Standard Code Library

Your TeamName

Your School

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Contents

一切的开始

宏定义

- 需要 C++11

```cpp
#include <bits/stdc++.h>
using namespace std;
using LL = long long;
#define FOR(i, x, y) for (decay<decltype(y)>::type i = (x), _##i = (y); i < _##i; ++i)
#define FORD(i, x, y) for (decay<decltype(x)>::type i = (x), _##i = (y); i > _##i; --i)
#ifdef DEBUG
#ifndef ONLINE_JUDGE
#define zerol
#endif
#endif
#ifdef zerol
#define dbg(x...) do { cout << "\033[32;1m" << #x << " -> "; err(x); } while (0)
void err() { cout << "\033[39;0m" << endl; }
template<template<typename...> class T, typename t, typename... A>
void err(T<t> a, A... x) { for (auto v: a) cout << v << ' '; err(x...); }
template<typename T, typename... A>
void err(T a, A... x) { cout << a << ' '; err(x...); }
#else
#define dbg(...)
#define err(...)
#endif
// -----------------------------------------------------------------------
```

- 调试时添加编译选项 `-DDEBUG`，提交时注释
- 注意检查判题系统编译选项，修改 `#ifndef ONLINE_JUDGE`
- `FOR` ++ 循环 `FOR(循环变量名称, 循环变量起始值, 循环变量结束值（不含）)`
- `FORD` -- 循环
- `err()` 调试时输出（支持单层迭代）
- `dbg()` 变色输出变量名和变量值（支持单层迭代）
- 黄色 `33`，蓝色 `34`，橙色 `31`

对拍

- Linux

```bash
#!/usr/bin/env bash
g++ -o r main.cpp -O2 -std=c++11
g++ -o std std.cpp -O2 -std=c++11
while true; do
    python gen.py > in
    ./std < in > stdout
    ./r < in > out
    if test $? -ne 0; then
        exit 0
    fi
    if diff stdout out; then
        printf "AC\n"
    else
        printf "GG\n"
        exit 0
    fi
done
```

- Windows

```bat
@echo off
setlocal enabledelayedexpansion

g++ -o r main.cpp -O2 -std=c++11
g++ -o std std.cpp -O2 -std=c++11

:loop
python gen.py > in
if !errorlevel! neq 0 exit /b
```

```
std.exe < in > stdout
if !errorlevel! neq 0 exit /b

r.exe < in > out
if !errorlevel! neq 0 exit /b

fc /b stdout out > nul
if !errorlevel! equ 0 (
    echo AC
) else (
    echo GG
    exit /b
)

goto loop
```

快速编译运行（配合无插件 VSC）

- Linux

```bash
#!/bin/bash
g++ $1.cpp -o $1 -O2 -std=c++14 -Wall -Dzerol -g
if $? -eq 0; then
    ./$1
fi
```

- Windows

```
@echo off
:: 参数为文件名（不含.cpp后缀）
g++ %1.cpp -o %1 -O2 -std=c++14 -Wall -Dzerol -g
if %errorlevel% equ 0 (
    %1.exe
)
```

数据结构

ST 表

- 一维

```cpp
#define M 10

struct RMQ {
    int f[22][M];
    inline int highbit(int x) { return 31 - __builtin_clz(x); }
    void init(int* v, int n) {
        FOR (i, 0, n) f[0][i] = v[i];
        FOR (x, 1, highbit(n) + 1)
        FOR (i, 0, n - (1 << x) + 1)
        f[x][i] = min(f[x - 1][i], f[x - 1][i + (1 << (x - 1))]);
    }
    int get_min(int l, int r) {
        assert(l <= r);
        int t = highbit(r - l + 1);
        return min(f[t][l], f[t][r - (1 << t) + 1]);
    }
};
```

- 二维

```cpp
#define maxn 10
LL n, m, a[maxn][maxn];

struct RMQ2D{
    int f[maxn][maxn][10][10];
    inline int highbit(int x) { return 31 - __builtin_clz(x); }
```

```
    inline int calc(int x, int y, int xx, int yy, int p, int q) {
        return max(
            max(f[x][y][p][q], f[xx - (1 << p) + 1][yy - (1 << q) + 1][p][q]),
            max(f[xx - (1 << p) + 1][y][p][q], f[x][yy - (1 << q) + 1][p][q])
            );
    }
    void init() {
        FOR (x, 0, highbit(n) + 1)
        FOR (y, 0, highbit(m) + 1)
        FOR (i, 0, n - (1 << x) + 1)
        FOR (j, 0, m - (1 << y) + 1) {
            if (!x && !y) { f[i][j][x][y] = a[i][j]; continue; }
            f[i][j][x][y] = calc(
                i, j,
                i + (1 << x) - 1, j + (1 << y) - 1,
                max(x - 1, 0), max(y - 1, 0)
                );
        }
    }
    inline int get_max(int x, int y, int xx, int yy) {
        return calc(x, y, xx, yy, highbit(xx - x + 1), highbit(yy - y + 1));
    }
};
```

线段树

朴素线段树

- 默认为最大值，可自行修改 `struct Q` `struct P` `P operator &`
- 注意建树时的下标问题 (1-based)

```
const LL INF = LONG_LONG_MAX;
#define maxn 10
LL n;

namespace SGT {
    struct Q {
        LL setv;
        explicit Q(LL setv = -1): setv(setv) {}
        void operator += (const Q& q) { if (q.setv != -1) setv = q.setv; }
    };
    struct P {
        LL max;
        explicit P(LL max = -INF): max(max) {}
        void up(Q& q) { if (q.setv != -1) max = q.setv; }
    };
    template<typename T>
    P operator & (T&& a, T&& b) {
        return P(max(a.max, b.max));
    }
    P p[maxn << 2];
    Q q[maxn << 2];
#define lson o * 2, l, (l + r) / 2
#define rson o * 2 + 1, (l + r) / 2 + 1, r
    void up(int o, int l, int r) {
        if (l == r) p[o] = P();
        else p[o] = p[o * 2] & p[o * 2 + 1];
        p[o].up(q[o]);
    }
    void down(int o, int l, int r) {
        q[o * 2] += q[o]; q[o * 2 + 1] += q[o];
        q[o] = Q();
        up(lson); up(rson);
    }
    template<typename T>
    void build(T&& f, int o = 1, int l = 1, int r = n) {
        if (l == r) q[o] = f(l);
        else { build(f, lson); build(f, rson); q[o] = Q(); }
        up(o, l, r);
    }
    P query(int ql, int qr, int o = 1, int l = 1, int r = n) {
```

```cpp
        if (ql > r || l > qr) return P();
        if (ql <= l && r <= qr) return p[o];
        down(o, l, r);
        return query(ql, qr, lson) & query(ql, qr, rson);
    }
    void update(int ql, int qr, const Q& v, int o = 1, int l = 1, int r = n) {
        if (ql > r || l > qr) return;
        if (ql <= l && r <= qr) q[o] += v;
        else {
            down(o, l, r);
            update(ql, qr, v, lson); update(ql, qr, v, rson);
        }
        up(o, l, r);
    }
}

// -----------------------------------------------------------------
void solve(){
    vector<LL> arr = {1, 5, 7, 4, 2, 8 ,3, 6, 10, 9};
    n = arr.size();
    SGT::build([&](int idx){
        return SGT::Q(arr[idx-1]);
    });
    for(LL i=1; i<=n; i++){
        dbg(SGT::query(1, i).max);
    }
    SGT::update(2, 4, SGT::Q(-3));
    cout << "MODIFIED\n";
    for(LL i=1; i<=n; i++){
        dbg(SGT::query(1, i).max);
    }
}
```

- 区间修改，区间累加，查询区间和、最大值、最小值。

```cpp
#define maxn 100005
#define INF LONG_LONG_MAX
LL a[maxn];

struct IntervalTree {
#define ls o * 2, l, m
#define rs o * 2 + 1, m + 1, r
    static const LL M = maxn * 4, RS = 1E18 - 1;
    LL addv[M], setv[M], minv[M], maxv[M], sumv[M];
    int n;
    void init() {
        memset(addv, 0, sizeof addv);
        fill(setv, setv + M, RS);
        memset(minv, 0, sizeof minv);
        memset(maxv, 0, sizeof maxv);
        memset(sumv, 0, sizeof sumv);
    }
    void maintain(LL o, LL l, LL r) {
        if (l < r) {
            LL lc = o * 2, rc = o * 2 + 1;
            sumv[o] = sumv[lc] + sumv[rc];
            minv[o] = min(minv[lc], minv[rc]);
            maxv[o] = max(maxv[lc], maxv[rc]);
        } else sumv[o] = minv[o] = maxv[o] = 0;
        if (setv[o] != RS) { minv[o] = maxv[o] = setv[o]; sumv[o] = setv[o] * (r - l + 1); }
        if (addv[o]) { minv[o] += addv[o]; maxv[o] += addv[o]; sumv[o] += addv[o] * (r - l + 1); }
    }
    void build(LL o, LL l, LL r) {
        if (l == r) addv[o] = a[l];
        else {
            LL m = (l + r) / 2;
            build(ls); build(rs);
        }
        maintain(o, l, r);
    }
    void pushdown(LL o) {
        LL lc = o * 2, rc = o * 2 + 1;
```

```cpp
        if (setv[o] != RS) {
            setv[lc] = setv[rc] = setv[o];
            addv[lc] = addv[rc] = 0;
            setv[o] = RS;
        }
        if (addv[o]) {
            addv[lc] += addv[o]; addv[rc] += addv[o];
            addv[o] = 0;
        }
    }
    void update(LL p, LL q, LL o, LL l, LL r, LL v, LL op) {
        if (p <= r && l <= q){
            if (p <= l && r <= q) {
                if (op == 2) { setv[o] = v; addv[o] = 0; }
                else addv[o] += v;
            } else {
                pushdown(o);
                LL m = (l + r) / 2;
                update(p, q, ls, v, op); update(p, q, rs, v, op);
            }
        }
        maintain(o, l, r);
    }
    void query(LL p, LL q, LL o, LL l, LL r, LL add, LL& ssum, LL& smin, LL& smax) {
        if (p > r || l > q) return;
        if (setv[o] != RS) {
            LL v = setv[o] + add + addv[o];
            ssum += v * (min(r, q) - max(l, p) + 1);
            smin = min(smin, v);
            smax = max(smax, v);
        } else if (p <= l && r <= q) {
            ssum += sumv[o] + add * (r - l + 1);
            smin = min(smin, minv[o] + add);
            smax = max(smax, maxv[o] + add);
        } else {
            LL m = (l + r) / 2;
            query(p, q, ls, add + addv[o], ssum, smin, smax);
            query(p, q, rs, add + addv[o], ssum, smin, smax);
        }
    }
    // 简化接口
    void build(int _n) {
        n = _n;
        build(1, 1, n);
    }

    void range_add(int l, int r, int val) {
        update(l, r, 1, 1, n, val, 1);
    }

    void range_set(int l, int r, int val) {
        update(l, r, 1, 1, n, val, 2);
    }

    void range_query(int l, int r, LL& sum, LL& min_val, LL& max_val) {
        sum = 0;
        min_val = INF;
        max_val = -INF;
        query(l, r, 1, 1, n, 0, sum, min_val, max_val);
    }
} IT;
// ---------------------------------------------------------------------
void solve(){
    IT.init();

    LL n = 5;
    vector<int> data = {1, 3, 5, 7, 9};
    for (int i = 0; i < n; i++) {
        a[i + 1] = data[i];  // 注意：线段树从 1 开始索引
    }

```

```
    IT.build(n);

    LL sum, min_val, max_val;
    IT.range_query(1, 5, sum, min_val, max_val);
    cout << " " << sum << " " << min_val << " " << max_val << endl;

    IT.range_add(2, 4, 2);
    IT.range_query(1, 5, sum, min_val, max_val);
    cout << " " << sum << " " << min_val << " " << max_val << endl;

    IT.range_set(3, 5, 10);
    IT.range_query(1, 5, sum, min_val, max_val);
    cout << " " << sum << " " << min_val << " " << max_val << endl;

    IT.range_query(2, 4, sum, min_val, max_val);
    cout << " " << sum << " " << min_val << " " << max_val << endl;
}
```

动态开点

```
namespace SGT{
    const LL N = 3e5 + 10, INF = LONG_LONG_MAX;
    LL sum[N << 2], lazy[N << 2];
//  LL minn[N << 2], lazy2[N << 2];
    LL lson[N << 2], rson[N << 2],  tot = 0, root = 0;

    inline void push_up(LL rt){
        sum[rt] = sum[lson[rt]] + sum[rson[rt]];
//      minn[rt] = min(minn[lson[rt]], minn[rson[rt]]);
    }
    inline void push_down(LL rt, LL m){
        if(!lazy[rt]) return;
        if(!lson[rt]){
            lson[rt] = ++tot;
//          minn[lson[rt]] = INF;
        }
        if(!rson[rt]){
            rson[rt] = ++tot;
//          minn[rson[rt]] = INF;
        }
        lazy[lson[rt]] += lazy[rt], lazy[rson[rt]] += lazy[rt];
        sum[lson[rt]] += lazy[rt] * (m - (m >> 1));
        sum[rson[rt]] += lazy[rt] * (m >> 1);
        lazy[rt] = 0;

//      lazy2[lson[rt]] = min(lazy2[lson[rt]], lazy2[rt]);
//      lazy2[rson[rt]] = min(lazy2[rson[rt]], lazy2[rt]);
//      minn[lson[rt]] = min(minn[lson[rt]], lazy2[rt]);
//      minn[rson[rt]] = min(minn[rson[rt]], lazy2[rt]);
//      lazy2[rt] = INF;
    }

    static void add_range(LL &rt ,LL l, LL r, LL L, LL R, LL val){
        if(!rt){
            rt = ++tot;
//          minn[rt] = INF;
//          lazy2[rt] = INF;
        }
        if(l >= L && r <= R){
            lazy[rt] += val;
            sum[rt] += val * (r - l + 1);
//          minn[rt] = min(minn[rt], val);
//          lazy2[rt] = min(lazy2[rt], val);
            return;
        }
        push_down(rt, r - l + 1);
        LL mid = l + r >> 1;
        if(mid >= L) add_range(lson[rt], l, mid, L, R, val);
        if(mid < R) add_range(rson[rt], mid + 1, r, L, R, val);
        push_up(rt);
    }
```

```cpp

    static void add_point(LL &rt, LL l, LL r, LL pos, LL val){
        if(!rt){
            rt = ++tot;
//          minn[rt] = INF;
        }
        if(l == r){
            sum[rt] += val;
//          minn[rt] = min(minn[rt], val);
            return;
        }
        LL mid = l + r >> 1;
        if(mid >= pos) add_point(lson[rt], l, mid, pos, val);
        else add_point(rson[rt], mid + 1, r, pos, val);
        push_up(rt);
    }

    static LL query(LL rt, LL l, LL r, LL L, LL R){
        if(!rt) return 0;
//      if(!rt) return INF;
        if(l >= L && r <= R) return sum[rt];
//      if(l >= L && r <= R) return minn[rt];
        push_down(rt, r - l + 1);
        LL mid = l + r >> 1;
        LL ans = 0;
//      LL ans = INF;
        if(mid >= L) ans += query(lson[rt], l, mid, L, R);
        if(mid < R) ans += query(rson[rt], mid + 1, r, L, R);
//      if(mid >= L) ans = min(ans, query(lson[rt], l, mid, L, R));
//      if(mid < R) ans = min(ans, query(rson[rt], mid + 1, r, L, R));
        return ans;
    }
};
// --------------------------- Template End ----------------------------
// !!!!!!!!!!!!!!!!!!!!!!!!!!! Tester Start !!!!!!!!!!!!!!!!!!!!!!!!!!!!

void solve(){
    LL rt = 0, l = -1e9, r = 1e9; // 根（必需为 0）和值域（操作范围）
    SGT::root = rt;
    SGT::add_point(rt, l, r, 1e4, 1e4);
    SGT::add_range(rt, l, r, 100000, 100010, 100);
    cout << SGT::query(rt, l, r, -1e9, 0) << '\n';
    cout << SGT::query(rt, l, r, 0, 1e5) << '\n';
    cout << SGT::query(rt, l, r, 1e5, 1e6) << '\n';
    cout << SGT::query(rt, l, r, 1e6, 1e10) << '\n';
}
```

树状数组

- 单点修改，区间查询
- 频次统计下的 k 小值
- 维护差分数组时的区间修改，单点查询

```cpp
#define M 100005

namespace BIT {
    LL c[M]; // 注意初始化开销
    inline int lowbit(int x) { return x & -x; }
    void add(int x, LL v) { // 单点加
        for (int i = x; i < M; i += lowbit(i))
            c[i] += v;
    }
    LL sum(int x) { // 前缀和
        LL ret = 0;
        for (int i = x; i > 0; i -= lowbit(i))
            ret += c[i];
        return ret;
    }
    int kth(LL k) { // 频次统计下从小到大第 k 个，详见应用
        int p = 0;
```

```cpp
        for (int lim = 1 << 20; lim; lim /= 2)
            if (p + lim < M && c[p + lim] < k) {
                p += lim;
                k -= c[p];
            }
        return p + 1;
    }
    LL sum(int l, int r) { return sum(r) - sum(l - 1); } // 区间和
    // 区间加（此时树状数组为差分数组，sum(x) 为第 x 个数的值）
    void add(int l, int r, LL v) { add(l, v); add(r + 1, -v); }
}
// ----------------------------------------------------------------------
void solve(){
    vector<LL> a={9, 9, 9, 9, 5, 3, 3, 3, 1, 1};
    LL n = a.size(), i;
    for(i=1; i<=n; i++) BIT::add(a[i-1], 1);
    // 1 1 3 3 3 5 9 9 9 9
    for(i=1; i<=n; i++) cout << BIT::kth(i) << ' ';
}
```

- 区间修改、区间查询

```cpp
#define maxn 100005

namespace BIT {
    int n;
    int c[maxn], cc[maxn];
    inline int lowbit(int x) { return x & -x; }
    void init(int siz){ // 初始化
        n = siz;
        for(LL i=0; i<=n; i++){
            c[i] = cc[i] = 0;
        }
    }
    void add(int x, int v) { // 不要用这个
        for (int i = x; i <= n; i += lowbit(i)) {
            c[i] += v; cc[i] += x * v;
        }
    }
    void add(int l, int r, int v) { add(l, v); add(r + 1, -v); } // 区间修改
    int sum(int x) { // 前缀和
        int ret = 0;
        for (int i = x; i > 0; i -= lowbit(i))
            ret += (x + 1) * c[i] - cc[i];
        return ret;
    }
    int sum(int l, int r) { return sum(r) - sum(l - 1); } // 区间和
}
// ----------------------------------------------------------------------
void solve(){
    LL i, n=8;
    BIT::init(n);
    BIT::add(2, 4, 2);
    for(i=1; i<=n; i++) cout << BIT::sum(i, i) << ' ';
    cout << '\n';
    cout << BIT::sum(5) << '\n';
    cout << BIT::sum(2, 3) << '\n';
}
```

- 三维

```cpp
#define maxn 105

namespace BIT{
    int n;
    LL c[maxn][maxn][maxn];
    inline int lowbit(int x) { return x & -x; }
    void init(int siz){
        n = siz;
        for(int i=0; i<=n; i++){
            for(int j=0; j<=n; j++){
                for(int k=0; k<=n; k++){
```

```
                    c[i][j][k] = 0;
                }
            }
        }
    }
    void update(int x, int y, int z, int d) {
        for (int i = x; i <= n; i += lowbit(i))
            for (int j = y; j <= n; j += lowbit(j))
                for (int k = z; k <= n; k += lowbit(k))
                    c[i][j][k] += d;
    }
    LL query(int x, int y, int z) {
        LL ret = 0;
        for (int i = x; i > 0; i -= lowbit(i))
            for (int j = y; j > 0; j -= lowbit(j))
                for (int k = z; k > 0; k -= lowbit(k))
                    ret += c[i][j][k];
        return ret;
    }
    LL solve(int x, int y, int z, int xx, int yy, int zz) {
        return    query(xx, yy, zz)
        - query(xx, yy, z - 1)
        - query(xx, y - 1, zz)
        - query(x - 1, yy, zz)
        + query(xx, y - 1, z - 1)
        + query(x - 1, yy, z - 1)
        + query(x - 1, y - 1, zz)
        - query(x - 1, y - 1, z - 1);
    }
}
```

数学

快速乘

```
LL mul(LL a, LL b, LL m) {
    LL ret = 0;
    while (b) {
        if (b & 1) {
            ret += a;
            if (ret >= m) ret -= m;
        }
        a += a;
        if (a >= m) a -= m;
        b >>= 1;
    }
    return ret;
}
```

- O(1)

```
LL mul(LL u, LL v, LL p) {
    return (u * v - LL((long double) u * v / p) * p + p) % p;
}
LL mul(LL u, LL v, LL p) { // 卡常
    LL t = u * v - LL((long double) u * v / p) * p;
    return t < 0 ? t + p : t;
}
```

高斯消元

- n 是方程个数，m 是未知量个数，a[n][m+1] 是增广矩阵
- x[m] 是每个未知量的解（如果有），free_x[m] 是每个未知量是否为自由变量。

```
typedef double LD;
const LD eps = 1E-10;
const int maxn = 2000 + 10;

int n, m;
```

```cpp
LD a[maxn][maxn], x[maxn];
bool free_x[maxn];

inline int sgn(LD x) { return (x > eps) - (x < -eps); }

int gauss(LD a[maxn][maxn], int n, int m) {
//int gauss() {
    memset(free_x, 1, sizeof free_x); memset(x, 0, sizeof x);
    int r = 0, c = 0;
    while (r < n && c < m) {
        int m_r = r;
        FOR (i, r + 1, n)
        if (fabs(a[i][c]) > fabs(a[m_r][c])) m_r = i;
        if (m_r != r)
            FOR (j, c, m + 1)
            swap(a[r][j], a[m_r][j]);
        if (!sgn(a[r][c])) {
            a[r][c] = 0;
            ++c;
            continue;
        }
        FOR (i, r + 1, n)
        if (a[i][c]) {
            LD t = a[i][c] / a[r][c];
            FOR (j, c, m + 1) a[i][j] -= a[r][j] * t;
        }
        ++r; ++c;
    }
    FOR (i, r, n)
    if (sgn(a[i][m])) return -1;
    if (r < m) {
        FORD (i, r - 1, -1) {
            int f_cnt = 0, k = -1;
            FOR (j, 0, m)
            if (sgn(a[i][j]) && free_x[j]) {
                ++f_cnt;
                k = j;
            }
            if(f_cnt > 0) continue;
            LD s = a[i][m];
            FOR (j, 0, m)
            if (j != k) s -= a[i][j] * x[j];
            x[k] = s / a[i][k];
            free_x[k] = 0;
        }
        return m - r;
    }
    FORD (i, m - 1, -1) {
        LD s = a[i][m];
        FOR (j, i + 1, m)
        s -= a[i][j] * x[j];
        x[i] = s / a[i][i];
    }
    return 0;
}
```

快速幂

- 如果模数是素数，则可在函数体内加上 `n %= MOD - 1;`（费马小定理）。

```cpp
LL bin(LL x, LL n, LL MOD) {
    LL ret = MOD != 1;
    for (x %= MOD; n; n >>= 1, x = x * x % MOD)
        if (n & 1) ret = ret * x % MOD;
    return ret;
}
```

- 防爆 LL
- 前置模板：快速乘

```cpp
LL bin(LL x, LL n, LL MOD) {
```

```cpp
    LL ret = MOD != 1;
    for (x %= MOD; n; n >>= 1, x = mul(x, x, MOD))
        if (n & 1) ret = mul(ret, x, MOD);
    return ret;
}
```

高精度

- https://github.com/Baobaobear/MiniBigInteger/blob/main/bigint_tiny.h，带有压位优化
- 按需实现

```cpp
#include <algorithm>
#include <cstdio>
#include <string>
#include <vector>

struct BigIntTiny {
    int sign;
    std::vector<int> v;

    BigIntTiny() : sign(1) {}
    BigIntTiny(const std::string &s) { *this = s; }
    BigIntTiny(int v) {
        char buf[21];
        sprintf(buf, "%d", v);
        *this = buf;
    }
    void zip(int unzip) {
        if (unzip == 0) {
            for (int i = 0; i < (int)v.size(); i++)
               v[i] = get_pos(i * 4) + get_pos(i * 4 + 1) * 10 + get_pos(i * 4 + 2) * 100 + get_pos(i * 4 + 3) * 1000;
        } else
            for (int i = (v.resize(v.size() * 4), (int)v.size() - 1), a; i >= 0; i--)
                a = (i % 4 >= 2) ? v[i / 4] / 100 : v[i / 4] % 100, v[i] = (i & 1) ? a / 10 : a % 10;
        setsign(1, 1);
    }
    int get_pos(unsigned pos) const { return pos >= v.size() ? 0 : v[pos]; }
    BigIntTiny &setsign(int newsign, int rev) {
        for (int i = (int)v.size() - 1; i > 0 && v[i] == 0; i--)
            v.erase(v.begin() + i);
        sign = (v.size() == 0 || (v.size() == 1 && v[0] == 0)) ? 1 : (rev ? newsign * sign : newsign);
        return *this;
    }
    std::string to_str() const {
        BigIntTiny b = *this;
        std::string s;
        for (int i = (b.zip(1), 0); i < (int)b.v.size(); ++i)
            s += char(*(b.v.rbegin() + i) + '0');
        return (sign < 0 ? "-" : "") + (s.empty() ? std::string("0") : s);
    }
    bool absless(const BigIntTiny &b) const {
        if (v.size() != b.v.size()) return v.size() < b.v.size();
        for (int i = (int)v.size() - 1; i >= 0; i--)
            if (v[i] != b.v[i]) return v[i] < b.v[i];
        return false;
    }
    BigIntTiny operator-() const {
        BigIntTiny c = *this;
        c.sign = (v.size() > 1 || v[0]) ? -c.sign : 1;
        return c;
    }
    BigIntTiny &operator=(const std::string &s) {
        if (s[0] == '-')
            *this = s.substr(1);
        else {
            for (int i = (v.clear(), 0); i < (int)s.size(); ++i)
                v.push_back(*(s.rbegin() + i) - '0');
            zip(0);
        }
        return setsign(s[0] == '-' ? -1 : 1, sign = 1);
```

```cpp
    }
    bool operator<(const BigIntTiny &b) const {
        return sign != b.sign ? sign < b.sign : (sign == 1 ? absless(b) : b.absless(*this));
    }
    bool operator==(const BigIntTiny &b) const { return v == b.v && sign == b.sign; }
    BigIntTiny &operator+=(const BigIntTiny &b) {
        if (sign != b.sign) return *this = (*this) - -b;
        v.resize(std::max(v.size(), b.v.size()) + 1);
        for (int i = 0, carry = 0; i < (int)b.v.size() || carry; i++) {
            carry += v[i] + b.get_pos(i);
            v[i] = carry % 10000, carry /= 10000;
        }
        return setsign(sign, 0);
    }
    BigIntTiny operator+(const BigIntTiny &b) const {
        BigIntTiny c = *this;
        return c += b;
    }
    void add_mul(const BigIntTiny &b, int mul) {
        v.resize(std::max(v.size(), b.v.size()) + 2);
        for (int i = 0, carry = 0; i < (int)b.v.size() || carry; i++) {
            carry += v[i] + b.get_pos(i) * mul;
            v[i] = carry % 10000, carry /= 10000;
        }
    }
    BigIntTiny operator-(const BigIntTiny &b) const {
        if (b.v.empty() || b.v.size() == 1 && b.v[0] == 0) return *this;
        if (sign != b.sign) return (*this) + -b;
        if (absless(b)) return -(b - *this);
        BigIntTiny c;
        for (int i = 0, borrow = 0; i < (int)v.size(); i++) {
            borrow += v[i] - b.get_pos(i);
            c.v.push_back(borrow);
            c.v.back() -= 10000 * (borrow >>= 31);
        }
        return c.setsign(sign, 0);
    }
    BigIntTiny operator*(const BigIntTiny &b) const {
        if (b < *this) return b * *this;
        BigIntTiny c, d = b;
        for (int i = 0; i < (int)v.size(); i++, d.v.insert(d.v.begin(), 0))
            c.add_mul(d, v[i]);
        return c.setsign(sign * b.sign, 0);
    }
    BigIntTiny operator/(const BigIntTiny &b) const {
        BigIntTiny c, d;
        BigIntTiny e=b;
        e.sign=1;

        d.v.resize(v.size());
        double db = 1.0 / (b.v.back() + (b.get_pos((unsigned)b.v.size() - 2) / 1e4) +
                           (b.get_pos((unsigned)b.v.size() - 3) + 1) / 1e8);
        for (int i = (int)v.size() - 1; i >= 0; i--) {
            c.v.insert(c.v.begin(), v[i]);
            int m = (int)((c.get_pos((int)e.v.size()) * 10000 + c.get_pos((int)e.v.size() - 1)) * db);
            c = c - e * m, c.setsign(c.sign, 0), d.v[i] += m;
            while (!(c < e))
                c = c - e, d.v[i] += 1;
        }
        return d.setsign(sign * b.sign, 0);
    }
    BigIntTiny operator%(const BigIntTiny &b) const { return *this - *this / b * b; }
    bool operator>(const BigIntTiny &b) const { return b < *this; }
    bool operator<=(const BigIntTiny &b) const { return !(b < *this); }
    bool operator>=(const BigIntTiny &b) const { return !(*this < b); }
    bool operator!=(const BigIntTiny &b) const { return !(*this == b); }
};
```

矩阵运算

```
#define MOD 998244353
#define M 10

struct Mat {
    LL m;
    LL v[M][M];
    Mat(int siz=2) {
        m = siz;
        for(int i=0; i<=m; i++){
            for(int j=0; j<=m; j++){
                v[i][j] = 0;
            }
        }
    }
    void eye() { FOR (i, 0, m) v[i][i] = 1; }
    LL* operator [] (LL x) { return v[x]; }
    const LL* operator [] (LL x) const { return v[x]; }
    Mat operator * (const Mat& B) {
        const Mat& A = *this;
        Mat ret;
        FOR (k, 0, m)
        FOR (i, 0, m) if (A[i][k])
            FOR (j, 0, m)
            ret[i][j] = (ret[i][j] + A[i][k] * B[k][j]) % MOD;
        return ret;
    }
    Mat pow(LL n) const {
        Mat A = *this, ret; ret.eye();
        for (; n; n >>= 1, A = A * A)
            if (n & 1) ret = ret * A;
        return ret;
    }
    Mat operator + (const Mat& B) {
        const Mat& A = *this;
        Mat ret;
        FOR (i, 0, m)
        FOR (j, 0, m)
        ret[i][j] = (A[i][j] + B[i][j]) % MOD;
        return ret;
    }
    void pprint() const {
        FOR (i, 0, m)
        FOR (j, 0, m)
        printf("%lld%c", (*this)[i][j], j == m - 1 ? '\n' : ' ');
    }
};
// -------------------------------------------------------------------
void solve(){
    Mat mat1, mat2;
    mat1.eye();
    mat1[1][0] = 2; // 0-based
    mat2.eye();
    mat2[1][1] = 4;
    Mat mat3 = mat1 * mat2;
    mat3.pprint();
}
```

数论分块

$f(i) = \lfloor \frac{n}{i} \rfloor = v$ 时 i 的取值范围是 $[l, r]$。

```
void sqrt_decomposition(LL n){
    for (LL l = 1, v, r; l <= n; l = r + 1) {
        v = n / l; r = n / v;
        printf("%lld / [%lld, %lld] = %lld\n", n, l, r, v);
    }
}
```

质数筛

- $\mathcal{O}(n)$

```cpp
const LL p_max = 1E6 + 100;
LL pr[p_max], p_sz;
void get_prime() {
    static bool vis[p_max];
    FOR (i, 2, p_max) {
        if (!vis[i]) pr[p_sz++] = i;
        FOR (j, 0, p_sz) {
            if (pr[j] * i >= p_max) break;
            vis[pr[j] * i] = 1;
            if (i % pr[j] == 0) break;
        }
    }
}
```

欧拉函数

朴素

```cpp
int phi(int x)
{
    int res = x;
    for (int i = 2; i <= x / i; i ++ )
        if (x % i == 0)
        {
            res = res / i * (i - 1);
            while (x % i == 0) x /= i;
        }
    if (x > 1) res = res / x * (x - 1);

    return res;
}
```

筛法求欧拉函数

- 前置模板：质数筛

```cpp
const LL p_max = 1E5 + 100;
LL phi[p_max];
void get_phi() {
    phi[1] = 1;
    static bool vis[p_max];
    static LL prime[p_max], p_sz, d;
    FOR (i, 2, p_max) {
        if (!vis[i]) {
            prime[p_sz++] = i;
            phi[i] = i - 1;
        }
        for (LL j = 0; j < p_sz && (d = i * prime[j]) < p_max; ++j) {
            vis[d] = 1;
            if (i % prime[j] == 0) {
                phi[d] = phi[i] * prime[j];
                break;
            }
            else phi[d] = phi[i] * (prime[j] - 1);
        }
    }
}
```

素性测试

试除法

- $\mathcal{O}(\sqrt{n})$

```cpp
bool is_prime(int x)
{
```

```
    if (x < 2) return false;
    for (int i = 2; i <= x / i; i ++ )
        if (x % i == 0)
            return false;
    return true;
}
```

Miller-Rabin

- 前置：快速幂
- $\mathcal{O}(k \times \log^3 n)$

```
bool miller_rabin(LL n) {
    static vector<LL> tester = {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37};
    if (n < 3 || n % 2 == 0) return n == 2;
    if (n % 3 == 0) return n == 3;
    LL u = n - 1, t = 0;
    while (u % 2 == 0) u /= 2, ++t;
    for (auto nt: tester) {
        if(nt >= n) continue;
        LL v = bin(nt, u, n);
        if (v == 1) continue;
        LL s;
        for (s = 0; s < t; ++s) {
            if (v == n - 1) break;
            v = v * v % n;
        }
        if (s == t) return false;
    }
    return true;
}
```

质因数分解

朴素质因数分解

- 前置模板：素数筛
- 带指数
- $\mathcal{O}(\frac{\sqrt{N}}{\ln N})$

```
LL factor[30], f_sz, factor_exp[30];
void get_factor(LL x) {
    f_sz = 0;
    LL t = sqrt(x + 0.5);
    for (LL i = 0; pr[i] <= t; ++i)
        if (x % pr[i] == 0) {
            factor_exp[f_sz] = 0;
            while (x % pr[i] == 0) {
                x /= pr[i];
                ++factor_exp[f_sz];
            }
            factor[f_sz++] = pr[i];
        }
    if (x > 1) {
        factor_exp[f_sz] = 1;
        factor[f_sz++] = x;
    }
}
```

- 不带指数

```
LL factor[30], f_sz;
void get_factor(LL x) {
    f_sz = 0;
    LL t = sqrt(x + 0.5);
    for (LL i = 0; pr[i] <= t; ++i)
        if (x % pr[i] == 0) {
            factor[f_sz++] = pr[i];
            while (x % pr[i] == 0) x /= pr[i];
        }
```

```
    if (x > 1) factor[f_sz++] = x;
}
```

Pollard-Rho

- 前置：素数测试

```
mt19937 mt(time(0));
LL pollard_rho(LL n, LL c) {
    LL x = uniform_int_distribution<LL>(1, n - 1)(mt), y = x;
    auto f = [&](LL v) { LL t = mul(v, v, n) + c; return t < n ? t : t - n; };
    while (1) {
        x = f(x); y = f(f(y));
        if (x == y) return n;
        LL d = gcd(abs(x - y), n);
        if (d != 1) return d;
    }
}

LL fac[100], fcnt;
void get_fac(LL n, LL cc = 19260817) {
    if (n == 4) { fac[fcnt++] = 2; fac[fcnt++] = 2; return; }
    if (miller_rabin(n)) { fac[fcnt++] = n; return; }
    LL p = n;
    while (p == n) p = pollard_rho(n, --cc);
    get_fac(p); get_fac(n / p);
}

void go_fac(LL n) { fcnt = 0; if (n > 1) get_fac(n); }
```

原根

- 前置模板：质因数分解、快速幂
- 要求 p 为质数
- 别忘了调用质因数分解的函数

```
LL find_smallest_primitive_root(LL p) {
    get_factor(p - 1);
    FOR (i, 2, p) {
        bool flag = true;
        FOR (j, 0, f_sz)
        if (bin(i, (p - 1) / factor[j], p) == 1) {
            flag = false;
            break;
        }
        if (flag) return i;
    }
//  assert(0);
    return -1;
}
```

欧几里得

- 朴素

```
int gcd(int a, int b)
{
    return b ? gcd(b, a % b) : a;
}
```

- 卡常

```
inline int ctz(LL x) { return __builtin_ctzll(x); }
LL gcd(LL a, LL b) {
    if (!a) return b; if (!b) return a;
    int t = ctz(a | b);
    a >>= ctz(a);
    do {
        b >>= ctz(b);
        if (a > b) swap(a, b);
```

```
        b -= a;
    } while (b);
    return a << t;
}
```

扩展欧几里得

- 求 $ax + by = gcd(a, b)$ 的一组解
- 如果 a 和 b 互素，那么 x 是 a 在模 b 下的逆元
- 注意 x 和 y 可能是负数

```
LL ex_gcd(LL a, LL b, LL &x, LL &y) {
    if (b == 0) { x = 1; y = 0; return a; }
    LL ret = ex_gcd(b, a % b, y, x);
    y -= a / b * x;
    return ret;
}
```

二次剩余

- 求解二次同余方程
- 给定 a, p，求一组 x 满足 $x^2 \equiv a \pmod p$
- 前置模板：快速幂

```
LL q1, q2, w;
struct P { // x + y * sqrt(w)
    LL x, y;
};

P pmul(const P& a, const P& b, LL p) {
    P res;
    res.x = (a.x * b.x + a.y * b.y % p * w) % p;
    res.y = (a.x * b.y + a.y * b.x) % p;
    return res;
}

P bin(P x, LL n, LL MOD) {
    P ret = {1, 0};
    for (; n; n >>= 1, x = pmul(x, x, MOD))
        if (n & 1) ret = pmul(ret, x, MOD);
    return ret;
}
LL Legendre(LL a, LL p) { return bin(a, (p - 1) >> 1, p); }

LL equation_solve(LL b, LL p) {
    if (p == 2) return 1;
    if ((Legendre(b, p) + 1) % p == 0)
        return -1;
    LL a;
    while (true) {
        a = rand() % p;
        w = ((a * a - b) % p + p) % p;
        if ((Legendre(w, p) + 1) % p == 0)
            break;
    }
    return bin({a, 1}, (p + 1) >> 1, p).x;
}
// !!!!!!!!!!!!!!!!!!!!!!!!!!! Tester Start !!!!!!!!!!!!!!!!!!!!!!!!!!!!
void solve(){
    LL a, p; cin >> a >> p;
    a = a % p;
    LL x = equation_solve(a, p);
    if (x == -1) {
        puts("No root");
    } else {
        LL y = p - x;
        if (x == y){
            cout << x << endl;
        }else{
```

```
            LL tx = min(x, y), ty = max(x, y);
            cout << tx << " " << ty << endl;
        }
    }
}
// !!!!!!!!!!!!!!!!!!!!!!!!!!!! Tester End !!!!!!!!!!!!!!!!!!!!!!!!!!!!!

```

中国剩余定理

- 求解线性同余方程组
- $$\begin{cases} x & \equiv r_1 \pmod{m_1} \\ x & \equiv r_2 \pmod{m_2} \\ & \vdots \\ x & \equiv r_k \pmod{m_k} \end{cases}$$
- 无解返回 -1
- 前置模板：扩展欧几里得

```
LL CRT(LL *m, LL *r, LL n) {
    if (!n) return 0;
    LL M = m[0], R = r[0], x, y, d;
    FOR (i, 1, n) {
        d = ex_gcd(M, m[i], x, y);
        if ((r[i] - R) % d) return -1;
        x = (r[i] - R) / d * x % (m[i] / d);
        // 防爆 LL
        // x = mul((r[i] - R) / d, x, m[i] / d);
        R += x * M;
        M = M / d * m[i];
        R %= M;
    }
    return R >= 0 ? R : R + M;
}
```

逆元

- 如果 p 是素数，使用快速幂（费马小定理）
- 前置模板：快速幂

```
inline LL get_inv(LL x, LL p) { return bin(x, p - 2, p); }
```

- 如果 p 不是素数，使用拓展欧几里得
- 前置模板：扩展欧几里得

```
LL get_inv(LL a, LL M) {
    static LL x, y;
    assert(exgcd(a, M, x, y) == 1);
    return (x % M + M) % M;
}
```

- 预处理 1~n 的逆元

```
LL inv[N];
void inv_init(LL n, LL p) {
    inv[1] = 1;
    FOR (i, 2, n)
        inv[i] = (p - p / i) * inv[p % i] % p;
}
```

- 预处理阶乘及其逆元

```
LL invf[M], fac[M] = {1};
void fac_inv_init(LL n, LL p) {
    FOR (i, 1, n)
        fac[i] = i * fac[i - 1] % p;
```

```
    invf[n - 1] = bin(fac[n - 1], p - 2, p);
    FORD (i, n - 2, -1)
        invf[i] = invf[i + 1] * (i + 1) % p;
}
```

组合数

组合数预处理（递推法）

```
LL C[M][M];
void init_C(int n) {
    FOR (i, 0, n) {
        C[i][0] = C[i][i] = 1;
        FOR (j, 1, i)
            C[i][j] = (C[i - 1][j] + C[i - 1][j - 1]) % MOD;
    }
}
```

预处理逆元法

- 如果数较小，模较大时使用逆元
- 前置模板：逆元-预处理阶乘及其逆元

```
inline LL C(LL n, LL m) { // n >= m >= 0
    return n < m || m < 0 ? 0 : fac[n] * invf[m] % MOD * invf[n - m] % MOD;
}
```

Lucas 定理

- 如果模数较小，数字较大，使用 Lucas 定理
- 前置模板可选 1：求组合数（如果使用阶乘逆元，需 `fac_inv_init(MOD, MOD);`）

```
LL C(LL n, LL m) { // m >= n >= 0
    if (m - n < n) n = m - n;
    if (n < 0) return 0;
    LL ret = 1;
    FOR (i, 1, n + 1)
        ret = ret * (m - n + i) % MOD * bin(i, MOD - 2, MOD) % MOD;
    return ret;
}
```

- 前置模板可选 2：模数不固定下使用，无法单独使用。

```
LL Lucas(LL n, LL m) { // m >= n >= 0
    return m ? C(n % MOD, m % MOD) * Lucas(n / MOD, m / MOD) % MOD : 1;
}
```

求具体值

- 分解质因数法

```
int primes[N], cnt;     // 存储所有质数
int sum[N];     // 存储每个质数的次数
bool st[N];     // 存储每个数是否已被筛掉

void get_primes(int n)      // 线性筛法求素数
{
    for (int i = 2; i <= n; i ++ )
    {
        if (!st[i]) primes[cnt ++ ] = i;
        for (int j = 0; primes[j] <= n / i; j ++ )
        {
            st[primes[j] * i] = true;
            if (i % primes[j] == 0) break;
        }
    }
}


int get(int n, int p)       // 求 n! 中的次数
```

```
{
    int res = 0;
    while (n)
    {
        res += n / p;
        n /= p;
    }
    return res;
}


vector<int> mul(vector<int> a, int b)       // 高精度乘低精度模板
{
    vector<int> c;
    int t = 0;
    for (int i = 0; i < a.size(); i ++ )
    {
        t += a[i] * b;
        c.push_back(t % 10);
        t /= 10;
    }

    while (t)
    {
        c.push_back(t % 10);
        t /= 10;
    }

    return c;
}

get_primes(a);  // 预处理范围内的所有质数

for (int i = 0; i < cnt; i ++ )     // 求每个质因数的次数
{
    int p = primes[i];
    sum[i] = get(a, p) - get(b, p) - get(a - b, p);
}

vector<int> res;
res.push_back(1);

for (int i = 0; i < cnt; i ++ )     // 用高精度乘法将所有质因子相乘
    for (int j = 0; j < sum[i]; j ++ )
        res = mul(res, primes[i]);
```

FFT & NTT & FWT

FFT

- 计算多项式乘法，可用于高精度乘法
- $\mathcal{O}(n \log n)$

```
typedef double LD;
const LD PI = acos(-1.0);

struct Complex {
    LD r, i;
    Complex(LD r = 0, LD i = 0) : r(r), i(i) {}
    Complex operator + (const Complex& other) const {
        return Complex(r + other.r, i + other.i);
    }
    Complex operator - (const Complex& other) const {
        return Complex(r - other.r, i - other.i);
    }
    Complex operator * (const Complex& other) const {
        return Complex(r * other.r - i * other.i, r * other.i + i * other.r);
    }
};

// 快速傅里叶变换，p=1 为正向，p=-1 为反向
```

```cpp
void FFT(vector<Complex>& x, int p) {
    int n = x.size();
    for (int i = 0, t = 0; i < n; ++i) {
        if (i > t) swap(x[i], x[t]);
        for (int j = n >> 1; (t ^= j) < j; j >>= 1);
    }
    for (int h = 2; h <= n; h <<= 1) {
        Complex wn(cos(p * 2 * PI / h), sin(p * 2 * PI / h));
        for (int i = 0; i < n; i += h) {
            Complex w(1, 0);
            for (int j = 0; j < h / 2; ++j) {
                Complex u = x[i + j];
                Complex v = x[i + j + h/2] * w;
                x[i + j] = u + v;
                x[i + j + h/2] = u - v;
                w = w * wn;
            }
        }
    }
    if (p == -1) {
        for (int i = 0; i < n; ++i) {
            x[i].r /= n;
        }
    }
}

// 计算两个多项式的卷积，返回结果多项式的系数向量
vector<LD> convolution(const vector<LD>& a, const vector<LD>& b) {
    int len = 1;
    int n = a.size(), m = b.size();
    while (len < n + m - 1) len <<= 1;
    vector<Complex> fa(len), fb(len);
    for (int i = 0; i < n; ++i) fa[i] = Complex(a[i], 0);
    for (int i = 0; i < m; ++i) fb[i] = Complex(b[i], 0);
    FFT(fa, 1);
    FFT(fb, 1);
    for (int i = 0; i < len; ++i) {
        fa[i] = fa[i] * fb[i];
    }
    FFT(fa, -1);
    vector<LD> res(n + m - 1);
    for (int i = 0; i < n + m - 1; ++i) {
        res[i] = fa[i].r;
    }
    return res;
}
```

NTT

- 用于大整数乘法时，位数不宜过高（在 MOD=998244353 的情况下，总位数不超过 $12324004(3510^2)$）
- 前置模板：快速幂、逆元

```cpp
const int N = 1e5+10;
const int MOD = 998244353; // 模数
const int G = 3; // 原根

LL wn[N << 2], rev[N << 2];
int NTT_init(int n_) {
    int step = 0; int n = 1;
    for ( ; n < n_; n <<= 1) ++step;
    FOR (i, 1, n)
    rev[i] = (rev[i >> 1] >> 1) | ((i & 1) << (step - 1));
    int g = bin(G, (MOD - 1) / n, MOD);
    wn[0] = 1;
    for (int i = 1; i <= n; ++i)
        wn[i] = wn[i - 1] * g % MOD;
    return n;
}

void NTT(vector<LL> &a, int n, int f) {
    FOR (i, 0, n) if (i < rev[i])
```

```
        std::swap(a[i], a[rev[i]]);
    for (int k = 1; k < n; k <<= 1) {
        for (int i = 0; i < n; i += (k << 1)) {
            int t = n / (k << 1);
            FOR (j, 0, k) {
                LL w = f == 1 ? wn[t * j] : wn[n - t * j];
                LL x = a[i + j];
                LL y = a[i + j + k] * w % MOD;
                a[i + j] = (x + y) % MOD;
                a[i + j + k] = (x - y + MOD) % MOD;
            }
        }
    }
    if (f == -1) {
        LL ninv = get_inv(n, MOD);
        FOR (i, 0, n)
        a[i] = a[i] * ninv % MOD;
    }
}

vector<LL> conv(vector<LL> a, vector<LL> b){
    int len_a = a.size(), len_b = b.size();
    int len = len_a + len_b - 1;
    int n = NTT_init(len);
    a.resize(n);
    b.resize(n);
    NTT(a, n, 1);
    NTT(b, n, 1);
    vector<LL> c(n);
    for (int i = 0; i < n; ++i) {
        c[i] = a[i] * b[i] % MOD;
    }
    NTT(c, n, -1);
    vector<LL> res(len);
    for (int i = 0; i < len; ++i) {
        res[i] = c[i];
    }
    return res;
}
```

FWT

```
const LL MOD = 998244353;

template<typename T>
void fwt(vector<LL> &a, int n, T f) {
    for (int d = 1; d < n; d *= 2)
        for (int i = 0, t = d * 2; i < n; i += t)
            FOR (j, 0, d)
            f(a[i + j], a[i + j + d]);
}

void AND(LL& a, LL& b) { a += b; }
void OR(LL& a, LL& b) { b += a; }
void XOR (LL& a, LL& b) {
    LL x = a, y = b;
    a = (x + y) % MOD;
    b = (x - y + MOD) % MOD;
}
void rAND(LL& a, LL& b) { a -= b; }
void rOR(LL& a, LL& b) { b -= a; }
void rXOR(LL& a, LL& b) {
    static LL INV2 = (MOD + 1) / 2;
    LL x = a, y = b;
    a = (x + y) * INV2 % MOD;
    b = (x - y + MOD) * INV2 % MOD;
}

int next_power_of_two(int n) {
    if (n <= 0) return 1;
    // __lg(n-1) 返回 n-1 的最高位所在位置（0-based）
```

```
    return 1 << (__lg(n - 1) + 1);
}

template<typename T, typename F>
vector<LL> conv(vector<LL> a, vector<LL> b, T f, F inv_f){
    LL len_a = a.size(), len_b = b.size(), len = max(len_a, len_b), n = next_power_of_two(len);
    a.resize(n), b.resize(n);
    fwt(a, n, f), fwt(b, n, f);
    vector<LL> c(n);
    for (int i = 0; i < n; i++) {
        c[i] = a[i] * b[i] % MOD;
    }
    fwt(c, n, inv_f);
    // 提取结果 (可选)
    c.resize(len);
    return c;
}
```

线性基

贪心法

可查询最大异或和

```
struct BasisGreedy{
    ULL p[64];
    BasisGreedy(){memset(p, 0, sizeof p);}
    void insert(ULL x) {
        for (int i = 63; ~i; --i) {
            if (!(x >> i))  // x 的第 i 位是 0
                continue;
            if (!p[i]) {
                p[i] = x;
                break;
            }
            x ^= p[i];
        }
    }
    ULL query_max(){
        ULL ans = 0;
        for (int i = 63; ~i; --i) {
            ans = std::max(ans, ans ^ p[i]);
        }
        return ans;
    }
};
```

高斯消元法

可查询任意大异或和

```
struct BasisGauss{
    vector<ULL> a;
    LL n, tmp, cnt;

    BasisGauss(){a = {0};}

    void insert(ULL x){
        a.push_back(x);
    }

    void init(){
        n = (LL)a.size() - 1;
        LL k=1;
        for(int i=63;i>=0;i--){
            int t=0;
            for(LL j=k;j<=n;j++){
                if((a[j]>>i)&1){
                    t=j;
                    break;
```

```
                }
            }
            if(t){
                swap(a[k],a[t]);
                for(LL j=1;j<=n;j++){
                    if(j!=k&&(a[j]>>i)&1) a[j]^=a[k];
                }
                k++;
            }
        }
        cnt = k-1;
        tmp = 1LL << cnt;
        if(cnt==n) tmp--;
    }

    LL query_xth(LL x){ // 从小到大，若 x 为负数，则查询倒数第几个
        if(x<0) x = tmp + x + 1;
        if(x>tmp) return -1;
        else{
            if(n>cnt) x--;
            LL ans=0;
            for(LL i=0; i<cnt; i++){
                if((x>>i)&1) ans^=a[cnt-i];
            }
            return ans;
        }
    }
};
```

性质与公式

低阶等幂求和

- $\sum_{i=1}^{n} i^1 = \frac{n(n+1)}{2} = \frac{1}{2}n^2 + \frac{1}{2}n$
- $\sum_{i=1}^{n} i^2 = \frac{n(n+1)(2n+1)}{6} = \frac{1}{3}n^3 + \frac{1}{2}n^2 + \frac{1}{6}n$
- $\sum_{i=1}^{n} i^3 = \left[\frac{n(n+1)}{2}\right]^2 = \frac{1}{4}n^4 + \frac{1}{2}n^3 + \frac{1}{4}n^2$
- $\sum_{i=1}^{n} i^4 = \frac{n(n+1)(2n+1)(3n^2+3n-1)}{30} = \frac{1}{5}n^5 + \frac{1}{2}n^4 + \frac{1}{3}n^3 - \frac{1}{30}n$
- $\sum_{i=1}^{n} i^5 = \frac{n^2(n+1)^2(2n^2+2n-1)}{12} = \frac{1}{6}n^6 + \frac{1}{2}n^5 + \frac{5}{12}n^4 - \frac{1}{12}n^2$

一些组合公式

- 错排公式（对于 $1 \sim n$ 的排列 P，满足 $P_i \neq i$）：$D_1 = 0, D_2 = 1, D_n = (n-1)(D_{n-1} + D_{n-2}) = n!(\frac{1}{2!} - \frac{1}{3!} + \cdots + (-1)^n \frac{1}{n!}) = \lfloor \frac{n!}{e} + 0.5 \rfloor$
- 卡塔兰数（n 对括号合法方案数，n 个结点二叉树个数，$n \times n$ 方格中对角线下方的单调路径数，凸 $n+2$ 边形的三角形划分数，n 个元素的合法出栈序列数）：$C_n = \frac{1}{n+1}\binom{2n}{n} = \frac{(2n)!}{(n+1)!n!}$

互质

若整数 a 与 m 互质（即 $\gcd(a, m) = 1$）

- 对于整数 $k = 0, 1, 2, \ldots, m-1$，$ak \mod m$ 的结果恰好是 $0, 1, 2, \ldots, m-1$ 的一个**排列**（每个数出现且仅出现一次）。
- 存在唯一的整数 b（$1 \leq b < m$），使得 $ab \equiv 1 \mod m$，此时 b 称为 a 在模 m 下的**乘法逆元**（记为 $a^{-1} \mod m$）。

图论

最短路

朴素 djikstra 算法

- 无负权边、稠密图

```
int g[N][N];  // 存储每条边
int dist[N];  // 存储 1 号点到每个点的最短距离
```

```
bool st[N];   // 存储每个点的最短路是否已经确定

// 求 1 号点到 n 号点的最短路，如果不存在则返回-1
int dijkstra(){
    memset(dist, 0x3f, sizeof dist);
    dist[1] = 0;
    for (int i = 0; i < n - 1; i ++ ){
        int t = -1;     // 在还未确定最短路的点中，寻找距离最小的点
        for (int j = 1; j <= n; j ++ )
            if (!st[j] && (t == -1 || dist[t] > dist[j]))
                t = j;
        // 用 t 更新其他点的距离
        for (int j = 1; j <= n; j ++ )
            dist[j] = min(dist[j], dist[t] + g[t][j]);
        st[t] = true;
    }
    if (dist[n] == 0x3f3f3f3f) return -1;
    return dist[n];
}
```

堆优化的 djikstra

- 无负权边、稀疏图

```
typedef pair<int, int> PII;

int n;      // 点的数量
int h[N], w[N], e[N], ne[N], idx;       // 邻接表存储所有边
int dist[N];        // 存储所有点到 1 号点的距离
bool st[N];     // 存储每个点的最短距离是否已确定

// 求 1 号点到 n 号点的最短距离，如果不存在，则返回-1
int dijkstra(){
    memset(dist, 0x3f, sizeof dist);
    dist[1] = 0;
    priority_queue<PII, vector<PII>, greater<PII>> heap;
    heap.push({0, 1});      // first 存储距离，second 存储节点编号
    while (heap.size()){
        auto t = heap.top();
        heap.pop();
        int ver = t.second, distance = t.first;
        if (st[ver]) continue;
        st[ver] = true;
        for (int i = h[ver]; i != -1; i = ne[i]){
            int j = e[i];
            if (dist[j] > distance + w[i]){
                dist[j] = distance + w[i];
                heap.push({dist[j], j});
            }
        }
    }
    if (dist[n] == 0x3f3f3f3f) return -1;
    return dist[n];
}
```

Bellman-Ford 算法

- 有负权边、可以处理负环

```
int n, m;       // n 表示点数，m 表示边数
int dist[N];        // dist[x] 存储 1 到 x 的最短路距离

struct Edge{     // 边，a 表示出点，b 表示入点，w 表示边的权重
    int a, b, w;
}edges[M];

// 求 1 到 n 的最短路距离，如果无法从 1 走到 n，则返回-1。
int bellman_ford(){
    memset(dist, 0x3f, sizeof dist);
    dist[1] = 0;

```

```
    // 如果第 n 次迭代仍然会松弛三角不等式，就说明存在一条长度是 n+1 的最短路径，由抽屉原理，路径中至少存在两个相同的点，说明图中存在负权回路。
    for (int i = 0; i < n; i ++ ){
        for (int j = 0; j < m; j ++ ){
            int a = edges[j].a, b = edges[j].b, w = edges[j].w;
            if (dist[b] > dist[a] + w)
                dist[b] = dist[a] + w;
        }
    }

    if (dist[n] > 0x3f3f3f3f / 2) return -1;
    return dist[n];
}
```

spfa 算法

- 有负权边、不能有负环，快

```
int n;      // 总点数
int h[N], w[N], e[N], ne[N], idx;       // 邻接表存储所有边
int dist[N];        // 存储每个点到 1 号点的最短距离
bool st[N];     // 存储每个点是否在队列中

// 求 1 号点到 n 号点的最短路距离，如果从 1 号点无法走到 n 号点则返回-1
int spfa(){
    memset(dist, 0x3f, sizeof dist);
    dist[1] = 0;
    queue<int> q;
    q.push(1);
    st[1] = true;
    while (q.size()){
        auto t = q.front();
        q.pop();
        st[t] = false;
        for (int i = h[t]; i != -1; i = ne[i]){
            int j = e[i];
            if (dist[j] > dist[t] + w[i]){
                dist[j] = dist[t] + w[i];
                if (!st[j]){     // 如果队列中已存在 j，则不需要将 j 重复插入
                    q.push(j);
                    st[j] = true;
                }
            }
        }
    }
    if (dist[n] == 0x3f3f3f3f) return -1;
    return dist[n];
}
```

spfa 判断负环

```
int n;      // 总点数
int h[N], w[N], e[N], ne[N], idx;       // 邻接表存储所有边
int dist[N], cnt[N];        // dist[x] 存储 1 号点到 x 的最短距离，cnt[x] 存储 1 到 x 的最短路中经过的点数
bool st[N];     // 存储每个点是否在队列中

// 如果存在负环，则返回 true，否则返回 false。
bool spfa(){
    // 不需要初始化 dist 数组
    // 原理：如果某条最短路径上有 n 个点（除了自己），那么加上自己之后一共有 n+1 个点，由抽屉原理一定有两个点相同，所以存在环。
    queue<int> q;
    for (int i = 1; i <= n; i ++ ){
        q.push(i);
        st[i] = true;
    }
    while (q.size()){
        auto t = q.front();
        q.pop();
        st[t] = false;
        for (int i = h[t]; i != -1; i = ne[i]){
            int j = e[i];
            if (dist[j] > dist[t] + w[i]){
```

```
                dist[j] = dist[t] + w[i];
                cnt[j] = cnt[t] + 1;
                if (cnt[j] >= n) return true;       // 如果从 1 号点到 x 的最短路中包含至少 n 个点（不包括自己），则说明存在环
                if (!st[j]){
                    q.push(j);
                    st[j] = true;
                }
            }
        }
    }
    return false;
}
```

floyd 算法

```
初始化:
    for (int i = 1; i <= n; i ++ )
        for (int j = 1; j <= n; j ++ )
            if (i == j) d[i][j] = 0;
            else d[i][j] = INF;

// 算法结束后，d[a][b] 表示 a 到 b 的最短距离
void floyd(){
    for (int k = 1; k <= n; k ++ )
        for (int i = 1; i <= n; i ++ )
            for (int j = 1; j <= n; j ++ )
                d[i][j] = min(d[i][j], d[i][k] + d[k][j]);
}
```

最小生成树

朴素 Prim 算法

- 稠密图（m 接近于 n^2 ）

```
int n;      // n 表示点数
int g[N][N];        // 邻接矩阵，存储所有边
int dist[N];        // 存储其他点到当前最小生成树的距离
bool st[N];     // 存储每个点是否已经在生成树中
// 如果图不连通，则返回 INF(值是 0x3f3f3f3f), 否则返回最小生成树的树边权重之和
int prim(){
    memset(dist, 0x3f, sizeof dist);
    int res = 0;
    for (int i = 0; i < n; i ++ ){
        int t = -1;
        for (int j = 1; j <= n; j ++ )
            if (!st[j] && (t == -1 || dist[t] > dist[j]))
                t = j;
        if (i && dist[t] == INF) return INF;
        if (i) res += dist[t];
        st[t] = true;
        for (int j = 1; j <= n; j ++ ) dist[j] = min(dist[j], g[t][j]);
    }
    return res;
}
```

Kruskal 算法

- 实现简单，稀疏图（m 接近 n ）

```
int n, m;       // n 是点数，m 是边数
int p[N];       // 并查集的父节点数组
struct Edge{    // 存储边
    int a, b, w;
    bool operator< (const Edge &W)const{
        return w < W.w;
    }
}edges[M];

int find(int x){    // 并查集核心操作
```

```cpp
    if (p[x] != x) p[x] = find(p[x]);
    return p[x];
}

int kruskal(){
    sort(edges, edges + m);
    for (int i = 1; i <= n; i ++ ) p[i] = i;    // 初始化并查集
    int res = 0, cnt = 0;
    for (int i = 0; i < m; i ++ ){
        int a = edges[i].a, b = edges[i].b, w = edges[i].w;
        a = find(a), b = find(b);
        if (a != b){    // 如果两个连通块不连通，则将这两个连通块合并
            p[a] = b;
            res += w;
            cnt ++ ;
        }
    }
    if (cnt < n - 1) return INF;
    return res;
}
```

拓扑排序

- 有向图
- 别忘了存储入度
- 当 `toporder(int n)` 返回值的长度不等于 n 时，不存在拓扑排序。

```cpp
const int N = 1e5+10;
vector<int> G[N];
int deg[N]; // 入度

vector<int> toporder(int n) {
    vector<int> orders;
    queue<int> q;
    for (int i = 1; i <= n; i++)
        if (!deg[i]) {
            q.push(i);
            orders.push_back(i);
        }
    while (!q.empty()) {
        int u = q.front(); q.pop();
        for (int v: G[u])
            if (!--deg[v]) {
                q.push(v);
                orders.push_back(v);
            }
    }
    return orders;
}
```

差分约束

一个系统 n 个变量和 m 个约束条件组成，每个约束条件形如 $x_j - x_i \le b_k$。可以发现每个约束条件都形如最短路中的三角不等式 $d_u - d_v \le w_{u,v}$。因此连一条边 (i, j, b_k) 建图。

若要判断解的存在性，使用 spfa 判断是否存在负环，有则无解。

若要使得所有量两两的值最接近，源点到各点的距离初始成 0，跑最远路。

若要使得某一变量与其他变量的差尽可能大，则源点到各点距离初始化成 ∞，跑最短路。

最近公共祖先

```cpp
const LL N = 5e5+10, SP = log2(N)+1;
vector<int> G[N];
int pa[N][SP], dep[N];

void dfs(int u, int fa) {
    pa[u][0] = fa; dep[u] = dep[fa] + 1;
```

```
    FOR (i, 1, SP) pa[u][i] = pa[pa[u][i - 1]][i - 1];
    for (int& v: G[u]) {
        if (v == fa) continue;
        dfs(v, u);
    }
}

int lca(int u, int v) {
    if (dep[u] < dep[v]) swap(u, v);
    int t = dep[u] - dep[v];
    FOR (i, 0, SP) if (t & (1 << i)) u = pa[u][i];
    FORD (i, SP - 1, -1) {
        int uu = pa[u][i], vv = pa[v][i];
        if (uu != vv) { u = uu; v = vv; }
    }
    return u == v ? u : pa[u][0];
}
```

树链剖分

- 将树上操作转化为区间操作，套用区间数据结构
- 别忘了选一种方法（取消注释）
- `fa[N]`：存储每个节点的父节点
- `dep[N]`：存储每个节点的深度
- `idx[N]`：存储每个节点在线段树中的索引（DFS 序）
- `out[N]`：存储每个节点子树在 DFS 序中的结束位置
- `ridx[N]`：存储 DFS 序到节点的反向映射
- `sz[N]`：存储每个节点的子树大小
- `son[N]`：存储每个节点的重儿子（子树最大的儿子）
- `top[N]`：存储每个节点所在重链的顶端节点
- `clk`：DFS 序计数器
- `init()`：初始化（先建图再调用）
- `go(u, v, f)`：`f` 是一个形如 `f(int l, int r)` 的函数。对树上节点 `u` 到节点 `v` 的简单路径，分解为 dfs 序中的区间 $[l, r]$，调用函数 `f`
- 子树操作：`u` 的子树的 dfs 序区间为 $[idx[u], out[u]]$

```
const int N = 3e4+10;

vector<int> G[N];
int fa[N], dep[N], idx[N], out[N], ridx[N];
namespace hld {
    int sz[N], son[N], top[N], clk;
    void predfs(int u, int d) {
        dep[u] = d; sz[u] = 1;
        int& maxs = son[u] = -1;
        for (int& v: G[u]) {
            if (v == fa[u]) continue;
            fa[v] = u;
            predfs(v, d + 1);
            sz[u] += sz[v];
            if (maxs == -1 || sz[v] > sz[maxs]) maxs = v;
        }
    }
    void dfs(int u, int tp) {
        top[u] = tp; idx[u] = ++clk; ridx[clk] = u;
        if (son[u] != -1) dfs(son[u], tp);
        for (int& v: G[u])
            if (v != fa[u] && v != son[u]) dfs(v, v);
        out[u] = clk;
    }
    void init(){
        clk = 0;
        predfs(1, 1);
        dfs(1, 1);
    }
    template<typename T>
    int go(int u, int v, T&& f = [](int, int) {}) {
```

```cpp
        int uu = top[u], vv = top[v];
        while (uu != vv) {
            if (dep[uu] < dep[vv]) { swap(uu, vv); swap(u, v); }
            f(idx[uu], idx[u]);
            u = fa[uu]; uu = top[u];
        }
        if (dep[u] < dep[v]) swap(u, v);
        // 下面两行代码选择一个
        // f(idx[v], idx[u]); // 包含 lca(u, v)
        // if (u != v) f(idx[v] + 1, idx[u]); // 不包含 lca(u, v)
        return v;
    }
    int up(int u, int d) { // 查询 u 节点向上走 d 步的节点编号
        while (d) {
            if (dep[u] - dep[top[u]] < d) {
                d -= dep[u] - dep[top[u]];
                u = top[u];
            } else return ridx[idx[u] - d];
            u = fa[u]; --d;
        }
        return u;
    }
    int finds(int u, int rt) { // 找 u 在 rt 的哪个儿子的子树中
        while (top[u] != top[rt]) {
            u = top[u];
            if (fa[u] == rt) return u;
            u = fa[u];
        }
        return ridx[idx[rt] + 1];
    }
}
```

网络流

- 最大流

```cpp
const LL INF = LONG_LONG_MAX;

struct E {
    LL to, cp;
    E(LL to, LL cp): to(to), cp(cp) {}
};

struct Dinic {
    static const LL M = 1E5 * 5;
    LL m, s, t;
    vector<E> edges;
    vector<LL> G[M];
    LL d[M];
    LL cur[M];

    void init(LL n, LL s, LL t) {
        this->s = s; this->t = t;
        for (LL i = 0; i <= n; i++) G[i].clear();
        edges.clear(); m = 0;
    }

    void addedge(LL u, LL v, LL cap) {
        edges.emplace_back(v, cap);
        edges.emplace_back(u, 0);
        G[u].push_back(m++);
        G[v].push_back(m++);
    }

    bool BFS() {
        memset(d, 0, sizeof d);
        queue<LL> Q;
        Q.push(s); d[s] = 1;
        while (!Q.empty()) {
            LL x = Q.front(); Q.pop();
            for (LL& i: G[x]) {
```

```
                E &e = edges[i];
                if (!d[e.to] && e.cp > 0) {
                    d[e.to] = d[x] + 1;
                    Q.push(e.to);
                }
            }
        }
        return d[t];
    }

    LL DFS(LL u, LL cp) {
        if (u == t || !cp) return cp;
        LL tmp = cp, f;
        for (LL& i = cur[u]; i < G[u].size(); i++) {
            E& e = edges[G[u][i]];
            if (d[u] + 1 == d[e.to]) {
                f = DFS(e.to, min(cp, e.cp));
                e.cp -= f;
                edges[G[u][i] ^ 1].cp += f;
                cp -= f;
                if (!cp) break;
            }
        }
        return tmp - cp;
    }

    LL go() {
        LL flow = 0;
        while (BFS()) {
            memset(cur, 0, sizeof cur);
            flow += DFS(s, INF);
        }
        return flow;
    }
} DC;
```

- 最小费用最大流

```
const LL M = 5e4+10;
const int INF = INT_MAX;

struct E {
    int from, to, cp, v;
    E() {}
    E(int f, int t, int cp, int v) : from(f), to(t), cp(cp), v(v) {}
};

struct MCMF {
    int n, m, s, t;
    vector<E> edges;
    vector<int> G[M];
    bool inq[M];
    int d[M], p[M], a[M];

    void init(int _n, int _s, int _t) {
        n = _n; s = _s; t = _t;
        FOR (i, 0, n + 1) G[i].clear();
        edges.clear(); m = 0;
    }

    void addedge(int from, int to, int cap, int cost) {
        edges.emplace_back(from, to, cap, cost);
        edges.emplace_back(to, from, 0, -cost);
        G[from].push_back(m++);
        G[to].push_back(m++);
    }

    bool BellmanFord(int &flow, int &cost) {
        FOR (i, 0, n + 1) d[i] = INF;
        memset(inq, 0, sizeof inq);
        d[s] = 0, a[s] = INF, inq[s] = true;
        queue<int> Q; Q.push(s);
```

```cpp
        while (!Q.empty()) {
            int u = Q.front(); Q.pop();
            inq[u] = false;
            for (int& idx: G[u]) {
                E &e = edges[idx];
                if (e.cp && d[e.to] > d[u] + e.v) {
                    d[e.to] = d[u] + e.v;
                    p[e.to] = idx;
                    a[e.to] = min(a[u], e.cp);
                    if (!inq[e.to]) {
                        Q.push(e.to);
                        inq[e.to] = true;
                    }
                }
            }
        }
        if (d[t] == INF) return false;
        flow += a[t];
        cost += a[t] * d[t];
        int u = t;
        while (u != s) {
            edges[p[u]].cp -= a[t];
            edges[p[u] ^ 1].cp += a[t];
            u = edges[p[u]].from;
        }
        return true;
    }

    pair<int, int> go() {
        int flow = 0, cost = 0;
        while (BellmanFord(flow, cost));
        return {flow, cost};
    }
} MM;
```

树上路径交

- 前置模板：最近公共祖先

```cpp
int intersection(int x1, int y1, int x2, int y2) {
    int t[4] = {lca(x1, x2), lca(x1, y2), lca(y1, x2), lca(y1, y2)};
    int p1 = 0,p2 = 0;
    FOR(j,0,4)
    if(dep[t[j]] > dep[p1]) p2 = p1, p1 = t[j];
    else if(dep[t[j]] > dep[p2]) p2 = t[j];
    int h1 = lca(x1,y1), h2 = lca(x2,y2);
    if(p1 == p2){
        if(dep[p1] < dep[h1] || dep[p1] < dep[h2]) return 0;
        else return 1;
    }
    else{
        int ans = dep[p1]+dep[p2]-2*dep[lca(p1,p2)]+1;
        return ans;
    }
}
```

树上点分治（树的重心）

```cpp
const LL N = 2e4+10, N2 = N * 2;

int h[N], e[N2], ne[N2], idx;

void add(int a, int b){
    e[idx] = b, ne[idx] = h[a], h[a] = idx++;
}

vector<bool> vis;

// 获取子树的重心（自动处理父子关系）(如果有两个重心，输出编号小的那个)
// 若重心为 u，则 mx[u] 为以 u 为重心子树大小的最大值
```

```cpp
int q[N], fa[N], sz[N], mx[N];
int get_rt(int u) {
    int p = 0, cur = -1;
    q[p++] = u; fa[u] = -1;
    while (++cur < p) {
        u = q[cur]; mx[u] = 0; sz[u] = 1;
        for (int i = h[u]; i!=-1; i=ne[i]){
            int j = e[i];
            if(vis[j] or j == fa[u]) continue;
            fa[q[p++] = j] = u;
        }
    }
    FORD (i, p - 1, -1) {
        u = q[i];
        mx[u] = max(mx[u], p - sz[u]);
        if (mx[u] * 2 <= p) return u;
        sz[fa[u]] += sz[u];
        mx[fa[u]] = max(mx[fa[u]], sz[u]);
    }
//  assert(0);
}

// 分治 dfs（起点任意）
void dfs(int u) {
    cout << "u: " << u;
    u = get_rt(u);
    vis[u] = true;
    // 处理子树逻辑
    cout << " centroid: " << u << '\n';
    // 如果在此处 DFS，会遍历整棵子树（if(vis[u]) return）
    // ...

    for(int i=h[u]; i!=-1; i=ne[i]){
        int j = e[i];
        if(vis[j]) continue;
        dfs(j);
    }
}
```

二分图

最大匹配

- 最小覆盖数 = 最大匹配数
- 最大独立集 = 顶点数 - 二分图匹配数
- DAG 最小路径覆盖数 = 结点数 - 拆点后二分图最大匹配数

```cpp
const int N = 500+10;

struct MaxMatch {
    int n;
    vector<int> G[N];
    int vis[N], left[N], clk;

    void init(int n) {
        this->n = n;
        FOR (i, 0, n + 1) G[i].clear();
        memset(left, -1, sizeof left);
        memset(vis, -1, sizeof vis);
    }

    bool dfs(int u) {
        for (int v: G[u])
            if (vis[v] != clk) {
                vis[v] = clk;
                if (left[v] == -1 || dfs(left[v])) {
                    left[v] = u;
                    return true;
                }
            }
```

```cpp
            return false;
        }

        int match() {
            int ret = 0;
            for (clk = 0; clk <= n; ++clk)
                if (dfs(clk)) ++ret;
            return ret;
        }
} MM;
// !!!!!!!!!!!!!!!!!!!!!!!!!!! Tester Start !!!!!!!!!!!!!!!!!!!!!!!!!!!!
void solve(){
    LL n1, n2, m, n, i, t1, t2;
    cin >> n1 >> n2 >> m;
    n = n1 + n2;
    MM.init(n);
    for(i=0; i<m; i++){
        cin >> t1 >> t2;
        MM.G[t1].push_back(n1+t2);
    }
    cout << MM.match() << '\n';
}
```

最大权匹配

- `py[j] = i` 表示右侧顶点 j 与左侧顶点 i 匹配

```cpp
namespace R {
    const int M = 400 + 5;
    const int INF = 2E9;
    int n;
    int w[M][M], kx[M], ky[M], py[M], vy[M], slk[M], pre[M];

    LL KM() {
        FOR (i, 1, n + 1)
        FOR (j, 1, n + 1)
        kx[i] = max(kx[i], w[i][j]);
        FOR (i, 1, n + 1) {
            fill(vy, vy + n + 1, 0);
            fill(slk, slk + n + 1, INF);
            fill(pre, pre + n + 1, 0);
            int k = 0, p = -1;
            for (py[k = 0] = i; py[k]; k = p) {
                int d = INF;
                vy[k] = 1;
                int x = py[k];
                FOR (j, 1, n + 1)
                if (!vy[j]) {
                    int t = kx[x] + ky[j] - w[x][j];
                    if (t < slk[j]) { slk[j] = t; pre[j] = k; }
                    if (slk[j] < d) { d = slk[j]; p = j; }
                }
                FOR (j, 0, n + 1)
                if (vy[j]) { kx[py[j]] -= d; ky[j] += d; }
                else slk[j] -= d;
            }
            for (; k; k = pre[k]) py[k] = py[pre[k]];
        }
        LL ans = 0;
        FOR (i, 1, n + 1) ans += kx[i] + ky[i];
        return ans;
    }
}
// !!!!!!!!!!!!!!!!!!!!!!!!!!! Tester Start !!!!!!!!!!!!!!!!!!!!!!!!!!!!
void solve(){
    LL n1, n2, i, t1, t2, t3, m, n, j;
    cin >> n1 >> n2 >> m;
    // 初始化
    n = max(n1, n2);
    R::n = n;
    for(i=0; i<=n; i++){
```

```
        for(j=0; j<=n; j++){
            R::w[i][j] = 0;
        }
    }
    // 读数据
    for(i=0; i<m; i++){
        cin >> t1 >> t2 >> t3;
        R::w[t1][t2] = t3;
    }
    // 计算
    LL maxx = R::KM();
    cout << maxx << '\n';
    // 结果转换
    vector<pair<LL, LL>> anss;
    for(i=1; i<=n; i++){ // 注意遍历最大范围
        if(R::w[R::py[i]][i]){
            anss.push_back({R::py[i], i});
        }else{
            // 未匹配
            anss.push_back({R::py[i], 0});
        }
    }
    sort(anss.begin(), anss.end());
    for(i=0; i<n1; i++){
        cout << anss[i].second << ' ';
    }
}
```

Tarjan

割点

- 判断割点（无向图）
- 注意原图可能不连通

```
int dfn[N], low[N], clk;
void init() { clk = 0; memset(dfn, 0, sizeof dfn); }
void tarjan(int u, int fa) {
    low[u] = dfn[u] = ++clk;
    int cc = fa != -1;
    for (int& v: G[u]) {
        if (v == fa) continue;
        if (!dfn[v]) {
            tarjan(v, u);
            low[u] = min(low[u], low[v]);
            cc += low[v] >= dfn[u];
        } else low[u] = min(low[u], dfn[v]);
    }
    if (cc > 1) // u 是割点
}
```

桥

- 无向图
- 注意原图不连通和重边

```
int dfn[N], low[N], clk;
void init() { memset(dfn, 0, sizeof dfn); clk = 0; }
void tarjan(int u, int fa) {
    low[u] = dfn[u] = ++clk;
    int _fst = 0;
    for (E& e: G[u]) {
        int v = e.to; if (v == fa && ++_fst == 1) continue;
        if (!dfn[v]) {
            tarjan(v, u);
            if (low[v] > dfn[u]) // (u, v) 是桥
            low[u] = min(low[u], low[v]);
        } else low[u] = min(low[u], dfn[v]);
```

```
        }
}
```

强连通分量缩点

- 有向图
- B：强连通分量的数量计数器
- bl[N]：记录每个顶点所属的强连通分量编号
- bcc[N]：存储每个强连通分量包含的顶点列表

```
int low[N], dfn[N], clk, B, bl[N];
vector<int> bcc[N];
void init() { B = clk = 0; memset(dfn, 0, sizeof dfn); }
void tarjan(int u) {
    static int st[N], p;
    static bool in[N];
    dfn[u] = low[u] = ++clk;
    st[p++] = u; in[u] = true;
    for (int& v: G[u]) {
        if (!dfn[v]) {
            tarjan(v);
            low[u] = min(low[u], low[v]);
        } else if (in[v]) low[u] = min(low[u], dfn[v]);
    }
    if (dfn[u] == low[u]) {
        while (1) {
            int x = st[--p]; in[x] = false;
            bl[x] = B; bcc[B].push_back(x);
            if (x == u) break;
        }
        ++B;
    }
}
```

点双连通分量 / 广义圆方树

- 数组开两倍
- 一条边也被计入点双了（适合拿来建圆方树），可以用点数 <= 边数过滤
- B：双连通分量的数量（编号从 0 开始）。
- bc[B]：存储第 B 个双连通分量包含的节点。
- be[B]：存储第 B 个双连通分量包含的边（索引）。
- bno[x]：标记节点 x 属于哪个双连通分量（用于去重）。

```
struct E { int to, nxt; } e[N];
int hd[N], ecnt;
void addedge(int u, int v) {
    e[ecnt] = {v, hd[u]};
    hd[u] = ecnt++;
}
int low[N], dfn[N], clk, B, bno[N];
vector<int> bc[N], be[N];
bool vise[N];
void init() {
    memset(vise, 0, sizeof vise);
    memset(hd, -1, sizeof hd);
    memset(dfn, 0, sizeof dfn);
    memset(bno, -1, sizeof bno);
    B = clk = ecnt = 0;
}

void tarjan(int u, int feid) {
    static int st[N], p;
    static auto add = [&](int x) {
        if (bno[x] != B) { bno[x] = B; bc[B].push_back(x); }
    };
    low[u] = dfn[u] = ++clk;
    for (int i = hd[u]; ~i; i = e[i].nxt) {
        if ((feid ^ i) == 1) continue;
        if (!vise[i]) { st[p++] = i; vise[i] = vise[i ^ 1] = true; }
```

```
            int v = e[i].to;
            if (!dfn[v]) {
                tarjan(v, i);
                low[u] = min(low[u], low[v]);
                if (low[v] >= dfn[u]) {
                    bc[B].clear(); be[B].clear();
                    while (1) {
                        int eid = st[--p];
                        add(e[eid].to); add(e[eid ^ 1].to);
                        be[B].push_back(eid);
                        if ((eid ^ i) <= 1) break;
                    }
                    ++B;
                }
            } else low[u] = min(low[u], dfn[v]);
        }
}
```

计算几何

字符串

最小表示法

- 寻找一个字符串的循环同构串中最小的那一个，输出偏移量

```
int min_string(string s){
    int k = 0, i = 0, j = 1, n = s.length();
    while (k < n && i < n && j < n) {
        if (s[(i + k) % n] == s[(j + k) % n]) {
            k++;
        } else {
            s[(i + k) % n] > s[(j + k) % n] ? i = i + k + 1 : j = j + k + 1;
            if (i == j) i++;
            k = 0;
        }
    }
    return min(i, j);
}
```

字符串哈希

```
// 双值哈希开关
#define ENABLE_DOUBLE_HASH

typedef long long LL;
typedef unsigned long long ULL;

const int x = 135;
const int N = 4e5 + 10;
const int p1 = 1e9 + 7, p2 = 1e9 + 9;
ULL xp1[N], xp2[N], xp[N];

void init_xp() {
    xp1[0] = xp2[0] = xp[0] = 1;
    for (int i = 1; i < N; ++i) {
        xp1[i] = xp1[i - 1] * x % p1;
        xp2[i] = xp2[i - 1] * x % p2;
        xp[i] = xp[i - 1] * x;
    }
}

struct String {
    string s;
    int length, subsize;
    bool sorted;
    ULL h[N], hl[N];

```

```
    // 预处理并返回全串哈希 O(n)
    ULL hash() {
        length = s.length();
        ULL res1 = 0, res2 = 0;
        h[length] = 0;  // ATTENTION!
        for (int j = length - 1; j >= 0; --j) {
#ifdef ENABLE_DOUBLE_HASH
            res1 = (res1 * x + s[j]) % p1;
            res2 = (res2 * x + s[j]) % p2;
            h[j] = (res1 << 32) | res2;
#else
            res1 = res1 * x + s[j];
            h[j] = res1;
#endif
            // printf("%llu\n", h[j]);
        }
        return h[0];
    }

    // 获取子串哈希，左闭右开区间 O(1)
    ULL get_substring_hash(int left, int right) const {
        int len = right - left;
#ifdef ENABLE_DOUBLE_HASH
        // get hash of s[left...right-1]
        unsigned int mask32 = ~(0u);
        ULL left1 = h[left] >> 32, right1 = h[right] >> 32;
        ULL left2 = h[left] & mask32, right2 = h[right] & mask32;
        return (((left1 - right1 * xp1[len] % p1 + p1) % p1) << 32) |
        (((left2 - right2 * xp2[len] % p2 + p2) % p2));
#else
        return h[left] - h[right] * xp[len];
#endif
    }

    void get_all_subs_hash(int sublen) {
        subsize = length - sublen + 1;
        for (int i = 0; i < subsize; ++i)
            hl[i] = get_substring_hash(i, i + sublen);
        sorted = 0;
    }

    void sort_substring_hash() {
        sort(hl, hl + subsize);
        sorted = 1;
    }

    bool match(ULL key) const {
//      if (!sorted) assert (0);
        if (!subsize) return false;
        return binary_search(hl, hl + subsize, key);
    }

    void init(string t) {
        length = t.length();
        s = t;
    }
};

String S, T; // 栈溢出

// 验证 S 中长度为 ans 的子串是否都存在于 T 中（是 0 否 1）
int check(String &S, String &T, int ans) {
    if (T.length < ans) return 1;
    T.get_all_subs_hash(ans); T.sort_substring_hash();
    for (int i = 0; i < S.length - ans + 1; ++i)
        if (!T.match(S.get_substring_hash(i, i + ans)))
            return 1;
    return 0;
}

// 返回是否匹配
```

```cpp
bool match_once(String &S, String &T){
    S.get_all_subs_hash(T.length);
    S.sort_substring_hash();
    return S.match(T.get_substring_hash(0, T.length));
}

// 返回匹配下标
vector<int> match_any(const String &text, const String &pattern) {
    vector<int> positions;
    int n = text.length;
    int m = pattern.length;

    if (m == 0 || m > n) return positions;

    ULL pattern_hash = pattern.get_substring_hash(0, m);

    for (int i = 0; i <= n - m; ++i) {
        ULL text_sub_hash = text.get_substring_hash(i, i + m);
        if (text_sub_hash == pattern_hash) {
            positions.push_back(i);
        }
    }
    return positions;
}

// 最长公共前缀 a[ai...] == b[bi...]
int LCP(const String &a, const String &b, int ai, int bi) {
    int l = 0, r = min(a.length - ai, b.length - bi);
    while (l < r) {
        int mid = (l + r + 1) / 2;
        if (a.get_substring_hash(ai, ai + mid) == b.get_substring_hash(bi, bi + mid))
            l = mid;
        else r = mid - 1;
    }
    return l;
}

// --------------------------- Template End ----------------------------
// !!!!!!!!!!!!!!!!!!!!!!!!!!! Tester Start !!!!!!!!!!!!!!!!!!!!!!!!!!!!

void solve(){
//  cout << "AA\n";
    init_xp();  // DON'T FORGET TO DO THIS!
//  cout << "BB\n";
    string s, t;
    cin >> s >> t;
    S.init(s), T.init(t);
    S.hash(), T.hash();
    cout << match_once(S, T) << '\n';

    vector<int> v = match_any(S, T);
    for(int ii: v) cout << ii << ' ';
    cout << '\n';

    cout << "LCP:" << LCP(S, T, 0, 0) << '\n';

    // S 中所有长度为 l 的子串均在 T 中出现, 且 l 最大
    LL l=0, r=S.length;
    while (l < r){
        int mid = l + r + 1 >> 1;
        if (!check(S, T, mid)) l = mid;
        else r = mid - 1;
    }
    cout << "check: " << l << '\n';
}
```

杂项

日期

```cpp
string day_of_week[] = {"Mo", "Tu", "We", "Th", "Fr", "Sa", "Su"};

// 格里高利历 (yyyy-mm-dd) 转儒略历 (整型/天)
int date_to_int(int y, int m, int d){
    return
    1461 * (y + 4800 + (m - 14) / 12) / 4 +
    367 * (m - 2 - (m - 14) / 12 * 12) / 12 -
    3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +
    d - 32075;
}

// 儒略历转格里高利历
void int_to_date(int jd, int &y, int &m, int &d){
    int x, n, i, j;
    x = jd + 68569;
    n = 4 * x / 146097;
    x -= (146097 * n + 3) / 4;
    i = (4000 * (x + 1)) / 1461001;
    x -= 1461 * i / 4 - 31;
    j = 80 * x / 2447;
    d = x - 2447 * j / 80;
    x = j / 11;
    m = j + 2 - 12 * x;
    y = 100 * (n - 49) + i + x;
}
```

随机

随机素数表

NTT 素数表

$p = r2^k + 1$，原根是 g。

$ (MOD, G, K, C) $ 满足：MOD 是质数，G 是 MOD 的原根，$MOD - 1 = C \times 2^K$

挑选方法:

- MOD 大于系数最大值的平方乘以多项式长度
- $2^m \le 2^K$，其中 2^m 为多项式长度

3, 1, 1, 2; 5, 1, 2, 2; 17, 1, 4, 3; 97, 3, 5, 5; 193, 3, 6, 5; 257, 1, 8, 3; 7681, 15, 9, 17; 12289, 3, 12, 11; 40961, 5, 13, 3; 65537, 1, 16, 3; 786433, 3, 18, 10; 5767169, 11, 19, 3; 7340033, 7, 20, 3; 23068673, 11, 21, 3; 104857601, 25, 22, 3; 167772161, 5, 25, 3; 469762049, 7, 26, 3; 1004535809, 479, 21, 3; 2013265921, 15, 27, 31; 2281701377, 17, 27, 3; 3221225473, 3, 30, 5; 75161927681, 35, 31, 3; 77309411329, 9, 33, 7; 206158430209, 3, 36, 22; 2061584302081, 15, 37, 7; 2748779069441, 5, 39, 3; 6597069766657, 3, 41, 5; 39582418599937, 9, 42, 5; 79164837199873, 9, 43, 5; 263882790666241, 15, 44, 7; 1231453023109121, 35, 45, 3; 1337006139375617, 19, 46, 3; 3799912185593857, 27, 47, 5; 4222124650659841, 15, 48, 19; 7881299347898369, 7, 50, 6; 31525197391593473, 7, 52, 3; 180143985094819841, 5, 55, 6; 1945555039024054273, 27, 56, 5; 4179340454199820289, 29, 57, 3.

注意事项

- `1LL << k`
- `(LL)v.size()`
- 输入要读完
- 不要把 `while` 写成 `if`
- 树链剖分/dfs 序，初始化或者询问不要忘记 idx, ridx
- 想清楚到底是要 `multiset` 还是 `set`
- 数据结构注意数组大小（2 倍，4 倍）
- 模意义下不要用除法