Introduction to NLP

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Introduction to Parsing

Parsing programming languages

```
#include <stdio.h>
int main()
  int n, reverse = 0;
 printf("Enter a number to reverse\n");
  scanf("%d", &n);
 while (n != 0)
    reverse = reverse * 10;
    reverse = reverse + n%10;
   n = n/10;
 printf("Reverse of entered number is = %d\n", reverse);
 return 0;
```

Parsing human languages

- Rather different than computer languages
 - Can you think in which ways?

Parsing human languages

- Rather different than computer languages
 - No types for words
 - No brackets around phrases
 - Ambiguity
 - Words
 - Parses
 - Implied information

The parsing problem

- Parsing means associating tree structures to a sentence, given a grammar (e.g., CFG)
 - There may be exactly one such tree structure
 - There may be many such structures
 - There may be none
- Grammars (e.g., CFG) are declarative
 - They don't specify how the parse tree will be constructed

Syntactic Issues

- PP attachment
 - I saw the man with the telescope
- Gaps
 - Mary likes Physics but hates Chemistry
- Coordination scope
 - Small boys and girls are playing
- Particles vs. prepositions
 - She ran up a large bill
- Gerund vs. adjective
 - Frightening kids can cause trouble

Applications of parsing

- Grammar checking
 - I want to return this shoes.
- Question answering
 - How many people in sales make \$40K or more per year?
- Machine translation
 - E.g., word order SVO vs. SOV
- Information extraction
 - Breaking Bad takes place in New Mexico.
- Speech generation
- Speech understanding

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Parsing Evaluation

Parsing Evaluation

- Parseval: precision and recall
 - get the proper constituents
- Labeled precision and recall
 - also get the correct non-terminal labels
- F1
 - harmonic mean of precision and recall
- Crossing brackets
 - (A (B C)) vs ((A B) C)
- PTB corpus
 - training 02-21, development 22, test 23

Evaluation Example

```
GOLD = (S (NP (DT The) (JJ Japanese) (JJ industrial) (NNS companies))

(VP (MD should) (VP (VB know) (ADVP (JJR better)))) (...)

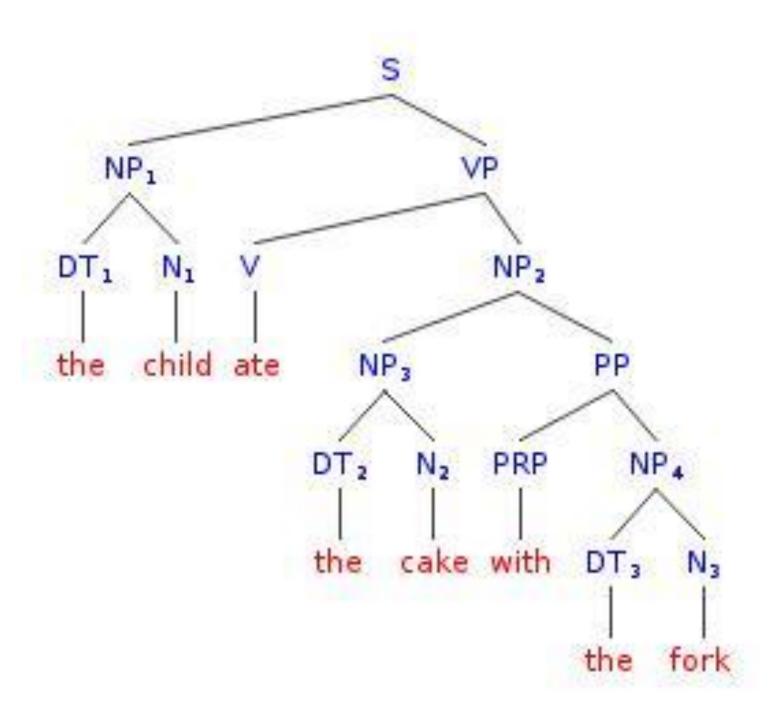
CHAR = (S (NP (DT The) (JJ Japanese) (JJ industrial) (NNS companies))

(VP (MD should) (VP (VB know)) ((ADVP (RBR better)))) (...))
```

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Classic parsing methods



```
S -> NP VP
NP -> DT N | NP PP
PP -> PRP NP
VP -> V NP | VP PP
DT -> 'a' | 'the'
N -> 'child' | 'cake' | 'fork'
PRP -> 'with' | 'to'
V -> 'saw' | 'ate'
```

Parsing as search

- There are two types of constraints on the parses
 - From the input sentence
 - From the grammar
- Therefore, two general approaches to parsing
 - Top-down
 - Bottom-up

```
S -> NP VP

NP -> DT N | NP PP

PP -> PRP NP

VP -> V NP | VP PP

DT -> 'a' | 'the'

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```

```
S -> NP VP

NP -> DT N | NP PP

PP -> PRP NP

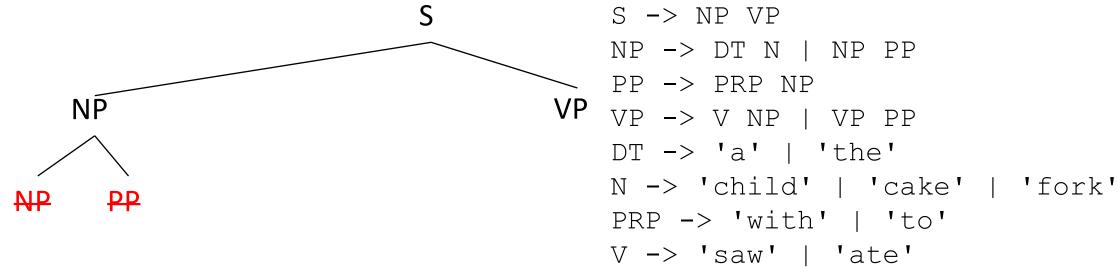
VP VP -> V NP | VP PP

DT -> 'a' | 'the'

N -> 'child' | 'cake' | 'fork'

PRP -> 'with' | 'to'

V -> 'saw' | 'ate'
```



```
S -> NP VP

NP -> DT N | NP PP

PP -> PRP NP

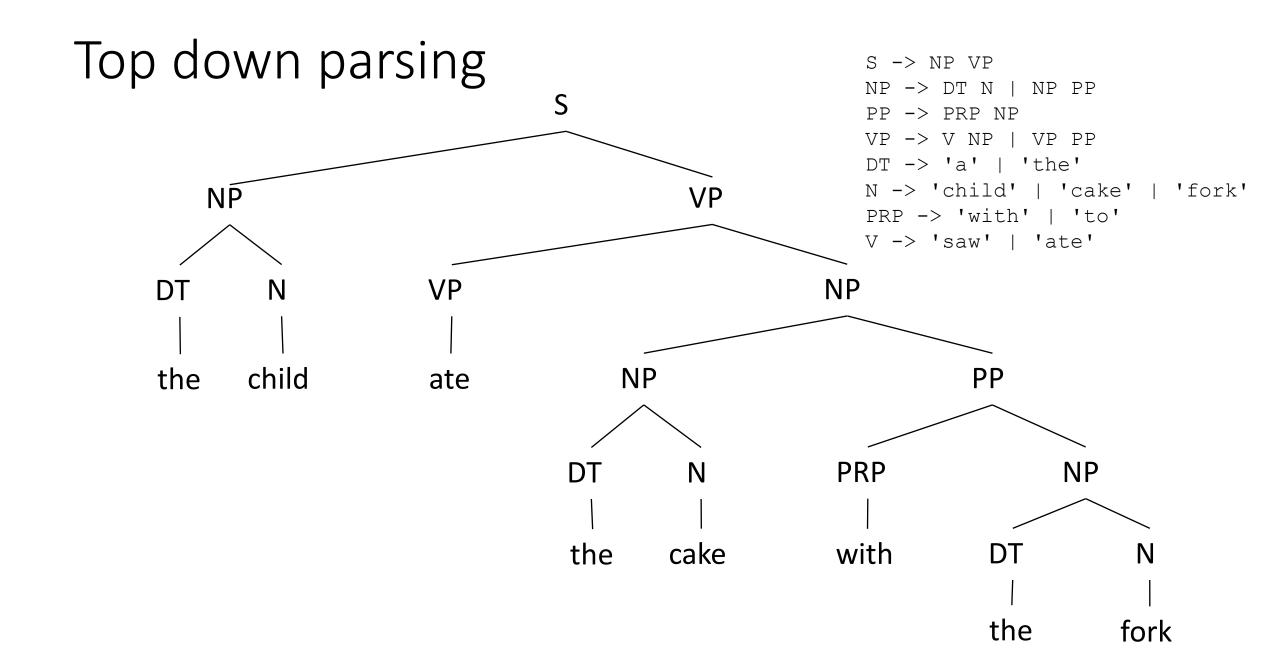
VP VP -> V NP | VP PP

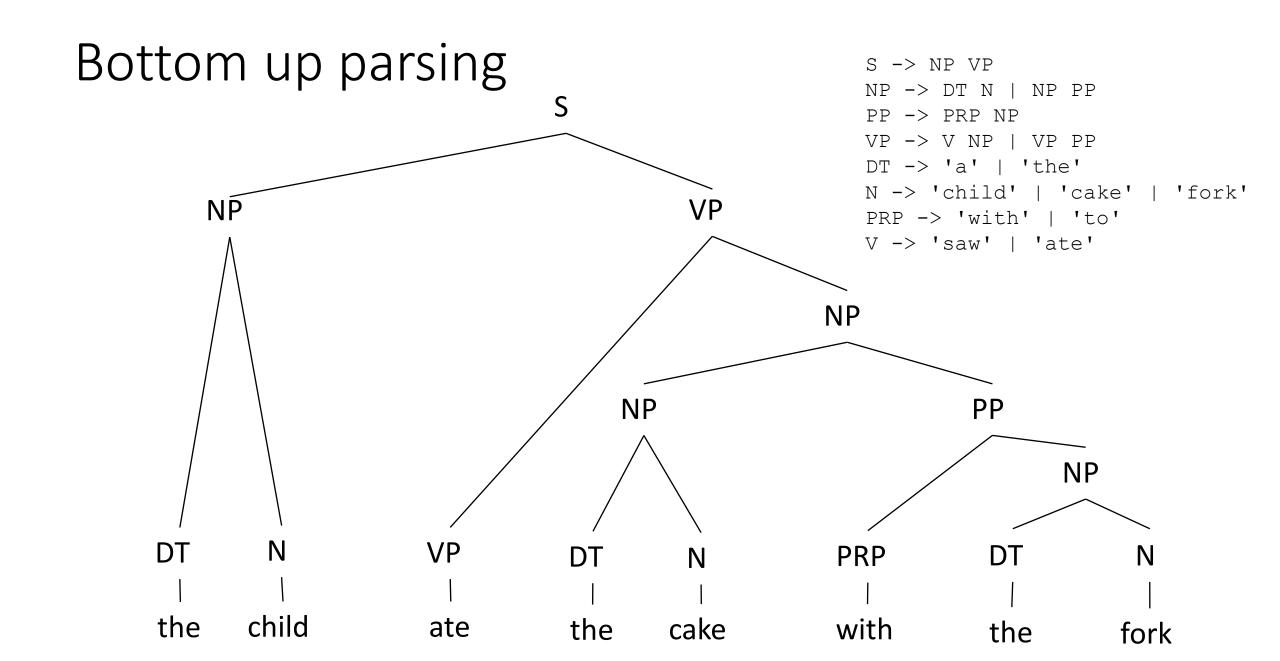
DT -> 'a' | 'the'

N -> 'child' | 'cake' | 'fork'

PRP -> 'with' | 'to'

V -> 'saw' | 'ate'
```





Bottom up vs. top down methods

- Bottom up
 - explores options that won't lead to a full parse
 - Example: shift-reduce (srparser in nltk)
 - Example: CKY (Cocke-Kasami-Younger)
- Top down
 - explores options that don't match the full sentence
 - Example: recursive descent (rdparser in nltk)
 - Example: Earley parser
- Dynamic programming
 - caches of intermediate results (memoization)

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Shift-Reduce Parsing

Shift-Reduce Parsing

- A bottom-up parser
 - Tries to match the RHS of a production until it can build an S
- Shift operation
 - Each word in the input sentence is pushed onto a stack
- Reduce-n operation
 - If the top *n* words on the top of the stack match the RHS of a production, then they are popped and replaced by the LHS of the production
- Breadth-first search
- Stopping condition
 - The process stops when the input sentence has been processed and S has been popped from the stack

Shift-Reduce Parsing Example

```
[ * the child ate the cake]
 S [ 'the' * child ate the cake]
 R [ DT * child ate the cake]
 S [ DT 'child' * ate the cake]
 R [ DT N * ate the cake]
 R [ NP * ate the cake]
 S [ NP 'ate' * the cake]
 R [ NP V * the cake]
 S [ NP V 'the' * cake]
 R [ NP V DT * cake]
 S [ NP V DT 'cake' * ]
 R [ NP V DT N * ]
 R [ NP V NP * ]
 R [ NP VP * ]
 R [ S * ]
(S (NP (DT the) (N child)) (VP (V ate) (NP (DT the) (N cake))))
```

Shift-Reduce Parsing

• In nltk

```
>>> from nltk.app import srparser;
>>> srparser())
```

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Cocke-Kasami-Younger (CKY) Parsing

Notes on Left Recursion

- Problematic for many parsing methods
 - Infinite loops when expanding
- But appropriate linguistically
 - NP -> DT N
 - NP -> PN
 - DT -> NP 's
 - Mary's mother's sister's friend

Chart Parsing

- Top-down parsers have problems with expanding the same non-terminal
 - In particular, pre-terminals such as POS
 - Bad idea to use top-down (recursive descent) parsing as is
- Bottom-up parsers have problems with generating locally feasible subtrees that are not viable globally
- Chart parsing will address these issues

Dynamic Programming

Motivation

- A lot of the work is repeated
- Caching intermediate results improves the complexity

Dynamic programming

• Building a parse for a substring [i,j] based on all parses [i,k] and [k, j] that are included in it.

Complexity

• $O(n^3)$ for recognizing an input string of length n

Dynamic Programming

- CKY (Cocke-Kasami-Younger)
 - bottom-up
 - requires a normalized (binarized) grammar
- Earley parser
 - top-down
 - more complicated
 - (separate lecture)

CKY Algorithm

```
function cky (sentence W, grammar G) returns table
  for i in 1..length(W) do
    table [i-1,i] = \{A \mid A->Wi \text{ in } G\}
  for j in 2...length(W) do
    for i in j-2 down to 0 do
       for k in (i+1) to (j-1) do
         table[i,j] = table[i,j] union \{A \mid A->BC \text{ in } G, B \text{ in } G\}
table [I,k], C in table [k,j]
```

If the start symbol S is in table [0,n] then W is in L(G)

Example

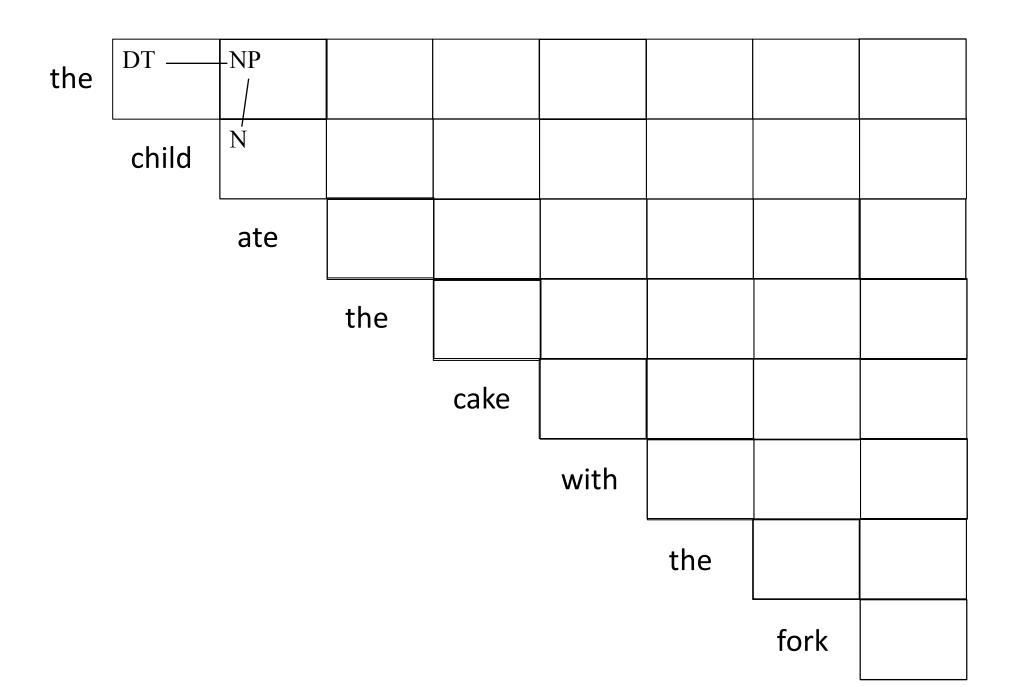
```
["the", "child", "ate", "the", "cake", "with", "the", "fork"]
      S \rightarrow NP VP
      NP -> DT N | NP PP
      PP -> PRP NP
      VP -> V NP | VP PP
      DT -> 'a' | 'the'
      N -> 'child' | 'cake' | 'fork'
      PRP -> 'with' | 'to'
      V -> 'saw' | 'ate'
```

the								
	child							
		ate						
			the					
				cake				
				·	with			
						the		
						'	fork	

the	DT							
	child							
		ate						
		·	the					
				cake				
				·	with			
						the		
						'	fork	

the	DT							
	child	N						
		ate						
			the					
				cake				
					with			
						the		
							fork	

the	DT	NP						
	child	N						
		ate						
			the					
			·	cake				
					with			
						the		
							fork	



the	DT	NP						
	child	N						
		ate	V					
			the					
			·	cake				
					with			
						the		
							fork	

the	DT	NP						
	child	N						
		ate	V					
			the	DT				
				cake				
				·	with			
						the		
						'	fork	

the	DT	NP						
	child	N						
		ate	V					
			the	DT				
				cake	N			
					with			
						the		
						,	fork	

the	DT	NP						
	child	N						
		ate	V					
			the	DT	NP			
				cake	N			
					with			
						the		
						'	fork	

the	DT	NP						
	child	N						
		ate	V					
			the	DT —	-NP			
				cake	N			
					with			
					'	the		
							fork	

the	DT	NP						
	child	N						
		ate	V		VP			
			the	DT	NP			
				cake	N			
					with			
						the		
						·	fork	

the	DT	NP						
	child	N						
		ate	V		VP 			
			the	DT	NP			
				cake	N			
					with			
						the		
						'	fork	

the	DT	NP			S			
	child	N						
		ate	V		VP			
			the	DT	NP			
				cake	N			
					with			
						the		
							fork	

the	DT	NP			-S			
	child	N						
		ate	V		VP			
			the	DT	NP			
				cake	N			
					with			
						the		
							fork	

the	DT	NP			S			
	child	N						
		ate	V		VP			
			the	DT	NP			
				cake	N			
					with	PRP		
						the		
							fork	

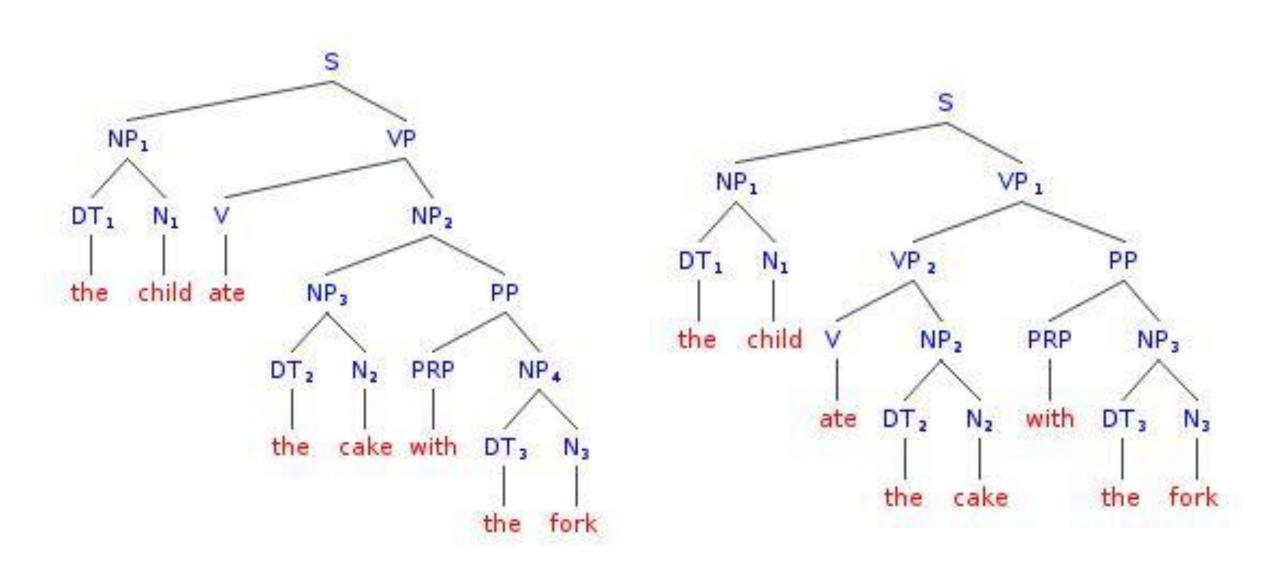
the	DT	NP			S			
	child	N						
		ate	V		VP			
			the	DT	NP			NP
				cake	N			
					with	PRP		PP
						the	DT	NP
							fork	N

the	DT	NP			S			
	child	N						
		ate	V		VP —			-VP
			the	DT	NP			NP
				cake	N			
					with	PRP		PP
						the	DT	NP
							fork	N

the	DT	NP			S			
	child	N						
		ate	V		VP			-VP
			the	DT	NP			NP
				cake	N			
					with	PRP		PP
						the	DT	NP
							fork	N

the	DT	NP			S			-S
	child	N						
		ate	V		VP			VP
			the	DT	NP			NP
				cake	N			
					with	PRP		PP
						the	DT	NP
							fork	N

	t	ne	DT		NP					S			S
		·	chi	ld	N								
					ate	9	V			VP			VP
							th	е	DT	NP			NP
[0]	DT	[1]			==>			[2]	cake	N			
[3][6][2][5]	DT DT V PRP	[4][7][3][6]	N N NP NP	[5] [8] [5] [8]	==>	[6] [2] [5]	NP NP VP PP	[5] [8] [5] [8]		with	PRP		PP
[0] [3] [2]	NP NP V	[2] [5]	VP PP NP	[5] [8]	==>	[0] [3]	S NP	[5] [8] [8]			the	DT	NP
[2] [0]	VP NP	[5]	PP	[8]	==>	[2]		[8]				fork	N



What is the *meaning* of each of these sentences?

```
(S
  (NP (DT the) (N child))
  (VP
     (VP (V ate) (NP (DT the) (N cake)))
     (PP (PRP with) (NP (DT the) (N fork)))))
```

```
(S
  (NP (DT the) (N child))
  (VP
        (V ate) (NP (DT the) (N cake)))
        (PRP with) (NP (DT the) (N fork))))
(S
  (NP (DT the) (N child))
  (VP
    (V ate)
    (NP
          (DT the) (N cake))
          (PRP with) (NP (DT the) (N fork)))))
```

Complexity of CKY

- Space complexity
 - There are $O(n^2)$ cells in the table
- Single parse
 - Each cell requires a linear lookup.
 - Total time complexity is $O(n^3)$
- All parses
 - Total time complexity is exponential

A longer example

```
["take", "this", "book"]
    S -> NP VP | Aux NP VP | VP
   NP -> PRON | Det Nom
   Nom -> N \mid Nom N \mid Nom PP
   PP -> PRP NP
   VP -> V | V NP | VP PP
   Det -> 'the' | 'a' | 'this'
   PRON -> 'he' | 'she'
   N -> 'book' | 'boys' | 'girl'
   PRP -> 'with' | 'in'
   V -> 'takes' | 'take'
```

Non-binary productions

```
["take", "this", "book"]
    S -> NP VP | Aux NP VP | VP
    NP -> PRON | Det Nom
    Nom -> N | Nom N | Nom PP
    PP -> PRP NP
    VP \rightarrow V \mid V NP \mid VP PP
    Det -> 'the' | 'a' | 'this'
    PRON -> 'he' | 'she'
    N -> 'book' | 'boys' | 'girl'
    PRP -> 'with' | 'in'
    V -> 'takes' | 'take'
```

Chomsky Normal Form (CNF)

- All rules have to be in binary form:
 - $X \rightarrow YZ$ or $X \rightarrow W$
- This introduces new non-terminals for
 - hybrid rules
 - n-ary rules
 - unary rules
 - epsilon rules (e.g., NP $\rightarrow \epsilon$)
- Any CFG can be converted to CNF
 - See Aho & Ullman p. 152

ATIS grammar

Original version

```
S \rightarrow NP VP
```

 $S \rightarrow Aux NP VP$

 $S \rightarrow VP$

 $NP \rightarrow Pronoun$

NP → Proper-Noun

 $NP \rightarrow Det Nominal$

Nominal \rightarrow Noun

Nominal → Nominal Noun

Nominal → Nominal PP

 $VP \rightarrow Verb$

 $VP \rightarrow Verb NP$

 $VP \rightarrow VP PP$

 $PP \rightarrow Prep NP$

ATIS grammar in CNF

Original version

```
S \rightarrow NP VP
```

 $S \rightarrow Aux NP VP$

$$S \rightarrow VP$$

 $NP \rightarrow Pronoun$

 $NP \rightarrow Proper-Noun$

 $NP \rightarrow Det Nominal$

Nominal → Noun

Nominal → Nominal Noun

Nominal → Nominal PP

 $VP \rightarrow Verb$

 $VP \rightarrow Verb NP$

 $VP \rightarrow VP PP$

 $PP \rightarrow Prep NP$

CNF version

 $S \rightarrow NP VP$

 $S \rightarrow X1 VP$

 $X1 \rightarrow Aux NP$

 $S \rightarrow book \mid include \mid prefer$

 $S \rightarrow Verb NP$

 $S \rightarrow VP PP$

 $NP \rightarrow I \mid he \mid she \mid me$

 $NP \rightarrow Houston \mid NWA$

 $NP \rightarrow Det Nominal$

Nominal → book | flight | meal | money

Nominal → Nominal Noun

Nominal → Nominal PP

 $VP \rightarrow book \mid include \mid prefer$

 $VP \rightarrow Verb NP$

 $VP \rightarrow VP PP$

 $PP \rightarrow Prep NP$

ATIS grammar in CNF

Original version

```
S \rightarrow NP VP
```

 $S \rightarrow Aux NP VP$

 $S \rightarrow VP$

NP → Pronoun

NP → **Proper-Noun**

 $NP \rightarrow Det Nominal$

Nominal → **Noun**

Nominal → Nominal Noun

Nominal → Nominal PP

 $VP \rightarrow Verb$

 $VP \rightarrow Verb NP$

 $VP \rightarrow VP PP$

 $PP \rightarrow Prep NP$

CNF version

 $S \rightarrow NP VP$

 $S \rightarrow X1 VP$

 $X1 \rightarrow Aux NP$

 $S \rightarrow book \mid include \mid prefer$

 $S \rightarrow Verb NP$

 $S \rightarrow VP PP$

 $NP \rightarrow I \mid he \mid she \mid me$

NP → Houston | NWA

 $NP \rightarrow Det Nominal$

Nominal → book | flight | meal | money

Nominal → Nominal Noun

Nominal → Nominal PP

 $VP \rightarrow book \mid include \mid prefer$

 $VP \rightarrow Verb NP$

 $VP \rightarrow VP PP$

 $PP \rightarrow Prep NP$

Chomsky Normal Form

- All rules have to be in binary form:
 - $X \rightarrow YZ$ or $X \rightarrow w$
- New non-terminals for hybrid rules, n-ary and unary rules:
 - INF-VP → to VP becomes
 - INF-VP → TO VP
 - TO → to
 - $S \rightarrow Aux NP VP$ becomes
 - $S \rightarrow R1 VP$
 - R1 \rightarrow Aux NP
 - $S \rightarrow VP$ $VP \rightarrow Verb$ $VP \rightarrow Verb$ $VP \rightarrow Verb$ PP becomes
 - $S \rightarrow book$
 - $S \rightarrow buy$
 - $S \rightarrow R2 PP$
 - $S \rightarrow Verb PP$
 - etc.

Issues with CKY

- Weak equivalence only
 - Same language, different structure
 - If the grammar had to be converted to CNF, then the final parse tree doesn't match the original grammar
 - However, it can be converted back using a specific procedure
- Syntactic ambiguity
 - (Deterministic) CKY has no way to perform syntactic disambiguation

Notes

- Demo:
 - http://lxmls.it.pt/2015/cky.html
- Recognizing vs. parsing
 - Recognizing just means determining if the string is part of the language defined by the CFG
 - Parsing is more complicated it involves producing a parse tree