

Types and Lambda Calculus Notes

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These notes are not necessarily correct, consistent, representative of the course as it stands today, or rigorous. Any result of the above is not the author's fault.

These notes are in progress.

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1 The Axioms of Lambda Calculus

We suppose that we have a countably infinite set of variables \mathbb{V} (we usually refer to elements of this set as x, y, z , etc.), from this we define the alphabet of lambda calculus $\mathbb{V} + \{\lambda, ., (,)\}$.

The set of terms of lambda calculus Λ is defined inductively by the axioms. A string is a term if and only if there is a proof tree using the following axioms that concludes it is a term.

1.1 Axiom of Variables

For some x in \mathbb{V} , we have the axiom:

$$\overline{x \in \Lambda},$$

which says that each variable is a term.

1.2 Axiom of Application

We have the axiom:

$$\frac{M \in \Lambda \quad N \in \Lambda}{(MN) \in \Lambda},$$

which says that for some terms M and N , (MN) is also a term.

1.3 Axiom of Abstraction

For some x in \mathbb{V} , we have the axiom:

$$\frac{M \in \Lambda}{(\lambda x.M) \in \Lambda},$$

which says that for some term M , $(\lambda x.M)$ is also a term.

1.4 Subterms

Subterms of a term M are substrings of M that are themselves terms and not captured by a λ (directly preceded by).

1.5 Syntactical Conventions

Parentheses allow our lambda calculus to be unambiguous, but for the sake of simplicity, we will construct conventions that will allow us to retain unique meaning with less parentheses:

- Omit outermost parentheses,
- Terms associate to the left, (MNP) parses as $((MN)P)$,
- Bodies of abstractions end at parentheses, $(\lambda x.MN)$ parses as $(\lambda x.(MN))$,
- Group repeated abstractions, $(\lambda xy.M)$ parses as $(\lambda x.(\lambda y.M))$.

2 Alpha

2.1 Free Variables

We define the function $FV : \Lambda \rightarrow \mathcal{P}(\mathbb{V})$, which returns the set of variables contained within a term M that are not bound. We define it recursively on the structure of terms:

$$\begin{aligned} FV(x) &= \{x\}, \\ FV(MN) &= FV(M) \cup FV(N), \\ FV(\lambda x.M) &= FV(M) \setminus \{x\}. \end{aligned}$$

If a term has no free variables we say it is closed, and if a term has at least one free variable then we say it is open. The set of all closed terms is denoted by Λ^0 .

2.2 Substitution

We define 'capture-avoiding' substitution of a term M for a variable x recursively on the structure of terms:

$$\begin{aligned} y[M/x] &= y && \text{if } y \neq x, \\ y[M/x] &= M && \text{if } y = x, \\ (PQ)[M/x] &= P[M/x] \cup Q[M/x], \\ (\lambda y.P)[M/x] &= \lambda y.P && \text{if } y = x, \\ (\lambda y.P)[M/x] &= \lambda y.P[M/x] && \text{if } y \neq x \text{ and } y \notin FV(M). \end{aligned}$$

On the final case, we stipulate that y cannot be a free variable of M because otherwise free variables in the substitution could be captured by the lambda.

2.3 Alpha Equivalence

Suppose we have a term $\lambda x.M$ and y in $\mathbb{V} \setminus FV(M)$, then substituting y for x is a change of bound variable name. If two terms can be made identical through changes of bound variable name, they are α -equivalent. The set of λ -terms is the set Λ under α -equivalence.

This equivalence is much more useful to us than string comparison, so for the remainder of the notes we will always be referring to λ -terms as terms.

2.4 The Variable Convention

For M_1, \dots, M_k terms occurring in the same scope, we assume each term has distinct bound variables. We can make this assumption as otherwise, we can use changes of bound variable names to make it so.

3 Beta

A term of the form $(\lambda x.M)N$ is called a β -redex and $M[N/x]$ is the contraction of the redex.

3.1 Beta Reductions

The one-step β -reduction relation, denoted by \rightarrow_β , is inductively defined by the redex rule:

$$\overline{(\lambda x.M)N \rightarrow_\beta M[N/x]},$$

the left and right application rules:

$$\frac{M \rightarrow_\beta M'}{MN \rightarrow_\beta M'N} \quad \frac{N \rightarrow_\beta N'}{MN \rightarrow_\beta MN'},$$

and the abstraction rule:

$$\frac{M \rightarrow_\beta N}{\lambda x.M \rightarrow_\beta \lambda x.N}.$$

A term M is said to be in β -normal form if there is no term N such that $M \rightarrow_\beta N$.

In general, β -reductions are sequences of consecutive one-step β -reductions:

$$M_0 \rightarrow_\beta M_1 \rightarrow_\beta \cdots \rightarrow_\beta M_k,$$

for some k in \mathbb{N}_0 . We say that M_0 β -reduces to M_k , denoted by $M_0 \rightarrow_\beta M_k$.

If $M \rightarrow_\beta N$, we say that N is a reduct of M , and is a proper reduct if $N \neq M$. If we can choose some β -normal N such that $M \rightarrow_\beta N$ then M is normalisable. If a term admits no infinite β -reductions then we say that it is strongly normalisable.

3.2 Standard Combinators

We have some interesting programs described below:

$\mathbf{I} = \lambda x.x$ $\mathbf{K} = \lambda xy.x$ $\mathbf{S} = \lambda xyz.xz(yz)$ $\omega = \lambda x.xx$ $\Omega = \omega\omega$ $\Theta = (\lambda xy.y(xxy))(\lambda xy.y(xxy))$	$\mathbf{I}M \rightarrow_\beta M$ $\mathbf{K}MN \rightarrow_\beta M$ $\mathbf{S}MNP \rightarrow_\beta MP(NP)$ $\omega M \rightarrow_\beta MM$ $\Omega \rightarrow_\beta \Omega$ $\Theta M \rightarrow_\beta M(\Theta M)$
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3.3 Confluence of Beta

For a term M , if $M \rightarrow_{\beta} P$ and $M \rightarrow_{\beta} Q$ then there exists some term N such that $P \rightarrow_{\beta} N$ and $Q \rightarrow_{\beta} N$.

3.4 Beta Convertibility

For M, N terms, if M and N have a common reduct then we say that M and N are β -convertible, denoted by $M =_{\beta} N$.