

# FUNCTIONAL AND CONCURRENT PROGRAMMING

SI4

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# **MORE ABOUT FUNCTIONS**

Closure, Currying, Tail Recursion

# CLOSURE

- **Functions have a context**
- Example
  - ```
Lst<Integer> greaterThan(Lst<Integer> l, int v) {  
    return filter(x -> x > v, l);  
}
```
  - $v$  is defined in the context of the lambda  $x \rightarrow x > v$
- **Closure = function + context**
- **What we use to call “function” is a closure!**
  - Function may use or not this context

# CLOSURE AND METHOD REFERENCES

- **Every functional interface has a context**

- Example

- ```
class Counter {  
    int v = 0;  
    public int next() {  
        return v++;  
    }  
}
```

- counterInstance::next has v in its context
- Counter::next does not

# CLOSURE AND VARIABLES

- Functions cannot modify local variables

- ```
static void print(Lst<T> l) {  
    int i = 0;  
    iter(x -> { System.out.println((i++) + ": " + x); }, l);  
}
```

- Accessed variables must be (implicitly) final

- ```
static void print(Lst<T> l) {  
    int i = 0; i += 1;  
    iter(x -> { System.out.println(i + ": " + x); }, l);  
}
```

## CLOSURE AND SIDE-EFFECT

- But closures can access mutable objects

- ```
static void print(Lst<T> l) {  
    final AtomicInteger i = new AtomicInteger(0);  
    iter(x -> { System.out.println(i.getAndIncrement() + ": " + x); }, l);  
}
```

- Or even modify (instance or class) attributes

- ```
static int i;  
static void print(Lst<T> l) {  
    i = 0;  
    iter(x -> { System.out.println((i++) + ": " + x); }, l);  
}
```

## NON-SAFE CLOSURE

- A closure that uses a mutable context is “NON-SAFE” (aka non-functional)
- No longer a pure function (something else may modify the context)
  - Non repeatable
  - Difficult to test
  - Very difficult to parallelise
  - ...

# NON-FUNCTIONALITY

- This print method is functional

- ```
static void print(Lst<T> l) {  
    final AtomicInteger i = new AtomicInteger(0);  
    iter(x -> { System.out.println(i.getAndIncrement() + ": " + x); }, l);  
}
```

- This one is not

- ```
static int i;  
static void print(Lst<T> l) {  
    i = 0;  
    iter(x -> { System.out.println((i++) + ": " + x); }, l);  
}
```



# CURRYING

- Translating the evaluation of a function that takes multiple arguments into evaluating a sequence of functions, each with a single argument
- Example
  - $f.\text{apply}(a, b) \quad \Rightarrow \quad f.\text{apply}(a).\text{apply}(b)$
  - $T \times R \rightarrow U \quad \Rightarrow \quad T \rightarrow (R \rightarrow U)$
  - `int add(int x, int y) { return x + y; }`
  - `Function<Integer,Integer> addc(int x) { return y -> x + y; }`

## CURRYING ADVANTAGE

- Make a closure functional
- ```
class Foo {  
    Integer b;  
    Lst<Integer> calc(Lst<Integer> l) {  
        return map(x -> x * b, l);  
    }  
  
    Lst<Integer> calcf(Lst<Integer> l, int a) {  
        Function<Integer, Integer> cf =  
            y -> x -> x * y;  
        return map(cf.apply(b), l);  
    }  
}
```
- calcf is thread-safe while calc is not

# CURRYING USAGE EXAMPLE

- Have more than 2 parameters (beside BiFunction)
  - $T_1 \times T_2 \times \dots \times T_N \rightarrow R \Rightarrow T_1 \rightarrow T_2 \rightarrow \dots \rightarrow T_N \rightarrow R$
- Builder Pattern
  - ```
class Letter {  
    static AddReturnAddress builder(){  
        return returnAddress  
            -> closing  
            -> dateOfLetter  
            -> insideAddress  
            -> salutation  
            -> body  
            -> new Letter(returnAddress, insideAddress, dateOfLetter,  
salutation, body, closing);  
    }  
  
    interface AddReturnAddress { Letter.AddClosing withReturnAddress(String  
returnAddress); }
```

# TAIL RECURSION

- **A function is tail recursive** if the recursive call is the **last statement** that is executed by the function

- ```
static void print(int n) {  
    if (n < 0) return;  
    System.out.print(" " + n);  
    print(n - 1);  
}
```

- But not

- ```
static int fact(int n) {  
    if (n == 0) return 1;  
    return n * fact(n - 1);  
}
```

## WHY TAIL RECURSION?

- A tail recursion fonction can be executed as a loop
  - So without call stack comsuption
  - Many languages compile/interpret tail recursion as loop (Tail Call Optimization)
  - But not (yet) Java ☹
    - Only with specific libraries (e.g. [https://github.com/judekeyser/tail\\_rec](https://github.com/judekeyser/tail_rec))
- Better control on time complexity
  - Double recursion :  $O(2^n)$  versus tail recursion :  $O(n \cdot c_{body}(n))$  - Usually
  - But still programming easiness of recursion

# TAIL ACCUMULATOR

- Tail recursive function often use helper functions with supplementary parameters for
  - Loop variables
  - Or accumulator (the result being calculated)
- Example
  - ```
static long power(int base, int exponent) {  
    return power(base, exponent, 1);  
}
```

  

```
static long power(int base, int exponent, int res) {  
    if (exponent == 0) return res;  
    return power(base, exponent - 1, res * base);  
}
```
- Helps memoization (dynamic programming approach)

## COMPLEXITY COMPARED

| Function      | Non-tail rec | Tail rec |
|---------------|--------------|----------|
| reverse list  | $O(n^2)$     | $O(n)$   |
| fibonnaci     | $O(2^n)$     | $O(n)$   |
| palindrome    | $O(n^2)$     | $O(n)$   |
| edit distance | $O(n^2)$     | $O(2^n)$ |
| ...           |              |          |

- Not always easy (e.g. tree traversals) or possible without stack simulation (e.g. quicksort)