

# Disjoint Types

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# Objectives

- Describe the syntax for disjoint data types in Haskell.
- Describe a few use-cases for them.

# Simple Type Definitions

## Disjoint Type Syntax

$$\text{data } Name = Name [type \dots] [ | Name [type \dots] \dots ]$$

- Note: Constructor names must be capitalized.
- Constructor names also must be unique.

```
1 data Contest = Rock | Scissors | Paper
2 data Velocity = MetersPerSecond Float
3               | FeetPerSecond Float
```

# Example of contest and velocity

```
1 winner Rock Scissors = "Player 1"
2 winner Scissors Paper = "Player 1"
3 winner Paper Rock = "Player 1"
4 winner Scissors Rock = "Player 2"
5 winner Paper Scissors = "Player 2"
6 winner Rock Paper = "Player 2"
7 winner _ _ = "Tie"
8
9 thrust (FeetPerSecond x) = x / 3.28
10 thrust (MetersPerSecond x) = x
```

# The Most Fun Datatypes are Recursive

## Our Own List Construct

```
1 data Mylist = Cons Int Mylist
2             | Nil
3   deriving Show
4 mklist [] = Nil
5 mklist (x:xs) = Cons x (mklist xs)
```

We can run it like this:

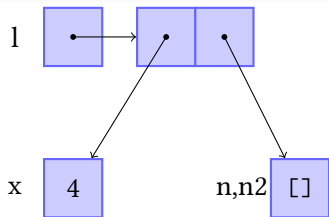
```
*Main> let l1 = mklist [2,3,4]
*Main> l1
Cons 2 (Cons 3 (Cons 4 Nil))
```

- 1 A recursive type without a recursive case is not really recursive.
- 2 A recursive type without a base case is dangerous, but using Haskell, it might even make sense.

# Type Constructors and Memory

- When a type constructor is invoked, it causes memory to be allocated.
  - Writing an integer
  - Writing [] or Nil
  - Using : or Cons
- Writing down a variable does not cause memory to be allocated.

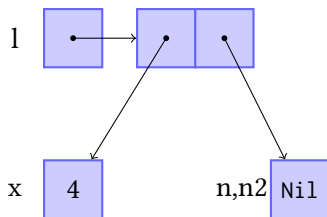
```
1 x = 4  -- allocates 4
2 n = [] -- allocates empty list
3 n2 = n -- does NOT allocate memory
4 l = x:n -- A cons cell is allocated, but not the 4 or the empty list
```



# Similarly...

```
1 x = 4
2 n = Nil
3 n2 = n
4 l = Cons x n
```

- Our own types do the same thing.



# Parameters

Haskell supports parametric polymorphism, like templates in C++ or generics in Java.

## Parametric Polymorphism

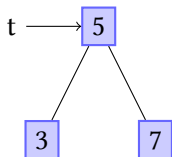
```
1 data Glist a = GCons a (Glist a)
2             | GNil
3 deriving Show

1 x1 = GCons 1 (GCons 2 (GCons 4 GNil))
2 x2 = GCons "hi" (GCons "there" GNil)
3 x3 = GCons Nil (GCons (Cons 5 Nil) GNil)
```



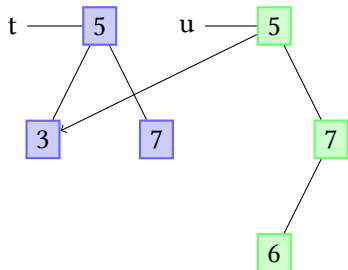
# Functional Updating

- It is important to understand functional updating.
- We don't update in place. We make copies, and share whatever we can.
  - Example: add 5,3,7 to a tree t.
  - `let u = add t 6`
  - `let v = add u 1`



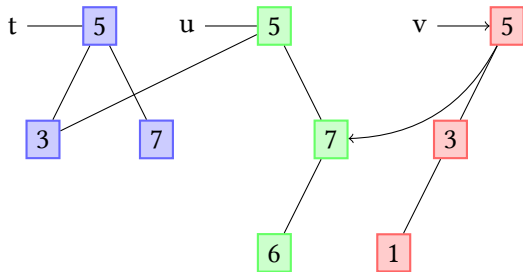
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# The Maybe Type

## The Maybe Type

```
1 data Maybe a = Just a | Nothing
```

We can use it in places where we want to return something, but we are not sure that the item exists.

```
1 getItem key [] = Nothing
2 getItem key ((k,v):xs) =
3     if key == k then Just v
4     else getItem key xs
```

Example:

```
*Main> getItem 3 [(2,"french hens"), (3,"turtle doves")]
Just "turtle doves"
*Main> getItem 5 [(2,"french hens"), (3,"turtle doves")]
Nothing
```