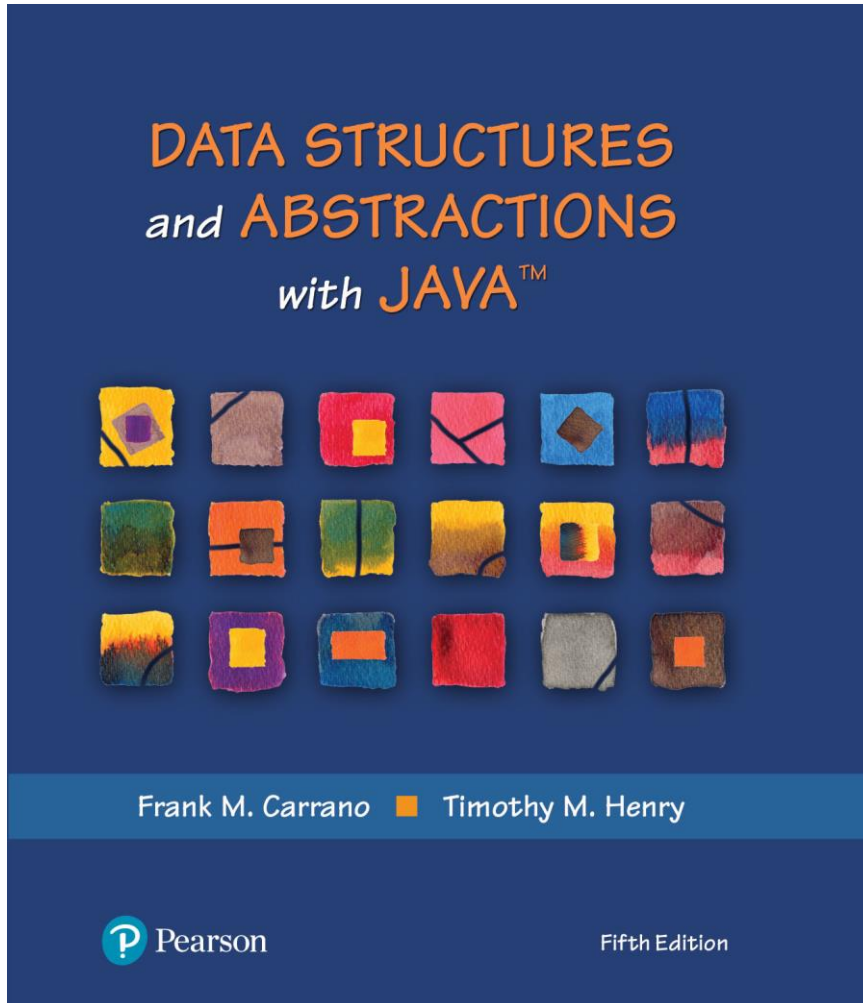


Data Structures and Abstractions with Java™

5th Edition



Chapter 12

A List Implementation That Links Data

Advantages of Linked Implementation

- Uses memory only as needed
- When entry removed, unneeded memory returned to system
- Avoids moving data when adding or removing entries

Adding a Node at Various Positions

- Possible cases:
 - Chain is empty
 - Adding node at chain's beginning
 - Adding node between adjacent nodes
 - Adding node to chain's end

Adding a Node

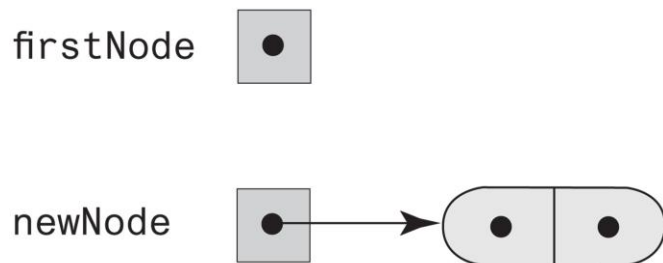
This pseudocode establishes a new node for the given data

newNode references a new instance of Node

Place newEntry in newNode

firstNode = address of newNode

(a) An empty chain and a new node



(b) After adding the new node to the chain

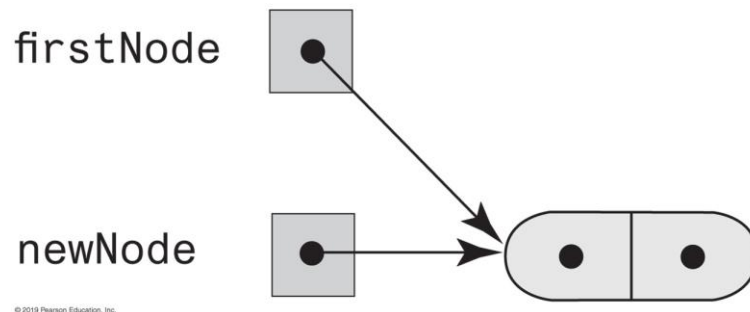


FIGURE 12-1 Adding a node to an empty chain

Adding a Node

This pseudocode describes the steps needed to add a node to the beginning of a chain.

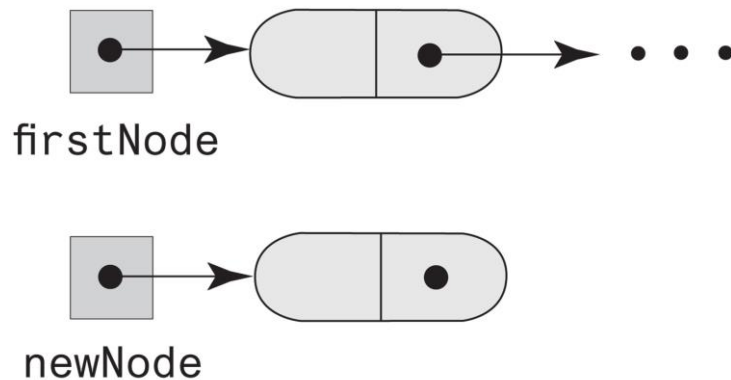
newNode references a new instance of Node

Place newEntry in newNode

Set newNode's link to firstNode

Set firstNode to newNode

(a) A chain of nodes and a new node



(b) After adding the new node to the beginning of the chain

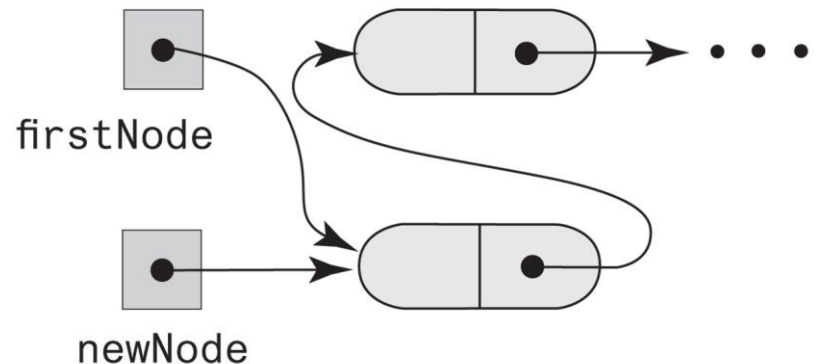


FIGURE 12-2 Adding a node to the beginning of a chain

Adding a Node

Pseudocode to add a node to a chain between two existing, consecutive nodes

newNode references the new node

Place newEntry in newNode

Let nodeBefore reference the node that will be before the new node

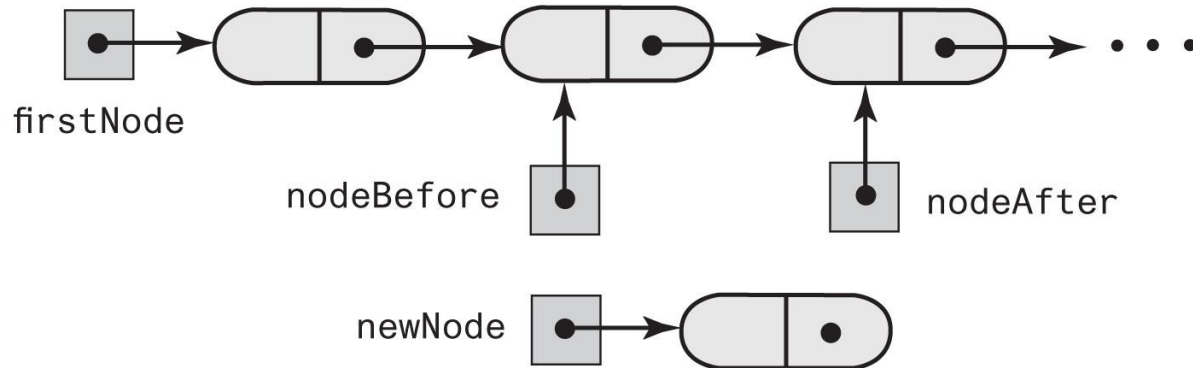
Set nodeAfter to nodeBefore's link

Set newNode's link to nodeAfter

Set nodeBefore's link to newNode

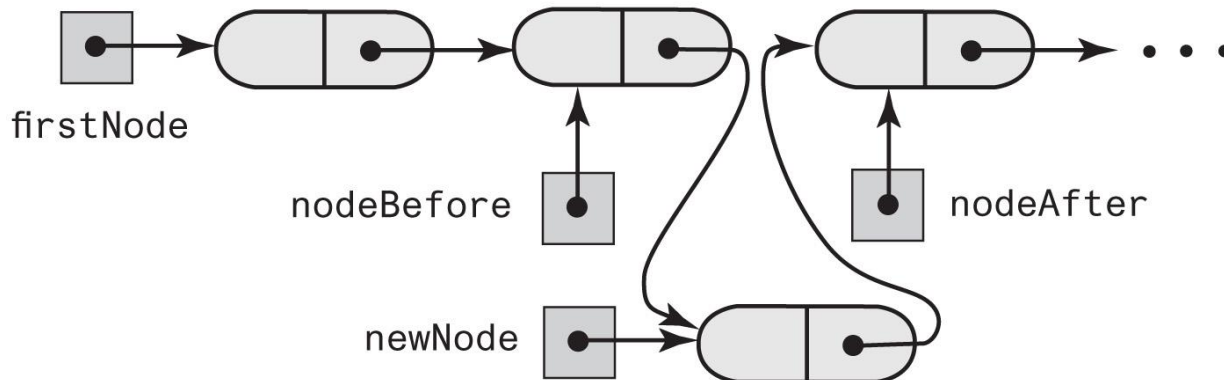
Adding a Node

(a) A chain of nodes and a new node



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(b) After adding the new node between adjacent nodes



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FIGURE 12-3 Adding a node between two adjacent nodes

Adding a Node

Steps to add a node at the end of a chain.

newNode references a new instance of Node

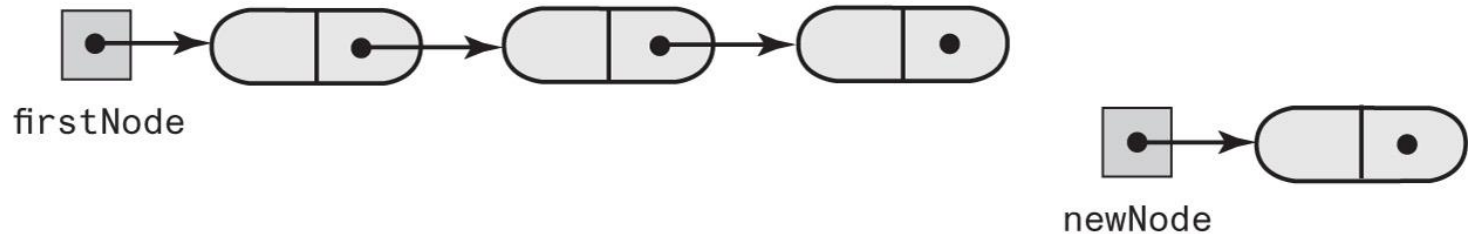
Place newEntry in newNode

Locate the last node in the chain

Place the address of newNode in this last node

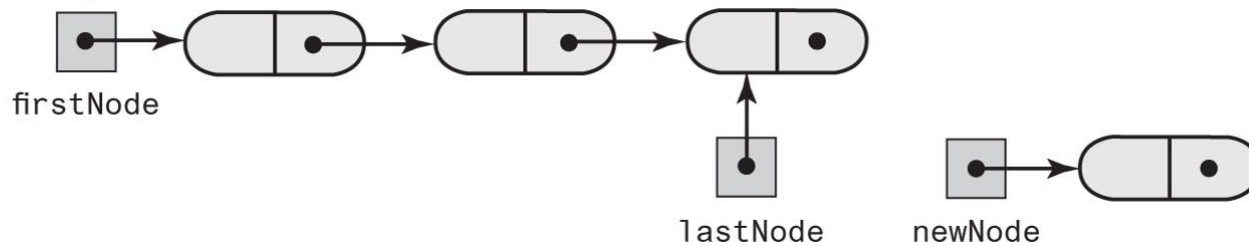
Adding a Node

(a) A chain of nodes and a new node



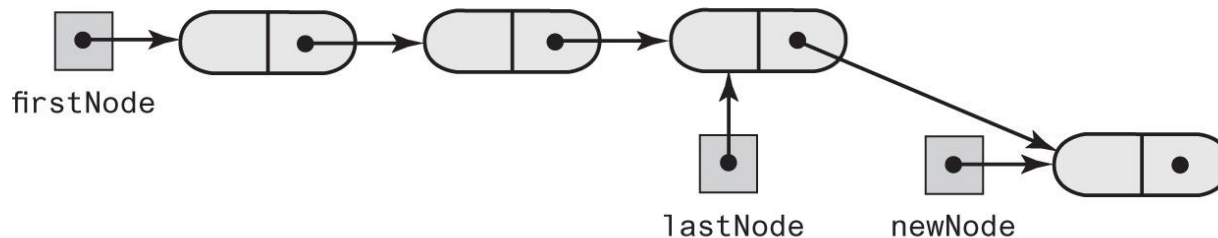
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(b) After locating the last node



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(c) After adding the new node to the end of the chain



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FIGURE 12-4 Adding a node to the end of a chain

Removing a Node

- Possible cases
 - Removing the first node
 - Removing a node other than first one

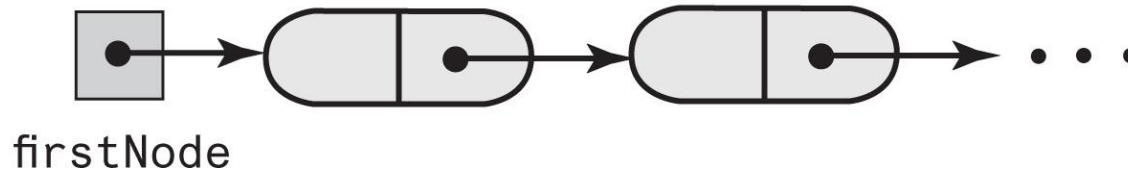
Removing a Node

Steps for removing the first node.

*Set **firstNode** to the link in the first node; **firstNode** now either references the second node or is null if the chain had only one node.*

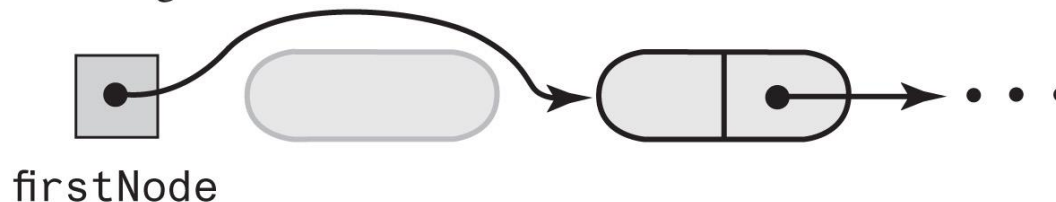
Since all references to the first node no longer exist, the system automatically recycles the first node's memory.

(a) A chain of nodes



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(b) After removing the first node



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FIGURE 12-5 Removing the first node from a chain

Removing a Node

Removing a node other than the first one.

Let nodeBefore reference the node before the one to be removed.

Set nodeToRemove to nodeBefore's link; nodeToRemove now references the node to be removed.

Set nodeAfter to nodeToRemove's link; nodeAfter now either references the node after the one to be removed or is null.

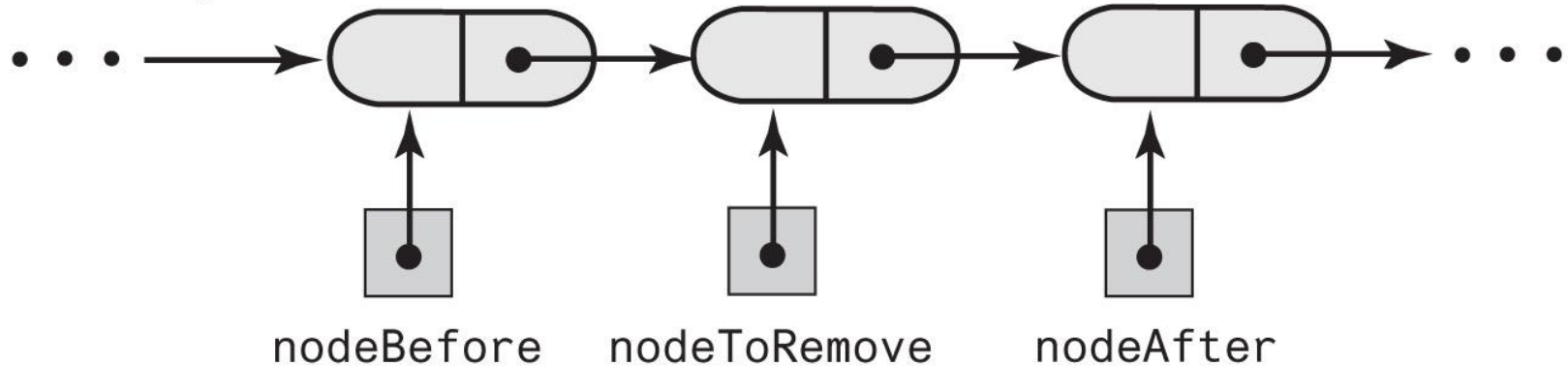
Set nodeBefore's link to nodeAfter. (nodeToRemove is now disconnected from the chain.)

Set nodeToRemove to null.

Since all references to the disconnected node no longer exist, the system automatically recycles the node's memory.

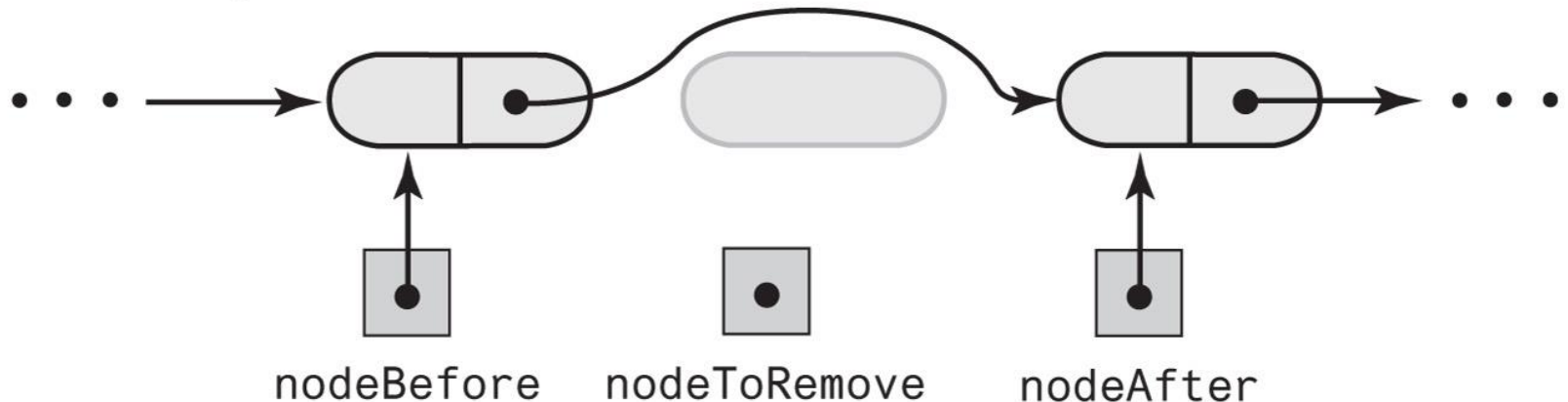
Removing a Node

(a) After locating the node to remove



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(b) After removing the node



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FIGURE 12-6 Removing an interior node from a chain

Removing a Node

```
// Returns a reference to the node at a given position.
// Precondition: The chain is not empty;
//      1 <= givenPosition <= numberOfEntries.
private Node getNodeAt(int givenPosition)
{
    // Assertion: (firstNode != null) &&
    //      (1 <= givenPosition) && (givenPosition <= numberOfEntries)
    Node currentNode = firstNode;
    // Traverse the chain to locate the desired node
    // (skipped if givenPosition is 1)
    for (int counter = 1; counter < givenPosition; counter++)
        currentNode = currentNode.getNextNode();
    // Assertion: currentNode != null
    return currentNode;
} // end getNodeAt
```

forloop alternatives

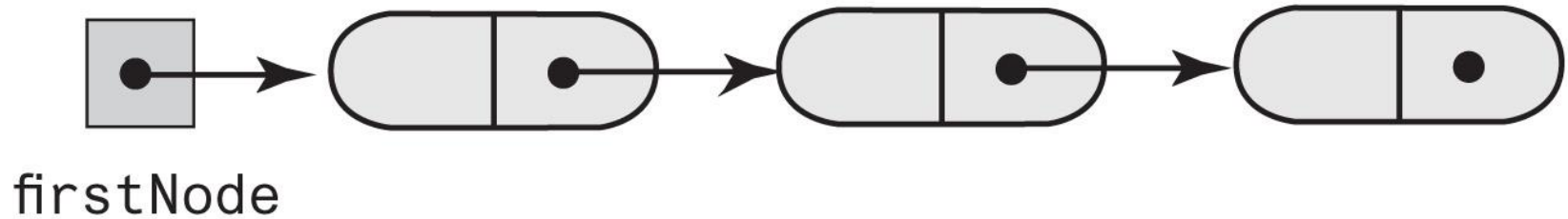
```
while(currentNode != null){
    if(givent position equals counter)
        break

    currentNode = currentNode.getNextNode()
    counter++;
}
```

Operations on a chain depended on the method `getNodeAt`

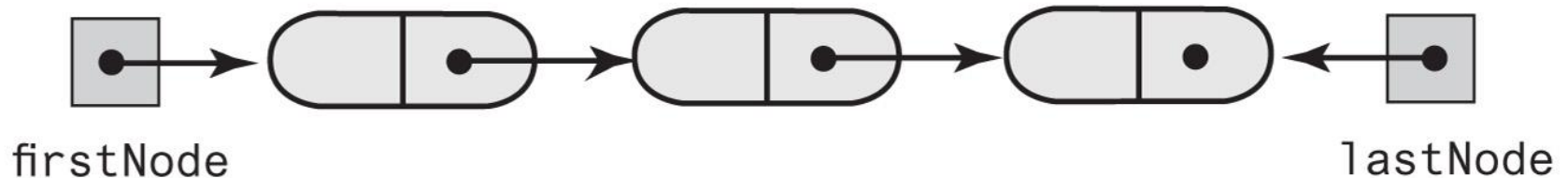
Using a Tail Reference

(a) With only a head reference



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(b) With both a head reference and a tail reference



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FIGURE 12-7 Two linked chains

Data Fields and Constructor (Part 1)

```
/** A linked implementation of the ADT list. */
public class LList<T> implements ListInterface<T>
{
    private Node firstNode;      // Reference to first node of chain
    private int numberOfEntries;

    public LList()
    {
        initializeDataFields();
    } // end default constructor

    public void clear()
    {
        initializeDataFields();
    } // end clear

    /* < Implementations of the public methods add, remove, replace, getEntry, contains,
       getLength, isEmpty, and toArray go here. >
       ... */

    // Initializes the class's data fields to indicate an empty list.
    private void initializeDataFields()
    {
        firstNode = null;
        numberOfEntries = 0;
    } // end initializeDataFields
}
```

LISTING 12-1 An outline of the class LList

Data Fields and Constructor (Part 2)

```
// Returns a reference to the node at a given position.
// Precondition: The chain is not empty;
//      1 <= givenPosition <= numberOfEntries.
private Node getNodeAt(int givenPosition)
{
    // Assertion: (firstNode != null) &&
    //      (1 <= givenPosition) && (givenPosition <= numberOfEntries)
    Node currentNode = firstNode;

    // Traverse the chain to locate the desired node
    // (skipped if givenPosition is 1)
    for (int counter = 1; counter < givenPosition; counter++)
        currentNode = currentNode.getNextNode();
    // Assertion: currentNode != null
    return currentNode;
} // end getNodeAt

private class Node
{
    // < Implementation of private inner class Node >
} // end Node

} // end LList
```

LISTING 12-1 An outline of the class LList

Adding to the End of the List

```
public void add(T newEntry)
{
    Node newNode = new Node(newEntry);

    if (isEmpty())
        firstNode = newNode;
    else // Add to end of nonempty list
    {
        Node lastNode = getNodeAt(numberOfEntries);
        lastNode.setNextNode(newNode); // Make last node reference new node
    } // end if

    numberOfEntries++;
} // end add
```

The method add assumes method getNodeAt

Adding at a Given Position

```
public void add(int givenPosition, T newEntry)
{
    if ((givenPosition >= 1) && (givenPosition <= numberOfEntries + 1))
    {
        Node newNode = new Node(newEntry);
        if (givenPosition == 1)           // Case 1
        {
            newNode.setNextNode(firstNode);
            firstNode = newNode;
        }
        else                               // Case 2: list is not empty
        {
            // and givenPosition > 1
            Node nodeBefore = getNodeAt(givenPosition - 1);
            Node nodeAfter = nodeBefore.getNextNode();
            newNode.setNextNode(nodeAfter);
            nodeBefore.setNextNode(newNode);
        } // end if
        numberOfEntries++;
    }
    else
        throw new IndexOutOfBoundsException(
            "Illegal position given to add operation.");
} // end add
```

add method.

Method isEmpty

```
public boolean isEmpty()
{
    boolean result;

    if (numberOfEntries == 0) // Or getLength() == 0
    {
        // Assertion: firstNode == null
        result = true;
    }
    else
    {
        // Assertion: firstNode != null
        result = false;
    } // end if

    return result;
} // end isEmpty
```

Method toArray

```
public T[] toArray()  
{  
    // The cast is safe because the new array contains null entries  
    @SuppressWarnings("unchecked")  
    T[] result = (T[])new Object[numberOfEntries];  
  
    int index = 0;  
    Node currentNode = firstNode;  
    while ((index < numberOfEntries) && (currentNode != null))  
    {  
        result[index] = currentNode.getData();  
        currentNode = currentNode.getNextNode();  
        index++;  
    } // end while  
  
    return result;  
} // end toArray
```

Traverses chain, loads an array.

Testing Core Methods

```
public static void main(String[] args)
{
    System.out.println("Create an empty list.");
    ListInterface<String> myList = new LList<String>();
    System.out.println("List should be empty; isEmpty returns " +
        myList.isEmpty() + ".");

    System.out.println("\nTesting add to end:");
    myList.add("15");
    myList.add("25");
    myList.add("35");
    myList.add("45");
    System.out.println("List should contain 15 25 35 45.");
    displayList(myList);
    System.out.println("List should not be empty; isEmpty() returns " +
        myList.isEmpty() + ".");
    System.out.println("\nTesting clear():");
    myList.clear();
    System.out.println("List should be empty; isEmpty returns " +
        myList.isEmpty() + ".");
} // end main
```

LISTING 12-2 A main method that tests part of the implementation of the ADT list

Testing Core Methods

Program Output

Create an empty list.

List should be empty; isEmpty returns true.

Testing add to end:

List should contain 15 25 35 45.

The list contains 4 entries, as follows:

15 25 35 45

List should not be empty; isEmpty() returns false.

Testing clear():

List should be empty; isEmpty returns true.

LISTING 12-2 A main method that tests part of the implementation of the ADT list

remove method returns entry it deletes from list

```
public T remove(int givenPosition)
{
    T result = null;           // Return value
    if ((givenPosition >= 1) && (givenPosition <= numberOfEntries))
    {
        // Assertion: !isEmpty()
        if (givenPosition == 1)    // Case 1: Remove first entry
        {
            result = firstNode.getData();    // Save entry to be removed
            firstNode = firstNode.getNextNode(); // Remove entry
        }
        else                        // Case 2: Not first entry
        {
            Node nodeBefore = getNodeAt(givenPosition - 1);
            Node nodeToRemove = nodeBefore.getNextNode();
            result = nodeToRemove.getData(); // Save entry to be removed
            Node nodeAfter = nodeToRemove.getNextNode();
            nodeBefore.setNextNode(nodeAfter); // Remove entry
        } // end if
        numberOfEntries--;           // Update count
        return result;               // Return removed entry
    }
    else
        throw new IndexOutOfBoundsException(
            "Illegal position given to remove operation.");
} // end remove
```


Continuing the Implementation

```
public T replace(int givenPosition, T newEntry)
{
    if ((givenPosition >= 1) && (givenPosition <= numberOfEntries))
    {
        // Assertion: !isEmpty()
        Node desiredNode = getNodeAt(givenPosition);
        T originalEntry = desiredNode.getData();
        desiredNode.setData(newEntry);
        return originalEntry;
    }
    else
        throw new IndexOutOfBoundsException(
            "Illegal position given to replace operation.");
} // end replace
```

Replacing a list entry requires us to replace the data portion of a node with other data.

Continuing the Implementation

```
public T getEntry(int givenPosition)
{
    if ((givenPosition >= 1) && (givenPosition <= numberOfEntries))
    {
        // Assertion: !isEmpty()
        return getNodeAt(givenPosition).getData();
    }
    else
        throw new IndexOutOfBoundsException(
            "Illegal position given to getEntry operation.");
} // end getEntry
```

Retrieving a list entry is straightforward.

Continuing the Implementation

```
public boolean contains(T anEntry)
{
    boolean found = false;
    Node currentNode = firstNode;

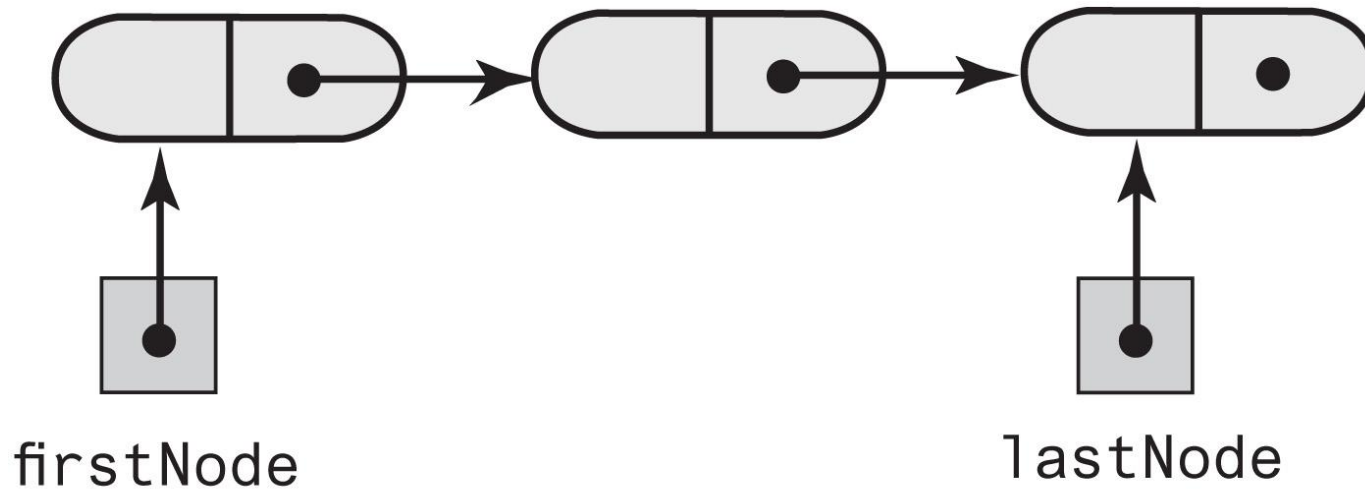
    while (!found && (currentNode != null))
    {
        if (anEntry.equals(currentNode.getData()))
            found = true;
        else
            currentNode = currentNode.getNextNode();
    } // end while

    return found;
} // end contains
```

Checking to see if an entry is in the list, the method contains.

A Refined Linked Implementation

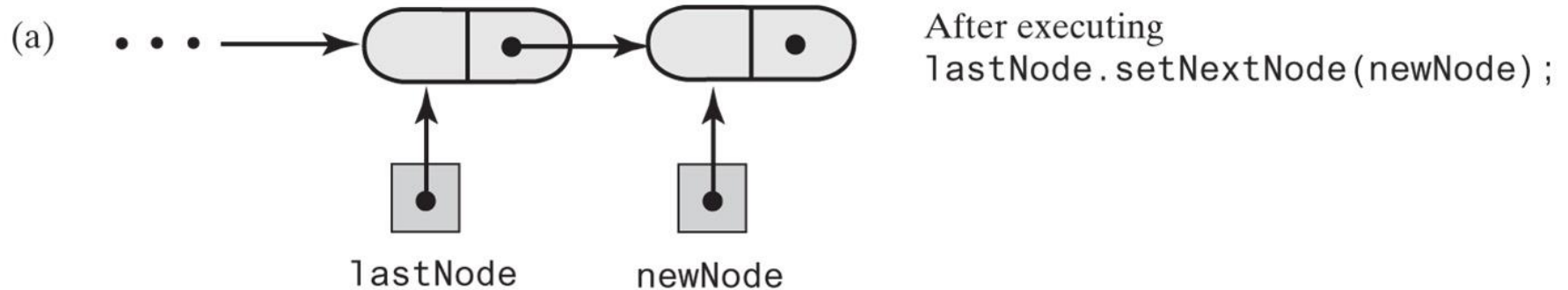
```
private Node firstNode;    // Head reference to first node  
private Node lastNode;    // Tail reference to last node  
private int numberOfEntries; // Number of entries in list
```



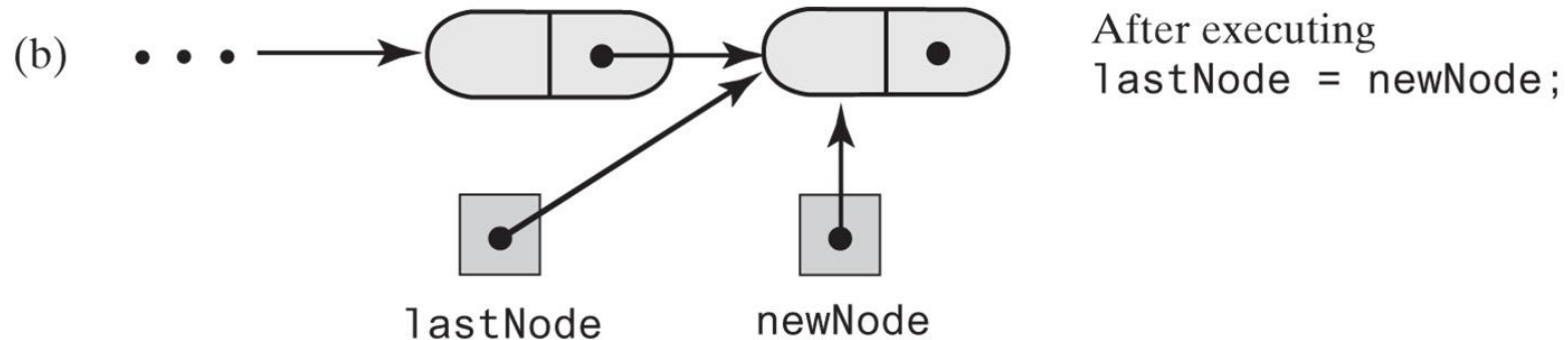
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FIGURE 12-8 A linked chain with both a head reference and a tail reference

A Refined Linked Implementation



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FIGURE 12-9 Adding a node to the end of a nonempty chain that has a tail reference

A Refined Linked Implementation

```
public void add(T newEntry)
{
    Node newNode = new Node(newEntry);

    if (isEmpty())
        firstNode = newNode;
    else
        lastNode.setNextNode(newNode);

    lastNode = newNode;
    numberOfEntries++;
} // end add
```

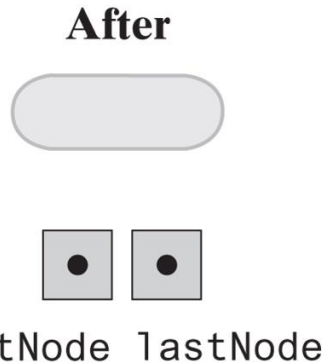
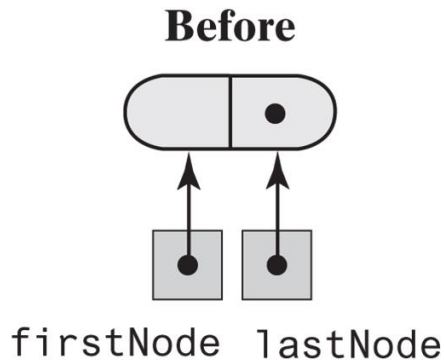
Revision of the first add method

A Refined Linked Implementation - refined add by position

```
public void add(int givenPosition, T newEntry) {
    if ((givenPosition >= 1) && (givenPosition <= numberOfEntries + 1))
    {
        Node newNode = new Node(newEntry);
        if (isEmpty())
        {
            firstNode = newNode;
            lastNode = newNode;
        }
        else if (givenPosition == 1)
        {
            newNode.setNextNode(firstNode);
            firstNode = newNode;
        }
        else if (givenPosition == numberOfEntries + 1)
        {
            lastNode.setNextNode(newNode);
            lastNode = newNode;
        }
        else {
            Node nodeBefore = getNodeAt(givenPosition - 1);
            Node nodeAfter = nodeBefore.getNextNode();
            newNode.setNextNode(nodeAfter);
            nodeBefore.setNextNode(newNode);
        } // end if
        numberOfEntries++;
    }
    else
        throw new IndexOutOfBoundsException(
            "Illegal position given to add operation.");
} // end add
```

A Refined Linked Implementation

(a) A one-node chain



(b) A chain of two or more nodes

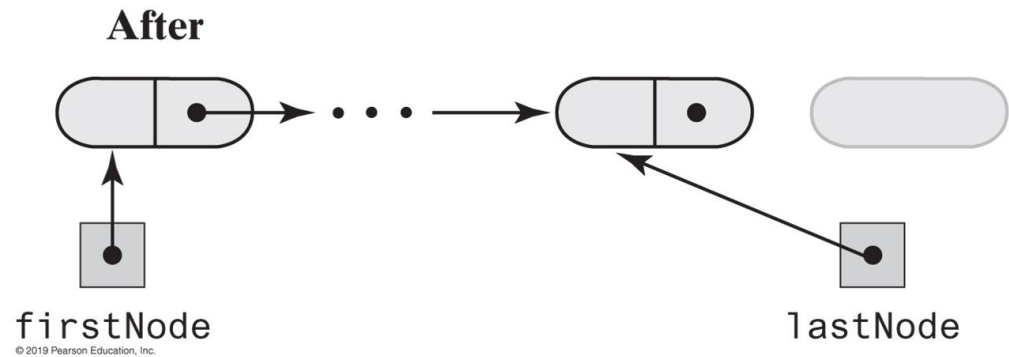
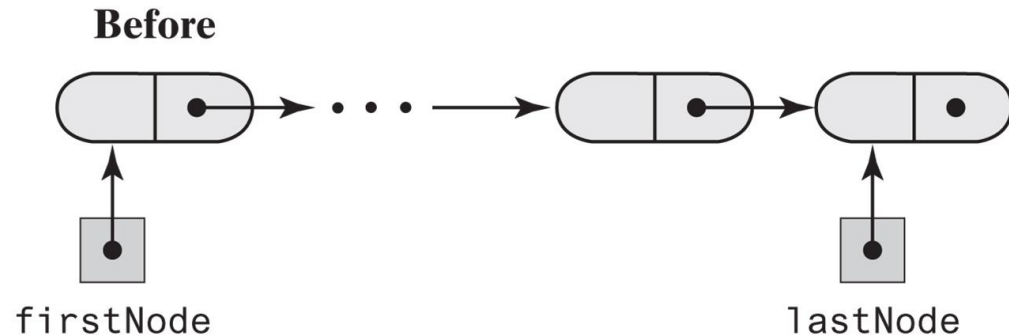


FIGURE 12-10 Before and after removing the last node from a chain that has both head and tail references and contains one or more nodes

A Refined Linked Implementation — refined remove

```
public T remove(int givenPosition) {
    T result = null;           // Return value
    if ((givenPosition >= 1) && (givenPosition <= numberOfEntries))
    {
        // Assertion: !isEmpty()
        if (givenPosition == 1)           // Case 1: Remove first entry
        {
            result = firstNode.getData(); // Save entry to be removed
            firstNode = firstNode.getNextNode();
            if (numberOfEntries == 1)
                lastNode = null;         // Solitary entry was removed
        }
        else                             // Case 2: Not first entry
        {
            Node nodeBefore = getNodeAt(givenPosition - 1);
            Node nodeToRemove = nodeBefore.getNextNode();
            Node nodeAfter = nodeToRemove.getNextNode();
            nodeBefore.setNextNode(nodeAfter);
            result = nodeToRemove.getData();
            if (givenPosition == numberOfEntries)
                lastNode = nodeBefore;    // Last node was removed
        } // end if
        numberOfEntries--;
    }
    else
        throw new IndexOutOfBoundsException(
            "Illegal position given to remove operation.");

    return result;              // Return removed entry
} // end remove
```

Efficiency of Using a Chain

best; average; worst

Operation	Alist	LList	LListWithTail
<code>add(newEntry)</code>	$O(1)$	$O(n)$	$O(1)$
<code>add(givenPosition, ^[SEP]newEntry)</code>	$O(n)$; $O(n)$; $O(1)$	$O(1)$; $O(n)$	$O(1)$; $O(n)$; $O(1)$
<code>toArray()</code>	$O(n)$	$O(n)$	$O(n)$
<code>remove(givenPosition)</code>	$O(n)$; $O(n)$; $O(1)$	$O(1)$; $O(n)$	$O(1)$; $O(n)$
<code>replace(givenPosition, newEntry)</code>	$O(1)$	$O(1)$; $O(n)$	$O(1)$; $O(n)$; $O(1)$
<code>getEntry(givenPosition)</code>	$O(1)$	$O(1)$; $O(n)$	$O(1)$; $O(n)$; $O(1)$
<code>contains(anEntry)</code>	$O(n)$	$O(n)$	$O(n)$
<code>clear()</code> , <code>getLength()</code> , <code>isEmpty()</code>	$O(1)$	$O(1)$	$O(1)$

FIGURE 12-11 The time efficiencies of the ADT list operations for three implementations, expressed in Big Oh notation

Java Class Library: The Class `LinkedList`

- Implements the interface `List`
- `LinkedList` defines more methods than are in the interface `List`
- You can use the class `LinkedList` as implementation of ADT
 - `queue`
 - `deque`
 - or `list`.

End

Chapter 12