Number-Theoretic Algorithms

Hengfeng Wei

hfwei@nju.edu.cn

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Cancellation in modular arithmetic

(TC 31.4.2)

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Cancellation in modular arithmetic

(TC 31.4.2)

$$ad \equiv bd \pmod{n} \not\Longrightarrow a \equiv b \pmod{n}$$

$$ad \equiv bd \pmod{n}, a \perp n \Longrightarrow a \equiv b \pmod{n}$$

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Changing the modulus

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$$ad \equiv bd \pmod n \iff a \equiv b \pmod{\frac{n}{(d,n)}}$$

Changing the modulus

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$$a \equiv b \pmod{100} \implies a \equiv b \pmod{20} \implies a \equiv b \pmod{5}$$

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$$a \equiv b \pmod{n_1}, a \equiv b \pmod{n_2} \iff a \equiv b \pmod{n_1 n_2}, \text{ if } n_1 \perp n_2$$

$$\forall 1 \le i \le k, a \equiv b \pmod{n_i} \iff a \equiv b \pmod{n}, \text{ if } n_i \perp n_j$$

Worst-case analysis of Euclid's algorithm

(TC 31.2–5)

1. If $a > b \geq 0$, $\textsc{Euclid}(a, b)$ makes $\leq 1 + \log_{\phi} b$ recursive calls.

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$$\text{To prove } b < F_{3+\log_\phi b}.$$

$$F_k = \frac{\phi^k - \hat{\phi}^k}{\sqrt{5}} = [\frac{\phi^k}{\sqrt{5}}] \ge \frac{\phi^k}{\sqrt{5}}$$

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$$\begin{aligned} &(16, 12) \\ &= (12, 4) \\ &= (4, 0) \\ &= 4 \end{aligned}$$

$$\begin{aligned} &(4, 3) \\ &= (3, 1) \\ &= (1, 0) \\ &= 1 \end{aligned}$$

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$$\textsc{Euclid}(a, b) \leftrightarrow \textsc{Euclid}(\frac{a}{(a,b)}, \frac{b}{(a,b)})$$

$$\textsc{Euclid}(b, a \bmod b) \leftrightarrow \textsc{Euclid}(\frac{b}{(a,b)}, \frac{a}{(a,b)} \bmod \frac{b}{(a,b)})$$

$$\frac{a}{(a,b)} \bmod \frac{b}{(a,b)} = \frac{a \bmod b}{(a,b)}$$

Worst-case analysis of Euclid's algorithm

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2. Improve this bound to $1 + \log_{\phi}(\frac{b}{(a,b)})$.

Lemma (Generalization of Lemma 31.10)

If $a > b \geq 1, d = (a, b)$ *and* $\text{Euclid}(a, b)$ *performs* $k \geq 1$ *recursive calls, then* $a \geq dF_{k+2}$ *and* $b \geq dF_{k+1}$.

Average-case analysis of Euclid's algorithm

$$T(m,0) = 0; \quad T(m,n) = 1 + T(n, m \bmod n) \; n \geq 1$$

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Assume that, for $0 \leq k < n$, $(n \bmod k)$ is "random":

$$T_n \approx 1 + \frac{1}{n}(T_0 + T_1 + \cdots + T_{n-1})$$

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Reference

"The Art of Computer Programming, Vol 2: Seminumerical Algorithms (Section 4.5.3)" by Donald E. Knuth, 3rd edition.

Pairwise relatively prime

(TC 31.2–9)

$$n_1, n_2, n_3, n_4 \text{ are pairwise relatively prime}$$

$$\iff$$

$$\gcd(n_1 n_2, n_3 n_4) = \gcd(n_1 n_3, n_2 n_4) = 1$$

Pairwise relatively prime

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a set of $\lceil \lg k \rceil$ pairs of numbers derived from the n_i are relatively prime.

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$$\gcd(\boxed{1_L}, \boxed{1_R}) = \gcd(\boxed{2_L}, \boxed{2_R}) = \cdots = \gcd(\boxed{\lceil \lg k \rceil_L}, \boxed{\lceil \lg k \rceil_R}) = 1$$

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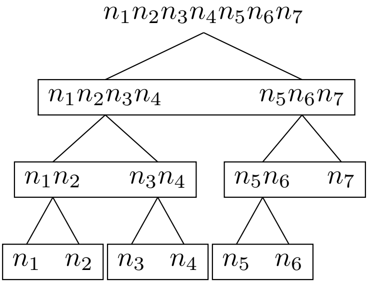
$$\gcd(\boxed{1_L}, \boxed{1_R}) = \gcd(\boxed{2_L}, \boxed{2_R}) = \cdots = \gcd(\boxed{\lceil \lg k \rceil_L}, \boxed{\lceil \lg k \rceil_R}) = 1$$

$$k = 3: \quad \gcd(n_1, n_2 n_3) = \gcd(n_2, n_3) = 1$$

$$k = 2: \quad \gcd(n_1, n_2) = 1$$

Pairwise relatively prime: divide-and-conquer



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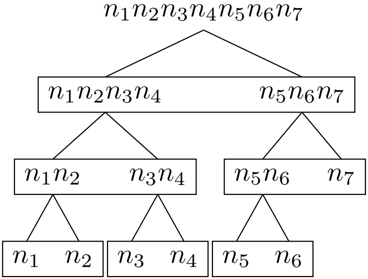
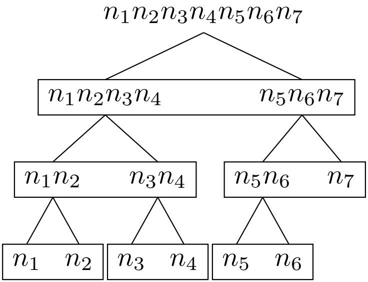


$$\begin{cases} T(1) = 0 \\ T(k) = 2T(\frac{k}{2}) + 1 \end{cases}$$

Pairwise relatively prime: divide-and-conquer



$$\begin{cases} T(1) = 0 \\ T(k) = 2T(\frac{k}{2}) + 1 \end{cases} \implies T(k) = k - 1$$

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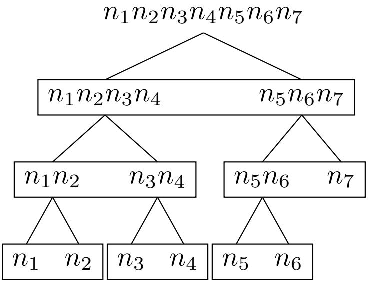

$$\begin{cases} T(1) = 0 \\ T(k) = 2T(\frac{k}{2}) + 1 \end{cases} \implies T(k) = k - 1$$

$$T_k = k - 1 : (n_i, n_{i+1} n_{i+2} \cdots n_k) \quad \forall 1 \le i < k$$

Pairwise relatively prime: smarter combination

$$\begin{cases} T(1) = 0 \\ T(k) = T(\frac{k}{2}) + 1 \end{cases}$$

Pairwise relatively prime: smarter combination

$$\begin{cases} T(1) = 0 \\ T(k) = T(\frac{k}{2}) + 1 \end{cases} \implies T(k) = \lceil \lg k \rceil$$

Pairwise relatively prime: the dividing pattern

$$k = 7 : \quad n_0, n_1, n_2, \ldots, n_6$$

$$\begin{array}{c} 000 \\ 001 \\ 010 \\ 011 \\ 100 \\ 101 \\ 110 \end{array}$$

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Prove by (strong) mathematical induction.

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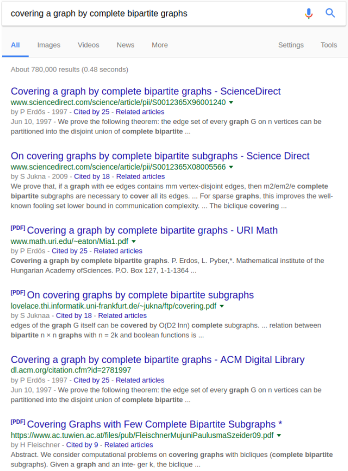
Prove by (strong) mathematical induction.

$$\begin{aligned} T(k) &\geq 1 + T(\lceil \frac{k}{2} \rceil) \\ &\geq 1 + \lceil \lg \lceil \frac{k}{2} \rceil \rceil \\ &= \lceil \lg k \rceil \end{aligned}$$

Biclique covering

Covering a complete graph with few complete bipartite subgraphs.



covering a graph by complete bipartite graphs

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Covering a graph by complete bipartite graphs - ScienceDirect
www.sciencedirect.com/science/article/pii/S0012365X96001240
by P Erdős - 1997 - Cited by 25 - Related articles
Jun 10, 1997 - We prove the following theorem: the edge set of every **graph** G on n vertices can be partitioned into the disjoint union of **complete bipartite** ...

On covering graphs by complete bipartite subgraphs - Science Direct
www.sciencedirect.com/science/article/pii/S0012365X08005566
by S Jukna - 2009 - Cited by 18 - Related articles
We prove that, if a **graph** with ee edges contains mm vertex-disjoint edges, then m2/em2/e **complete bipartite** subgraphs are necessary to **cover** all its edges. ... For sparse **graphs**, this improves the well-known fooling set lower bound in communication complexity. ... The biclique **covering** ...

[PDF] Covering a graph by complete bipartite graphs - URI Math
www.math.uri.edu/~eaton/Mia1.pdf
by P Erdős - Cited by 25 - Related articles
Covering a graph by complete bipartite graphs. P. Erdos, L. Pyber,*. Mathematical institute of the Hungarian Academy ofSciences. P.O. Box 127, 1-1-1364 ...

[PDF] On covering graphs by complete bipartite subgraphs
lovelace.thi.informatik.uni-frankfurt.de/~jukna/ftp/covering.pdf
by S Juknaa - Cited by 18 - Related articles
edges of the **graph** G itself can be **covered** by O(D2 lnn) **complete** subgraphs. ... relation between **bipartite** n × n **graphs** with n = 2k and boolean functions is ...

Covering a graph by complete bipartite graphs - ACM Digital Library
dl.acm.org/citation.cfm?id=2781997
by P Erdős - 1997 - Cited by 25 - Related articles
Jun 10, 1997 - We prove the following theorem: the edge set of every **graph** G on n vertices can be partitioned into the disjoint union of **complete bipartite** ...

[PDF] Covering Graphs with Few Complete Bipartite Subgraphs *
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by H Fleischner - Cited by 9 - Related articles
Abstract. We consider computational problems on **covering graphs** with bicliques (**complete bipartite** subgraphs). Given a **graph** and an inte- ger k, the biclique ...

Biclique covering: rethinking the first divide-and-conquer

$$T(k) = k - 1$$

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edge-disjoint biclique partition

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Reference for $T(k) \geq k - 1$

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Biclique covering: rethinking the first divide-and-conquer

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edge-disjoint biclique partition

Reference for $T(k) \geq k - 1$

"On the Addressing Problem for Loop Switching" by Graham and Pollak, 1971.

Reference for *weighted* biclique partition

"Covering a Graph by Complete Bipartite Graphs" by P. Erdős and L. Pyber, 1997.

Number-Theoretic Algorithms

Chinese Remainder Theorem (CRT)

Theorem (CRT)

$$n_1, \ldots, n_k; \quad a_1, \ldots, a_k$$

$$n_i \perp n_j \quad i \neq j, \quad n = n_1 n_2 \cdots n_k$$

$$\exists! a \ (0 \leq a < n) : a \equiv a_i \pmod{n_i}.$$

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$$a \leftrightarrow (a_1, a_2, \ldots, a_k)$$

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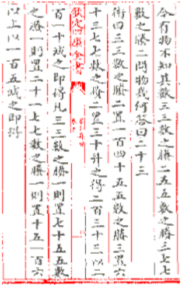
$$a \leftrightarrow (a_1, a_2, \ldots, a_k)$$

Proof for uniqueness.

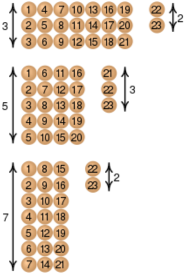
$$a \equiv a' \pmod{n_i} \implies n \mid a - a'.$$

$\square$

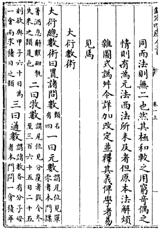
History of CRT



“孙子算经”



“物不知数”



秦九韶 “数书九章”
大衍求一术

Proof of CRT (1)

Nonconstructive proof.

$$f : [0, n) \to \prod_{1 \le i \le k} [0, a_i)$$

$$f : a \mapsto (a \bmod n_1, \ldots, a \bmod n_k)$$

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- f is one-to-one.

$\square$

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Nonconstructive proof.

$$f : [0, n) \to \prod_{1 \le i \le k} [0, a_i)$$

$$f : a \mapsto (a \bmod n_1, \ldots, a \bmod n_k)$$

- f is one-to-one.
- f is onto.

$$\exists a : f(a) = (a_1, \ldots, a_k).$$

□

Proof of CRT (2)

Constructive proof by induction.

$$a \equiv a_1 \pmod{n_1} \tag{1}$$

$$a \equiv a_2 \pmod{n_2} \tag{2}$$

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Constructive proof by induction.

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$$a \equiv a_2 \pmod{n_2} \tag{2}$$

$$(1) \implies a = a_1 + n_1 y$$

$$x = a_1 + n_1 n_1^{-1}(a_2 - a_1) \pmod{n_1 n_2}$$

□

Proof of CRT (3)

Constructive proof by induction.

$$a \equiv a_1 \pmod{n_1} \tag{3}$$
$$a \equiv a_2 \pmod{n_2} \tag{4}$$

$$n_1 \perp n_2 \implies n_1 n_1' + n_2 n_2' = 1$$

Proof of CRT (3)

Constructive proof by induction.

$$a \equiv a_1 \pmod{n_1} \tag{3}$$
$$a \equiv a_2 \pmod{n_2} \tag{4}$$

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$$x = a_1 n_1 n_1' + a_2 n_2 n_2' \pmod{n_1 n_2}$$

□

Proof of CRT (4)

Constructive proof.

1. $x \equiv 1 \pmod{n_i}, \quad x \equiv 0 \pmod{n_j}\ (i \neq j)$

$$x = M_i(M_i^{-1} \bmod n_i) \implies x = M_i M_i^{-1} \pmod{n}$$

Proof of CRT (4)

Constructive proof.

1. $x \equiv 1 \pmod{n_i}, \quad x \equiv 0 \pmod{n_j}\ (i \neq j)$

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2. $x \equiv a_i \pmod{n_i}, \quad x \equiv 0 \pmod{n_j}\ (i \neq j)$

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2. $x \equiv a_i \pmod{n_i}, \quad x \equiv 0 \pmod{n_j}\ (i \neq j)$

$$x = a_i M_i M_i^{-1} \pmod{n}$$

3. $a \equiv a_i \pmod{n_i}, \forall 1 \leq i \leq k$

$$a = \sum_{1 \leq i \leq k} a_i M_i M_i^{-1} \pmod{n}$$

□

Proof of CRT (5)

More efficient constructive proof.

Reference

"The Residue Number System" by Garner, 1959.

Reference

"The Art of Computer Programming, Vol 2: Seminumerical Algorithms (Section 4.3.2)" by Donald E. Knuth, 3rd edition.

□

Operations over CRT

$$a \leftrightarrow (a_1, a_2, \ldots, a_n)$$

$$a \pm b \leftrightarrow (a_1 \pm b_1, a_2 \pm b_2, \ldots, a_n \pm b_n)$$
$$a \times b \leftrightarrow (a_1 \times b_1, a_2 \times b_2, \ldots, a_n \times b_n)$$

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TC 31.5–3

$$a \leftrightarrow (a_1, a_2, \dots, a_n), (a, n) = 1 \implies a^{-1} \leftrightarrow (a_1^{-1}, a_2^{-2}, \dots, a_n^{-1})$$

Operations over CRT

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$$a^{-1} \equiv a_i^{-1} \pmod{n_i} \impliedby \begin{cases} a \equiv a_i \pmod{n_i} \\ (a, n) = 1 \end{cases}$$

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Theorem (The ϕ function)

$$\phi(p) = p - 1$$
$$\phi(p^k) = p^k - p^{k-1}$$

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"We shall not prove this formula here." — *CLRS (Section 31.3)*

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$$U_{mn} = \{a \bmod mn, (a, mn) = 1\}$$
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$$f : U_{mn} \to U_m \times U_n$$

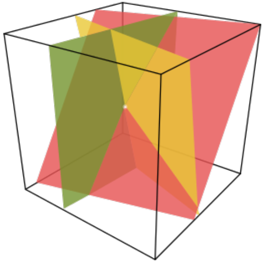
$$f(a \bmod mn) = (a \bmod m, a \bmod n).$$

□

Secret sharing using the CRT

Definition ((k, n)-threshold secret sharing scheme)

$(2, 3)$-secret sharing:



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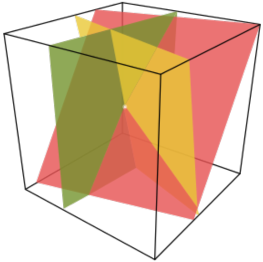

Reference

"How to Share a Secret" by Maurice Mignotte, 1982.

Secret sharing using the CRT

1. Choose m_i:

$$m_1 < m_2 < \cdots < m_n, \quad m_i \perp m_j, \quad \prod_{i=n-k+2}^{n} m_i < \prod_{i=1}^{k} m_i$$

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3. Compute the shares:

$$s_i = S \bmod m_i$$

Solving the system of congruences

(TC 31.5–2)

$$\begin{cases} x \equiv 1 \pmod{9} \\ x \equiv 2 \pmod{8} \\ x \equiv 3 \pmod{7} \end{cases}$$

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$$\begin{cases} 19x \equiv 556 \pmod{3} \\ 19x \equiv 556 \pmod{5} \\ 19x \equiv 556 \pmod{7} \\ 19x \equiv 556 \pmod{11} \end{cases}$$

$$\begin{cases} x \equiv 1 \pmod{3} \\ x \equiv 4 \pmod{5} \\ x \equiv 2 \pmod{7} \\ x \equiv 9 \pmod{11} \end{cases}$$

Solving the system of congruences

CRT with non-pairwisely co-prime moduli

$$\begin{cases} x \equiv 3 \pmod{8} \\ x \equiv 11 \pmod{20} \\ x \equiv 1 \pmod{15} \end{cases}$$

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CRT with non-pairwisely co-prime moduli

$$\begin{cases} x \equiv 3 \pmod 8 \\ x \equiv 11 \pmod{20} \\ x \equiv 1 \pmod{15} \end{cases}$$

$$\begin{cases} x \equiv 3 \pmod{2^3} \end{cases} \qquad \begin{cases} x \equiv 3 \pmod{2^2} \\ x \equiv 1 \pmod 5 \end{cases} \qquad \begin{cases} x \equiv 1 \pmod 3 \\ x \equiv 1 \pmod 5 \end{cases}$$

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Solving the system of congruences

Theorem (CRT with non-pairwisely coprime moduli)

$$a_i \equiv a_j \pmod{(n_i, n_j)}$$

$$0 \leq a < \mathrm{lcm}(n_1, n_2, \ldots, n_k)$$


THANK YOU