Treasure Hunting

Zheng Zangwei*

April 15, 2018

Problem 1

1. Analyze the worst case W(n) and the best case B(n) time complexity of mergesort as accurately as possible.

Explore the relation between them and the binary representations of numbers.

Plot W(n) and B(n) and explain what you observe.

2. Analyze the average case A(n) time complexity of mergesort.

Plot A(n) and explain what you observe.

3. Prove that: The minimum number of comparisons needed to merge two sorted arrays of equal size m is 2m-1.

Algorithm 1

```
 \begin{aligned} \textit{Subroutine} \quad & \textit{Mergesort}\left(A,p,r\right) \\ & \textit{if} \quad (p \! < \! r) \\ & q = \lfloor (p \! + \! r)/2 \rfloor; \\ & \textit{Mergesort}\left(A,p,q\right); \\ & \textit{Mergesort}\left(A,q \! + \! 1,r\right); \\ & \textit{Merge}\left(A,p,q,r\right); \end{aligned}
```

1 Problem 1.1

We know from the algorithm, W(0) = 0, $W(\lfloor \frac{n}{2} \rfloor) + W(\lceil \frac{n}{2} \rceil) + n - 1$; We have following lemma which I omit the proof:

^{*}E-mail: zzw@smail.nju.edu.com Student ID: 171860658

Lemma 1

$$\lceil \frac{n}{2} \rceil = \lfloor \frac{n+1}{2} \rfloor$$

Lemma 2

If $T(n) = 1 + T(\lfloor \frac{n}{k} \rfloor)$ we have $T(n) = \lceil \log_k(n+1) \rceil$. This can be derived from k-base number.

Now consider W(n+1)-W(n), we have

$$W(n+1) - W(n) = W(\lfloor \frac{n}{2} \rfloor + 1) - W(\lfloor \frac{n}{2} \rfloor) + 1.$$

Let f(n) = W(n+1) - W(n) we have $f(n) = 1 + f(\lfloor \frac{n}{2} \rfloor)$ so that $W(n+1) - W(n) = f(n) = \lceil \log_k(n+1) \rceil$. By summation, $W(n) = \sum_{k=1}^n \lceil \log_2 k \rceil$.

This introduced an interesting property which I want but cannot find a intuitive proof to it.

Corollary 1

If n is even, W(n) denotes the number of bits in binary representations, then

$$\sum_{k=0}^{\frac{n}{2}-1} W(k) + n - 1 = \sum_{k=\frac{n}{2}+1}^{n} W(k)$$

Then from the great *Concrete Mathematics* problem 3.34 we have the following calculation. Let $m = \lceil log_2 n \rceil$,

$$W(n) + (2^{m} - n)m = \sum_{k=1}^{2^{m}} \lceil \log_{2} k \rceil$$

$$= \sum_{j,k} j[j = \lceil \log_{2} k \rceil] [1 \le k \le 2^{m}]$$

$$= \sum_{j,k} j[2^{j-1} < k \le 2^{j}] [1 \le jm]$$

$$= \sum_{j=1}^{m} j * 2^{j-1} = 2^{m} (m-1) + 1$$

Thus $W(n) = n\lceil \log_2 n \rceil - 2^{\lceil \log_2 n \rceil} + 1$.

2 Problem 1.2

Still, we know $B(0) = 0, B(\lfloor \frac{n}{2} \rfloor) + B(\lceil \frac{n}{2} \rceil) + \lfloor \frac{n}{2} \rfloor$