

Synchronization: Reader Writer Problem

Harrison Vuong (hvuong@ucsc.edu)

David Zou (dzou@ucsc.edu)

Derek Frank (dmfrank@ucsc.edu)

1 Goal:

The goal of reader writer is to use semaphores to eliminate starvation of the writers.

2 Available Resources:

Minix Documentations

Interprocess Communications and Synchronization Slides

3 Design:

void reader (int mutex, int writing, int priority, int readerid)

note: shared variables

number of reader file

time file for reader

open the time file for readers

generate a random number based on time

close the time file for readers

if the random number is below 20,000 + 20,000

pause for the given random number

semdown the priority

semdown the mutex

open then number of reader file

variable to keep number of readers

close the number of readers

increment the number of readers

open the number of reader file

write the new count of readers if any

close the number of reader file

if number of readers is 1

 semdown writing

 writer goes to sleep since there is a reader

semup the priority

print the reader is currently reading

generate a new random time

if new random time is less than 20,000 add 20,000 to it

- open time file for readers
- write the new time in file then close the file
- decrease the number of readers
- after opening and recording old number of readers
- open and write new number of readers since reader has decreased

- if there are currently no readers
 - semup for the writers

- semup the mutex
- print that the reader is finished reading

int main (int argc, char **argv) → reader.c

- check for correct usage
- open the sem file
- check for errors while opening
- scan the contents of the semfile (which keeps track of currently used semaphores)
- call reader with the retrieved semaphores

void writer (int mutex, int writing, int priority, int writerid)

- create a writer time file to keep track of wait time
- create a random number using random while using an algorithm to ensure that its between the given time for wait
- usleep using the calculated random number
- sem down on the priority
- sem down on writing
 - this is where the writing will take place
 - show the writer is writing
- sem up on writing
- sem up on priority
- the writer is now done writing

int main (int argc, char **argv) → writer.c

- check the usage of writer by checking options
- set values for wid, mutex, writing and priority
- open the semaphore file
- do necessary error checking while opening
- get values for writing and priority from the sem file
- pass variables to writer for execution

int main (int argc, char **argv) → rwinit.c

- create two file pointers
- initialize mutex, priority and writing to 0
- create the mutex semaphore
 - do the necessary error checking using a switch to ensure that the semaphore is valid

- create the writing semaphore
 - do the necessary error checking using a switch to ensure that the semaphore is valid
- create the priority semaphores
 - do the necessary error checking using a switch to ensure validity again

- open the sem file and write the semaphores to the file
- close the file

- open the number of readers and write 0
- close the file

int main (int argc, char **argv) → rtfree.c

- check to see if there is an error when opening sem file
 - otherwise open the semfile
- get the variables from the sem file
- close the sem file
- if semfree of the mutex is below 0, return error
- if semfree of writing is 0 or below 0, return error
- if semfree of priority is 0 or below 0, return error
- delete number of readers file
- delete reader time file
- delete the semaphore file

4 Testing:

The main purpose of this test is to see if the usage of semaphores correctly. To prove this, writers are not supposed to starve, meaning, if there are many writers to readers, writers will eventually get priority to go over readers that entered after the writers. To do this, the test was to create multiple writer and reader processes running as background processes. In order to determine if the semaphores are working is to trace through which reader and which writer are getting access to which critical regions at a time.