COMP9334 Project Report Z5004850 Tengyu Ma

In this project, a simulation of a multi-server system was implemented in Python. The code was generate under the Python 3.5 with pylab installed.

## 1. Introduction of Design

According to the instruction of the multi-server system, preprocessor and subserves obey the first-come-first-serve queueing discipline and all of them has infinite buffer, which means during the simulation period, no task would be decline. So a simple way to implement the simulation would be just handle each arrival request and calculate the finish time. If any request go over the hole simulation time, the request will not take into consideration as it is not finished. In other word, if we handle the previous finish time of preprocessor and subserves, each request can be calculate independently.

## **Formula Design**

Internal arrive time is defined as the sum of two part. The sequence a1k is exponentially distributed with a mean arrival rate 0.85 requests per time unit. The sequence a2k is uniformly distributed in the interval [0.05, 0.25]. I use the formula in sim\_mm1:

#### -math.log(1-random.random())/rate

to generate exponentially distributed sequence where random.random() cloud generate random number in range[0,1):

#### random.uniform(0.05,0.25)

Previous function could generate uniformly distributed sequence, where uniform could generate uniform distribute in specific range. Inter\_arrival time is sum of those two random number

Subserves process time is the inverse function of the The probability density function f(ts). First the integral of the function should be calculated and then we could write the formula as following:

Here random.uniform(0,1) is used to simulate Pareto distribution in python, as there is no build function to generate it. But it could be implement by uniform distribution between 0 and 1. Thus, I use random.uniform(0,1).

## **Work flow Design**

When a request come into the system, a parameter is used to record the arrival time. While the arrival time is smaller than whole simulation period, we keep handle the request. Then we made a judgement whether the preprocessor is busy. Preprocessor was designed as an integer to record previous task's finish time. If the preprocessor is greater than current request's arrival time, preprocessor is busy. So current task will finish its preprocess in

## preprocessor+preprocess\_time

Otherwise current task will finish in

#### Arrival\_time + preprocess time

After preprocess, the whole task is distribute into n sub\_tasks and send to each subserves. Here a python build in function **random.sample((0,10),n)** is used to generate a randomly list of subserves distribution without repeat(we assume preprocessor do not repeatedly forward to subserves for one request). Subserves are designed as Array in order to record their previous task's finish time:

### $sub\_servers = [0]*10$

In each subserves we did same thing as what we had done for preprocessor. Comparing the value of subserves and renew the subserves. The finish time of a request is the longest finish time in each subserves. In this way, each request could be calculate independently.

The result of the simulation is stored in a list. The graph of the simulation or mean response time could be generate from the list.

In order to make the simulation reproducible, I use fake random which could generate same random number in different test.

# 2. Output Analysis

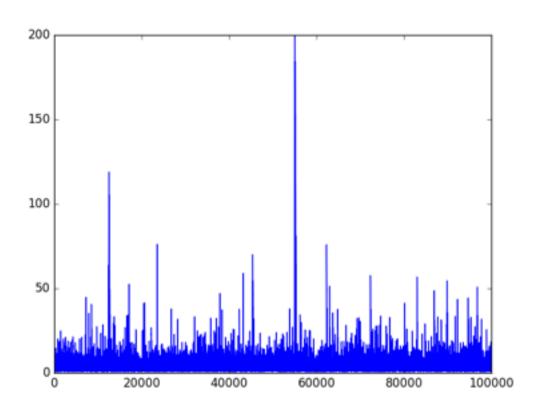
When set the whole simulation time as 10000, 80000, 90000, 100000and 200000 time unit, we could get following data set:

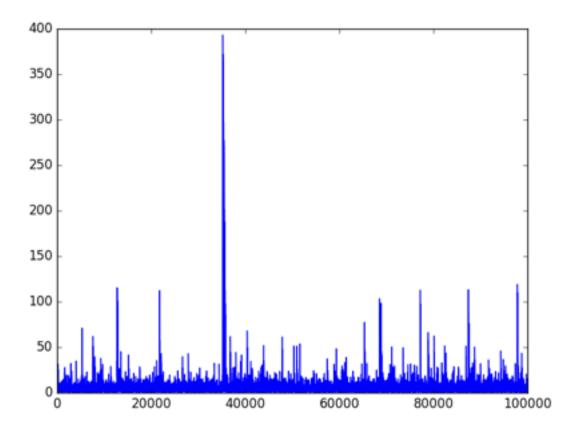
| ı | n | 10000                 | 80000                  | 90000                  | 100000                | 200000                | 300000        |
|---|---|-----------------------|------------------------|------------------------|-----------------------|-----------------------|---------------|
| - | 1 | 1584.1545225<br>50526 | 13646.969993<br>770357 | 15397.080720<br>752654 | 17192.30246<br>064073 | 34323.36268<br>422478 | 51349.6708340 |

| 2  | 142.91670579          | 300.32539670          | 273.62034990           | 258.0213212           | 280.3624003            | 328.2582478           |
|----|-----------------------|-----------------------|------------------------|-----------------------|------------------------|-----------------------|
|    | 648658                | 816595                | 62273                  | 923813                | 812613                 | 097392                |
| 3  | 28.987195470<br>80968 | 31.587639270<br>24156 | 31.014511132<br>401942 | 31.65077172<br>174367 | 1.34015238353(         | 33.31181835<br>073347 |
| 4  | 14.369633845          | 15.931379480          | 16.040043226           | 17.07701258           | 16.18986494            | 17.19878589           |
|    | 52062                 | 487541                | 215605                 | 357501                | 608148                 | 5210353               |
| 5  | 10.571561664          | 18.517094002          | 9.9846227384           | 9.743526070           | 12.70911726            | 28.49081981           |
|    | 603573                | 257593                | 16883                  | 379309                | 5624542                | 3913117               |
| 6  | 7.6021382108<br>54852 | 8.1375035206<br>67755 | .5591854178997         | 8.358498062<br>145816 | 10.84551657<br>9330193 | 8.772087243<br>548873 |
| 7  | 7.1779874906          | 8.8735829476          | 7.1936460243           | 7.445295893           | 8.975120302            | 8.981218331           |
|    | 45033                 | 36786                 | 97184                  | 778244                | 817741                 | 644696                |
| 8  | 6.0424016658<br>03292 | 7.0025784368<br>98714 | 6.9499175056<br>44183  | 7.252334583<br>016475 | 3.668591299962         | 6.879488979<br>911506 |
| 9  | .6933075440542        | 6.7698513984<br>25083 | 6.1956693247<br>03803  | 7.042908680<br>561254 | 8.000783066<br>363867  | 6.860541980<br>789675 |
| 10 | 6.3242722927          | 7.3289622749          | 8.7063468825           | 8.512071879           | 7.594867188            | 7.348871239           |
|    | 59348                 | 90749                 | 06097                  | 948845                | 841812                 | 605482                |

It is clear that although there are some fluctuate, when test time is long enough, the simulation will generate the solution that when break into 8 or 9 sub-tasks, the mean response time would be least. Generally, when distribute tasks into to many subtasks, there is a bottle neck on preprocessor as it will spend more time when breaking task. So there is a balance point for breaking tasks, too many subtask or too small separation will lead to longer response time.

In order to make a solution of which distribution is better, we should do further test.





Take n=8 and simulation time = 100000 time unit

We could find there are very few results which are extremely huge. According to test result diagram we can set the confident interval as results which is less than 50 as most of the data distribute in this area. So we could generate new data set for n=8 and n=9 to analysis which choice is better:

| n | 10000       | 80000       | 900000      | 100000      | 200000      | 300000      |
|---|-------------|-------------|-------------|-------------|-------------|-------------|
| 8 | 5.824452698 | 5.742732056 | 5.737556093 | 5.791103215 | 5.748870647 | 5.867649837 |
|   | 099063      | 222329      | 064794      | 949062      | 473025      | 149568      |
| 9 | 5.838954234 | 6.079313128 | 5.946730086 | 6.039459659 | 5.908259323 | 6.013586517 |
|   | 828516      | 693893      | 436914      | 539555      | 114307      | 2902525     |

New test result shows that when we ignore result greater than 50, each test case shows that the n=8 is a better choice. In this case we could say that distribute request into 8 sub-tasks will be a better choice