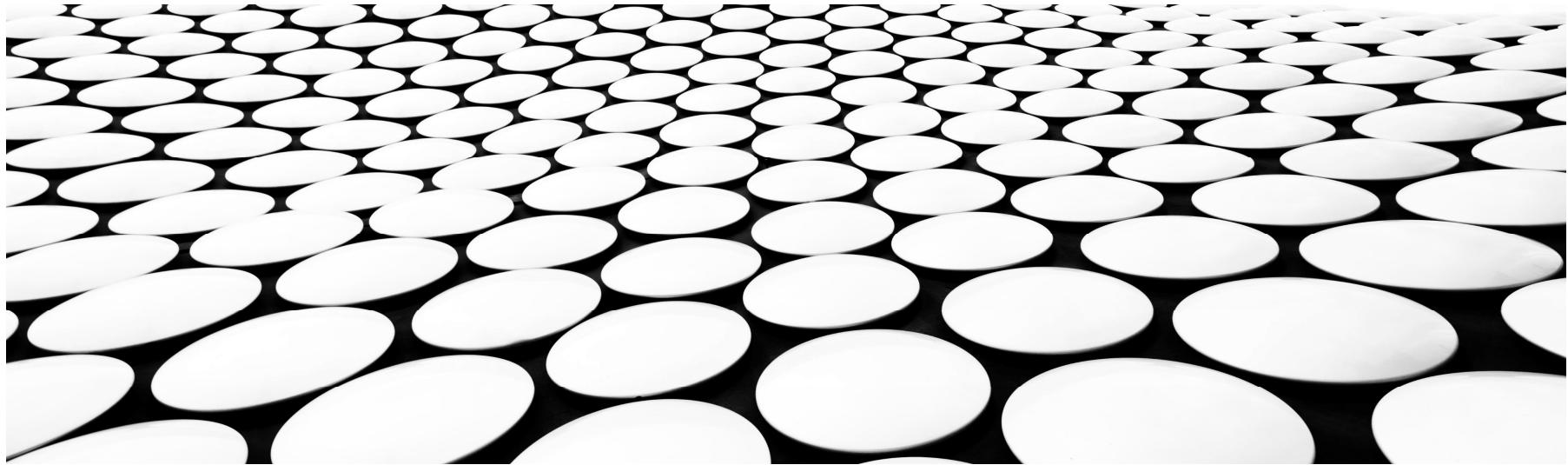

FINITE AUTOMATA

IF 2124 TEORI BAHASA FORMAL OTOMATA



Motivation

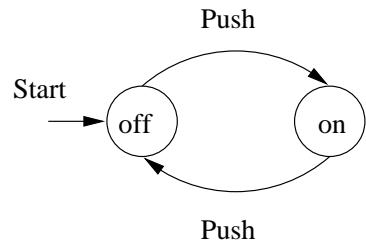
- Automata = abstract computing devices
- Turing studied Turing Machines (= computers) before there were any real computers
- We will also look at simpler devices than Turing machines (Finite State Automata, Push-down Automata, . . .), and specification means, such as grammars and regular expressions.
- NP-hardness = what cannot be efficiently computed

Finite Automata

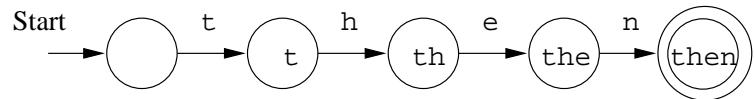
Finite Automata are used as a model for

- Software for designing digital circuits
- Lexical analyzer of a compiler
- Searching for keywords in a file or on the web.
- Software for verifying finite state systems, such as communication protocols.

- Example: Finite Automaton modelling an on/off switch



- Example: Finite Automaton recognizing the string `then`



Central Concepts

Alphabet: Finite, nonempty set of symbols

Example: $\Sigma = \{0, 1\}$ binary alphabet

Example: $\Sigma = \{a, b, c, \dots, z\}$ the set of all lower case letters

Example: The set of all ASCII characters

Strings: Finite sequence of symbols from an alphabet Σ , e.g. 0011001

Empty String: The string with zero occurrences of symbols from Σ

- The empty string is denoted ϵ

Length of String: Number of positions for symbols in the string.

$|w|$ denotes the length of string w

$$|0110| = 4, |\epsilon| = 0$$

Powers of an Alphabet: Σ^k = the set of strings of length k with symbols from Σ

Example: $\Sigma = \{0, 1\}$ $\Sigma = \{0, 1\}$

$$\Sigma^1 = \{0, 1\} \quad \Sigma^1 = \{0, 1\}$$

$$\Sigma^2 = \{00, 01, 10, 11\}$$

$$\Sigma^0 = \{\epsilon\} \quad \Sigma^3 = \{000, 001, 010, 011, \dots, 111\}.$$

Question: How many strings are there in Σ^3

The set of all strings over Σ is denoted Σ^*

$$\Sigma^* = \Sigma^0 \cup \Sigma^1 \cup \Sigma^2 \cup \dots$$

Also:

$$\Sigma^+ = \Sigma^1 \cup \Sigma^2 \cup \Sigma^3 \cup \dots$$

$$\Sigma^* = \Sigma^+ \cup \{\epsilon\}$$

Concatenation: If x and y are strings, then xy is the string obtained by placing a copy of y immediately after a copy of x

$$x = a_1 a_2 \dots a_i, y = b_1 b_2 \dots b_j$$

$$xy = a_1 a_2 \dots a_i b_1 b_2 \dots b_j$$

Example: $x = 01101, y = 110, xy = 01101110$

Note: For any string x

$$x\epsilon = \epsilon x = x$$

Languages:

If Σ is an alphabet, and $L \subseteq \Sigma^* = \{\epsilon, 01, 00, 01, 10, 11, 000, 001, \dots \dots\}$,
then L is a language

Examples of languages:

- The set of legal English words
- The set of legal C programs
- The set of strings consisting of n 0's followed by n 1's

$$\{\epsilon, 01, 0011, 000111, \dots\}$$

- The set of strings with equal number of 0's and 1's

$\{\epsilon, 01, 10, 0011, 0101, 1001, \dots\}$

- L_P = the set of binary numbers whose value is prime

$\{10, 11, 101, 111, 1011, \dots\}$

- The empty language \emptyset
- The language $\{\epsilon\}$ consisting of the empty string

Note: $\emptyset \neq \{\epsilon\}$

Note2: The underlying alphabet Σ is always finite