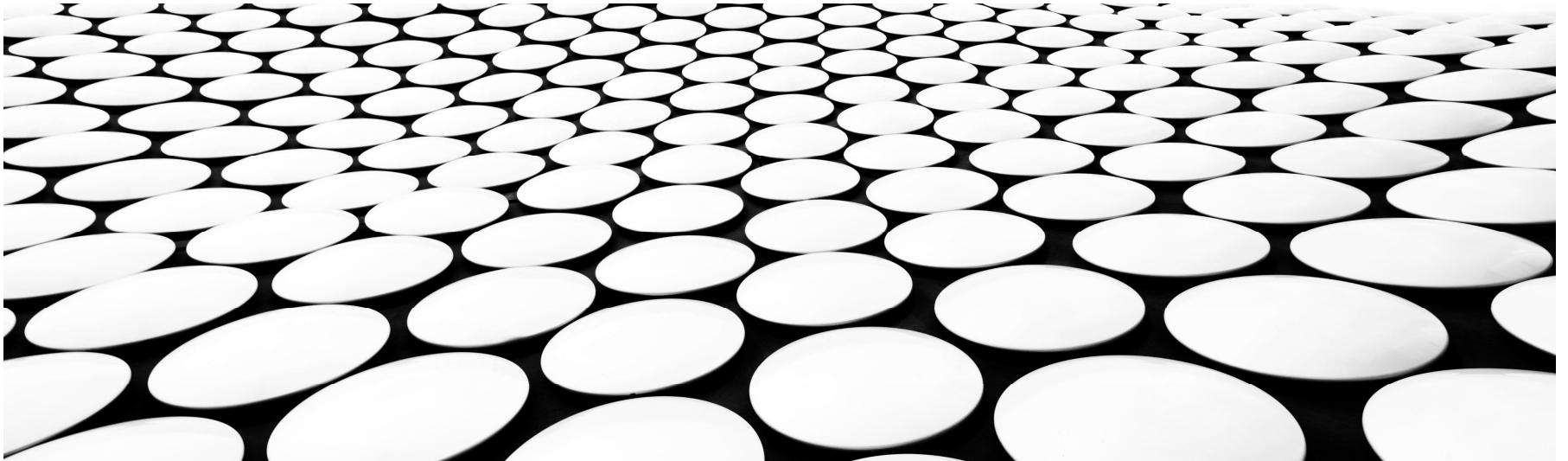

REGULAR EXPRESSIONS AND CONVERSIONS

IF 2124 TEORI BAHASA FORMAL OTOMATA



Regular expressions

A FA (NFA or DFA) is a “blueprint” for constructing a machine recognizing a regular language.

A *regular expression* is a “user-friendly,” declarative way of describing a regular language.

Example: $01^* + 10^*$ = $\{0, 1, 01, 10, 011, 100, 0111, \dots\}$.
Handwritten notes: $10^* = \{1, 10, 100, 1000, \dots\}$
 $01^* = \{0, 01, 011, 0111, \dots\}$

Regular expressions are used in e.g.

1. UNIX `grep` command
2. UNIX Lex (Lexical analyzer generator) and Flex (Fast Lex) tools.

Operations on languages

Union:

$$L \cup M = \{w : w \in L \text{ or } w \in M\}$$

Concatenation:

$$L.M = \{w : w = xy, x \in L, y \in M\}$$

Powers:

$$L^0 = \{\epsilon\}, L^1 = L, L^{k+1} = L.L^k$$

Kleene Closure:

$$L^* = \bigcup_{i=0}^{\infty} L^i \quad \{\epsilon, L, L^2, L^3, \dots\}$$

Question: What are \emptyset^0 , \emptyset^i , and \emptyset^*

$$\emptyset^0 = \{\epsilon\} \quad \emptyset^i = \{\} \quad \emptyset^* = \{\epsilon\}$$

Building regex's

Inductive definition of regex's:

Basis: ϵ is a regex and \emptyset is a regex.

$$L(\epsilon) = \{\epsilon\}, \text{ and } L(\emptyset) = \emptyset.$$

If $a \in \Sigma$, then a is a regex.

$$L(a) = \{a\}.$$

Induction:

If E is a regex's, then (E) is a regex.

$$L((E)) = L(E).$$

If E and F are regex's, then $E + F$ is a regex.

$$L(E + F) = L(E) \cup L(F).$$

If E and F are regex's, then $E.F$ is a regex.

$$L(E.F) = L(E).L(F).$$

If E is a regex's, then E^* is a regex.

$$L(E^*) = (L(E))^*$$

$$W \subset \{0,1\}^* = \{\epsilon, 0, 1, 00, \dots\}$$

Example: Regex for

$$L = \{w \in \{0,1\}^* : 0 \text{ and } 1 \text{ alternate in } w\}$$

$$\begin{aligned}(01)^* &= \{\epsilon, 01, 0101, 010101, \dots\} \\ (10)^* &= \{\epsilon, 10, 1010, 101010, \dots\} \\ 0(10)^* &= \{0, 010, 01010, \dots\} \\ 1(01)^* &= \{1, 101, 10101, \dots\}\end{aligned}$$

$$(01)^* + (10)^* + 0(10)^* + 1(01)^*$$

or, equivalently,

$$(\epsilon + 1)(01)^*(\epsilon + 0) \quad \hookrightarrow \quad \{\epsilon, 0, 1, 01, 10, 1010, 010, \dots\}$$

ϵ $\{ \}$

Order of precedence for operators:

1. Star

2. Dot 01^*

3. Plus $+$, union.

$$\begin{aligned}\phi &= \{ \} \neq \{ \epsilon \} \\ \phi^0 &= \{ \epsilon \}\end{aligned}$$

Example: $01^* + 1$ is grouped $(0(1)^*) + 1$