# **Algebras and Fixpoints: Recursion Non-Recursively!**

#### Who am I?

David Barri. Hi!

Github: japgolly

Twitter: @japgolly

Programming 28 years.

Scala & FP: 4 years. (14%)

#### What'cha workin' on?

- ShipReq: Software requirements startup.
- Open Source: github.com/japgolly
  - scalajs-react
  - scalajs-benchmark
  - ScalaCSS
  - Test State
  - Nyaya
  - univeq
  - o microlibs

## **FP = Amazing**

Great experience.

My startup & OSS is all FP.

Humbling. Can do so much with so little.

$$S \Rightarrow (S, A)$$

 $F[A] \Rightarrow A$ 

## **Quick Foreward**

Who knows how to write a red-black tree?

## **Quick Foreward**

- Who knows how to use Map or Set?
- Who finds they get benefit when using Map or Set?

## **Quick Foreward**

Of the things I'll present tonight,

if you leave not understanding **how/why it works**, that's ok. It might take a few attempts; explore later at your leisure.

More important to understand how to use it for your benefit.

#### **Structure**

- 1. When/where should I use this stuff?
- 2. Really awesome thing.
- 3. Really awesome thing backwards!
- 4. Two extensions.
- 5. Example & live demo.
- 6. Summary & resources.

#### **Recursive Data**

Often pops up in my experience.

What is it?

What are some examples?

Where is this talk applicable?

• Use cases

```
1.0. Create account.
1.0.1. System prompts for username & password.
1.0.2. User enters details.
1.0.2.a. User says "der I'm a user"
1.0.2.b. User mashes keyboard.
1.0.3. System crashes cos actors have no type safety.
Goto 1.0.1.
```

• Filter expression

```
type != blah and (active || foo == bar)
```

• Trie (Prefix tree)

```
com.blah.cool_library
com.blah.cool_library.dao
com.blah.cool_library.lib
com.blah.cool_library.util
```

#### modeled as:

```
com
blah
cool_library
dao
lib
util
```

• Self-referential many-to-many relationship.

Focus on a node, get a recursive tree of children or parents.

#### **Recursive Data**

• Nyaya (OSS): properties

```
eg. P = A \wedge \neg B \wedge (C \rightarrow D)
```

• Test-State (OSS): actions, properties, assertions.

```
eg. getMilk = fridge.open >> take(milk) >>= drink >> fridge.close
```

That's me just recently. Quite common.

## **Everyday Recursive Data Types**

- Linked lists
- Maps, sets
- JSON

## Why is recursive data useful?

- Models reality
  - o FriendBook: My friends have friends who have friends who have friends...
  - File system: directories have directories which have directories...

## Why is recursive data useful?

- Efficiency
  - Time hashmap, binary search
  - Space trie, interval trees

# **Typical Recursive Model**

Calculator example.

```
sealed trait Calc
case class Number (i: Int) extends Calc
case class Add (a: Calc, b: Calc) extends Calc
case class Multiply(a: Calc, b: Calc) extends Calc
// 1 + 2
Add(
  Number(1),
  Number(2))
// 3 * (1 + 2)
Multiply(
  Number(3),
  Add(
   Number(1),
   Number(2)))
```

#### Problem #1

No control over the recursion.

It's hardcoded to be always recursive, and always infinitely.

Can't abstract, define depth.

#### Problem #2

Recursion is usually hard. Often requires practice.

Non-recursion easier than recursion.

## Problem #3

Can't annotate nodes.

#### No annotations

#### Annotated with IDs

Calculator annotated with logging:

How would we allow **any** annotation to our calculator?

```
sealed abstract class Calc[A] {
 val ann: A
case class Add [A](ann: A, a: Calc[A], b: Calc[A]) extends Calc[A]
case class Multiply[A](ann: A, a: Calc[A], b: Calc[A]) extends Calc[A]
// 3 * (1 + 2)
Multiply("3 * (1 + 2)",
 Number("3", 3),
 Add("(1 + 2)",
   Number("1", 1),
   Number("2", 2)))
```

Is there a better way? 🦻

## **Step 1: Generalise the recursive-type.**

Before:

```
sealed trait Calc
case class Number (i: Int) extends Calc
case class Add (a: Calc, b: Calc) extends Calc
case class Multiply(a: Calc, b: Calc) extends Calc
```

After:

#### Wait...

```
val WHY_YOU_NO_WORK_?! : Calc[Calc] =
   Multiply(
    Number(3),
   Add(
      Number(1),
      Number(2)))
```

# **Type constructors**

- List
- Map
- Future

## **Types**

- Int
- String
- Unit

## **Type constructors**

- List
- Map
- Future

## **Types**

- List[Int]
- Map[Int, String]
- Future[List[Int]]

## **Type constructors**

Calc

#### **Types**

- Calc[Unit]
- Calc[List[Int]]
- Calc[Nothing]
- Calc[Calc[Nothing]]
- Calc[Calc[Nothing]]]
- Calc[Calc[Calc[Nothing]]]]
- Calc[Calc[Calc[Calc[Nothing]]]]]

This works but isn't nice.

```
val MaxDepthOf3: Calc[Calc[Nothing]]] =
   Number(456)

val MaxDepthOf3: Calc[Calc[Nothing]]] =
   Multiply(
     Number(3),
     Add(
         Number(1),
         Number(2)))
```

How do we get unlimited recursion back?

## **Step 2: Fixpoints!**

```
final case class Fix[F[_]](unfix: F[Fix[F]])
```

```
final case class Fix[F[_]](unfix: F[Fix[F]])
```

```
Fix[F] = F[ Fix[F] ]

//

= F[F[ Fix[F] ]]

//

= F[F[F[ Fix[F] ]]]

//
```

#### Before:

```
val calc: Calc =
    Multiply(
        Number(3),
        Add(
            Number(1),
            Number(2)))
```

#### After:

```
val calc: Fix[Calc] =
  Fix(Multiply(
    Fix(Number(3)),
    Fix(Add(
       Fix(Number(1)),
       Fix(Number(2))))))
```

It goes back and forth.

```
def gimmeFix (c: Calc[Fix[Calc]]):     Fix[Calc] = Fix[Calc](c)
def gimmeCalc(f: Fix[Calc]): Calc[Fix[Calc]] = f.unfix
```

```
Fix[Calc]

| ↑
unfix | Fix

↓ |

Calc[Fix[Calc]]
```

## **ALL YOU NEED TO KNOW IS:**

- Generalise (use an A) in place of a self-reference.
- Wrap your stuff in Fix when you want recursion.

# **Algebra**

F A -> A

type Algebra[F[\_], A] = F[A] => A

Scala has many ways of representing the same thing...

```
val listSumAlg: Algebra[List, Int] =
   _.sum

val listSumFn: List[Int] => Int =
   _.sum

def listSumMethod(list: List[Int]): Int =
   list.sum
```

They're all algebras.

```
listSumAlg : Algebra[List, Int]
listSumMethod: Algebra[List, Int]
listSumFn : Algebra[List, Int]
```

## Ok...and?

A special function exists...

```
def awesome[F[_]: Functor, A](data: Fix[F], alg: Algebra[F, A]): A
```

```
def awesome[F[_]: Functor, A](data: Fix[F], alg: Algebra[F, A]): A
```

Algebra needs a F[A].

No A in Fix[F].

Like magic!

# Let's try it out!

Well, first there's a prerequisite:

```
import scalaz.Functor
implicit val functor: Functor[Calc] =
  new Functor[Calc] {
    override def map[A, B](c: Calc[A])(f: A => B): Calc[B] = ???
}
```

s/Calc/YourDataType/g

```
import scalaz.Functor
implicit val functor: Functor[Calc] =
  new Functor[Calc] {
    override def map[A, B](c: Calc[A])(f: A => B): Calc[B] =
    c match {
```

case Number (i) => Number (i)

case Add  $(a, b) \Rightarrow Add$  (f(a), f(b))

case Multiply(a, b) => Multiply(f(a), f(b))

```
}
```

# Let's try it out!

```
val eval: Algebra[Calc, Int] = {
  case Number (i) => i
  case Add (a, b) => a + b
  case Multiply(a, b) => a * b
}
```

```
def awesome[F[_]: Functor, A](data: Fix[F], alg : Algebra[F, A]): A

val eval: Algebra[Calc, Int] = {
   case Number (i) => i
   case Add (a, b) => a + b
   case Multiply(a, b) => a * b
}
```

```
val c: Fix[Calc] = ...
awesome(c, eval) // returns 9
```

```
val explain: Algebra[Calc, String] = {
  case Number (i) => i.toString
  case Add (a, b) => s"($a + $b)"
  case Multiply(a, b) => s"($a * $b)"
}
```

```
val explain: Algebra[Calc, String] = ...
val eval : Algebra[Calc, Int] = ...

val explainAndEval: Algebra[Calc, (String, Int)] =
   explain zip eval
```

```
val c: Fix[Calc] = ...
awesome(c, explainAndEval) // returns ("3 * (1 + 2)", 9)
```

# **Amazing!**

Real name of awesome is catamorphism.

It goes to the ends of a tree, then calculates its way back up.

Bottom-up.

## How does that work?!



```
def cata: Fix[Calc] => A =
   fix => {
     val calc : Calc[Fix[Calc]] = fix.unfix
     val calcA: Calc[A] = calc.map(cata)
     val a : A = alg(calcA)
     a
   }
}
```

```
cata :: Functor f => (Algebra f a) -> Fix f -> a
cata alg = alg . fmap (cata alg) . unfix
```

## **ALL YOU NEED TO KNOW IS:**

- Create a functor.
- Create an algebra when you want to fold/reduce a tree.
- Call cata . Done!

**Next up: CoAlgebras** 

Everything in category theory can go backwards - this is called a dual.

### **Algebra**

• F[A] => A

### Coalgebra

• A => F[A]

### **Algebra**

- F[A] => A
- squash a structure (F) into a single value (A)

### Coalgebra

- A => F[A]
- expand a single value (A) into a structure (F)

### **Algebra**

- F[A] => A
- fold a structure (F) into a single value (A)

### Coalgebra

- A => F[A]
- unfold a single value (A) into a structure (F)

```
type CoAlgebra[F[_], A] = A => F[A]
```

```
val factors: CoAlgebra[Calc, Int] = i =>
   if (i > 2 && i % 2 == 0)
      Multiply(2, i / 2)
   else
      Number(i)
```

# **Another magic function**

```
def emosewa[F[_]: Functor, A](data: A, alg: CoAlgebra[F, A]): Fix[F]
```

Notice, there's no A in the result.

Real name is anamorphism.

Dual of catamorphism.

```
cata: Fix f -> (f a -> a ) -> a
ana : a -> (a -> f a) -> Fix f
```

It starts with the at the root, then calculates its way down to the nodes until complete. Top-down.

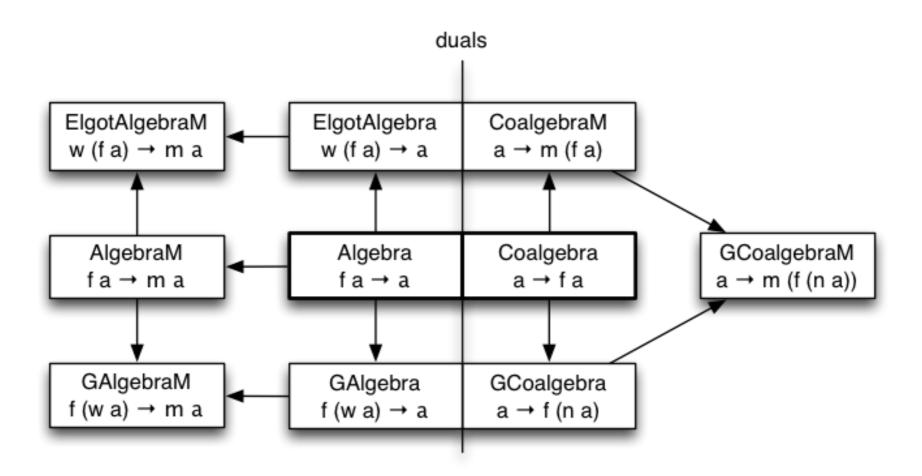
### **ALL YOU NEED TO KNOW IS:**

- Create a coalgebra when you want to build up a structure, using the seed/instruction/specification as the input.
- Call ana . Done!

## **Recursion Scheme Basics**

- Fixpoints
- (Co)Algebra
- {cata,ana}morphism

Most everything else builds on the above.



### Recursion Schemes unfolds (tear down a structure)

algebra  $f a \rightarrow Fix f \rightarrow a$ 

coalgebra  $f a \rightarrow a \rightarrow Fix f$ 

	catamorphism f a → a  prepromorphism* after applying a NatTrans	anamorphism a → f a  postpromorphism* before applying a NatTrans	<b>g</b> eneralized $(m f \rightarrow f m) \rightarrow (a \rightarrow f (m \beta))$
<b>g</b> eneralized $(f w \rightarrow w f) \rightarrow (f (w a) \rightarrow \beta)$	(f a → a) → (f → f) <b>para</b> morphism*  with primitive recursion  f (Fix f × a) → a	(a → f a) → (f → f)  apomorphism*  returning a branch or single level a → f (Fix f ∨ a)	
	zygomorphism* with a helper function $(f b \rightarrow b) \rightarrow (f (b \times a) \rightarrow a)$	<b>g apo</b> morphism $(b \rightarrow f b) \rightarrow (a \rightarrow f (b \lor a))$	
<b>g histo</b> morphism $(f h \rightarrow h f) \rightarrow (f (w a) \rightarrow a)$	histomorphism with prev. answers it has given f (w a) → a	futumorphism multiple levels at a time a → f (m a)	<b>g futu</b> morphism $(h f \rightarrow f h) \rightarrow (a \rightarrow f (m a))$

refolds (build up then tear down a structure) algebra  $a b \rightarrow (f \rightarrow g) \rightarrow coalgebra f a \rightarrow a \rightarrow b$ 

#### others **synchro**morphism ???

exomorphism

???

mutumorphism

???

algebra g b ' (1 -> g)		coalgebra ra · a	- 0	
<b>hylo</b> mo		rphism		
	cata	ana		
<b>dyna</b> morphism		codynamorphism		<b>generalized</b> apply the generalizations for both
histo	; ana	cata; futu		the relevant fold and unfold
chronom		norphism		the relevant lold and diffold
	histo; futu			
Elgot algebra		coElgot algebra		
may short-circuit while building		may short-circuit while tearing		
cata; a → b ∨ f a		$a \times g b \rightarrow b$ ; ana		
				•

reunfolds (tear down then build up a structure)

coalgebra  $g b \rightarrow (a \rightarrow b) \rightarrow algebra f a \rightarrow Fix f \rightarrow Fix g$ 

metamorphism	<b>g</b> eneralized	
ana; cata	apply both [un]fold	

combinations (combine two structures)

algebra  $f a \rightarrow Fix f \rightarrow Fix f \rightarrow a$ 

#### zippamorphism fa → a

#### mergamorphism

... which may fail to combine  $(f(Fix f) \times f(Fix f)) \vee fa \rightarrow a$ 

These can be combined in various ways. For example, a "zygohistomorphic prepromorphism" combines the zygo, histo, and prepro aspects into a signature like  $(f b \rightarrow b) \rightarrow (f \rightarrow f) \rightarrow (f (w (b \times a)) \rightarrow a) \rightarrow Fix f \rightarrow a$ 

Stolen from Edward Kmett's http://comonad.com/reader/ 2009/recursion-schemes/

\* This gives rise to a family of related recursion schemes, modeled in recursion-schemes with distributive law combinators

Two very useful extensions...

## 1. Operation Fusion

Combine catamorphism & anamorphism into a single operation (called a hylomorphism).

Hylomorphism sounds scary...

...but it's really simple. Pass the same arguments to 1 method instead of 2.

```
def ana[F[_]: Functor, A]
      (coalg: Coalgebra[F, A])(a: A): Fix[F]

def cata[F[_]: Functor, B]
      (alg: Algebra[F, B])(f: Fix[F]): B
```

```
def hylo[F[_]: Functor, A, B]
    (coalg: Coalgebra[F, A], alg: Algebra[F, B])(a: A): B
```

...but it's really simple. Pass the same arguments to 1 method instead of 2.

```
def unfold[F[_]: Functor, A]
      (coalg: Coalgebra[F, A])(a: A): Fix[F]

def fold[F[_]: Functor, B]
      (alg: Algebra[F, B])(f: Fix[F]): B
```

```
def unfoldIntoFold[F[_]: Functor, A, B]
     (coalg: Coalgebra[F, A], alg: Algebra[F, B])(a: A): B
```

Build-up & tear-down in one pass.

Generate & consume without actually creating the whole tree.

 $\Theta(n)$  instead of  $\Theta(2n)$ .

### 2. Monadic versions

Algebras can return monads.

#### What's a monad?

Oh god...

That's a separate talk *but*, speaking *extremely* loosely:

- A composable wrapper around data or intent.
- Something with map and flatMap methods.
- Something you can use in for comprehensions.

Some monads you've probably already used:

- Option[A]
- List[A]
- Future[A]
- Either[A, B] / Scalaz's disjunction A \/ B

```
val eval: Calc[Int] => Either[String, Int] = {
  case Num(i) => Right(i)
  case Div(_, 0) => Left("Division by zero: Australia says no!")
  case Div(a, b) => Right(a / b)
}
```

```
type Algebra [F[_], A] = F[A] => A
type CoAlgebra[F[_], A] = A => F[A]
```

```
type AlgebraM [M[_], F[_], A] = F[A] => M[A]
type CoAlgebraM[M[_], F[_], A] = A => M[F[A]]
```

```
def unfold[F[_]: Functor, A]
        (coalg: Coalgebra[F, A])(a: A): Fix[F]

def fold[F[_]: Functor, B]
        (alg: Algebra[F, B])(f: Fix[F]): B
```

```
def monadicUnfold[M[_]: Monad, F[_]: Traverse, A]
        (coalg: CoalgebraM[M, F, A])(a: A): M[Fix[F]]

def monadicFold[M[_]: Monad, F[_]: Traverse, B]
        (alg: AlgebraM[M, F, B])(f: Fix[F]): M[B]
```

- (Co)Algebra -> (Co)AlgebraM
- Result is now M[\_]
- Functor -> Traverse

```
type AlgebraM [M[_], F[_], A] = F[A] => M[A]
type CoAlgebraM[M[_], F[_], A] = A => M[F[A]]
```

```
val eval: Calc[Int] => String \/ Int = {
   case Num(i) => \/-(i)
   case Div(_, 0) => -\/("Division by zero detected.")
   case Div(a, b) => \/-(a / b)
}
```

```
// Calc[Int] => String \/ Int
// F [A] => M [A]

type StringOr[A] = String \/ A

val eval: AlgebraM[StringOr, Calc, Int] = {
```

Great! We've just added error handling and short-circuiting.

# **Real Example**

Random JSON generator.

### **Random Data**

Generating random data is important.

It's the secret sauce of property testing that ensures that, (asymptotically), you test your code with every possible, legal or desirable value.

#### Also useful for:

- load testing
- stress testing
- benchmarking (if generator supports determinism)

## Nyaya

```
import nyaya.gen._
val genWhatever: Gen[Whatever] =
  for {
    enabled <- Gen.boolean
    position <- Gen.chooseInt(-128, 128).pair
    stuff <- Gen.long.mapTo(Gen.string.option)
  } yield Whatever(enabled, position, stuff)</pre>
```

```
println(genWhatever.sample())

// Whatever(
// true,
// (-30, 83),
// Map(
// 2340946662719216224 -> Some("!XM91u"),
// 7161527918171176759 -> None
// )
// )
```

## Hint: Gen is a monad.

```
for {
   enabled <- Gen.boolean
   position <- Gen.chooseInt(-128, 128).pair
   stuff <- Gen.long.mapTo(Gen.string.option)
} yield Whatever(enabled, position, stuff)</pre>
```

Our imaginary app uses Play JSON.

Play JSON (like everything else) has hardcoded recursion.

We need to abstract over recursion and use fixpoints...

## **Fixpoint JSON**

## **Fixpoint JSON**

### Traverse (skeleton)

s/JsonF/YourDataType/g

## Traverse (body)

1x A: Use G.apply .
2x A: Use G.apply2 .
3x A: Use G.apply3 .
...
Collection of A s: Use .traverse then .map .

#### Done:

- Custom JSON data type
- Traverse & Functor instance

#### TODO:

- Generate JSON (using custom data type)
- Convert custom JSON into Play JSON

## Live code time...

# In summary...

#### Recursive data types:

- Trees, indented lists
- Tries
- DSLs, expression languages
- Anything self-referential
- Logical propositions, assertions
- Composable actions/data

- Allows you to not think about, or implement recursion.
- Where there'd normally be recursion, just declare what you want.
   Magic will grant your wish.
- Traditionally high barrier to entry. It's needless.
   You don't need to be a pilot to catch a plane.
   You don't need to understand Greek or category theory to use this.
- It wasn't discovered simply but now that it is, it's simple to use.

### **ALL YOU NEED TO KNOW TO BENEFIT IS:**

- Use a library or copy-and-paste little snippets that you need. They're tiny.
- Parametise your data; use Fix instead of self-referencing.
- When writing your logic, use plain old functions--just use the right shapes.



### **Libraries:**

#### My recursion micro-library

- https://github.com/japgolly/microlibs-scala
- Minimalistic: Few algebras & morphisms. Only supports Fix.
- Fast.
- Beginner focused.
  - Contains an EasyRecursion module with English instead of Greek.
  - Hopefully less intimidating for non-FP teams.
  - Nice stepping-stone before graduating to Matryoshka.

### **Libraries:**

#### Matryoshka

- https://github.com/slamdata/matryoshka
- Very comprehensive. Kitchen-sink of recursion.
- Will become part of Typelevel suite.
- Currently undergoing a lot of change.
- Super-smart people seriously working on it.
- Going to be awesome.

### Resources

- Youtube: Any talks by Greg Pfeil.
- Youtube: "Pure functional database programming with fixpoint types" by Rob Norris.
- Google: "recursion scheme cheatsheets".
- Paper: "Functional Programming with Bananas, Lenses, Envelopes and Barbed Wire"

## All done; Thank you!

Huge thanks to Rob Norris & Greg Pfeil for teaching me!

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