

MIPS Architecture

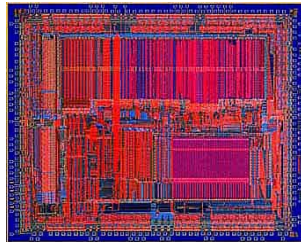
Many sheets copied from Andreas Klappenecker.
Additional information "MIPS Assembly Language programming",
Robert L. Britton (Prentice Hall, ISBN 0-13-142044-5) and
many sources on internet

MIPS Design Paradigms

- **Simplicity favors regularity**
 - all instructions single size
 - three register operands in arithmetic instr.
 - keep register fields in the same place
- **Smaller is faster**
 - 32 registers
- **Make the common case fast**
 - PC-relative addressing for conditional branches

MIPS R2000

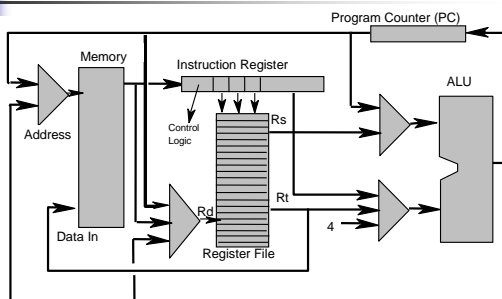
- **Several firsts:**
 - First RISC microprocessor
 - First microprocessor to provide integrated support for instruction & data cache
 - First pipelined microprocessor
- **Implemented in 1985**
 - 125,000 transistors
 - 5-8 MIPS



Basic Functional Components

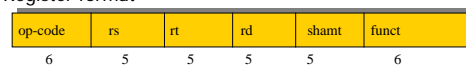
- Control unit
- Register file
- Arithmetic logic unit (ALU)
- Program counter (PC)
- Memory
- Instruction register (IR)

Simplified Datapath Diagram



Instruction Word Formats

- **Register format**



- **Immediate format**

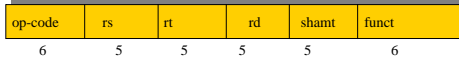


- **Jump format**



Register Format (R-Format)

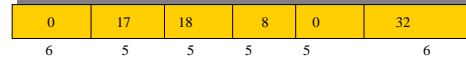
- Register format



- op: *basic operation of instruction*
- funct: *variant of instruction*
- rs: *first register source operand*
- rt: *second register source operand*
- rd: *register destination operand*
- shamt: *shift amount*

R-Format Example

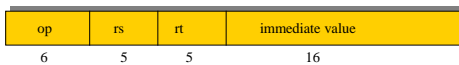
- Register format



- (op, funct)=(0,32): add
- rs=17: first source operand is \$s1
 - Note: de 32 register addresses have names starting with '\$'
- rt=18: second source operand is \$s2
- Rd=8: register destination is \$t0
- add \$t0, \$s1, \$s2

Immediate Format (I-Format)

- Immediate format



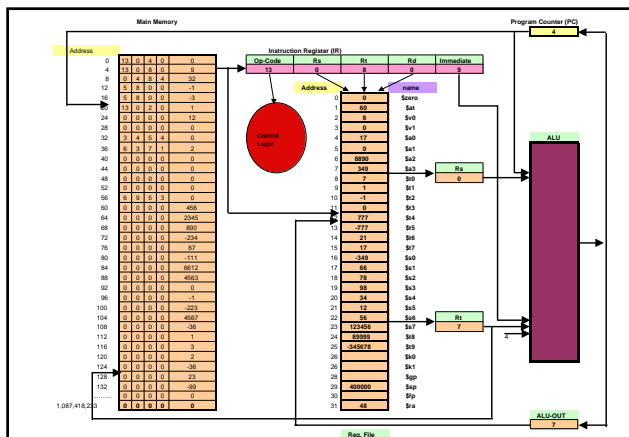
- op determines the instruction
- rs is the source register
- rt is the destination register
- 16bit immediate value

I-Format Example

- Immediate format



- op=8 means addi
- rs=29 means source register is \$sp
- rt=29 means \$sp is destination register
- immediate value = 4
- addi \$sp, \$sp, 4



Registers

Number	Value	Name
0		\$zero
1		\$at
2		\$v0
3		\$v1
4		\$a0
5		\$a1
6		\$a2
7		\$a3
8		\$t0
9		\$t1
10		\$t2
11		\$t3
12		\$t4
13		\$t5
14		\$t6
15		\$t7
16		\$s0
17		\$s1
18		\$s2
19		\$s3
20		\$s4
21		\$s5
22		\$s6
23		\$s7
24		\$t8

return values from functions
\$s0-\$s7 are used to hold temporary values. E.g. can be used in functions
\$s0-\$s7 are used to hold values that need to be saved (not modified) while functions are being called

Some problems

- Immediate field is 16 bits. How can we load a register with 32 bits?
- The result of multiplication of two 32 bits is 64 bits. Where is the result?

How can we load a 32bit constant?

- How can we load \$s0 with
 - 0000 0000 0011 1101 0000 1001 0000 0000
- Load upper 16bits with
 - `lui $s0, 0x003d` so that \$s0 contains (16 bits on the right)
 - 0000 0000 0011 1101 0000 0000 0000 0000
- Add immediate 16bit value to complete load
 - `addi $s0, $s0, 0x0900` so that \$s0 contains value
 - 0000 0000 0011 1101 0000 1001 0000 0000

Where are my multiply result?

- Multiple 32 bits x 32 bits → 64 bits
 - `mult $a1, $s1`
 - `mfhi $v0`
 - `mflo $v1`

MIPS Addressing Modes

- Register addressing
- Base displacement addressing
- Immediate addressing
- PC-relative addressing
 - address is the sum of the PC and a constant in the instruction
- Pseudo-direct addressing
 - jump address is 26bits of instruction concatenated with upper bits of PC

MIPS Assembly Language

MIPS Assembly Instructions

- `add $t0, $t1, $t2` # $\$t0 = \$t1 + \$t2$
- `sub $t0, $t1, $t2` # $\$t0 = \$t1 - \$t2$
- `lw $t1, a_addr` # $\$t1 = \text{Mem}[a_addr]$
- `sw $s1, a_addr` # $\text{Mem}[a_addr] = \$t1$

Assembler directives

- `.text` assembly instructions follow
- `.data` data follows
- `.globl` globally visible label
= symbolic address

Hello World!

```
.text                # code section
.globl main
main: li $v0, 4        # system call for print string
      la $a0, str      # load address of string to print
      syscall          # print the string
      li $v0, 10       # system call for exit
      syscall          # exit
.data
str:   .asciz "Hello world!\n" # NUL terminated string, as in C
```

Addressing modes

`lw $s1, addr` # load \$s1 from addr
`lw $s1, 8($s0)` # \$s1 = Mem[\$s0+8]
register `$s0` contains the base address
access the address (`$s0`)
possibly add an offset `8($s0)`

Load and move instructions

`la $a0, addr` # load address addr into \$a0 (macro)
`li $a0, 12` # load immediate \$a0 = 12
`lbu $a0, c($s1)` # load \$a0 = zero_extended(Mem[\$s1+c])
`lb $a0, c($s1)` # load \$a0 = sign_extended(Mem[\$s1+c])
`lw $a0, c($s1)` # load word
`move $s0, $s1` # \$s0 = \$s1

Control Structures

Assembly language has very few control structures:

- Branch instructions `if cond then goto label`
- Jump instructions `goto label`

We can build while loops, for loops, repeat-until loops,
if-then-else structures from these primitives

Branch instructions

<code>beqz \$s0, label</code>	<code>if \$s0==0</code>	<code>goto label</code>
<code>bnez \$s0, label</code>	<code>if \$s0!=0</code>	<code>goto label</code>
<code>bge \$s0, \$s1, label</code>	<code>if \$s0>=\$s1</code>	<code>goto label</code>
<code>ble \$s0, \$s1, label</code>	<code>if \$s0<=\$s1</code>	<code>goto label</code>
<code>blt \$s0, \$s1, label</code>	<code>if \$s0<\$s1</code>	<code>goto label</code>
<code>beq \$s0, \$s1, label</code>	<code>if \$s0==\$s1</code>	<code>goto label</code>
<code>bgez \$s0, \$s1, label</code>	<code>if \$s0>=0</code>	<code>goto label</code>

if-then-else structures

```
if ($t0==$t1) then /* blockA */ else /* blockB */
```

```
    beq $t0, $t1, blockA
```

```
    j blockB
```

```
blockA: ... instructions of then block ...
```

```
    j exit
```

```
blockB: ... instructions of else block ...
```

```
exit:  ... subsequent instructions ...
```

repeat-until loop

```
repeat ... until $t0>$t1
```

```
loop: ... instructions of loop ...
```

```
    ble $t0, $t1, loop    # if $t0<=$t1 goto loop
```

Other loop structures are similar...

System calls

- load argument registers
- load call code
- syscall

```
li $a0, 10    # load argument $a0=10
```

```
li $v0, 1     # call code to print integer
```

```
syscall       # print $a0
```

SPIM system calls

procedure	code \$v0	argument
print int	1	\$a0 contains number
print float	2	\$f12 contains number
print double	3	\$f12 contains number
print string	4	\$a0 address of string

SPIM system calls

procedure	code \$v0	result
read int	5	res returned in \$v0
read float	6	res returned in \$f0
read double	7	res returned in \$f0
Exit	10	

Example programs

- Loop printing integers 1 to 10

```
1  
2  
3
```

- Increasing array elements by 5

```
for(i=0; i<len; i++) {  
    a[i] = a[i] + 5;  
}
```

Print numbers 1 to 10

```
main: li $s0, 1          # $s0 = loop counter
      li $s1, 10        # $s1 = upper bound of loop
loop: move $a0, $s0      # print loop counter $s0
      li $v0, 1
      syscall
      li $v0, 4          # print "\n"
      la $a0, linebrk    # linebrk: .asciz "\n"
      syscall
      addi $s0, $s0, 1    # increase counter by 1
      ble $s0, $s1, loop # if ($s0 <= $s1) goto loop
      li $v0, 10        # exit
      syscall
      .data
linebrk: .asciz "\n"
```

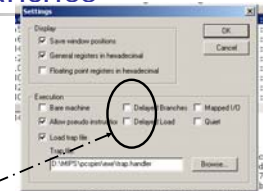
Increase array elements by 5

```
.text
.globl main
main: la $t0, Aaddr      # $t0 = pointer to array A
      lw $t1, len        # $t1 = length (of array A)
      sll $t1, $t1, 2     # $t1 = 4*length (all=shift left logical)
      add $t1, $t1, $t0   # $t1 = address(A)+4*length
loop: lw $t2, 0($t0)      # $t2 = A[i]
      addi $t2, $t2, 5    # $t2 = $t2 + 5
      sw $t2, 0($t0)      # A[i] = $t2
      addi $t0, $t0, 4    # i = i+1
      bne $t0, $t1, loop # if $t0 < $t1 goto loop
      .data
Aaddr: .word 0,2,1,4,5    # array with 5 elements
len:   .word 5
```

Procedures

- jal addr
 - (store address + (2x4)) into \$ra
 - Jump to the **second** instruction after the branch (see page 139 of MIPS32™ Architecture For Programmers Volume II, Revision 2.00)
 - jump to address addr
- jr \$ra
 - allows subroutine to jump back
 - Execute the instruction following the jump, in the branch delay slot, before jumping (page 145 of document above)
 - care must be taken to preserve \$ra!
 - more work for non-leaf procedures

(No) delayed branches



- PCPSIM: Simulator → settings
 - If these options have "no tick" then jump is to next address (also the next address is not executed in case of a branch).

Procedures

- one of the few means to structure your assembly language program
- small entities that can be tested separately
- can make an assembly program more readable
- recursive procedures

Write your own procedures

```
# prints the integer contained in $a0
print_int:
    li $v0, 1          # system call to
    syscall            # print integer
    jr $ra             # return
    nop                # this instruction is executed
                       # required in case of a
                       # 'delayed' branch.

main: . . .
    li $a0, 10         # we want to print 10
    jal print_int      # print integer in $a0
```

Write your own procedures

```
.data
linebrk: .asciiz "\n"
.text
print_eol:                # prints "\n"
    li $v0, 4             #
    la $a0, linebrk       #
    syscall               #
    jr $ra                # return
                          # this instruction is executed
                          # required in case of a
                          # 'delayed' branch.
    nop
main: . . .
    jal print_eol         # printf("\n")
```

Write your own procedures

```
.data
main:
    li $s0, 1             # $s0 = loop ctr
    li $s1, 10            # $s1 = upperbnd
loop: move $a0, $s0        # print loop ctr
    jal print_int         #
    jal print_eol         # print "\n"
    addi $s0, $s0, 1      # loop ctr +1
    ble $s0, $s1, loop    # unless $s0>$s1...
```

Non-leaf procedures

- Suppose that a procedure `procA` calls another procedure `jal procB`
- Problem:** `jal` stores return address of procedure `procB` and destroys return address of procedure `procA`
- Save `$ra` and all necessary variables onto the stack, call `procB`, and restore

Stack

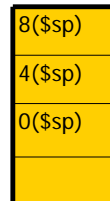
The stack can be used for

- parameter passing
- storing return addresses
- storing result variables
- stack pointer `$sp`

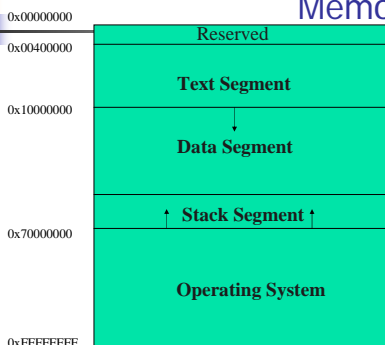
high address

stack pointer `$sp -->`

low address



Memory layout



Factorial

$$0! = 1$$

$$1! = 1$$

$$N! = N \times (N-1)!$$

Factorial is a recursive procedure with one argument `$v0`

Solution based on description in "MIPS Assembly Language programming"

```

        .data
prompt: .asciiz "\n\n Give me a value for 'N': "
msg:    .asciiz " N factorial is: "
bye:    .asciiz "\n### Good-bye ###"
        .text
        .globl main
main:    addiu    $sp, $sp, -8      # allocate space
mloop:
        li      $v0, 4
        la      $a0, prompt
        syscall
        li      $v0, 5           # Get value for N
        syscall
        bltz    $v0, quit        # Quit if <0
        sw      $v0, 0($sp)
        jal     fac              # Call factorial
        nop
        li      $v0, 4
        la      $a0, msg
        syscall
        lw      $a0, 4($sp)      # Get result
        li      $v0, 1
        syscall
        b       mloop
quit:
        addiu    $sp, $sp, 8      # Deallocate space
        li      $v0, 4
        la      $a0, bye
        syscall
        li      $v0, 10
        syscall

```

```

fac:
        lw      $a0, 0($sp)
        bltz    $a0, problem
        addi    $t1, $a0, -13
        bgtz    $t1, problem     # 13 is largest value we can accept
        addiu   $sp, $sp, -16    # Allocate
        sw      $ra, 12($sp)     # Save return address
        sw      $a0, 8($sp)
        slti    $t0, $a0, 2      # If N is 1 or 0 then return the value 1
        beqz    $t0, go
        li      $v0, 1
        b       facret
        # If value is 0 or 1
        # no recursive call
        # Check if value is
        # Between 0 and 13

```

```

go:      addi    $a0, $a0, -1
        sw      $a0, 0($sp)      # Pass N-1 to factorial function
        jal     fac              # Recursive call
        nop
        lw      $v0, 4($sp)      # Get (N-1)! back
        lw      $ra, 12($sp)
        lw      $a0, 8($sp)
        mult    $v0, $a0         # N*(N-1)!
        mflo    $v0
facret:  addiu   $sp, $sp, 16     # Deallocate
        sw      $v0, 4($sp)
        jr      $ra
problem: sw      $v0, 4($sp)
        jr      $ra

```