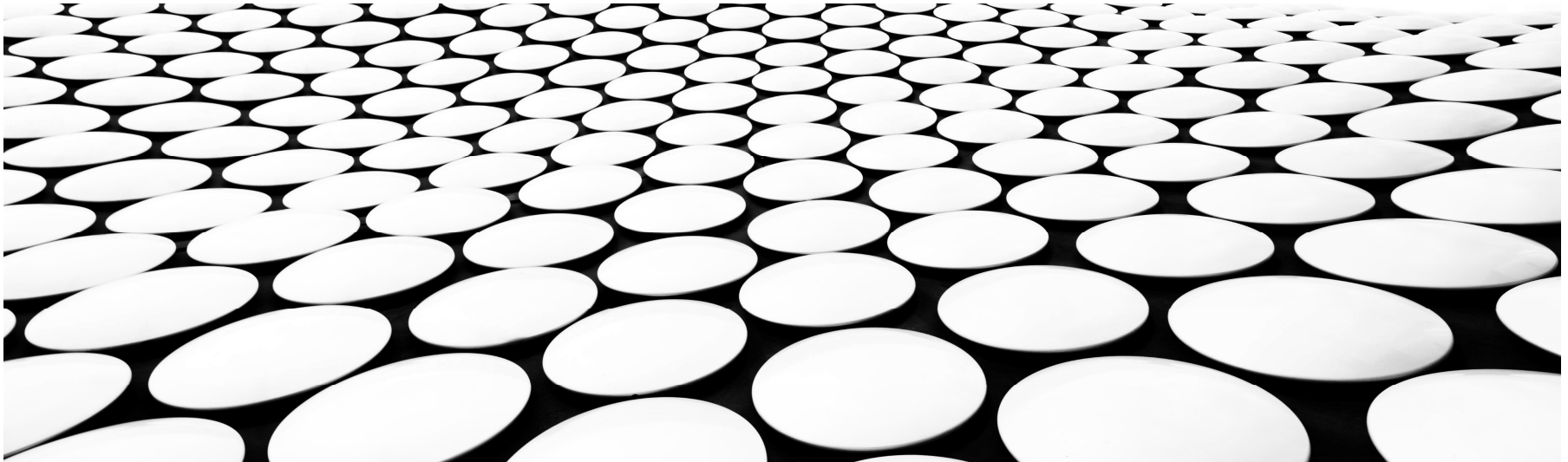


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# FINITE AUTOMATA

IF 2124 TEORI BAHASA FORMAL OTOMATA



## Motivation

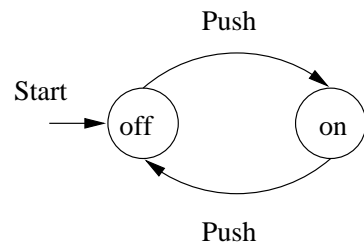
- Automata = abstract computing devices
- Turing studied Turing Machines (= computers) before there were any real computers
- We will also look at simpler devices than Turing machines (Finite State Automata, Push-down Automata, . . . ), and specification means, such as grammars and regular expressions.
- NP-hardness = what cannot be efficiently computed

## **Finite Automata**

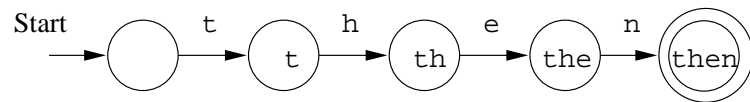
Finite Automata are used as a model for

- Software for designing digital circuits
- Lexical analyzer of a compiler
- Searching for keywords in a file or on the web.
- Software for verifying finite state systems, such as communication protocols.

- Example: Finite Automaton modelling an on/off switch



- Example: Finite Automaton recognizing the string then



## Central Concepts

**Alphabet:** Finite, nonempty set of symbols

Example:  $\Sigma = \{0, 1\}$  binary alphabet

Example:  $\Sigma = \{a, b, c, \dots, z\}$  the set of all lower case letters

Example: The set of all ASCII characters

**Strings:** Finite sequence of symbols from an alphabet  $\Sigma$ , e.g. 0011001

**Empty String:** The string with zero occurrences of symbols from  $\Sigma$

- The empty string is denoted  $\epsilon$

**Length of String:** Number of positions for symbols in the string.

$|w|$  denotes the length of string  $w$

$$|0110| = 4, |\epsilon| = 0$$

**Powers of an Alphabet:**  $\Sigma^k$  = the set of strings of length  $k$  with symbols from  $\Sigma$

Example:  $\Sigma = \{0, 1\}$        $\Sigma = \{0, 1\}$

$\Sigma^1 = \{0, 1\}$        $\Sigma^1 = \{0, 1\}$

$\Sigma^2 = \{00, 01, 10, 11\}$

$\Sigma^0 = \{\epsilon\}$        $\Sigma^3 = \{000, 001, 010, 011, \dots, 111\}$

**Question:** How many strings are there in  $\Sigma^3$

The set of all strings over  $\Sigma$  is denoted  $\Sigma^*$

$$\Sigma^* = \Sigma^0 \cup \Sigma^1 \cup \Sigma^2 \cup \dots$$

Also:

$$\Sigma^+ = \Sigma^1 \cup \Sigma^2 \cup \Sigma^3 \cup \dots$$

$$\Sigma^* = \Sigma^+ \cup \{\epsilon\}$$

**Concatenation:** If  $x$  and  $y$  are strings, then  $xy$  is the string obtained by placing a copy of  $y$  immediately after a copy of  $x$

$$x = a_1a_2 \dots a_i, y = b_1b_2 \dots b_j$$

$$xy = a_1a_2 \dots a_ib_1b_2 \dots b_j$$

Example:  $x = 01101, y = 110, xy = 01101110$

**Note:** For any string  $x$

$$x\epsilon = \epsilon x = x$$

## Languages:

If  $\Sigma$  is an alphabet, and  $L \subseteq \Sigma^* = \{\epsilon, 0, 1, 00, 01, 10, 11, 000, 001, \dots\}$ , then  $L$  is a language

Examples of languages:

- The set of legal English words
- The set of legal C programs
- The set of strings consisting of  $n$  0's followed by  $n$  1's

$\{\epsilon, 01, 0011, 000111, \dots\}$



- The set of strings with equal number of 0's and 1's

$$\{\epsilon, 01, 10, 0011, 0101, 1001, \dots\}$$

- $L_P$  = the set of binary numbers whose value is prime

$$\{10, 11, 101, 111, 1011, \dots\}$$

- The empty language  $\emptyset$
- The language  $\{\epsilon\}$  consisting of the empty string

**Note:**  $\emptyset \neq \{\epsilon\}$

**Note2:** The underlying alphabet  $\Sigma$  is always finite