i < V; i++) {
        for (auto it : adj[i]) {
            indegree[it]++;
        }
    }
    // Queue to store vertices with indegree 0
    queue<int> q;
    for (int i = 0; i < V; i++) {
        if (indegree[i] == 0) {
            q.push(i);
        }
    }
    vector<int> result;
    while (!q.empty()) {
        int node = q.front();
        q.pop();
        result.push_back(node);

        // Decrease indegree of adjacent vertices as the
        // current node is in topological order
        for (auto it : adj[node]) {
            indegree[it]--;

            // If indegree becomes 0, push it to the queue
            if (indegree[it] == 0)
                q.push(it);
        }
    }

    // Check for cycle
    if (result.size() != V) {
        cout << "Graph contains cycle!" << endl;
        return {};
    }

    return result;
}

int main()
{
    // Number of nodes

```



```
int V = 6;

// Edges
vector<vector<int> > edges
    = {{0, 1}, {1, 2}, {2, 3},
        {4, 5}, {5, 1}, {5, 2}};

vector<int> result = topologicalSort(V, edges);

// Displaying result
for (auto i : result) {
    cout << i << " ";
}

return 0;
}
```