

Higher order unification for free!

Reusing the meta-language unification for the object language

Davide Fissore

davide.fissore@inria.fr

Université Côte d'Azur, Inria
France

Enrico Tassi

enrico.tassi@inria.fr

Université Côte d'Azur, Inria
France

ABSTRACT

Specifying and implementing a proof system from scratch requires significant effort. Logical Frameworks and Higher Order Logic Programming Languages provide dedicated, high-level Meta Languages (ML) to facilitate this task in two key ways: 1) variable binding and substitution are simplified when ML binders represent object logic ones; 2) proof construction, and even proof search, is greatly simplified by leveraging the unification procedure provided by the ML. Notable examples of ML are Elf [13], Twelf [15], λ Prolog [9] and Isabelle [21] which have been utilized to implement various formal systems such as First Order Logic [3], Set Theory [12], Higher Order Logic [11], and even the Calculus of Constructions [2].

The object logic we are interested in is Coq's [19] Calculus of Inductive Constructions (CIC), for which we aim to implement a unification procedure \approx_o using the ML Elpi [1], a dialect of λ Prolog. Elpi's equational theory comprises $\eta\beta$ equivalence and comes equipped with a higher order unification procedure \approx_λ restricted to the pattern fragment [8]. We want \approx_o to be as powerful as \approx_λ but on the object logic CIC. Elpi also comes with an encoding for CIC that works well for meta-programming [18, 17, 6, 4]. Unfortunately this encoding, which we refer to as \mathcal{F}_o , "underuses" \approx_λ by restricting it to first-order unification problems only. To address this issue, we propose a better-behaved encoding, \mathcal{H}_o , demonstrate how to map unification problems in \mathcal{F}_o to related problems in \mathcal{H}_o , and illustrate how to map back the unifiers found by \approx_λ , effectively implementing \approx_o on top of \approx_λ for the encoding \mathcal{F}_o .

KEYWORDS

Logic Programming, Meta-Programming, Higher-Order Unification, Proof Automation

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1 INTRODUCTION

Meta languages such as Elf [13], Twelf [15], λ Prolog [9] and Isabelle [21] have been utilized to specify various logics [3, 11, 12,

2]. The use of these meta languages facilitated this task in two key ways. The first and most well know one is that variable binding and substitution come for free. The second one is that these meta languages come equipped with some form of unification, a cornerstone in proof construction and proof search.

The object logic we are interested in is Coq's [19] Calculus of Inductive Constructions (CIC) and we want to implement a form of proof search known as type-class [20, 16] resolution. In particular we want to leverage the Elpi [18] meta programming language, a dialect of λ Prolog already used to extend Coq in various ways [18, 17, 6, 4]. Type-class solvers are unification based proof search procedures reminiscent of Prolog that back-chain lemmas taken from a designated database of "type class instances", hence we can expect that Elpi is a good fit for implementing such as form of automation. In this paper we focus on one aspect of this work, namely *how to reuse the higher order unification procedure of the meta language in order to implement a type-class solver for the object language*. As it turns out, re-using the unification of the meta language is not a trivial task.

We take as an example the Decision and Finite type classes from the Stdpp [7] library. The class Decision identifies predicates equipped with a decision procedure, while Finite the types whose inhabitants can be enumerated in a (finite) list. The following three designated Instances state that: 1) the type `fin n`, of natural numbers smaller than `n`, is finite; 2) the predicate `nfact n nf`, linking a natural number `n` to the number of its prime factors `nf`, is decidable; 3) the universal closure of a predicate has a decision procedure if its domain is finite and if the predicate is decidable.

```
Instance fin_fin:  $\forall n$ , Finite (fin n). (* r1 *)
Instance nfact_dec:  $\forall n$  nf, Decision (nfact n nf). (* r2 *)
Instance forall_dec:  $\forall A$  P, Finite A  $\rightarrow$  (* r3 *)
 $\forall x:A$ , Decision (P x)  $\rightarrow$  Decision ( $\forall x:A$ , P x).
```

Given this database a type-class solver is expected to prove the following statement automatically:

```
Decision ( $\forall x$ : fin 7, nfact x 3) (* g *)
```

The proof found by the solver back-chains on rule 3 (the only rule about the \forall quantifier), and then solves the premises with rules rules 1 and 2 respectively. Note that rule 3 features a second order parameter `P` that stands for a function of type `A \rightarrow Prop` (a predicate over `A`). The solver has to infer a value for `P` by unifying the conclusion of rule 3 with the goal, and in particular by solving the unification problem `P x = nfact x 3`. This higher order problem falls in the so called pattern-fragment \mathcal{L}_λ [8] and hence admits a unique solution `P = λx . nfact x 3`.

In order to implement such a search in Elpi we shall describe the encoding of CIC terms and then the encoding of rules. Elpi comes

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with an Higher Order Abstract Syntax [14] datatype of CIC terms, called `tm`, that features (among others) the following constructors:

```

type lam  tm -> (tm -> tm) -> tm.    % lambda abstraction
type app  list tm -> tm.             % n-ary application
type all  tm -> (tm -> tm) -> tm.    % forall quantifier
type con  string -> tm.              % constants

```

Following λ Prolog [9]’s standard syntax, the meta level binding of a variable x in an expression e is written $\llbracket x \rrbracket e$, and square brackets denote a list of terms separated by comma. For example the term $\llbracket \forall y:t, \text{nfact } y \ 3 \rrbracket$ is encoded as follows:

```
all (con "t") y\ app [con "nfact", y, con "3"]
```

We now illustrate the encoding of the three instances above as higher-order logic-programming rules: capital letters denote rule parameters; `:-` separates the rule’s head from the premises and `pi w\` introduces a fresh nominal constant w for the premise p .

```
finite (app [con "fin", N]).           (r1)
```

```
decision (app [con "nfact", N, NF]).   (r2)
```

```
decision (all A x\ app [P, x]) :- finite A, (r3)
```

```
pi w\ decision (app [P, w]).
```

Unfortunately this intuitive encoding of rule (r3) does not work, since it uses the predicate P as a first order term: for the meta language its type is `tm`. If we try to back-chain the rule (r3) on the encoding of the goal g given below

```
decision (all (app [con "fin", con "7"]) y\ (g)
  app [con "nfact", y, con "3"]].
```

we obtain an unsolvable unification problem (p): the two lists of terms have different lengths!

```
app [con "nfact", y, con "3"] = app [P, y] (p)
```

In this paper we study a more sophisticated encoding of rules that, on a first approximation, would shape (r3) as follows:

```
decision (all A x\ Pm x) :- link Pm P A, finite A, (r3')
pi x\ decision (app [P, x]).
```

Since Pm is an higher-order unification variable of type `tm -> tm`, with x in its scope, the unification problem (p') admits one solution:

```
app [con "nfact", y, con "3"] = Pm y (p')
Pm = x\ app [con "nfact", x, con "3"] (rho)
```

Once the head of rule (r3') unifies with the goal (g) the premise $\llbracket \text{link } Pm \ A \ P \rrbracket$ brings the assignment (ρ) back to the domain `tm` of Coq terms:

```
P = lam A a\ app [con "nfact", a, con "3"]
```

This simple example is sufficient to show that the encoding we seek is not trivial and does not only concern the head of rules, but the entire sequence of unification problems that constitute the execution of a logic program. In fact the solution for P above generates a (Coq) β -redex in the second premise (the predicate under the `pi w\`). We show below the premise before and after the instantiation of P :

```
decision (app [P, w])
decision (app [lam A (a\ app [con "nfact", a, con "3"]) , w])
```

In turn this redex prevents the rule (r2) to backchain properly since the following unification problem has no solution:

```
app [lam A (a\ app [con "nfact", a, con "3"]) , x] =
app [con "nfact" , N, NF]
```

The root cause of the problems we sketched in this example is that the unification procedure \approx_λ of the meta language is not aware of the equational theory of the object logic, even if both theories include $\eta\beta$ -conversion and admit most general unifiers for unification problems in the pattern-fragment.

Contributions. In this paper we identify a minimal language \mathcal{F}_0 in which the problems sketched in the introduction can be formally described. We detail an encoding of a logic program on \mathcal{F}_0 to a strongly related logic program in \mathcal{H}_0 (the language of the meta language) and we show that the unification procedure of the meta language \approx_λ can be effectively used to simulate a unification procedure \approx_o for the object language that features $\eta\beta$ -conversion in the pattern-fragment.

section 2 formally states the problem and gives the intuition behind our solution. section 3 discusses alternative term encodings and related works. section 4 introduces the languages \mathcal{F}_0 and \mathcal{H}_0 , section 5 describes a basic simulation of higher order logic programs. sections 6 and 7 completes its equational theory with support for η -conversion. section 8 deals with the practical necessity of “tolerating” terms outside of the pattern-fragment and discusses how heuristic can be applied. Finally section 9 discusses the implementation in Elpi.

The λ Prolog code discussed in the paper can be accessed at the URL: <https://github.com/FissoreD/paper-ho>.

2 PROBLEM STATEMENT AND SOLUTION

The equational theory of Coq’s Dependent Type Theory is very rich. In addition to the usual $\eta\beta$ -equivalence for functions, terms (hence types) are compared up to proposition unfolding and fix-point unrolling. Still, for efficiency and predictability reasons, most form of automatic proof search employ a unification procedure that captures a simpler one, just $\eta\beta$, and that solves higher-order problems restricted to the pattern fragment \mathcal{L}_λ . We call this unification procedure \approx_o .

The equational theory of the meta language Elpi is strikingly similar, since it it comprises $\eta\beta$ (for the meta language functions), and the unification procedure \approx_λ solves problems in \mathcal{L}_λ as well.

In spite of the similarity the link between \approx_λ and \approx_o is not trivial, since the abstraction and application term constructors the two unification procedures deal with are different. For example:

$x\ f\ x$	$\approx_\lambda\ f$	$P\ x$	$\approx_\lambda\ x$
$\text{lam } A\ x\ \text{app}[con "f", x]$	$\approx_o\ con "f"$	$\text{app}[P, x]$	$\approx_o\ x$
$\text{lam } A\ x\ \text{app}[con "f", x]$	$\neq_\lambda\ con "f"$	$\text{app}[P, x]$	$\neq_\lambda\ x$

To state precisely the problem we solve we need a \mathcal{F}_0 representation of CIC terms and a \mathcal{H}_0 one. We extend the equality over ground terms to open ones with the reflexivity of unification variables (that is unification variable X is equal to itself, not to another one such as Y). We call \approx_o that equality in \mathcal{F}_0 and \approx_λ its counterpart in \mathcal{H}_0 . We call \approx_o the unification procedure we want to implement and \approx_λ the one provided by the meta language. We assume that the unification of our meta language is correct and complete in \mathcal{L}_λ .

We write $t_1 \approx_\lambda t_2 \mapsto \sigma$ when t_1 and t_2 unify with substitution σ ; we write σt for the application of the substitution to t , and $\sigma X = \{\sigma t \mid t \in X\}$ when X is a set; we write $\sigma \subseteq \sigma'$ when σ is more general than σ' and we use set union to denote the concatenation of two substitutions. We shall use ρ for \mathcal{F}_0 substitutions, and σ for the \mathcal{H}_0 ones.

We illustrate a compilation $\langle s \rangle \mapsto (t, m, l)$ that maps a term s in \mathcal{F}_0 to a term t in \mathcal{H}_0 and produces a variable mapping m and list of links l . The variable map connects unification variables in \mathcal{H}_0 with variables in \mathcal{F}_0 and is used to “decompile” the assignment, $\langle \sigma, m, l \rangle^{-1} \mapsto \rho$. Links are an accessory piece of information whose description is deferred to section 2.1.

We represent a logic program *run* in \mathcal{F}_0 as a list *steps* p of length N . Each step p represents a unification problem between two terms \mathbb{P}_{p_l} and \mathbb{P}_{p_r} , taken from the set of all terms \mathbb{P} . The composition of these steps starting from the empty substitution ρ_0 produces the final substitution ρ_N .¹

$$\begin{aligned} \text{fstep}(\mathbb{P}, p, \rho) &\mapsto \rho'' \stackrel{\text{def}}{=} \rho \mathbb{P}_{p_l} \approx_o \rho \mathbb{P}_{p_r} \mapsto \rho' \wedge \rho'' = \rho \cup \rho' \\ \text{frun}(\mathbb{P}, N) &\mapsto \rho_N \stackrel{\text{def}}{=} \bigwedge_{p=1}^N \text{fstep}(\mathbb{P}, p, \rho_{p-1}) \mapsto \rho_p \end{aligned}$$

We simulate each run in \mathcal{F}_0 with a run in \mathcal{H}_0 as follows.

$$\begin{aligned} \text{hstep}(\mathbb{T}, p, \sigma, \mathbb{L}) &\mapsto (\sigma'', \mathbb{L}') \stackrel{\text{def}}{=} \\ &\sigma \mathbb{T}_{p_l} \approx_\lambda \sigma \mathbb{T}_{p_r} \mapsto \sigma' \wedge \text{progress}(\mathbb{L}, \sigma \cup \sigma') \mapsto (\mathbb{L}', \sigma'') \\ \text{hrun}(\mathbb{P}, N) &\mapsto \rho_N \stackrel{\text{def}}{=} \\ &\mathbb{T} \times \mathbb{M} \times \mathbb{L}_0 = \{(t_j, m_j, l_j) \mid s_j \in \mathbb{P}, \langle s_j \rangle \mapsto (t_j, m_j, l_j)\} \\ &\bigwedge_{p=1}^N \text{hstep}(\mathbb{T}, p, \sigma_{p-1}, \mathbb{L}_{p-1}) \mapsto (\sigma_p, \mathbb{L}_p) \\ &\langle \sigma_N, \mathbb{M}, \mathbb{L}_N \rangle^{-1} \mapsto \rho_N \end{aligned}$$

Here *hstep* is made of two sub-steps: a call to the meta language unification and a check for progress on the set of links. We claim:

PROPOSITION 2.1 (SIMULATION). $\forall \mathbb{P}, \forall N$, if $\mathbb{P} \subseteq \mathcal{L}_\lambda$

$$\text{frun}(\mathbb{P}, N) \mapsto \rho_N \Leftrightarrow \text{hrun}(\mathbb{P}, N) \mapsto \rho_N$$

That is, the two executions give the same result. Moreover:

PROPOSITION 2.2 (SIMULATION FIDELITY). *In the context of *hrun*, if $\mathbb{P} \subseteq \mathcal{L}_\lambda$ we have that $\forall p \in 1 \dots N$,*

$$\text{fstep}(\mathbb{P}, p, \rho_{p-1}) \mapsto \rho_p \Leftrightarrow \text{hstep}(\mathbb{T}, p, \sigma_{p-1}, \mathbb{L}_{p-1}) \mapsto (\sigma_p, \mathbb{L}_p)$$

In particular this property guarantees that a *failure* in the \mathcal{F}_0 run is matched by a failure in \mathcal{H}_0 at the same step. We consider this property very important from a practical point of view since it guarantees that the execution traces are strongly related and in turn this enables a user to debug a logic program in \mathcal{F}_0 by looking at its execution trace in \mathcal{H}_0 .

We can define $s_1 \approx_o s_2$ as follows:

$$\begin{aligned} s_1 \approx_o s_2 &\mapsto \rho \stackrel{\text{def}}{=} \\ \langle s_1 \rangle &\mapsto (t_1, m_1, l_1) \wedge \langle s_2 \rangle \mapsto (t_2, m_2, l_2) \\ t_1 &\approx_\lambda t_2 \mapsto \sigma' \wedge \text{progress}(\{l_1, l_2\}, \sigma') \mapsto (L, \sigma'') \wedge \\ \langle \sigma'', \{m_1, m_2\}, L \rangle^{-1} &\mapsto \rho \end{aligned}$$

¹If the same rule is used multiple time in a run we just consider as many copies as needed of the terms composing the rules, with fresh unification variables each time

PROPOSITION 2.3 (PROPERTIES OF \approx_o IN \mathcal{L}_λ).

$$s_i \in \mathcal{L}_\lambda \Rightarrow s_1 \approx_o s_2 \mapsto \rho \Rightarrow \rho s_1 =_o \rho s_2 \quad (1)$$

$$s_i \in \mathcal{L}_\lambda \Rightarrow \rho s_1 =_o \rho s_2 \Rightarrow \exists \rho', s_1 \approx_o s_2 \mapsto \rho' \wedge \rho' \subseteq \rho \quad (2)$$

In addition to correctness and completeness in \mathcal{L}_λ , we claim that *hrun* handles terms outside \mathcal{L}_λ in the following sense:

PROPOSITION 2.4 (PROPERTIES OF *hrun* OUTSIDE \mathcal{L}_λ).

$$\exists \rho, \rho s_1 =_o \rho s_2 \Rightarrow \langle s_i \rangle \mapsto (t_i, m_i, l_i) \Rightarrow t_1 \approx_\lambda t_2 \mapsto \sigma \quad (3)$$

$$\sigma_{p-1} \mathbb{T}_p \in \mathcal{L}_\lambda \Rightarrow \text{hstep}(\mathbb{T}, p, \sigma_{p-1}, \mathbb{L}_{p-1}) \mapsto (\sigma_p, \mathbb{L}_p) \mapsto \sigma \quad (4)$$

Property 3 states that two terms for which there is a unifier (given by an oracle, not necessarily a most general one), then the compiler generates two terms that unify in \mathcal{H}_0 . In other words the compiler moves out of the way all problematic terms.

Property 4 says that if the two terms involved in a step re-enter \mathcal{L}_λ , then *hstep* succeeds. This is a typical example in which the order of the unification problems in a logic program run do matter. The simplest example is the sequence $F = \lambda x.a$ and $F a = a$: the second problem is not in \mathcal{L}_λ and has two unifiers, namely $F = \lambda x.x$ and $F = \lambda x.a$. The first problem picks one of the two, making the second problem re-enter \mathcal{L}_λ .

2.1 The intuition in a nutshell

A term s is compiled in a term t where every “problematic” sub term p is replaced by a fresh unification variable h and an accessory *link* that represent a suspended unification problem $h \approx_\lambda p$. As a result \approx_λ is “well behaved” on t , that is it does not contradict $=_o$ as it would otherwise do on “problematic” terms. We now define “problematic” and “well behaved” more formally.

Definition 2.5 ($\diamond \eta$). $\diamond \eta = \{t \mid \exists \rho, \rho t \text{ is an eta expansion}\}$

An example of term t in $\diamond \eta$ is $\lambda x.\lambda y.F y x$ since the substitution $\rho = \{F \mapsto \lambda a.\lambda b.fba\}$ makes $\rho t = \lambda x.\lambda y.fxy$ that is the eta long form of f . This term is problematic since its rigid part, the λ -abstractions, cannot justify a unification failure against a constant or an application.

Definition 2.6 ($\diamond \beta_0$). $\diamond \beta_0 = \{X x_1 \dots x_n \mid x_1 \dots x_n \text{ are distinct names}\}$

An example of term t in $\diamond \beta_0$ is the application $F x$. This term is problematic since the application node of its syntax tree cannot be used to justify a unification failure, i.e. by properly instantiating F the term head constructor may become a λ , or a constant or stay an application.

Definition 2.7 ($\overline{\mathcal{L}_\lambda}$). $\overline{\mathcal{L}_\lambda} = \{X t_1 \dots t_n \mid X t_1 \dots t_n \notin \mathcal{L}_\lambda\}$.

These terms are problematic for the very same reason terms in $\diamond \beta_0$ are, but cannot be handled directly by the unification of the meta language, that is only required to handle terms in \mathcal{L}_λ .

We write $\mathcal{P}(t)$ the set of sub-terms of t , and we write $\mathcal{P}(X) = \bigcup_{t \in X} \mathcal{P}(t)$ when X is a set of terms.

Definition 2.8 (Well behaved set). Given a set of terms $X \subseteq \mathcal{H}_0$,

$$\mathcal{W}(X) \Leftrightarrow \forall t \in \mathcal{P}(X), t \notin (\overline{\mathcal{L}_\lambda} \cup \diamond \eta \cup \diamond \beta_0)$$

PROPOSITION 2.9 (\mathcal{W} -PRESERVATION OF *hstep*). $\forall \mathbb{T}, \forall \mathbb{L}, \forall p, \forall \sigma, \forall \sigma'$

$$\mathcal{W}(\sigma \mathbb{T}) \wedge \sigma \mathbb{T}_{p_l} \approx_\lambda \sigma \mathbb{T}_{p_r} \mapsto \sigma' \Rightarrow \mathcal{W}(\sigma' \mathbb{T})$$

$$\mathcal{W}(\sigma \mathbb{T}) \wedge \text{progress}(\mathbb{L}, \sigma) \mapsto (_, \sigma') \Rightarrow \mathcal{W}(\sigma' \mathbb{T})$$

A less formal way to state 2.9 is that hstep and progress never “commit” an unneeded λ -abstraction in σ (a λ that could be erased by an η -contraction), nor put in σ an application whose head is flexible (an application node that could be erased by a β -reduction).

Note that proposition 2.9 does not hold for \approx_o as a whole since decompilation can introduce (actually restore) terms in $\diamond\eta$, $\diamond\beta_0$ or $\overline{\mathcal{L}_\lambda}$ that were moved out of the way during compilation.

3 OTHER ENCODINGS AND RELATED WORK

One could ignore the similarity between \approx_o and \approx_λ and “just” describe the object language unification procedure in the meta language by crafting a unif routine and using it as follows in rule (r3):

```
decision X :- unif X (all A x\ app [P, x]), finite A,
pi x\ decision (app [P, x]).
```

This choice would underuse the logic programming engine provided by the meta language since, by removing any datum from the head of rules, indexing degenerates. Moreover the unification procedure unif programmed in the meta language is likely to be an order of magnitude slower than one that is built-in.

Another possibility is to avoid having the application and abstraction nodes in the syntax tree, and use the ones of meta language, as in the following:

```
finite (fin N).
decision (nfact N NF).
decision (all A x\ P x) :- finite A, pi x\ decision (P x).
```

There are two reasons for dismissing this encoding. The first one is that in CIC it is not always possible to adopt it since the type system of the meta language is too weak to accommodate for the one of the object language. In CIC the lambda abstraction has to carry a type in order to make type checking decidable. Moreover CIC allows for functions with a variable arity, like the following example:

```
Fixpoint arr T n := if n is S m then T -> arr T m else T.
Definition sum n : arr nat n := ...
Check sum 2 7 8 : nat.
Check sum 3 7 8 9 : nat.
```

The type system of the λ Prolog is too stringent to accept this terms. The second reason is that the CIC encoding provided by Elpi is used for meta programming (extending) the Coq system, hence it must accommodate the manipulation of terms that are now known in advance (not even defined in Coq) without using introspection primitives such as Prologs’s functor and arg. In this sense constants have to live in an open world, like the string data type used in the examples so far.

In the literature we could find related encoding of the Calculus of Constructions [2]. The goal of that work was to exhibit a logic program performing proof checking in CC and hence relate the proof system of intuitionistic higher-order logic (that animates λ Prolog programs) with the Calculus of Constructions. The encoding is hence tailored toward a different goal, and utilizes three relations to represent the equational theory of CC. Section 6 contains a discussion about the use of the unification procedure of the

meta language in presence of non ground goals, but the authors do not aim at exploiting it to the degree we want.

4 PRELIMINARIES: \mathcal{F}_0 AND \mathcal{H}_0

In order to reason about unification we provide a description of the \mathcal{F}_0 and \mathcal{H}_0 languages where unification variables are first class terms, i.e. they have a concrete syntax. We keep these languages minimal, for example, we omit the all quantifier of CIC we used in the example in Section 1 together with the type notation of terms carried by the lam constructor.

```
kind fm type.          kind tm type.
type fapp list fm -> fm. type app list tm -> tm.
type flam (fm -> fm) -> fm. type lam (tm -> tm) -> tm.
type fcon string -> fm.   type con string -> tm.
type fuva addr -> fm.     type uva addr -> list tm -> tm.
```

Figure 1: The \mathcal{F}_0 and \mathcal{H}_0 languages

Unification variables (fuva term constructor) in \mathcal{F}_0 have no explicit scope: the arguments of an higher order variable are given via the fapp constructor. For example the term $P \ x$ is represented as $fapp[fuva \ N, \ x]$, where N is a memory address and x is a bound variable.

In \mathcal{H}_0 the representation of $P \ x$ is instead $uva \ N \ [x]$, since unification variables come equipped with an explicit scope. We say that the unification variable occurrence $uva \ N \ L$ is in \mathcal{L}_λ if and only if L is made of distinct names. The predicate to test this condition is called pattern-fragment:

```
type pattern-fragment list A -> o.
```

Natural numbers represent the memory addresses that identify unification variables in both languages. The memory and its associated operations are described below:

```
kind addr type.
type addr nat -> addr.
typeabbrev (mem A) (list (option A)).

type set? addr -> mem A -> A -> o.
type unset? addr -> mem A -> o.
type assign addr -> mem A -> A -> mem A -> o.
type new mem A -> addr -> mem A -> o.
```

If a memory cell is none, then the corresponding unification variable is not set. assign sets an unset cell to the given value, while new finds the first unused address and sets it to none.

E:is new used?

Since in \mathcal{H}_0 unification variables have a scope, their solution needs to be abstracted over it to enable the instantiation of a single solution to different scopes. This is obtained via the inctx container, and in particular via its abs binding constructor. On the contrary a solution to a \mathcal{F}_0 variable is a plain term.

```
typeabbrev fsubst (mem fm).
```

```
kind inctx type -> type.          (· ⊢ ·)
type abs (tm -> inctx A) -> inctx A.
type val A -> inctx A.
typeabbrev assignment (inctx tm).
typeabbrev subst (mem assignment).
```


We call $fsubst$ the memory of \mathcal{F}_0 , while we call $subst$ the one of \mathcal{H}_0 . Both have the invariant that they are not cyclic, TODO: explain.

```

kind arity type.
type arity nat -> arity.

kind fvariable type.
type fv addr -> fvariable.

kind hvariable type.
type hv addr -> arity -> hvariable.

kind mapping type.
type mapping fvariable -> hvariable -> mapping.
typeabbrev mmap (list mapping).
```

INVARIANT 1 (UNIFICATION VARIABLE ARITY). *Each variable A in \mathcal{H}_0 has a (unique) arity N and each occurrence $(uva\ A\ L)$ is such that $(len\ L\ N)$ holds*

The compiler establishes a mapping between variables of the two languages. In order to preserve invariant 1 we store the arity of each hvariable in the mapping and we reuse an existing mapping only if the arity matches.

```

type m-alloc fvariable -> hvariable -> mmap -> mmap ->
  subst -> subst -> o. (malloc)
m-alloc Fv Hv M M S S :- mem M (mapping Fv Hv), !.
m-alloc Fv Hv M [mapping Fv Hv|M] S S1 :- Hv = hv N _,
  alloc S N S1.
```

When a single fvariable occurs multiple times with different numbers of arguments the compiler generates multiple mappings for it, on a first approximation, and then makes the mapping bijective by introducing η -link; this detail is discussed in section 6.

As we mentioned in section 2.1 the compiler replaces terms in $\diamond\eta$ and \mathcal{L}_λ with fresh variables linked to the problematic terms. Each class of problematic terms has a dedicated link.

```

kind baselink type.
type link-eta tm -> tm -> baselink.
type link-beta tm -> tm -> baselink.
typeabbrev link (inctx baselink).
typeabbrev links (list link).
```

The right hand side of a link, the problematic term, can occur under binders. To accommodate this situation the compiler wraps baselink using the inctx container (see, $\cdot \vdash \cdot$).

INVARIANT 2 (LINK LEFT HAND SIDE). *The left hand side of a suspended link is a variable.*

New links are suspended by construction. If the left hand side variable is assigned during a step, then the link is considered for progress and possibly eliminated. This is discussed in section 6 and section 8.

4.1 Notational conventions

When we write \mathcal{H}_0 terms outside code blocks we follow the usual λ -calculus notation, reserving f, g, a, b for constants, x, y, z for bound variables and X, Y, Z, F, G, H for unification variables. However we

need to distinguish between the “application” of a unification variable to its scope and the application of a term to a list of arguments. We write the scope of unification variables in subscript while we use juxtaposition for regular application. Here a few examples:

```

f a      app[con "f", con "a"]
λx.λy.Fxy lam x\ lam y\ uva F [x, y]
λx.Fx a  lam x\ app[uva F [x], con "a"]
λx.Fx x  lam x\ app[uva F [x], x]
```

When variables x and y can occur in term t we shall write t_{xy} to stress this fact.

We write $\sigma = \{ A_{xy} \mapsto y \}$ for the assignment $abs\ x \backslash abs\ y \backslash y$ and $\sigma = \{ A \mapsto \lambda x. \lambda y. y \}$ for $lam\ x \backslash lam\ y \backslash y$.

When detailing examples we write links as equations between two terms under a context. The equality sign is subscripted with kind of baselink. For example $x \vdash A_x =_\beta F_x a$ corresponds to:

```
abs x\ val (link-beta (uva A [x]) (app[uva F [x], con "a"])))
```

When it is clear from the context we shall use the same syntax for \mathcal{F}_0 terms (although we never subscript unification variables).

4.2 Equational theory and Unification

In order to express properties ?? we need to equip \mathcal{F}_0 and \mathcal{H}_0 with term equality, substitution application and unification.

Term equality: $=_o$ vs. $=_\lambda$. We extend the equational theory over ground terms to the full languages by adding the reflexivity of unification variables (a variable is equal to itself).

The first four rules are common to both equalities and just define the usual congruence over terms, and since we use an HOAS encoding they also capture α -equivalence. In addition to that $=_o$ has rules for η and β -equivalence.

```

type (=o) fm -> fm -> o. (=o)
fcon X =o fcon X.
fapp A =o fapp B :- forall2 (=o) A B.
flam F =o flam G :- pi x\ x =o x => F x =o G x.
fuva N =o fuva N.
flam F =o T :- (=η)
  pi x\ beta T [x] (T' x), x =o x => F x =o T' x.
T =o flam F :- (=ηr)
  pi x\ beta T [x] (T' x), x =o x => T' x =o F x.
fapp [flam X | L] =o T :- beta (flam X) L R, R =o T. (βl)
T =o fapp [flam X | L] :- beta (flam X) L R, T =o R. (βr)
```

```

type (=λ) tm -> tm -> o.
con C =λ fcon C.
app A =λ fapp B :- forall2 (=λ) A B.
lam F =λ flam G :- pi x\ x =λ x => F x =λ G x.
uva N A =λ fuva N B :- forall2 (=λ) A B.
```

The main point in showing these equality tests is to remark how weaker $=_\lambda$ is, and to identify the four rules that need special treatment in the implementation of \approx_o .

For reference, $(beta\ T\ A\ R)$ reduces away lam nodes in head position in T whenever the list A provides a corresponding argument.

```

type beta fm -> list fm -> fm -> o.
beta A [] A.
beta (flam Bo) [H | L] R :- napp (Bo H) F, beta F L R.
```

```

581  beta (fapp A) L (fapp X) :- append A L X.
582  beta (fuva N) L (fapp [fuva N | L]).
583  beta (fcon H) L (fapp [fcon H | L]).
584  beta N L (fapp [N | L]) :- name N.
585
586  type napp fm -> fm -> o.
587  napp (fcon C) (fcon C).
588  napp (fuva A) (fuva A).
589  napp (flam F) (flam G) :- pi x\ napp (F x) (G x).
590  napp (fapp [fapp L1 | L2]) T :- !,
591    append L1 L2 L3, napp (fapp L3) T.
592  napp (fapp L) (fapp L1) :- map napp L L1.
593  napp N N :- name N.

```

The name predicate holds only on nominal constants (i.e. bound variables).² The choice of using n-ary application, rather than binary, is to make it easy to access the application's head. The price we pay is that substituting an application in the head of an application should be amended by "flattening" fapp nodes, that is the job of napp.³ Finally note that the cut operator is inessential, it could be removed at the cost of a verbose test on the head of L in the second rule about fapp: L's head can be fcon, flam or a name.

Substitution application: ρ_s and σ_t . Applying the substitution corresponds to dereferencing a term with respect to the memory. To ease the comparison we split \mathcal{F}_0 dereferencing into a fder step and a napp one. The former step replaces references to memory cells that are set with their values, and has a corresponding operation in \mathcal{H}_0 , namely deref. On the contrary napp has no corresponding operation in \mathcal{H}_0 . The reasons for this asymmetry is that an fapp node with a flexible head is always mapped to a uva (as per section 5.1 and section 8), preventing nested applications to materialize.

```

612  type fder fsubst -> fm -> fm -> o.
613  fder _ (fcon C) (fcon C).
614  fder S (fapp A) (fapp B) :- map (fder S) A B.
615  fder S (flam F) (flam G) :-
616    pi x\ fder S x x => fder S (F x) (G x).
617  fder S (fuva N) R :- set? N S T, fder S T R.
618  fder S (fuva N) (fuva N) :- unset? N S.
619
620  type fderef fsubst -> fm -> fm -> o. (ρs)
621  fderef S T T2 :- fder S T T1, napp T1 T2.

```

Applying the substitution in \mathcal{H}_0 is very similar, with the caveat that assignments have to be moved to the current scope, i.e. renaming the abs-bound variables with the names in the scope of the unification variable occurrence.

```

627  type deref subst -> tm -> tm -> o. (σt)
628  deref _ (con C) (con C).
629  deref S (app A) (app B) :- map (deref S) A B.
630  deref S (lam F) (lam G) :-
631    pi x\ deref S x x => deref S (F x) (G x).
632  deref S (uva N L) R :- set? N S A,
633    move A L T, deref S T R.

```

²Elpi provides it as a builtin, but one could implement it by systematically loading the hypothetical rule `name x` every time a nominal constant is postulated via `pi x\`

³Note that napp is an artefact of formalization of \mathcal{F}_0 we do in this presentation and, as we explain later, no equivalent of napp is needed in \mathcal{H}_0 .

```

deref S (uva N A) (uva N B) :- unset? N S,
map (deref S) A B.

```

Note that move strongly relies on invariant 1: the length of the arguments of all occurrences of a unification variable and the number of abstractions in its assignment have to match. In turn this grants that move never fails.

```

type move assignment -> list tm -> tm -> o.
move (abs Bo) [H|L] R :- move (Bo H) L R.
move (val A) [] A.

```

Term unification: \approx_o vs. \approx_λ . In this paper we assume to have an implementation of \approx_λ that satisfies properties ?? and ??. Although we provide an implementation in the appendix (that we used for testing purposes) we only describe its signature here. Elpi is expected to provide this brick, as well as any other implementation of λ Prolog.

```

type (≈λ) tm -> tm -> subst -> subst -> o.

```

The only detail worth discussing is the fact that the procedure updates a substitution, rather than just crafting one as presented in section 2. The reason is that the algorithm folds over a term, updating a substitution while it traverses it.

E:explain better

5 BASIC SIMULATION OF \mathcal{F}_0 IN \mathcal{H}_0

In this section we describe a basic compilation scheme that we refine later, in the following sections. This scheme is sufficient to implement an \approx_o that respects β -conversion for terms in \mathcal{L}_λ . The extension to $\eta\beta$ -conversion is described in Section 6 and the support for terms outside \mathcal{L}_λ in Section 8.

5.1 Compilation

E:manca beta normal in entrata

The main task of the compiler is to recognize \mathcal{F}_0 variables standing for functions and map them to higher order variables in \mathcal{H}_0 . In order to bring back the substitution from \mathcal{H}_0 to \mathcal{F}_0 the compiler builds a "memory map" connecting the the kind of variables using routine (*malloc*).

The signature of the comp predicate below allows for the generation of links (suspended unification problems) that play no role in this section but play a major role in section 6 and section 8. With respect to section 2 the signature also allows for updates to the substitution.

```

type comp fm -> tm -> mmap -> mmap -> links -> links ->
subst -> subst -> o.
comp (fcon C) (con C) M M L L S S.
comp (flam F) (lam F1) M1 M2 L1 L2 S1 S2 :- (cλ)
  comp-lam F F1 M1 M2 L1 L2 S1 S2.
comp (fuva A) (uva B [J]) M1 M2 L L S1 S2 :-
  m-alloc (fv A) (hv B (arity z)) M1 M2 S1 S2.
comp (fapp A) (app A1) M1 M2 L1 L2 S1 S2 :- (c@)
  fold6 comp A A1 M1 M2 L1 L2 S1 S2.

```

The code above uses that possibility in order to allocate space for the variables, i.e. sets their memory address to none (a details not worth mentioning in the previous sections).

```

type comp-lam (fm -> fm) -> (tm -> tm) ->
  mmap -> mmap -> links -> links -> subst -> subst -> o.
comp-lam F G M1 M2 L1 L3 S1 S2 :-
  pi x y\ (pi M L S\ comp x y M M L L S S) =>
    comp (F x) (G y) M1 M2 L1 (L2 y) S1 S2,
  close-links L2 L3.

```

In the code above the syntax `pi x y\..` is syntactic sugar for iterated pi abstraction, as in `pi x\ pi y\..`.

The auxiliary function `close-links` tests if the bound variable `v` really occurs in the link. If it is the case the link is wrapped into an additional `abs` node binding `v`. In this way links generated deep inside the compiled terms can be moved outside their original context of binders.

```

type close-links (tm -> links) -> links -> o.
close-links (v\ [X | L v]) [X|R] :- !, close-links L R.
close-links (v\ [X v | L v]) [abs X|R] :- close-links L R.
close-links (_\ []) [].

```

Note that we could remove the first rule, whose solve purpose is to make links more readable by pruning unused context entries.

5.1.1 Detection of $\diamond \beta_0$. The following rule is inserted just before rule ($c_{@}$).

```

comp (fapp [fuva A|Ag]) (uva B Ag1) M1 M2 L L S1 S2 :-
  pattern-fragment Ag, !,
  fold6 comp Ag Ag1 M1 M1 L L S1 S1,
  len Ag Arity,
  m-alloc (fv A) (hv B (arity Arity)) M1 M2 S1 S2.

```

Note that compiling `Ag` cannot create new mappings nor links, since `Ag` is made of bound variables and the hypothetical rule loaded by `comp-lam` (see below) grants this property. Also note that this rule generates no links.

5.2 Execution

A step in \mathcal{H}_0 consists in unifying two terms and reconsidering all links for progress. If any of the two tasks fail we say that the entire step fails, and it is at this granularity that we can relate steps in the two languages.

```

type hstep tm -> tm -> links -> links -> subst -> subst -> o.
hstep T1 T2 L1 L2 S1 S3 :-
  (T1  $\approx_\lambda$  T2) S1 S2,
  progress L1 L2 S2 S3.

```

Note that the infix notation $((A \approx_\lambda B) C D)$ is syntactic sugar for $((\approx_\lambda) A B C D)$.

Reconsidering links is a fixpoint, since the progress of a link can update the substitution and in turn enable another link to progress.

```

type progress links -> links -> subst -> subst -> o.
progress L L2 S1 S3 :-
  progress1 L L1 S1 S2,
  occur-check-links L1,
  if (L = L1, S1 = S2)

```

```

(L2 = L1, S3 = S1)
(progress L1 L2 S2 S3).

```

In the base compilation scheme `progress1` is the identity on both the links and the substitution, so the fixpoint trivially terminates. Sections 6 and 8 add rules to `progress1` and justify why the don't hinder termination. For brevity we omit the code that applies the substitution `S1` to all terms in \mathbb{L} .

Since compilation moves problematic terms out of the sigh of \approx_λ , that procedure can only perform a partial occur check. For example the unification problem $X \approx_\lambda f Y$ cannot generate a cyclic substitution alone, but should be disallowed if a \mathbb{L} contains a link like $\vdash Y =_\eta \lambda z.X_z$: We don't know yet if Y will feature a lambda in head position, but we surely know it contains X , hence $f Y$ and that fails the occur check. The procedure `occur-check-links` is in charge of ensuring that each link does not represent a (suspended) unification problem doomed to fail because of occur check. This check is needed in order to guarantee proposition 2.2 (simulation fidelity).

5.3 Substitution decompilation

Decompiling the substitution requires to first force the progress of links and then allocating new unassigned variables in the substitution for \mathcal{F}_0 and finally decompiling all assignments. Note that invariant 2 and the occur check allows us to update the subst.

```

type decompile mmap -> links -> subst ->
  fsubst -> fsubst -> o.
fsubst -> fsubst -> o.
decompile M1 L S F1 F3 :-
  commit-links L S S1,
  complete-mapping S1 S1 M1 M2 F1 F2,
  decompM M2 M2 S1 F2 F3.

```

Decompiling an assignment requires to turn abstractions into lambdas. For aesthetic purposes we also eta-contract the result (not needed since \mathcal{F}_0 equality can do that)

```

type decompM mmap -> mmap -> subst -> fsubst -> fsubst -> o.
decompM _ [] _ F F.
decompM M [mapping (fv V) (hv H _)]MS S F1 F3 :- set? H S A,
  deref-assmt S A A1,
  abs->lam A1 T, decompM M T T1,
  eta-contract T1 T2,
  assign V F1 T2 F2,
  decompM M MS S F2 F3.
decompM M [mapping _ (hv H _)]MS S F1 F2 :- unset? H S,
  decompM M MS S F1 F2.

```

Finally decompiling a term is trivial, now that we have an extended mapping containing all unassigned variables \approx_λ may have introduced.

```

type decomp mmap -> tm -> fm -> o.
decomp _ (con C) (fcon C).
decomp M (app A) (fapp B) :- map (decomp M) A B.
decomp M (lam F) (flam G) :-
  pi x y\ (pi M\ decomp M x y) => decomp M (F x) (G y).
decomp M (uva Hv Ag) R :-
  mem M (mapping (fv Fv) (hv Hv _)),
  map (decomp M) Ag Bg,
  beta (fuva Fv) Bg R.

```

E:What is commit-links and complete-mapping?

Note that we use beta to build fapp nodes when needed (if Ag is empty no fapp node should appear).

INVARIANT 3. *TODO: dire che il mapping è bijective*

5.4 Definition of \approx_o and its properties

```
type ( $\approx_o$ ) fm -> fm -> fsubst -> o.
(A  $\approx_o$  B) F :-
  comp A A' [] M1 [] [] S1,
  comp B B' M1 M2 [] [] S1 S2,
  hstep A' B' [] [] S2 S3,
  decomp M2 M2 S3 [] F.
```

The code given so far applies to terms in $\beta\eta$ -normal form where unification variables in \mathcal{F}_o can occur non linearly but always with the same number of arguments, and where their arguments are distinct names (as per \mathcal{L}_λ).

LEMMA 5.1 (COMPILATION ROUND TRIP). *If comp S T [] M [] - [] - then decomp M T S*

PROOF SKETCH. trivial, since the terms are beta normal beta just builds an app. \square

LEMMA 5.2. *Properties 1 and 2 hold for the implementation of \approx_o above*

PROOF SKETCH. In this setting \approx_λ is as strong as \approx_o on ground terms. What we have to show is that whenever two different \mathcal{F}_o terms can be made equal by a substitution ρ (plus the β_l and β_r if needed) we can find this ρ by finding a σ via \approx_λ on the corresponding \mathcal{H}_o terms and by decompiling it. If we look at the \mathcal{F}_o terms, there are two interesting cases:

- fuva $X \approx_o s$. In this case after comp we have $Y \approx_\lambda t$ that succeeds with $\sigma = \{Y \mapsto t\}$ and σ is decompiled to $\rho = \{Y \mapsto s\}$.
- fapp[fuva $X|L$] $\approx_o s$. In this case we have $Y_x \approx_\lambda t$ that succeeds with $\sigma = \{\vec{y} \mapsto Y \mapsto t[\vec{x}/\vec{y}]\}$ that in turn is decompiled to $\rho = \{Y \mapsto \lambda \vec{y}.s[\vec{x}/\vec{y}]\}$. Thanks to β_l ($\lambda \vec{y}.s[\vec{x}/\vec{y}]) \vec{x} \approx_o s$.

Since the mapping is a bijection occur check in \mathcal{H}_o corresponds to occur check in \mathcal{F}_o . \square

LEMMA 5.3. *Properties simulation (2.1) and fidelity (2.2) hold*

PROOF SKETCH. Since progress1 is trivial fstep and hstep are the same, that is in this context where input terms are $\beta\eta$ -normal and we disregard η -equivalence \approx_λ is equivalent to \approx_o . \square

5.5 Limitations of by this basic scheme

The basic compilation scheme is not about to deal with the following problem:

$$\mathbb{P} = \{ \lambda xy.X y x \approx_o \lambda xy.x \quad \lambda x.f (X x) x \approx_o Y \}$$

Note that here X is used with different arities, moreover in the second problem the left hand side happens to be an eta expansion (of $f(\lambda y.y)$) only after we discover (at run time) that $X = \lambda x \lambda y.y$ (i.e. that X discards the x argument). Both problems are addressed in the next two sections.

6 HANDLING OF $\diamond\eta$

η -reduction is an equivalence relation where a term of the form $\lambda x.t x$ can be converted to t any time x does not occur as a free variable in t . We call t the η -contraction of $\lambda x.t x$.

Following the compilation scheme of section 5.1 the unification problem \mathbb{P} is compiled as follows:

$$\begin{aligned} \mathbb{P} &= \{ \lambda x.X x \approx_o f \} \\ \mathbb{T} &= \{ \lambda x.A_x \approx_\lambda f \} \\ \mathbb{M} &= \{ X \mapsto A^1 \} \end{aligned}$$

While $\lambda x.X x \approx_o f$ does admit the solution $\rho = \{X \mapsto f\}$, the corresponding problem in \mathbb{T} does not: $\text{lam } x \backslash \text{uva } A [x]$ and $\text{con } "f"$ start with different, rigid, term constructors hence \approx_λ fails.

In order to guarantee proposition 2.1 we detect lambdas that can disappear by eta contraction (section 6.1) and we modify the compiled terms by putting fresh unification variables in their place: the problematic term is moved from \mathbb{T} to \mathbb{L} (section 6.2). The compilation of the problem \mathbb{P} above is refined to:

$$\begin{aligned} \mathbb{P} &= \{ \lambda x.X x \approx_o f \} \\ \mathbb{T} &= \{ A \approx_\lambda f \} \\ \mathbb{M} &= \{ X \mapsto B^1 \} \\ \mathbb{L} &= \{ \vdash A =_\eta \lambda x.B_x \} \end{aligned}$$

As per invariant 2 the term on the left is a variable, and its right counterpart is the term in $\diamond\eta$. That term has the following property:

INVARIANT 4 (η -link rhs). *The rhs of any η -link has the shape $\lambda x.t$ and t is not a lambda.*

η -link are kept in the link store \mathbb{L} during execution and activated when some conditions hold on lhs or rhs. Link activation is implemented by extending the progress1 predicate (defined in section 5.2).

6.1 Detection of $\diamond\eta$

When compiling a term t we need to determine if any subterm $s \in \mathcal{P}(t)$ that is of the form $\lambda x.r$, where x occurs in r , can be a η -expansion, i.e. if there exists a substitution ρ such that $\rho(\lambda x.r) =_o s$. The detection of lambda abstractions that can “disappear” is not as trivial as it may seem, here a few examples:

$$\begin{aligned} \lambda x.f (A x) &\in \diamond\eta \quad \rho = \{A \mapsto \lambda x.x\} \\ \lambda x.f (A x) x &\in \diamond\eta \quad \rho = \{A \mapsto \lambda x.a\} \\ \lambda x.f x (A x) &\notin \diamond\eta \\ \lambda x \lambda y.f (A x) (B y x) &\in \diamond\eta \quad \rho = \{A \mapsto \lambda x.x, B \mapsto \lambda y.\lambda x.y\} \end{aligned}$$

The first two examples are easy, and show how a unification variable can expose or erase a variable in their scope and turn the resulting term in an η -expansion or not.

The third example shows that when a variable occurs outside the scope of a unification variable it cannot be erased and can hence prevent a term from being an η -expansion.

The last example shows the recursive nature of the check we need to implement. The term starts with a spine of two lambdas hence the whole term is in $\diamond\eta$ iff the inner term $\lambda y.f (A x) (B y x)$ is in $\diamond\eta$ itself. If it is, it could η -contract to $f (A x)$ making $\lambda x.f (A x)$ a potential η -expansion.

We can now define more formally how $\diamond\eta$ terms are detected together with its auxiliary functions:

Definition 6.1 (may-contract-to). A β -normal term s may-contract-to a name x if there exists a substitution ρ such that $\rho s =_o x$.

LEMMA 6.2. A β -normal term $s = \lambda x_1 \dots x_n. t$ may-contract-to x only if one of the following three conditions holds:

- (1) $n = 0$ and $t = x$;
- (2) t is the application of x to a list of terms l and each l_i may-contract-to x_i (e.g. $\lambda x_1 \dots x_n. x \ x_1 \dots x_n =_o x$);
- (3) t is a unification variable with scope W , and for any $v \in \{x, x_1 \dots x_n\}$, there exists a $w_i \in W$, such that w_i may-contract-to v (if $n = 0$ this is equivalent to $x \in W$).

PROOF SKETCH. Since our terms are in β -normal form there is only one rule that can play a role (namely η_l), hence if the term s is not exactly x (case 1) it can only be an η -expansion of x , or a unification variable that can be assigned to x , or a combination of both. If s begins with a lambda, then the lambda can only disappear by η contraction. In that case the term t is under the spine of binders $x_1 \dots x_n$, t can either be x applied to terms that can may-contract-to these variables (case 2), or a unification variable that can be assigned to that application (case 3). \square

Definition 6.3 (occurs-rigidly). A name x occurs-rigidly in a β -normal term t , if $\forall \rho, x \in \mathcal{P}(\rho t)$

In other words x occurs-rigidly in t if it occurs in t outside of the scope of a unification variable X , otherwise an instantiation of X can make x disappear from t . Moreover, note that η -contracting t cannot make x disappear, since x is not a locally bound variable inside t .

We can now derive the implementation for $\Diamond\eta$ detection:

Definition 6.4 (maybe-eta). Given a β -normal term $s = \lambda x_1 \dots x_n. t$, maybe-eta s holds if any of the following holds:

- (1) t is a constant or a name applied to the arguments $l_1 \dots l_m$ such that $m \geq n$ and for every i such that $m - n < i \leq m$ the term l_i may-contract-to x_i , and no x_i occurs-rigidly in $l_1 \dots l_{m-n}$;
- (2) t is a unification variable with scope W and for each x_i there exists a $w_j \in W$ such that w_j may-contract-to x_i .

LEMMA 6.5 ($\Diamond\eta$ DETECTION). If t is a β -normal term and maybe-eta t holds, then $t \in \Diamond\eta$.

PROOF SKETCH. Follows from definition 6.3 and lemma 6.2 \square

Remark that the converse of lemma 6.5 does not hold: there exists a term t satisfying the criteria (1) of definition 6.4 that is not in $\Diamond\eta$, i.e. there exists no substitution ρ such that ρt is an η -expansion. A simple counter example is $\lambda x. f \ (A \ x) \ (A \ x)$ since x does not occur-rigidly in the first argument of f , and the second argument of f may-contract-to x . In other words $A \ x$ may either use or discard x , but our analysis does not take into account that the same term cannot have two contrasting behaviors.

As we will see in the rest of this section this is not a problem since it does not break proposition 2.1 nor proposition 2.2.

6.2 Compilation and decompilation

Compilation. The following rule is inserted just before rule (c_λ) from the code in section 5.1.

```
comp (flam F) (uva A Scope) M1 M2 L1 L3 S1 S3 :-
  maybe-eta (flam F) [], !,
  alloc S1 A S2,
  comp-lam F F1 M1 M2 L1 L2 S2 S3,
  get-scope (lam F1) Scope,
  L3 = [val (link-eta (uva A Scope) (lam F1)) | L2].
```

The rule triggers when the input term `flam F` is in $\Diamond\eta$. It compiles `flam F` to `lam F1` but puts the fresh variable `A` in its place. This variable sees all the names free in `lam F1`. The critical part of this rule is the creation of the η -link, which relates the variable `A` with `lam F1`. This link clearly validates invariant 2.

COROLLARY 6.6. The rhs of any η -link has exactly one lambda abstraction, hence the rule above respects invariant 4.

PROOF SKETCH. By contradiction, suppose that the rule above is triggered and that the rhs of the link is $\lambda x. \lambda y. t_{xy}$. If maybe-eta $\lambda y. t_{xy}$ holds the recursive call to `comp` (made by `comp-lam`) must have put a fresh variable in its place, so this case is impossible. Otherwise, if maybe-eta $\lambda y. t_{xy}$ does not hold, also maybe-eta $\lambda x. \lambda y. t_{xy}$ does not hold, contradicting the assumption that the rule triggered. \square

Decompilation. Decompilation of the remaining η -link (i.e. the η -link that have been activated) is performed by iterating over them and unifying lhs and rhs. Note that this unification never fails, since lhs is a flexible term not appearing in any other η -link (by definition 6.9).

6.3 Progress

η -link are meant to delay the unification of “problematic” terms until we know for sure if the term has to be η -contracted or not.

Definition 6.7 (progress- η -left). A link $\Gamma \vdash X =_\eta T$ is removed from \mathbb{L} when X becomes rigid. Let $y \in \Gamma$, there are two cases:

- (1) if $X = a$ or $X = y$ or $X = f \ a_1 \dots a_n$ we unify the η -expansion of X with T , that is we run $\lambda x. X \ x \simeq_\lambda T$
- (2) if $X = \lambda x. t$ we run $X \simeq_\lambda T$.

Definition 6.8 (progress- η -right). A link $\Gamma \vdash X =_\eta T$ is removed from \mathbb{L} when either 1) maybe-eta T does not hold (anymore) or 2) by η -contracting T to T' , T' is a term not starting with the `lam` constructor. In the first case, X is unified with T and in the second one, X is unified with T' (under the context Γ).

There is a third case in which a link is removed from \mathbb{L} , namely when the lhs is assigned to a variable that is the lhs of another η -link.

Definition 6.9 (progress- η -deduplicate). A link $\Gamma \vdash X_{\vec{s}} =_\eta T$ is removed from \mathbb{L} when another link $\Delta \vdash X_{\vec{r}} =_\eta T'$ is in \mathbb{L} . By invariant 1 the length of \vec{s} and \vec{r} is the same hence we can move the term T' from Δ to Γ by renaming its bound variables, i.e. $T'' = T' [\vec{r}/\vec{s}]$. We then run $T \simeq_\lambda T''$ (under the context Γ).

LEMMA 6.10. Given a η -link l , the unification done by progress- η -left is between terms in \mathcal{W}

PROOF SKETCH. Let σ be the substitution, which is $\mathcal{W}(\sigma)$ (by proposition 2.9). $\text{lhs} \in \sigma$, therefore $\mathcal{W}(\text{lhs})$. By definition 6.7, if 1) lhs is a name, a constant of an application, then, $\lambda x.\text{lhs } x$ is unified with rhs . By invariant 4, $\text{rhs} = \lambda x.t$, therefore $\mathcal{W}(t)$. Otherwise, 2) lhs has lam as functor, rhs should not be an η -expansion, so, $\mathcal{W}(\text{rhs})$. In both cases, unification is performed between terms in \mathcal{W} . \square

LEMMA 6.11. *Given a η -link l , the unification done by progress- η -right is between terms in \mathcal{W} .*

PROOF SKETCH. lhs is variable, and, by definition 6.8, rhs is either no more a $\diamond\eta$, i.e. rhs is not a η -expansion and, so, $\mathcal{W}(\text{rhs})$. Otherwise, rhs can reduce to a term which cannot be a η -expansion, and, so, $\mathcal{W}(\text{rhs})$. In both cases, unification is done between terms in \mathcal{W} . \square

LEMMA 6.12. *Given a η -link l , the unification done by progress- η -deduplicate is between terms in \mathcal{W} .*

PROOF. The unification is done between the rhs of two η -link. Both rhs has the shape $\lambda x.t$, with $\mathcal{W}(t)$, therefore, the unification is done between well-behaved terms. \square

LEMMA 6.13. *The introduction of η -link guarantees proposition 2.9 (\mathcal{W} -preservation of hstep)*

PROOF SKETCH. By lemmas 6.10 to 6.12, every unification performed by the activation of a η -link is done between terms in \mathcal{W} , therefore, the substitution remains \mathcal{W} .

D: Bisogna aggiungere un lemma nella section 2.1 che dice che unificare due termini in \mathcal{W} , in una σ , tale che $\mathcal{W}(\sigma)$, non invalida \mathcal{W}

LEMMA 6.14. *progress terminates.*

PROOF SKETCH. Rules definitions 6.7 and 6.8 and definition 6.9 remove one link from \mathbb{L} , hence they cannot be applied indefinitely. Moreover each rule only relies on terminating operations such as \simeq_λ , η -contraction, η -expansion, relocation (a recursive copy of a finite term). \square

LEMMA 6.15 (FIDELITY WITH η -link). *The introduction of η -link guarantees proposition 2.2 (Simulation fidelity)*

PROOF SKETCH. *progress- η -left* and *progress- η -deduplicate* activate a η -link when, in the original unification problem, a $\diamond\eta$ term is unified with respectively a well-behaved term or another $\diamond\eta$ term. In both cases, the links trigger a unification which succeeds iff the same unification in \mathcal{F}_0 succeeds, guaranteeing proposition 2.2. *progress- η -right* never fails, in fact, this progression refines a variable to a rigid term and plays no role in proposition 2.2. \square

Example of progress- η -left. The example at the beginning of section 6, once $\sigma = \{A \mapsto f\}$, triggers this rule since the link becomes $\vdash f =_\eta \lambda x.B_x$ and the lhs is a constant. In turn the rule runs $\lambda x.f x \simeq_\lambda \lambda x.B_x$, resulting in $\sigma = \{A \mapsto f; B_x \mapsto f\}$. Decomposition generates $\rho = \{X \mapsto f\}$, since X is mapped to B and f is the η -contracted version of $\lambda x.f x$.

Example of progress- η -deduplicate. A very basic example of η -link deduplication, is given below:

$$\begin{aligned} \mathbb{P} &= \{ \lambda x.(X x) \simeq_o \lambda x.(Y x) \} \\ \mathbb{T} &= \{ A \simeq_\lambda C \} \\ \mathbb{M} &= \{ X \mapsto B^1 \quad Y \mapsto D^1 \} \\ \mathbb{L} &= \{ \vdash A =_\eta \lambda x.B_x \quad \vdash C =_\eta \lambda x.D_x \} \end{aligned}$$

The result of $A \simeq_\lambda C$ is that the two η -link share the same lhs . By unifying the two rhs we get $\sigma = \{A \mapsto C, B \mapsto D\}$. In turn, given the map \mathbb{M} , this second assignment is decomposed to $\rho = \{X \mapsto Y\}$ as expected.

We delay at the end of next section an example of η -link progression due to *progress- η -right*

7 ENFORCING INVARIANT 1

We report here the problem given in section 5.5 where X is used with two different arities and the output of the compilation does not respect invariant 3 (merging the two mappings for s would break invariant 1). In this section we explain how to replace the duplicate mapping with some η -link in order to restore the invariants.

$$\begin{aligned} \mathbb{P} &= \{ \lambda x.\lambda y.(X y x) \simeq_o \lambda x.\lambda y.x \quad \lambda x.(f (X x) x) \simeq_o Y \} \\ \mathbb{T} &= \{ A \simeq_\lambda \lambda x.\lambda y.x \quad D \simeq_\lambda F \} \\ \mathbb{M} &= \{ X \mapsto E^1 \quad Y \mapsto F^0 \quad X \mapsto C^2 \} \\ \mathbb{L} &= \{ \vdash D =_\eta \lambda x.(f E_x x) \quad \vdash A =_\eta \lambda x.B_x \} \\ &\quad \{ x \vdash B_x =_\eta \lambda y.C_{yx} \} \end{aligned}$$

We see that the maybe-eta as identified $\lambda xy.X y x$ and $\lambda x.f (X x) x$ and the compiler has replaced them with A and D respectively. However, the mapping \mathbb{M} breaks invariant 3: the \mathcal{F}_0 variable X is mapped to two different \mathcal{H}_0 variables. To address this problem we adjust the compiler's output with a map-deduplication procedure.

Definition 7.1 (align-arity). Given two mappings $m_1 : X \mapsto A^m$ and $m_2 : X \mapsto C^n$ where $m < n$ and $d = n - m$, *align-arity* $m_1 m_2$ generates the following d links, one for each i such that $0 \leq i < d$,

$$x_0 \dots x_{m+i} \vdash B_{x_0 \dots x_{m+i}}^i =_\eta \lambda x_{m+i+1}. B_{x_0 \dots x_{m+i+1}}^{i+1}$$

where B^i is a fresh variable of arity $m + i$, and $B^0 = A$ as well as $B^d = C$.

The intuition is that we η -expand the occurrence of the variable with lower arity to match the higher arity. Since each η -link can add exactly one lambda, we need as many links as the difference between the two arities.

Definition 7.2 (map-deduplication). For all mappings $m_1, m_2 \in \mathbb{M}$ such that $m_1 : X \mapsto A^m$ and $m_2 : X \mapsto C^n$ and $m < n$ we remove m_1 from \mathbb{M} and add to \mathbb{L} the result of *align-arity* $m_1 m_2$.

If we look back the example give at the beginning of this section, we can deduplicate $X \mapsto E^1, X \mapsto C^2$ by removing the first mapping and adding the auxiliary η -link: $x \vdash E_x =_\eta \lambda y.C_{xy}$. After deduplication the compiler output is as follows:

$$\begin{aligned} \mathbb{P} &= \{ \lambda x.\lambda y.(X y x) \simeq_o \lambda x.\lambda y.x \quad \lambda x.(f (X x) x) \simeq_o Y \} \\ \mathbb{T} &= \{ A \simeq_\lambda \lambda x.\lambda y.x \quad D \simeq_\lambda F \} \\ \mathbb{M} &= \{ Y \mapsto F^0 \quad X \mapsto C^2 \} \\ \mathbb{L} &= \{ x \vdash E_x =_\eta \lambda y.C_{xy} \quad \vdash D =_\eta \lambda x.(f E_x x) \} \\ &\quad \{ \vdash A =_\eta \lambda x.B_x \quad x \vdash B_x =_\eta \lambda y.C_{yx} \} \end{aligned}$$

In this example, \mathbb{T}_1 assigns A which triggers \mathbb{L}_3 and then \mathbb{L}_4 by definition 6.7. C_{yx} is therefore assigned to x (the second variable of its scope). We can finally see the *progress- η -right* of \mathbb{L}_1 : its rhs is now $\lambda y.y (C_{xy} \text{ gives } y)$. Since it is no more in $\Diamond\eta$, $\lambda y.y$ is unified with E_x . Moreover, \mathbb{L}_2 is also triggered due to definition 6.8: $\lambda x.(f \cdot (\lambda y.y) \cdot x)$ is η -reducible to $f \cdot (\lambda y.y)$ which is a term not starting with the lam constructor.

8 HANDLING OF $\overline{\mathcal{L}_\lambda}$

D:I've rewritten it, it is clearer?

Until now, we have only dealt we unification of terms in \mathcal{L}_λ . However, we want the unification relation to be more robust so that it can work with terms in $\overline{\mathcal{L}_\lambda}$. In general, unification in $\overline{\mathcal{L}_\lambda}$ admits more then one solution and committing one of them in the substitution does not guarantee prop. 2. For instance, $X \ a \approx_o a$ is a unification problem admits two different substitutions: $\rho_1 = \{X \mapsto \lambda x.x\}$ and $\rho_2 = \{X \mapsto \lambda_.a\}$. Prefer one over the other may break future unifications.

It is the case that, given a list of unification problems, $\mathbb{P}_1 \dots \mathbb{P}_n$ with \mathbb{P}_n in $\overline{\mathcal{L}_\lambda}$, the resolution of $\bigwedge_{i=0}^{n-1} \mathbb{P}_i$ gives a partial substitution ρ , such that $\rho \mathbb{P}_n$ falls again in \mathcal{L}_λ .

$$\begin{aligned} \mathbb{P} &= \{ X \approx_o \lambda x.Y \quad (X \cdot a) \approx_o a \} \\ \mathbb{T} &= \{ A \approx_\lambda \lambda x.B \quad (A \cdot a) \approx_\lambda a \} \\ \mathbb{M} &= \{ Y \mapsto B^0 \quad X \mapsto A^0 \} \end{aligned}$$

In the example above, we see that \mathbb{P}_1 instantiates X so that \mathbb{P}_2 , can be solved in \mathcal{L}_λ .

E:it is even a ground term, there is no unification left to perform actually

D:i don't understand the note

On the other hand, we see that, \approx_λ can't solve the compiled problems \mathbb{T} . In fact, the resolution of \mathbb{T}_1 gives the substitution $\sigma = \{A \mapsto \lambda x.B\}$, but the dereferencing of \mathbb{T}_2 gives the non-unifiable problem $(\lambda x.B) \cdot a \approx_\lambda a$.

To address this unification problem, term compilation should capture the terms t in $\overline{\mathcal{L}_\lambda}$ and replace them with fresh variables X . The variables X and the terms t are linked through a β -link.

β -link guarantees invariant 2 and the term on the rhs has the following property:

D:Is it clearer?

INVARIANT 5 (β -link rhs). *The rhs of any β -link has the shape $X_{s_1 \dots s_n} t_1 \dots t_m$ such that X is a unification variable with scope $s_1 \dots s_n$ and $t_1 \dots t_m$ is a list of terms. This is equivalent to $\text{app}[\text{uva } X \ S \mid L]$, where $S = s_1 \dots s_n$ and $L = t_1 \dots t_m$.*

LEMMA 8.1 (β -link WITH RIGID lhs). *If the lhs of a β -link is instantiated to a rigid term and its rhs counterpart is still in $\overline{\mathcal{L}_\lambda}$, the original unification problem is not in \mathcal{L}_λ and the unification fails.*

PROOF SKETCH. Given $X \ t_1 \dots t_n \approx_\lambda t$ where t is a rigid term and $t_1 \dots t_n$ is not in \mathcal{L}_λ . By construction, $X \ t_1 \dots t_n$ is replaced with a variable Y , and the β -link $\Gamma \vdash Y =_\beta X \ t_1 \dots t_n$ is created. The unification instantiates Y to t , making the lhs of the link a rigid term, while rhs is still in $\overline{\mathcal{L}_\lambda}$. The original problem is in fact outside \mathcal{L}_λ . \square

8.1 Compilation and decompilation

Detection of $\overline{\mathcal{L}_\lambda}$ is quite simple to implement in the compiler, since it is sufficient to detect applications with flexible head and argument that are not in \mathcal{L}_λ .

```
comp (fapp [fuva A|Ag]) (uva B Scope) M1 M3 L1 L3 S1 S4 :- !,
  pattern-fragment-prefix Ag Pf Extra,
  len Pf Arity,
  alloc S1 B S2,
  m-alloc (fv A) (hv C (arity Arity)) M1 M2 S2 S3,
  fold6 comp Pf Pf1 M2 M2 L1 L1 S3 S3,
  fold6 comp Extra Extra1 M2 M3 L1 L2 S3 S4,
  Beta = app [uva C Pf1 | Extra1],
  get-scope Beta Scope,
  L3 = [val (link-beta (uva B Scope) Beta) | L2].
```

The list Ag is split into the list Pf and Extra such that $\text{append Pf Extra Ag}$ and Pf is the largest prefix of Ag such that Pf is in \mathcal{L}_λ . The rhs of the β -link is the application of a fresh variable C having in scope all the free variables appearing in the compiled version of Pf and Extra . The variable B , returned has the compiled term, is a fresh variable having in scope all the free variables occurring in Pf1 and Extra1 .

INVARIANT 6. *The rhs of a β -link has the shape $X_{s_1 \dots s_n} t_1 \dots t_m$.*

COROLLARY 8.2. *Let $X_{s_1 \dots s_n} t_1 \dots t_m$ be the rhs of a β -link, then $m > 0$.*

PROOF SKETCH. Assume we have a β -link, by contradiction, if $m = 0$, then the original \mathcal{F}_0 term has the shape $\text{fapp}[\text{fuva } M \mid \text{Ag}]$ where Ag is a list of distinct names (i.e. the list Extra is empty). This case is however captured by rule (c_λ) (from section 5.1) and no β -link is produced which contradicts our initial assumption. \square

COROLLARY 8.3. *Let $X_{s_1 \dots s_n} t_1 \dots t_m$ be the rhs of a β -link, then t_1 either appears in $s_1 \dots s_n$ or it is not a name.*

PROOF SKETCH. By construction, the lists $s_1 \dots s_n$ and $t_1 \dots t_m$ are built by splitting the list Ag from the original term $\text{fapp}[\text{fuva } A \mid \text{Ag}]$. $s_1 \dots s_n$ is the longest prefix of the compiled terms in Ag which is in \mathcal{L}_λ . Therefore, by definition of \mathcal{L}_λ , t_1 must appear in $s_1 \dots s_n$, otherwise $s_1 \dots s_n$ is not the longest prefix in \mathcal{L}_λ , or it is a term with a constructor of tm as functor. \square

E:Dire che maybe eta fa il detect anche su termini che non sono il llambda, oppure dirlo in section of maybeeta + dare un esempio?

Decompilation. During progress, as claimed in invariant 5, the decompilation can only have β -link with not instantiated lhs. In this case, lhs is unified with rhs.

D:not really sure of this, we can have $F \ a = \lambda x.Gx$. In this case when do we fail: for sure in decompile. But to respect fidelity, we should fail immediately: we have a β -link and a η -link with same lhs

8.2 Progress

The activation of a β -link is performed when its rhs falls under \mathcal{L}_λ under a given substitution.

Definition 8.4 (progress-beta- \mathcal{L}_λ). Given a substitution σ and a β -link $\Gamma \vdash T =_\beta X_{s_1 \dots s_n} t_1 \dots t_m$ such that σt_1 is a name, say t , and $t \notin s_1 \dots s_n$. If $m = 0$, then the β -link is removed and lhs is unified with $X_{s_1 \dots s_n}$. If $m > 0$, then the β -link is replaced by a refined version $\Gamma \vdash T =_\beta Y_{s_1 \dots s_n, t} t_2 \dots t_m$ with reduced list of arguments and Y being a fresh variable. Moreover, the new link $\Gamma \vdash X_{s_1 \dots s_n} =_\eta \lambda x. Y_{s_1 \dots s_n, x}$ is added to \mathbb{L} .

Definition 8.5 (progress-beta-rigid-head). A link $\Gamma \vdash X =_\beta X_{s_1 \dots s_n} t_1 \dots t_m$ is removed from \mathbb{L} if $X_{s_1 \dots s_n}$ is instantiated to a term t and the β -reduced term t' obtained from the application of t to $t_1 \dots t_m$ is in \mathcal{L}_λ . Moreover, X is unified to t .

Definition 8.6 (progress-beta-dedup). Given a β -link l_1 and second link $l_2 \in \mathbb{L}$, such that they share the same lhs. In this case, the two rhs are unified and a l_2 is removed from \mathbb{L} .

LEMMA 8.7. progress terminates

PROOF SKETCH. Let l a β -link in the store \mathbb{L} . If l is activated by progress-beta-rigid-head, then it disappears from \mathbb{L} and progress terminates. Otherwise, the rhs of l is made by a variable applied to m arguments. At each activation of progress-beta- \mathcal{L}_λ , l is replaced by a new β -link l^1 having $m - 1$ arguments. At the m^{th} iteration, the β -link l^m has no more arguments and is removed from \mathbb{L} . Note that at the m^{th} iteration, m new η -link have been added to \mathbb{L} , however, by lemma 6.14, the algorithm terminates. Finally progress-beta-dedup also guarantees termination since it makes a unification \mathbb{U} and if \mathbb{U} fails, then progress terminates and if \mathbb{U} succeeds, the recursive calls to progress have a specialized β -link ...

E:funziona. per essere più precisi io parlerei di ordine lessicografico (tipico ordine ben fondato usato per dimostrare terminazione). Nl nostro caso è la tripla (argomenti extra dei beta, numero di beta, numero di eta).

COROLLARY 8.8. Given a β -link, the variables occurring in its rhs are in \mathcal{L}_λ .

D:is it clearer?

PROOF SKETCH. By construction, the rhs of β -link has the shape $X_{s_1 \dots s_n} t_1 \dots t_m$, $s_1 \dots s_n$ is in \mathcal{L}_λ and all the terms $t_1 \dots t_n$ are in \mathcal{L}_λ , too. If a β -link is triggered by progress-beta-rigid-head, then, by definition 8.5, that link is removed by \mathbb{L} , and the property is satisfied. If the η -link is activated by progress-beta- \mathcal{L}_λ , then, by definition 8.4, the new β -link as a variable as a scope which is still in \mathcal{L}_λ .

LEMMA 8.9 (FIDELITY WITH β -link). The introduction of β -link guarantees proposition 2.2

PROOF SKETCH. Let \mathbb{T} a unification problem and σ a substitution such that $\mathbb{T} \in \overline{\mathcal{L}_\lambda}$. If $\sigma \mathbb{T}$ is in \mathcal{L}_λ , then by definitions 8.4 and 8.5, the β -link associated to the subterm of \mathbb{T} have been solved and removed. The unification is done between terms in \mathcal{L}_λ and by lemma 6.15 fidelity is guaranteed. If $\sigma \mathbb{T}$ is in $\overline{\mathcal{L}_\lambda}$, then, by lemma 8.1, the unification fails, as per the corresponding unification in \mathcal{F}_0 . \square

Example of progress-beta- \mathcal{L}_λ . Consider the β -link below:

$$\begin{aligned} \mathbb{P} &= \{ X \approx_o \lambda x.x \quad \lambda x.(Y (X x)) \approx_o f \} \\ \mathbb{T} &= \{ A \approx_\lambda \lambda x.x \quad B \approx_\lambda f \} \\ \mathbb{M} &= \{ Y \mapsto D^0 \quad X \mapsto A^0 \} \\ \mathbb{L} &= \left\{ \begin{array}{l} \vdash A =_\eta \lambda x.E_x \quad \vdash B =_\eta \lambda x.C_x \\ x \vdash C_x =_\beta (D E_x) \end{array} \right\} \end{aligned}$$

Initially the β -link rhs is a variable D applied to the E_x . The first unification problem results in $\sigma = \{A \mapsto \lambda x.x\}$. In turn this instantiation triggers \mathbb{L}_1 by progress- η -left and E_x is assigned to x . Under this substitution the β -link becomes $x \vdash C_x =_\beta (D x)$, and by progress-beta- \mathcal{L}_λ it is replaced with the link: $\vdash E =_\eta \lambda x.D_x$, while C_x is unified with D_x . The second unification problem assigns f to B , that in turn activates the second η -link (f is assigned to C), and then all the remaining links are solved. The final \mathcal{H}_0 substitution is $\sigma = \{A \mapsto \lambda x.x, B \mapsto f, C_x \mapsto (f x), D \mapsto f, E_x \mapsto x, F_x \mapsto C_x\}$ and is decompiled into $\rho = \{X \mapsto \lambda x.x, Y \mapsto f\}$.

Example of progress-beta-rigid-head. We can take the example provided in section 8. The problem is compiled into:

$$\begin{aligned} \mathbb{P} &= \{ X \approx_o \lambda x.Y \quad (X a) \approx_o a \} \\ \mathbb{T} &= \{ A \approx_\lambda \lambda x.B \quad C \approx_\lambda a \} \\ \mathbb{M} &= \{ Y \mapsto B^0 \quad X \mapsto A^0 \} \\ \mathbb{L} &= \{ \vdash C =_\beta (A a) \} \end{aligned}$$

The first unification problems is solved by the substitution $\sigma = \{A \mapsto \lambda x.B\}$. The β -link becomes $\vdash C =_\beta ((\lambda x.B) a)$ whose rhs can be β -reduced to B . B is in \mathcal{L}_λ and is unified with C . The resolution of the second unification problem gives the final substitution $\sigma = \{A \mapsto \lambda x.B, B \mapsto C, C \mapsto a\}$ which is decompiled into $\rho = \{X \mapsto \lambda x.a, Y \mapsto a\}$.

8.3 Relaxing lemma β -link with rigid lhs

Working with terms in \mathcal{L}_λ is sometime too restrictive. There exists systems such as λ Prolog [10], Abella [5], which delay the resolution of $\overline{\mathcal{L}_\lambda}$ unification problems if the substitution is not able to put them in \mathcal{L}_λ .

$$\mathbb{P} = \{ (X a) \approx_o a \quad X \approx_o \lambda x.Y \}$$

In the example above, \mathbb{P}_1 is in $\overline{\mathcal{L}_\lambda}$ and the object language cannot solve it, and, by proposition 2.2, the meta language neither. However, we can be more permissive, and relax lemma 8.1. This modification is quite simple to manage: we are introducing a new $\overline{\mathcal{L}_\lambda}$ progress rule, say progress-beta- $\overline{\mathcal{L}_\lambda}$, by which, if lhs is rigid and rhs is flexible, the considered β -link is kept in the store and no progression is done⁴. progress-beta- $\overline{\mathcal{L}_\lambda}$ makes occur-check-links partial, since the check is possible only on links with a variable on the lhs. This means that we can have two links $\vdash X =_\beta Y a$ and $\vdash f X =_\beta Y a$ where the occur check does not throw an error. Note however, that the decompilation of the two links will force the unification of X to $Y a$ and then the unification of $f (Y a)$ to $Y a$, which fails by the occur check of \approx_λ .

A second strategy to deal with problem that are in $\overline{\mathcal{L}_\lambda}$ is to make some approximation. This is the case for example of the unification algorithm of Coq used in its type class solver [16]. The approximation consists in forcing a choice (among the others) when

⁴This new rule trivially guarantees the termination of progress

the unification problem is in $\overline{\mathcal{L}_\lambda}$. For instance, in $X a b = Y b$, the last argument of the two terms is the same, therefore Y is assigned to Xa . Note that this is of course an approximation, since $\sigma = \{X = \lambda x.Y, Y = _ \}$ is another valid substitution for the original problem. This approximation can be easily introduced in our unification procedure, by adding new custom β -link progress rules.

The commit-link predicate can be extended to add heuristics if during the decompilation phase β -link remain. For example, the same approximation explained above can be delayed and applied only if the terms in $\diamond \beta_0$ never falls in \mathcal{L}_λ after the execution of all the unification problems. We want to point out, that we call this *approximation*, since we are making a choice among all the possible unifiers and therefore, we can pick the wrong one.

9 ACTUAL IMPLEMENTATION IN ELPI

In this paper we show a minimized example. The full code is there. But we also have to code things in Coq-Elpi.

The main difference between the presentation in the previous sections and the actual implementation for Coq is that the main loop `hrun` is replaced by the one of Prolog that chains calls to the unification procedure. In order implement the store of links we resort to Elpi's CLP engine and use constraints (suspended goals) to represent links, and constraint handling rules to implement progress operations involving more than one link.

about the progress of 1 link:

```
link-eta L R :- suspend-condition L R Holes, !,
  declare_constraint (link-eta L R) Holes.
link-eta L R :-
  progress. % e.g. L = R.
```

about the progress of 2 links:

```
constraint link-eta {
  rule (N1 > G1 ?- link-eta (uvar X LX1) T1) % match
    / (N2 > G2 ?- link-eta (uvar X LX2) T2) % remove
    | (relocate LX1 LX2 T2 T2') % condition
    <=> (N1 > G1 ?- T1 = T2'). % new goal
}
```

Remark how the invariant about `uvar` arity makes this easy, since `LX1` and `LX2` have the same length. Also note that `N1` only contains the names of the first link (while `relocate` runs in the disjoint union) and Elpi ensures that `T2'` can live in `N1`.

10 CONCLUSION

Benefits: less work, reuse efficient ho unif (3x faster), indexing,

Future: tabling and static analysis (reuse for ML again).

Very little is Coq specific. Applies to all OL that are not a sub-system of HOL, or for ML that are used for meta programming.

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APPENDIX

This appendix contains the entire code described in this paper. The code can also be accessed at the URL: <https://github.com/FissoreD/paper-ho>

Note that (a infix b) c d de-sugars to (infix) a b c d.

Explain builtin name (can be implemented by loading name after each pi)

11 THE MEMORY

```

kind addr type.
type addr nat -> addr.
typeabbrev (mem A) (list (option A)).

type set? addr -> mem A -> A -> o.
set? (addr A) Mem Val :- get A Mem Val.

type unset? addr -> mem A -> o.
unset? Addr Mem :- not (set? Addr Mem _).

type assign-aux nat -> mem A -> A -> mem A -> o.
assign-aux z (none :: L) Y (some Y :: L).
assign-aux (s N) (X :: L) Y (X :: L1) :- assign-aux N L Y L1.

type assign addr -> mem A -> A -> mem A -> o.
assign (addr A) Mem1 Val Mem2 :- assign-aux A Mem1 Val Mem2.

type get nat -> mem A -> A -> o.
get z (some Y :: _) Y.
get (s N) (_ :: L) X :- get N L X.

type alloc-aux nat -> mem A -> mem A -> o.
alloc-aux z [] [none] :- !.
alloc-aux z L L.
alloc-aux (s N) [] [none | M] :- alloc-aux N [] M.
alloc-aux (s N) [X | L] [X | M] :- alloc-aux N L M.

type alloc addr -> mem A -> mem A -> o.
alloc (addr A as Ad) Mem1 Mem2 :- unset? Ad Mem1,
  alloc-aux A Mem1 Mem2.

type new-aux mem A -> nat -> mem A -> o.
new-aux [] z [none].
new-aux [A | As] (s N) [A | Bs] :- new-aux As N Bs.

type new mem A -> addr -> mem A -> o.
new Mem1 (addr Ad) Mem2 :- new-aux Mem1 Ad Mem2.

```

12 THE OBJECT LANGUAGE

```

kind fm type.
type fapp list fm -> fm.
type flam (fm -> fm) -> fm.
type fcon string -> fm.
type fuva addr -> fm.

typeabbrev fsubst (mem fm).

```

```

type fder fsubst -> fm -> fm -> o.
fder _ (fcon C) (fcon C).
fder S (fapp A) (fapp B) :- map (fder S) A B.
fder S (flam F) (flam G) :-
  pi x\ fder S x x => fder S (F x) (G x).
fder S (fuva N) R :- set? N S T, fder S T R.
fder S (fuva N) (fuva N) :- unset? N S.

type fderef fsubst -> fm -> fm -> o. (ρs)
fderef S T T2 :- fder S T T1, napp T1 T2.

```

```

type (=o) fm -> fm -> o. (=o)
fcon X =o fcon X.
fapp A =o fapp B :- forall2 (=o) A B.
flam F =o flam G :- pi x\ x =o x => F x =o G x.
fuva N =o fuva N.
flam F =o T :- (ηl)
  pi x\ beta T [x] (T' x), x =o x => F x =o T' x.
T =o flam F :- (ηr)
  pi x\ beta T [x] (T' x), x =o x => T' x =o F x.
fapp [flam X | L] =o T :- beta (flam X) L R, R =o T. (βl)
T =o fapp [flam X | L] :- beta (flam X) L R, T =o R. (βr)

```

```

type extend-subst fm -> fsubst -> fsubst -> o.
extend-subst (fuva N) S S' :- mem.alloc N S S'.
extend-subst (flam F) S S' :-
  pi x\ (pi S\ extend-subst x S S) => extend-subst (F x) S S'.
extend-subst (fcon _) S S.
extend-subst (fapp L) S S1 :- fold extend-subst L S S1.

type beta fm -> list fm -> fm -> o.
beta A [] A.
beta (flam Bo) [H | L] R :- napp (Bo H) F, beta F L R.
beta (fapp A) L (fapp X) :- append A L X.
beta (fuva N) L (fapp [fuva N | L]).
beta (fcon H) L (fapp [fcon H | L]).
beta N L (fapp [N | L]) :- name N.

```

```

type napp fm -> fm -> o.
napp (fcon C) (fcon C).
napp (fuva A) (fuva A).
napp (flam F) (flam G) :- pi x\ napp (F x) (G x).
napp (fapp [fapp L1 | L2]) T :- !,
  append L1 L2 L3, napp (fapp L3) T.
napp (fapp L) (fapp L1) :- map napp L L1.
napp N N :- name N.

```

```

type beta-reduce fm -> fm -> o.
beta-reduce (uvar _ _) _ :- halt "Passed uvar to beta-reduce".
beta-reduce A A :- name A.
beta-reduce (fcon A) (fcon A).
beta-reduce (fuva A) (fuva A).
beta-reduce (flam A) (flam B) :-
  pi x\ beta-reduce (A x) (B x).
beta-reduce (fapp [flam B | L]) T2 :- !,

```

```

1741   beta (flam B) L T1, beta-reduce T1 T2.
1742   beta-reduce (fapp L) (fapp L1) :-
1743     map beta-reduce L L1.
1744
1745   type mk-app fm -> list fm -> fm -> o.
1746   mk-app T L S :- beta T L S.
1747
1748   type eta-contract fm -> fm -> o.
1749   eta-contract (fcon X) (fcon X).
1750   eta-contract (fapp L) (fapp L1) :- map eta-contract L L1.
1751   eta-contract (flam F) T :- eta-contract-aux [] (flam F) T.
1752   eta-contract (flam F) (flam F1) :-
1753     pi x\ eta-contract x x => eta-contract (F x) (F1 x).
1754   eta-contract (fuva X) (fuva X).
1755   eta-contract X X :- name X.
1756
1757   type eta-contract-aux list fm -> fm -> fm -> o.
1758   eta-contract-aux L (flam F) T :-
1759     pi x\ eta-contract-aux [x|L] (F x) T. % also checks H Prefix does not x
1760   eta-contract-aux L (fapp [H|Args] T) :-
1761     rev L LRev, append Prefix LRev Args,
1762     if (Prefix = []) (T = H) (T = fapp [H|Prefix]).
1763
1764
1765 13 THE META LANGUAGE
1766   kind inctx type -> type.
1767   type abs (tm -> inctx A) -> inctx A.
1768   type val A -> inctx A.
1769   typeabbrev assignment (inctx tm).
1770   typeabbrev subst (mem assignment).
1771
1772   kind tm type.
1773   type app list tm -> tm.
1774   type lam (tm -> tm) -> tm.
1775   type con string -> tm.
1776   type uva addr -> list tm -> tm.
1777
1778   type ( $\approx_\lambda$ ) tm -> tm -> subst -> subst -> o.
1779   (con C  $\approx_\lambda$  con C) S S.
1780   (app L1  $\approx_\lambda$  app L2) S S1 :- fold2 ( $\approx_\lambda$ ) L1 L2 S S1.
1781   (lam F1  $\approx_\lambda$  lam F2) S S1 :-
1782     pi x\ (pi S\ (x  $\approx_\lambda$  x) S S) => (F1 x  $\approx_\lambda$  F2 x) S S1.
1783   (uva N Args  $\approx_\lambda$  T) S S1 :-
1784     set? N S F,!, move F Args T1, (T1  $\approx_\lambda$  T) S S1.
1785   (T  $\approx_\lambda$  uva N Args) S S1 :-
1786     set? N S F,!, move F Args T1, (T  $\approx_\lambda$  T1) S S1.
1787   (uva M A1  $\approx_\lambda$  uva N A2) S1 S2 :- !,
1788     pattern-fragment A1, pattern-fragment A2,
1789     prune! M A1 N A2 S1 S2.
1790   (uva N Args  $\approx_\lambda$  T) S S1 :- not_occ N S T, pattern-fragment Args,
1791     bind T Args T1, assign N S T1 S1.
1792   (T  $\approx_\lambda$  uva N Args) S S1 :- not_occ N S T, pattern-fragment Args,
1793     bind T Args T1, assign N S T1 S1.
1794
1795   type prune! addr -> list tm -> addr ->
1796     list tm -> subst -> subst -> o.
1797   /* no pruning needed */
1798

```

```

1799   prune! N A N A S S :- !.
1800   prune! M A N A S1 S2 :- !, bind (uva M A) A Ass,
1801     assign N S1 Ass S2.
1802   /* prune different arguments */
1803   prune! N A1 N A2 S1 S3 :- !,
1804     new S1 W S2, prune-same-variable W A1 A2 [] Ass,
1805     assign N S2 Ass S3.
1806   /* prune to the intersection of scopes */
1807   prune! N A1 M A2 S1 S4 :- !,
1808     new S1 W S2, prune-diff-variables W A1 A2 Ass1 Ass2,
1809     assign N S2 Ass1 S3,
1810     assign M S3 Ass2 S4.
1811
1812   type prune-same-variable addr -> list tm -> list tm ->
1813     list tm -> assignment -> o.
1814   prune-same-variable N [] [] ACC (val (uva N Args)) :-
1815     rev ACC Args.
1816   prune-same-variable N [X|XS] [X|YS] ACC (abs F) :-
1817     pi x\ prune-same-variable N XS YS [x|ACC] (F x).
1818   prune-same-variable N [_|XS] [_|YS] ACC (abs F) :-
1819     pi x\ prune-same-variable N XS YS ACC (F x).
1820
1821   type permute list nat -> list tm -> list tm -> o.
1822   permute [] _ [].
1823   permute [P|PS] Args [T|TS] :-
1824     nth P Args T,
1825     permute PS Args TS.
1826
1827   type build-perm-assign addr -> list tm -> list bool ->
1828     list nat -> assignment -> o.
1829   build-perm-assign N ArgsR [] Perm (val (uva N PermutedArgs)) :-
1830     rev ArgsR Args, permute Perm Args PermutedArgs.
1831   build-perm-assign N Acc [tt|L] Perm (abs T) :-
1832     pi x\ build-perm-assign N [x|Acc] L Perm (T x).
1833   build-perm-assign N Acc [ff|L] Perm (abs T) :-
1834     pi x\ build-perm-assign N Acc L Perm (T x).
1835
1836   type keep list A -> A -> bool -> o.
1837   keep L A tt :- mem L A, !.
1838   keep _ _ ff.
1839
1840   type prune-diff-variables addr -> list tm -> list tm ->
1841     assignment -> assignment -> o.
1842   prune-diff-variables N Args1 Args2 Ass1 Ass2 :-
1843     map (keep Args2) Args1 Bits1,
1844     map (keep Args1) Args2 Bits2,
1845     filter Args1 (mem Args2) ToKeep1,
1846     filter Args2 (mem Args1) ToKeep2,
1847     map (index ToKeep1) ToKeep1 IdPerm,
1848     map (index ToKeep1) ToKeep2 Perm21,
1849     build-perm-assign N [] Bits1 IdPerm Ass1,
1850     build-perm-assign N [] Bits2 Perm21 Ass2.
1851
1852   type beta tm -> list tm -> tm -> o.
1853   beta A [] A :- !.
1854   beta (lam Bo) [H | L] R :- beta (Bo H) L R1, beta-aux R1 R.
1855   beta (app A) L (app X) :- append A L X.
1856

```



```

1857 beta (con H) L (app [con H | L]).
1858 beta X L (app[X|L]) :- name X.
1859
1860 type beta-aux tm -> tm -> o.
1861 beta-aux (app [HD|TL]) R :- !, beta HD TL R.
1862 beta-aux A A.
1863
1864 /* occur check for N before crossing a functor */
1865 type not_occ addr -> subst -> tm -> o.
1866 not_occ N S (uva M Args) :- set? M S F,
1867     move F Args T, not_occ N S T.
1868 not_occ N S (uva M Args) :- unset? M S, not (M = N),
1869     forall1 (not_occ_aux N S) Args.
1870 not_occ _ _ (con _).
1871 not_occ N S (app L) :- not_occ_aux N S (app L).
1872 /* Note: lam is a functor for the meta language! */
1873 not_occ N S (lam L) :- pi x\ not_occ_aux N S (L x).
1874 not_occ _ _ X :- name X.
1875 /* finding N is ok */
1876 not_occ N _ (uva N _).
1877
1878 /* occur check for X after crossing a functor */
1879 type not_occ_aux addr -> subst -> tm -> o.
1880 not_occ_aux N S (uva M _) :- unset? M S, not (N = M).
1881 not_occ_aux N S (uva M Args) :- set? M S F,
1882     move F Args T, not_occ_aux N S T.
1883 not_occ_aux N S (app L) :- forall1 (not_occ_aux N S) L.
1884 not_occ_aux N S (lam F) :- pi x\ not_occ_aux N S (F x).
1885 not_occ_aux _ _ (con _).
1886 not_occ_aux _ _ X :- name X.
1887 /* finding N is ko, hence no rule */
1888
1889 /* copy T T' fails if T contains a free variable, i.e. it
1890     performs scope checking for bind */
1891 type copy tm -> tm -> o.
1892 copy (con C) (con C).
1893 copy (app L) (app L') :- map copy L L'.
1894 copy (lam T) (lam T') :- pi x\ copy x x => copy (T x) (T' x).
1895 copy (uva A L) (uva A L') :- map copy L L'.
1896
1897 type bind tm -> list tm -> assignment -> o.
1898 bind T [] (val T') :- copy T T'.
1899 bind T [X | TL] (abs T') :- pi x\ copy X x => bind T TL (T' x).
1900
1901 type deref subst -> tm -> tm -> o. (σt)
1902 deref _ (con C) (con C).
1903 deref S (app A) (app B) :- map (deref S) A B.
1904 deref S (lam F) (lam G) :-
1905     pi x\ deref S x x => deref S (F x) (G x).
1906 deref S (uva N L) R :- set? N S A,
1907     move A L T, deref S T R.
1908 deref S (uva N A) (uva N B) :- unset? N S,
1909     map (deref S) A B.
1910
1911 type move assignment -> list tm -> tm -> o.
1912 move (abs Bo) [H|L] R :- move (Bo H) L R.
1913 move (val A) [] A.

```

```

type deref-assmt subst -> assignment -> assignment -> o.

```

```

deref-assmt S (abs T) (abs R) :- pi x\ deref-assmt S (T x) (R x).

```

```

deref-assmt S (val T) (val R) :- deref S T R.

```

14 THE COMPILER

```

kind arity type.

```

```

type arity nat -> arity.

```

```

kind fvariable type.

```

```

type fv addr -> fvariable.

```

```

kind hvariable type.

```

```

type hv addr -> arity -> hvariable.

```

```

kind mapping type.

```

```

type mapping fvariable -> hvariable -> mapping.

```

```

typeabbrev mmap (list mapping).

```

```

typeabbrev scope (list tm).

```

```

typeabbrev inctx ho.inctx.

```

```

kind baselink type.

```

```

type link-eta tm -> tm -> baselink.

```

```

type link-beta tm -> tm -> baselink.

```

```

typeabbrev link (inctx baselink).

```

```

typeabbrev links (list link).

```

```

macro @val-link-eta T1 T2 :- ho.val (link-eta T1 T2).

```

```

macro @val-link-beta T1 T2 :- ho.val (link-beta T1 T2).

```

```

type get-lhs link -> tm -> o.

```

```

get-lhs (val (link-beta A _)) A.

```

```

get-lhs (val (link-eta A _)) A.

```

```

type get-rhs link -> tm -> o.

```

```

get-rhs (val (link-beta _ A)) A.

```

```

get-rhs (val (link-eta _ A)) A.

```

```

type occurs-rigidly fm -> fm -> o.

```

```

occurs-rigidly N N.

```

```

occurs-rigidly _ (fapp [fuva _|_]) :- !, fail.

```

```

occurs-rigidly N (fapp L) :- exists (occurs-rigidly N) L.

```

```

occurs-rigidly N (flam B) :- pi x\ occurs-rigidly N (B x).

```

```

type reducible-to list fm -> fm -> fm -> o.

```

```

reducible-to _ N N :- !.

```

```

reducible-to L N (fapp [fuva _|Args]) :- !,

```

```
    forall1 (x\ exists (reducible-to [] x) Args) [N|L].

```

```

reducible-to L N (flam B) :- !,

```

```
    pi x\ reducible-to [x | L] N (B x).

```

```

reducible-to L N (fapp [N|Args]) :-

```

```
    last-n {len L} Args R,

```

```
    forall2 (reducible-to []) R {rev L}.

```

```

1973
1974 type maybe-eta fm -> list fm -> o. (◇η)
1975 maybe-eta (fapp[fuva _]Args] L :- !,
1976   forall1 (x\ exists (reducible-to [] x) Args) L, !.
1977 maybe-eta (flam B) L :- !, pi x\ maybe-eta (B x) [x | L].
1978 maybe-eta (fapp [T]Args] L :- (name T; T = fcon _),
1979   split-last-n {len L} Args First Last,
1980   none (x\ exists (y\ occurs-rigidly x y) First) L,
1981   forall2 (reducible-to []) {rev L} Last.
1982
1983
1984 type locally-bound tm -> o.
1985 type get-scope-aux tm -> list tm -> o.
1986 get-scope-aux (con _) [].
1987 get-scope-aux (uva _ L) L1 :-
1988   forall2 get-scope-aux L R,
1989   flatten R L1.
1990 get-scope-aux (lam B) L1 :-
1991   pi x\ locally-bound x => get-scope-aux (B x) L1.
1992 get-scope-aux (app L) L1 :-
1993   forall2 get-scope-aux L R,
1994   flatten R L1.
1995 get-scope-aux X [X] :- name X, not (locally-bound X).
1996 get-scope-aux X [] :- name X, (locally-bound X).
1997
1998 type names1 list tm -> o.
1999 names1 L :-
2000   names L1,
2001   new_int N,
2002   if (1 is N mod 2) (L1 = L) (rev L1 L).
2003
2004 type get-scope tm -> list tm -> o.
2005 get-scope T Scope :-
2006   get-scope-aux T ScopeDuplicata,
2007   undup ScopeDuplicata Scope.
2008 type rigid fm -> o.
2009 rigid X :- not (X = fuva _).
2010
2011 type comp-lam (fm -> fm) -> (tm -> tm) ->
2012   mmap -> mmap -> links -> links -> subst -> subst -> o.
2013 comp-lam F G M1 M2 L1 L3 S1 S2 :-
2014   pi x y\ (pi M L S\ comp x y M M L L S S) =>
2015   comp (F x) (G y) M1 M2 L1 (L2 y) S1 S2,
2016   close-links L2 L3.
2017
2018 type close-links (tm -> links) -> links -> o.
2019 close-links (v\ [X | L v]) [X|R] :- !, close-links L R.
2020 close-links (v\ [X v | L v]) [abs X|R] :- close-links L R.
2021 close-links (_[]) [].
2022 type comp fm -> tm -> mmap -> mmap -> links -> links ->
2023   subst -> subst -> o.
2024 comp (fcon C) (con C) M M L L S S.
2025 comp (flam F) (uva A Scope) M1 M2 L1 L3 S1 S3 :-
2026   maybe-eta (flam F) [], !,
2027   alloc S1 A S2,
2028   comp-lam F F1 M1 M2 L1 L2 S2 S3,
2029   get-scope (lam F1) Scope,
2030
2031   L3 = [val (link-eta (uva A Scope) (lam F1)) | L2].
2032 comp (flam F) (lam F1) M1 M2 L1 L2 S1 S2 :- (cλ)
2033   comp-lam F F1 M1 M2 L1 L2 S1 S2.
2034 comp (fuva A) (uva B []) M1 M2 L L S1 S2 :-
2035   m-alloc (fv A) (hv B (arity z)) M1 M2 S1 S2.
2036 comp (fapp [fuva A]Ag] (uva B Ag1) M1 M2 L L S1 S2 :-
2037   pattern-fragment Ag, !,
2038   fold6 comp Ag Ag1 M1 M1 L L S1 S1,
2039   len Ag Arity,
2040   m-alloc (fv A) (hv B (arity Arity)) M1 M2 S1 S2.
2041 comp (fapp [fuva A]Ag] (uva B Scope) M1 M3 L1 L3 S1 S4 :- !,
2042   pattern-fragment-prefix Ag Pf Extra,
2043   len Pf Arity,
2044   alloc S1 B S2,
2045   m-alloc (fv A) (hv C (arity Arity)) M1 M2 S2 S3,
2046   fold6 comp Pf Pf1 M2 M2 L1 L1 S3 S3,
2047   fold6 comp Extra Extra1 M2 M3 L1 L2 S3 S4,
2048   Beta = app [uva C Pf1 | Extra1],
2049   get-scope Beta Scope,
2050   L3 = [val (link-beta (uva B Scope) Beta) | L2].
2051 comp (fapp A) (app A1) M1 M2 L1 L2 S1 S2 :- (c@)
2052   fold6 comp A A1 M1 M2 L1 L2 S1 S2.
2053
2054 type alloc mem A -> addr -> mem A -> o.
2055 alloc S N S1 :- mem.new S N S1.
2056
2057 type compile-terms-diagnostic
2058   triple diagnostic fm fm ->
2059   triple diagnostic tm tm ->
2060   mmap -> mmap ->
2061   links -> links ->
2062   subst -> subst -> o.
2063 compile-terms-diagnostic (triple D F01 F02) (triple D H01 H02) M1 M3 M1 L
2064   fo.beta-reduce F01 F01',
2065   fo.beta-reduce F02 F02',
2066   comp F01' H01 M1 M2 L1 L2 S1 S2,
2067   comp F02' H02 M2 M3 L2 L3 S2 S3.
2068
2069 type compile-terms
2070   list (triple diagnostic fm fm) ->
2071   list (triple diagnostic tm tm) ->
2072   mmap -> links -> subst -> o.
2073 compile-terms T H M L S :-
2074   fold6 compile-terms-diagnostic T H [] M_ [] L_ [] S_,
2075   print-compil-result T H L_ M_,
2076   deduplicate-map M_ M S_ S L_ L.
2077
2078 type make-eta-link-aux nat -> addr -> addr ->
2079   list tm -> links -> subst -> subst -> o.
2080 make-eta-link-aux z Ad1 Ad2 Scope1 L H1 H1 :-
2081   rev Scope1 Scope, eta-expand (uva Ad2 Scope) T1,
2082   L = [val (link-eta (uva Ad1 Scope) T1)].
2083 make-eta-link-aux (s N) Ad1 Ad2 Scope1 L H1 H3 :-
2084   rev Scope1 Scope, alloc H1 Ad H2,
2085   eta-expand (uva Ad Scope) T2,
2086   (pi x\ make-eta-link-aux N Ad Ad2 [x|Scope1] (L1 x) H2 H3),
2087   close-links L1 L2,
2088

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2089   L = [val (link-eta (uva Ad1 Scope) T2) | L2].
2090
2091   type make-eta-link nat -> nat -> addr -> addr ->
2092         list tm -> links -> subst -> subst -> o.
2093   make-eta-link (s N) z Ad1 Ad2 Vars L H H1 :-
2094     make-eta-link-aux N Ad2 Ad1 Vars L H H1.
2095   make-eta-link z (s N) Ad1 Ad2 Vars L H H1 :-
2096     make-eta-link-aux N Ad1 Ad2 Vars L H H1.
2097   make-eta-link (s N) (s M) Ad1 Ad2 Vars Links H H1 :-
2098     (pi x\ make-eta-link N M Ad1 Ad2 [x|Vars] (L x) H H1),
2099     close-links L Links.
2100
2101   type deduplicate-map mmap -> mmap ->
2102         subst -> subst -> links -> links -> o.
2103   deduplicate-map [] [] H H L L.
2104   deduplicate-map [(mapping (fv 0) (hv M (arity LenM)) as X1) | Map1]
2105     take-list Map1 (mapping (fv 0) (hv M' (arity LenM'))), !,
2106     std.assert! (not (LenM = LenM')) "Deduplicate map, there is a bug"
2107     print "arity-fix links:" {ppmapping X1} "~!~" {ppmapping (mapping (fv 0) (hv M' (arity LenM')))},
2108     make-eta-link LenM LenM' M M' [] New H1 H2,
2109     print "new eta link" {pplinks New},
2110     append New L1 L2,
2111     deduplicate-map Map1 Map2 H2 H3 L2 L3.
2112   deduplicate-map [A|As] [A|Bs] H1 H2 L1 L2 :-
2113     deduplicate-map As Bs H1 H2 L1 L2, !.
2114   deduplicate-map [A|_] _ H _ _ :-
2115     halt "deduplicating mapping error" {ppmapping A} {ho.ppsubst H}.

```

15 THE PROGRESS FUNCTION

```

2119   macro @one :- s z.
2120
2121   type contract-rigid list ho.tm -> ho.tm -> ho.tm -> o.
2122   contract-rigid L (ho.lam F) T :-
2123     pi x\ contract-rigid [x|L] (F x) T. % also checks H Prefix does not
2124   contract-rigid L (ho.app [H|Args]) T :-
2125     rev L LRev, append Prefix LRev Args,
2126     if (Prefix = []) (T = H) (T = ho.app [H|Prefix]).
2127
2128   type progress-eta-link ho.tm -> ho.tm -> ho.subst -> ho.subst -> link-eta
2129   progress-eta-link (ho.app _ as T) (ho.lam x\ _ as T1) H H1 [] :- !, not (T1 = ho.uva _), !, fail.
2130     ({eta-expand T @one} ==1 T1) H H1.
2131   progress-eta-link (ho.con _ as T) (ho.lam x\ _ as T1) H H1 [] :- !
2132     ({eta-expand T @one} ==1 T1) H H1.
2133   progress-eta-link (ho.lam _ as T) T1 H H1 [] :- !,
2134     (T ==1 T1) H H1.
2135   progress-eta-link (ho.uva _ _ as X) T H H1 [] :-
2136     contract-rigid [] T T1, !, (X ==1 T1) H H1.
2137   progress-eta-link (ho.uva Ad _ as T1) T2 H H [@eval-link-eta T1 T2] :-
2138     if (ho.not_occ Ad H T2) true fail.
2139
2140   type is-in-pf ho.tm -> o.
2141   is-in-pf (ho.app [ho.uva _ _ | _]) :- !, fail.
2142   is-in-pf (ho.lam B) :- !, pi x\ is-in-pf (B x).
2143   is-in-pf (ho.con _).
2144   is-in-pf (ho.app L) :- forall1 is-in-pf L.
2145   is-in-pf N :- name N.

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2147   is-in-pf (ho.uva _ L) :- pattern-fragment L.
2148
2149   type arity ho.tm -> nat -> o.
2150   arity (ho.con _) z.
2151   arity (ho.app L) A :- len L A.
2152
2153   type occur-check-err ho.tm -> ho.tm -> ho.subst -> o.
2154   occur-check-err (ho.con _) _ _ :- !.
2155   occur-check-err (ho.app _) _ _ :- !.
2156   occur-check-err (ho.lam _) _ _ :- !.
2157   occur-check-err (ho.uva Ad _) T S :-
2158     not (ho.not_occ Ad S T).
2159
2160   type progress-beta-link-aux ho.tm -> ho.tm ->
2161         ho.subst -> ho.subst -> links -> o.
2162   progress-beta-link-aux T1 T2 S1 S2 [] :- is-in-pf T2, !,
2163     (T1 ==1 T2) S1 S2.
2164   progress-beta-link-aux T1 T2 S S [@eval-link-beta T1 T2] :- !.
2165   type progress-beta-link ho.tm -> ho.tm -> ho.subst ->
2166         ho.subst -> links -> o.
2167   progress-beta-link T (ho.app[ho.uva V Scope | L] as T2) S S2 [@eval-link-
2168     arity T Arity, len L ArgsNb, ArgsNb >n Arity, !,
2169     minus ArgsNb Arity Diff, mem.new S V1 S1,
2170     eta-expand (ho.uva V1 Scope) Diff T1,
2171     ((ho.uva V Scope) ==1 T1) S1 S2.
2172   progress-beta-link (ho.uva _ _ as T) (ho.app[ho.uva Ad1 Scope1 | L1] as
2173     append Scope1 L1 Scope1L,
2174     pattern-fragment-prefix Scope1L Scope2 L2,
2175     not (Scope1 = Scope2), !,
2176     mem.new S1 Ad2 S2,
2177     len Scope1 Scope1Len,
2178     len Scope2 Scope2Len,
2179     make-eta-link Scope1Len Scope2Len Ad1 Ad2 [] LinkEta S2 S3,
2180     if (L2 = []) (NewLinks = LinkEta, T2 = ho.uva Ad2 Scope2)
2181     (T2 = ho.app [ho.uva Ad2 Scope2 | L2],
2182     NewLinks = [@eval-link-beta T T2 | LinkEta]).
2183   progress-beta-link T1 (ho.app[ho.uva _ _ | _] as T2) _ _ _ :-
2184     progress-beta-link (ho.uva _ _ as T) (ho.app[ho.uva _ _ | _] as T2) S1
2185     occur-check-err T T2 S1, !, fail.
2186   progress-beta-link T1 (ho.app[ho.uva _ _ | _] as T2) H H [@eval-link-beta
2187     progress-beta-link T1 (ho.app [Hd | T1]) S1 S2 B :-
2188     ho.lam beta Hd T1 T3,
2189     progress-beta-link-aux T1 T3 S1 S2 B.
2190
2191   type solve-link-abs link -> links -> ho.subst -> ho.subst -> o.
2192   solve-link-abs (ho.abs X) R H H1 :-
2193     pi x\ ho.copy x x => (pi S\ ho.deref S x x) =>
2194     solve-link-abs (X x) (R' x) H H1,
2195     close-links R' R.

```

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2205 solve-link-abs (@val-link-eta A B) NewLinks S S1 :- !,
2206   progress-eta-link A B S S1 NewLinks.
2207
2208 solve-link-abs (@val-link-beta A B) NewLinks S S1 :- !,
2209   progress-beta-link A B S S1 NewLinks.
2210
2211 type take-link link -> links -> link -> links -> o.
2212 take-link A [B|XS] B XS :- link-abs-same-lhs A B, !.
2213 take-link A [L|XS] B [L|YS] :- take-link A XS B YS.
2214
2215 type link-abs-same-lhs link -> link -> o.
2216 link-abs-same-lhs (ho.abs F) B :-
2217   pi x\ link-abs-same-lhs (F x) B.
2218 link-abs-same-lhs A (ho.abs G) :-
2219   pi x\ link-abs-same-lhs A (G x).
2220 link-abs-same-lhs (@val-link-eta (ho.uva N _) _) (@val-link-eta (ho.uva N S2) B) H H1 :-
2221
2222 type same-link-eta link -> link -> ho.subst -> ho.subst -> o.
2223 same-link-eta (ho.abs F) B H H1 :- !, pi x\ same-link-eta (F x) B H H1 -> fm L (ho.uva VM TL) T :- mem L (mapping (fv V0) (hv VM _)),
2224 same-link-eta A (ho.abs G) H H1 :- !, pi x\ same-link-eta A (G x) H H1 -> fm L (ho.uva VM TL) T :- mem L (mapping (fv V0) (hv VM _)),
2225 same-link-eta (@val-link-eta (ho.uva N S1) A)
2226   (@val-link-eta (ho.uva N S2) B) H H1 :-
2227   std.map2 S1 S2 (x\y\r\ r = ho.copy x y) Perm,
2228   Perm => ho.copy A A',
2229   (A' ==l B) H H1.
2230
2231 type progress1 links -> links -> ho.subst -> ho.subst -> o.
2232 progress1 [] [] X X.
2233 progress1 [A|L1] [A|L3] S S2 :- take-link A L1 B L2, !,
2234   same-link-eta A B S S1,
2235   progress1 L2 L3 S1 S2.
2236 progress1 [L0|L1] L3 S S2 :- deref-link S L0 L,
2237   solve-link-abs L R S S1, !,
2238   progress1 L1 L2 S1 S2, append R L2 L3.
2239
2240
2241
2242 type abs->lam ho.assignment -> ho.tm -> o.
2243 abs->lam (ho.abs T) (ho.lam R) :- !, pi x\ abs->lam (T x) (R x).
2244 abs->lam (ho.val A) A.
2245
2246 type commit-links-aux link -> ho.subst -> ho.subst -> o.
2247 commit-links-aux (@val-link-eta T1 T2) H1 H2 :-
2248   ho.deref H1 T1 T1', ho.deref H1 T2 T2',
2249   (T1' ==l T2') H1 H2.
2250 commit-links-aux (@val-link-beta T1 T2) H1 H2 :-
2251   ho.deref H1 T1 T1', ho.deref H1 T2 T2',
2252   (T1' ==l T2') H1 H2.
2253 commit-links-aux (ho.abs B) H H1 :-
2254   pi x\ commit-links-aux (B x) H H1.
2255
2256 type commit-links links -> links -> ho.subst -> ho.subst -> o.
2257 commit-links [] [] H H.
2258 commit-links [Abs | Links] L H H2 :-
2259   commit-links-aux Abs H H1, !, commit-links Links L H1 H2.
2260
2261 type decomp1-subst map -> map -> ho.subst ->
2262
2263   fo.fsubst -> fo.fsubst -> o.
2264 decomp1-subst _ [A|_] _ _ _ :- fail.
2265 decomp1-subst _ [] _ F F.
2266 decomp1-subst Map [mapping (fv V0) (hv VM _)|T1] H F F2 :-
2267   mem.set? VM H T, !,
2268   ho.deref-assmt H T TTT,
2269   abs->lam TTT T', tm->fm Map T' T1,
2270   fo.eta-contract T1 T2, mem.assign V0 F T2 F1,
2271   decomp1-subst Map T1 H F1 F2.
2272 decomp1-subst Map [mapping _ (hv VM _)|T1] H F F2 :-
2273   mem.unset? VM H, decomp1-subst Map T1 H F F2.
2274
2275 type tm->fm map -> ho.tm -> fo.fm -> o.
2276 tm->fm _ (ho.con C) (fo.fcon C).
2277 tm->fm L (ho.lam B1) (fo.flam B2) :-
2278   tm->fm L (ho.app N1 y) tm->fm _ x y => tm->fm L (B1 x) (B2 y).
2279 tm->fm L (ho.app L1) T :- map (tm->fm L) L1 [Hd|T1],
2280   fo.mk-app Hd T1 T.
2281 tm->fm L (ho.uva VM TL) T :- mem L (mapping (fv V0) (hv VM _)),
2282   map (tm->fm L) TL T1, fo.mk-app (fo.fuva V0) T1 T.
2283
2284 type add-new-map-aux ho.subst -> list ho.tm -> map ->
2285   map -> fo.fsubst -> fo.fsubst -> o.
2286 add-new-map-aux _ [] _ [] S S.
2287 add-new-map-aux H [T|Ts] L L2 S S2 :-
2288   add-new-map H T L L1 S S1,
2289   add-new-map-aux H Ts L1 L2 S1 S2.
2290
2291 type add-new-map ho.subst -> ho.tm -> map ->
2292   map -> fo.fsubst -> fo.fsubst -> o.
2293 add-new-map _ (ho.uva N _) Map [] F1 F1 :-
2294   mem Map (mapping _ (hv N _)), !.
2295 add-new-map H (ho.uva N L) Map [Map1 | MapL] F1 F3 :-
2296   mem.new F1 M F2,
2297   len L Arity, Map1 = mapping (fv M) (hv N (arity Arity)),
2298   add-new-map H (ho.app L) [Map1 | MapL] MapL F2 F3.
2299 add-new-map H (ho.lam B) Map NewMap F1 F2 :-
2300   pi x\ add-new-map H (B x) Map NewMap F1 F2.
2301 add-new-map H (ho.app L) Map NewMap F1 F3 :-
2302   add-new-map-aux H L Map NewMap F1 F3.
2303 add-new-map _ (ho.con _) _ [] F F :- !.
2304 add-new-map _ N _ [] F F :- name N.
2305
2306 type complete-mapping-under-ass ho.subst -> ho.assignment ->
2307   map -> map -> fo.fsubst -> fo.fsubst -> o.
2308 complete-mapping-under-ass H (ho.val Val) Map1 Map2 F1 F2 :-
2309   add-new-map H Val Map1 Map2 F1 F2.
2310 complete-mapping-under-ass H (ho.abs Abs) Map1 Map2 F1 F2 :-
2311   pi x\ complete-mapping-under-ass H (Abs x) Map1 Map2 F1 F2.
2312
2313 type complete-mapping ho.subst -> ho.subst ->
2314   map -> map -> fo.fsubst -> fo.fsubst -> o.
2315 complete-mapping _ [] L L F F.
2316 complete-mapping H [none | T1] L1 L2 F1 F2 :-
2317   complete-mapping H T1 L1 L2 F1 F2.
2318 complete-mapping H [some T0 | T1] L1 L3 F1 F3 :-
2319   ho.deref-assmt H T0 T,
2320

```



```

2321   complete-mapping-under-ass H T L1 L2 F1 F2,
2322   append L1 L2 LAll,
2323   complete-mapping H T1 LAll L3 F2 F3.
2324

```

```

2325   type decompile map -> links -> ho.subst ->
2326   fo.fsubst -> fo.fsubst -> o.
2327   decompile Map1 L H0 F0 F02 :-
2328   commit-links L L1_ H0 H01, !,
2329   complete-mapping H01 H01 Map1 Map2 F0 F01,
2330   decomp1-subst Map2 Map2 H01 F01 F02.
2331

```

17 AUXILIARY FUNCTIONS

```

2333   type fold4 (A -> A1 -> B -> B -> C -> C -> o) -> list A ->
2334   list A1 -> B -> B -> C -> C -> o.
2335   fold4 _ [] [] A A B B.
2336   fold4 F [X|XS] [Y|YS] A A1 B B1 :- F X Y A A0 B B0,
2337   fold4 F XS YS A0 A1 B0 B1.
2338
2339   type len list A -> nat -> o.
2340   len [] z.
2341   len [_|L] (s X) :- len L X.
2342

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