COMP261 (2018) — Assignment 5

String Search and Compression

In this assignment you will implement two string search and two compression algorithms which have been discussed in lectures, and explore their performance.

Resources

The assignment webpage also contains:

- An archive of the template code and data.
- A provisional marksheet.

Please read the marking schedule, as it lists the marks for each part of the assignment.

To Submit

You should submit:

- All the source code for your program, including the template code. Please make sure you do this, without it your assignment cannot be marked.
- Any other files your program needs to run that *aren't* the data files provided.
- A report which answers the questions asked in this handout. A significant portion of the marks are for the report.

Your report should be submitted as a txt or pdf file. No fancy formatting of the report is required.

Please **do not** submit the provided data files, as they are quite large.

Note that you will need to sign up for a 15 minute slot with the markers.

Handout code

The handout code constructs a small GUI with two panes, selected by tabs at the top of the window:

- The first is a very small text editor with one interesting function: string search. This will call a search method which you need to fill in.
- The second pane provides a GUI for selecting a file to compress with either Huffman coding or Lempel-Ziv compression.

You are free to change or ignore the handout code, but there's not much reason to.

Handout data

Several data files have been provided to search and compress.

- War and Peace. A novel by Tolstoy that is often used as sample text in machine learning. It is roughly 3 MB.
- **Taisho.** A ninth century dictionary from the Chinese Buddhist Canon, a collection of texts written in Classical Chinese. It uses a very large alphabet and is about 3 MB in size.
- **Pi.** The first million digits of Pi, totalling about 1 MB.
- **Lenna.** A hexdump of the famous image of Lenna, often used as an example in image processing. It is small, only 300 kB. This makes it good for quickly testing your algorithms.
- **Apollo.** A text version of the "the eagle has landed" sound recording, from Apollo 11. There are two channels, and hence two numbers at each time step. It is about 6 MB.

What to do

The assignment has four parts, worth approx. 30%, 30%, 30% and 10% (see marksheet for a more detailed breakdown of marks). Each part has some programming and some questions to be answered in your report.

Part 1: String search

In this section, your task is to implement the brute force and KMP string search algorithms to enable searching in the text editor.

A method stub is provided in the Search class which you should fill in. The search method takes two arguments: the text to search through and the string to search for. The method returns an integer as follows:

- The starting index of the first match in the text if one exists.
- -1 if no match exists.

You should implement two versions of this method, one using brute force search and one using the KMP search algorithm.

Note that KMP has two stages: computing the match table for input string, and performing the string search itself. You should write a separate method for each of these.

Once you have both algorithms implemented, experiment with the provided files to see what differences you can observe in their performance. To do this, you can add code to measure the time taken, or to count the number of steps taken, or both.

By default, the provided code tries to load the War and Peace code initially, assuming that you have the data files in a directory called data. You may change this behaviour and/or location if you wish.

Question 1: Write a short summary of the performance you observed using the two search algorithms.

Part 2: Huffman coding

Your task in this part is to implement the Huffman coding and decoding algorithm, as described in lectures, and answer Questions 2 and 3.

A full implementation does three things:

- Create a tree of binary codes for each character in the input text.
- Encode input text using that tree.
- Decode previously encoded text with that tree.

You will need to write methods to do all three of these steps for a particular text. Remember, you need to dynamically generate the tree from a given input text, and *not* use a fixed tree that you supply manually.

Here are some implementation notes:

- The HuffmanCoding class has three methods that you should fill in.
- The encode method should return a *binary string*, a string containing only 1's and 0's. Similarly, decode takes a binary string as its argument.
- You could store the binary codes for each character in a Map<Character, String>.

• To debug your code, you could manually make an encoding tree and then generate a text using its frequencies.

Question 2: Report the binary tree of codes your algorithm generates, and the final size of War and Peace after Huffman coding.

Question 3: Consider the Huffman coding of war_and_peace.txt, taisho.txt, and pi.txt. Which of these achieves the best compression, i.e. the best reduction in size? What makes some of the encodings better than others?

Part 3: Lempel-Ziv compression

In this part, your task is to implement the Lempel-Ziv 77 compression and decompression algorithms, as described in lectures, and answer Questions 4 and 5.

Implementation notes:

- The LempelZiv class has two methods you should fill in.
- None of the provided data files include the characters [,], or |. This means you can use them to start, end, and delimit your tuples respectively.
- To debug your code, you could make some small files containing carefully constructed strings (all one character, one repetition, etc.) and check you get the expected result.

Question 4: The Lempel-Ziv algorithm has a parameter: the size of the sliding window. On a text of your choice, how does changing the window size affect the quality of the compression?

Question 5: What happens if you Huffman encode War and Peace *before* applying Lempel-Ziv compression to it? Do you get a smaller file size (in characters) overall?

Part 4: Challenge — Better string search or coding

You have a choice for the challenge, pick **one** of the following.

Better coding. Implement a more sophisticated coding scheme. Several of these exist, but two good candidates are:

- Adaptive Huffman coding. This uses a coding tree that changes throughout the text to get better performance in the case where the distribution of characters changes throughout the text.
- Arithmetic coding. This takes a different approach to coding, where strings are encoded into a single floating point number between 0 and 1.

Boyer-Moore. Implement the Boyer-Moore string search algorithm, which is faster but more complex than KMP. It improves efficiency by doing a more complex search starting from the end of a pattern instead of the start.

Lastly, there are a few spare marks available for an interesting question about coding.

Question 6: Suppose Alice has two binary strings (made only of 1's and 0's), X and Y. Make a pair of algorithms so that:

- Alice can encode X and Y into a single binary string Z, which she sends to Bob.
- Bob receives Z and can unambiguously decode it back into X and Y.

For example, a simple solution would be to concatenate X and Y together. This works if X = 101 and Y = 011. Bob then receives Z = 101011 and splits it in half to retrieve X and Y. But what if X = 10 and Y = 1011? The goal is to craft your encoding and decoding algorithms so that Z is as short as possible. Start by making algorithms so that |Z| = 2*(|X| + |Y|) or (where |Q| is the length of string Q).

You get the marks for this question if you find an algorithm such that |Z| is significantly less than 2*(|X| + |Y|).

Writing it yourself

Make sure that you write the code for the data structures yourself – you will not learn what you need to learn if you use code from somewhere else. You can build on code examples from somewhere else, but do not simply copy large segments of code and make sure that you acknowledge the source appropriately. If we identify any plagiarism, we **will** penalise it.