

Lab10-Turing Machine

CS214-Algorithm and Complexity, Xiaofeng Gao & Lei Wang, Spring 2021.

* If there is any problem, please contact TA Yihao Xie.

* Name: Beichen Yu Student ID: 519030910245 Email: polarisybc@sjtu.edu.cn

1. Design a one-tape TM M that computes the function $f(x, y) = \lfloor x/y \rfloor$, where x and y are positive integers ($x > y$). The alphabet is $\{1, 0, \square, \triangleright, \triangleleft\}$, and the inputs are x "1"s, \square and y "1"s. Below is the initial configuration for input $x = 7$ and $y = 3$. The result $z = f(x, y)$ should also be represented in the form of z "1"s on the tape with pattern of $\triangleright 111 \cdots 111 \triangleleft$, which is $\triangleright 11 \triangleleft$ for the example.

Initial Configuration

\triangleright	1	1	1	1	1	1	1	\square	1	1	1	\triangleleft
\uparrow												
q_s												

- (a) Please describe your design and then write the specifications of M in the form like $\langle q_s, \triangleright \rangle \rightarrow \langle q_1, \triangleright, R \rangle$. Explain the transition functions in detail.
- (b) Please draw the state transition diagram.
- (c) Show briefly and clearly the whole process from initial to final configurations for input $x = 7$ and $y = 3$. You may start like this:

$$(q_s, \underline{\triangleright} 1111111 \square 111 \triangleleft) \vdash (q_1, \underline{\triangleright} 1111111 \square 111 \triangleleft) \vdash^* (q_1, \triangleright 1111111 \underline{\square} 111 \triangleleft) \vdash (q_2, \triangleright 1111111 \square \underline{1} 11 \triangleleft)$$

(Note that for simplicity, we write $(q_1, \underline{\triangleright} 1111111 \square 111 \triangleleft) \vdash^* (q_1, \triangleright 1111111 \underline{\square} 111 \triangleleft)$ if the corresponding transaction repeats on multiple inputs with the same state.)

Solution. (a) Start state:

$$\langle q_s, \triangleright \rangle \rightarrow \langle q_1, \triangleright, R \rangle$$

Go to the first "1" of y :

$$\langle q_1, 1 \rangle \rightarrow \langle q_1, 1, R \rangle$$

$$\langle q_1, \square \rangle \rightarrow \langle q_2, \square, R \rangle$$

Change all "1" of y into 0, while delete the same number of "1" in x :

$$\langle q_2, 1 \rangle \rightarrow \langle q_3, 0, L \rangle$$

$$\langle q_3, 0 \rangle \rightarrow \langle q_3, 0, L \rangle$$

$$\langle q_3, \square \rangle \rightarrow \langle q_3, \square, L \rangle$$

$$\langle q_3, 1 \rangle \rightarrow \langle q_2, \square, R \rangle$$

$$\langle q_2, \square \rangle \rightarrow \langle q_2, \square, R \rangle$$

$$\langle q_2, 0 \rangle \rightarrow \langle q_2, 0, R \rangle$$

Write a "1" in the right side of \triangleleft , and turn all "0" back to "1":

$$\langle q_2, \triangleleft \rangle \rightarrow \langle q_4, \triangleleft, R \rangle$$

$$\langle q_4, 1 \rangle \rightarrow \langle q_4, 1, R \rangle$$

$$\langle q_4, \square \rangle \rightarrow \langle q_5, 1, L \rangle$$

$$\langle q_5, 1 \rangle \rightarrow \langle q_5, 1, L \rangle$$

$$\langle q_5, \triangleleft \rangle \rightarrow \langle q_5, \triangleleft, L \rangle$$

$$\langle q_5, 0 \rangle \rightarrow \langle q_5, 1, L \rangle$$

$$\langle q_5, \square \rangle \rightarrow \langle q_2, \square, R \rangle$$

If the remain “1” in x is less than the number of “1” in y , end the loop and get the answer:

$$\begin{aligned}
\langle q_3, \triangleright \rangle &\rightarrow \langle q_6, \square, R \rangle \\
\langle q_6, 1 \rangle &\rightarrow \langle q_6, \square, R \rangle \\
\langle q_6, \square \rangle &\rightarrow \langle q_6, \square, R \rangle \\
\langle q_6, 0 \rangle &\rightarrow \langle q_6, \square, R \rangle \\
\langle q_6, \triangleleft \rangle &\rightarrow \langle q_7, \triangleright, R \rangle \\
\langle q_7, 1 \rangle &\rightarrow \langle q_7, 1, R \rangle \\
\langle q_7, \square \rangle &\rightarrow \langle q_H, \triangleleft, R \rangle
\end{aligned}$$

(b) Below is the state transition diagram.

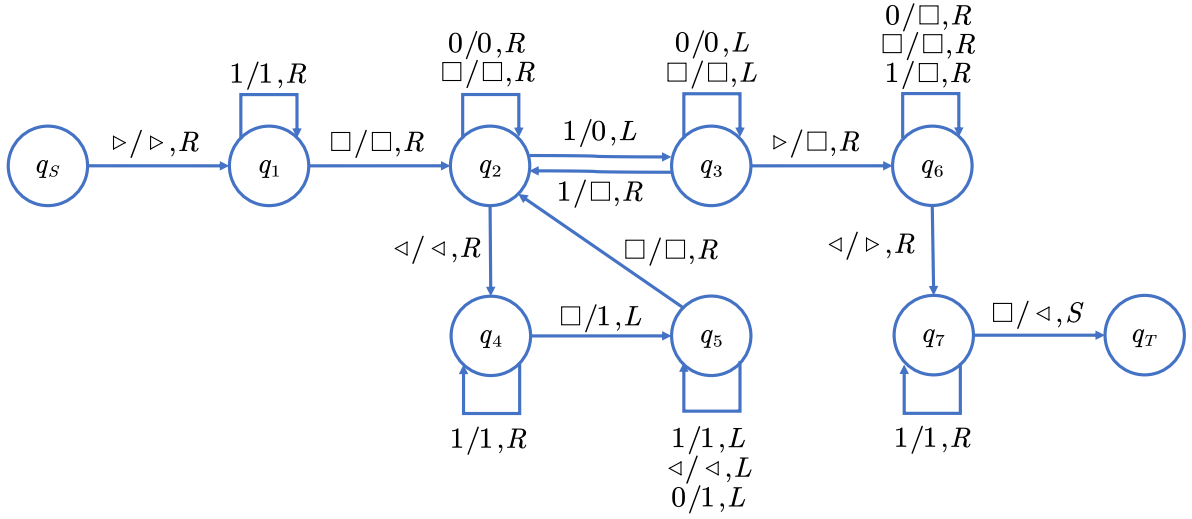


Figure 1: The state transition diagram

(c)

$$\begin{aligned}
&(q_s, \underline{\triangleright}1111111\square111\triangleleft) \\
&\vdash (q_1, \underline{\triangleright}1111111\square111\triangleleft) \\
&\vdash^* (q_1, \triangleright1111111\underline{\square}111\triangleleft) \\
&\vdash (q_2, \triangleright1111111\square\underline{1}11\triangleleft) \\
&\vdash (q_3, \triangleright1111111\underline{\square}011\triangleleft) \\
&\vdash (q_3, \triangleright111111\underline{1}\square011\triangleleft) \\
&\vdash (q_2, \triangleright111111\square\underline{\square}011\triangleleft) \\
&\vdash (q_2, \triangleright111111\square\square\underline{0}11\triangleleft) \\
&\vdash (q_2, \triangleright111111\square\square\underline{0}11\triangleleft) \\
&\vdash^* (q_2, \triangleright1111\square\square\square\underline{0}00\triangleleft)
\end{aligned}$$

$$\begin{aligned}
&\vdash(q_2, \triangleright 1111 \square \square \square \square 000 \triangleleft) \\
&\vdash(q_4, \triangleright 1111 \square \square \square \square 000 \triangleleft \square) \\
&\vdash(q_5, \triangleright 1111 \square \square \square \square 000 \triangleleft 1) \\
&\vdash(q_5, \triangleright 1111 \square \square \square \square 00 \square \triangleleft 1) \\
&\vdash(q_5, \triangleright 1111 \square \square \square \square 0 \square 1 \triangleleft 1) \\
&\vdash^*(q_5, \triangleright 1111 \square \square \square \square \square 111 \triangleleft 1) \\
&\vdash(q_2, \triangleright 1111 \square \square \square \square \square 11 \triangleleft 1) \\
&\vdash^*(q_2, \triangleright 1 \square \square \square \square \square \square \square 111 \triangleleft 11) \\
&\vdash^*(q_3, \triangleright \square \square \square \square \square \square \square \square 111 \triangleleft 11) \\
&\vdash(q_6, \square \square \square \square \square \square \square \square 111 \triangleleft 11) \\
&\vdash^*(q_6, \square 11 \triangleleft 11) \\
&\vdash(q_6, \square 1 \triangleleft 11) \\
&\vdash^*(q_6, \triangleleft 11) \\
&\vdash(q_7, \triangleright 11) \\
&\vdash(q_7, \triangleright 1 \square) \\
&\vdash(q_7, \triangleright 11 \square) \\
&\vdash(q_T, \triangleright 11 \triangleleft)
\end{aligned}$$

□

2. Given the alphabet $\{1, 0, \square, \triangleright, \triangleleft\}$, design a time efficient 3-tape TM M to compute $f : \{0, 1\}^* \rightarrow \{0, 1\}$ which verifies whether the number of 0 and the number of 1 are the same in an input consisting of only 0's and 1's. M should output 1 if the numbers are the same, and 0 otherwise. For example, for the input tape $\triangleright 001101 \triangleleft$, M should output 1
- (a) Please describe your design and then write the specifications of M in the form like $\langle q_S, \triangleright, \triangleright, \triangleright \rangle \rightarrow \langle q_1, \triangleright, \triangleright, R, R, S \rangle$. Explain the transition functions in detail.
- (b) Show the time complexity for one-tape TM M' to compute the same function f with n symbols in the input and give a brief description of such M' .

Solution. (a) Start state:

$$\langle q_S, \triangleright, \triangleright, \triangleright \rangle \rightarrow \langle q_1, \triangleright, \triangleright, R, R, S \rangle$$

Go right, write a “1” in the second tape if it reads a “1” in the first tape, and do nothing if reads “0” in the first tape:

$$\begin{aligned}
\langle q_1, 0, \square, \triangleright \rangle &\rightarrow \langle q_1, \square, \triangleright, R, S, S \rangle \\
\langle q_1, 1, \square, \triangleright \rangle &\rightarrow \langle q_1, 1, \triangleright, R, R, S \rangle \\
\langle q_1, \triangleleft, \square, \triangleright \rangle &\rightarrow \langle q_2, \square, \triangleright, L, L, S \rangle
\end{aligned}$$

Go left, move to left in the second tape if it reads a “0” in the first tape, and do nothing if reads “1” in the first tape:

$$\begin{aligned}
\langle q_2, 1, 1, \triangleright \rangle &\rightarrow \langle q_2, 1, \triangleright, L, S, S \rangle \\
\langle q_2, 0, 1, \triangleright \rangle &\rightarrow \langle q_2, 1, \triangleright, L, L, S \rangle
\end{aligned}$$

If the first tape and the second tape meet \triangleright together, then the number of 0 and the number of 1 are the same:

$$\begin{aligned}\langle q_2, \triangleright, \triangleright, \triangleright \rangle &\rightarrow \langle q_4, \triangleright, \triangleright, S, S, R \rangle \\ \langle q_4, \triangleright, \triangleright, \square \rangle &\rightarrow \langle q_5, \triangleright, 1, S, S, R \rangle \\ \langle q_5, \triangleright, \triangleright, \square \rangle &\rightarrow \langle q_H, \triangleright, \triangleleft, S, S, S \rangle\end{aligned}$$

Otherwise, then the number of 0 and the number of 1 are not the same:

$$\begin{aligned}\langle q_2, 0, \triangleright, \triangleright \rangle &\rightarrow \langle q_6, \triangleright, \triangleright, S, S, R \rangle \\ \langle q_2, 1, \triangleright, \triangleright \rangle &\rightarrow \langle q_6, \triangleright, \triangleright, S, S, R \rangle \\ \langle q_2, \triangleright, 1, \triangleright \rangle &\rightarrow \langle q_6, 1, \triangleright, S, S, R \rangle \\ \langle q_6, 0, \triangleright, \square \rangle &\rightarrow \langle q_7, \triangleright, 0, S, S, R \rangle \\ \langle q_6, 1, \triangleright, \square \rangle &\rightarrow \langle q_7, \triangleright, 0, S, S, R \rangle \\ \langle q_6, \triangleright, 1, \square \rangle &\rightarrow \langle q_7, 1, 0, S, S, R \rangle \\ \langle q_7, 0, \triangleright, \square \rangle &\rightarrow \langle q_H, \triangleright, \triangleleft, S, S, S \rangle \\ \langle q_7, 1, \triangleright, \square \rangle &\rightarrow \langle q_H, \triangleright, \triangleleft, S, S, S \rangle \\ \langle q_7, \triangleright, 1, \square \rangle &\rightarrow \langle q_H, 1, \triangleleft, S, S, S \rangle\end{aligned}$$

- (b) The M' can use the similar algorithm with the M . Each time the reading head reads a "1", it moves to the right side and write a flag; after it goes to the rightmost side, then turn left. In this turn, each time the reading head reads a "0", it moves to the right side and delete a flag. Then it can make a compare according to the number of flags.

Because the time complexity of moving right is $O(n)$ and there are n symbols in the input, the final time complexity is $O(n^2)$.

□

3. Define the corresponding decision or search problem of the following problems and give the "certificate" and "certifier" for each decision problem provided in the subquestions or defined by yourself.

- (a) *3-Dimensional Matching*. Given disjoint sets X, Y, Z all with the size of n , and a set $M \subseteq X \times Y \times Z$. Is there a subset M' of M of size n where no two elements of M' agree in any coordinate?
- (b) *Travelling Salesman Problem*. Given a list of cities and the distances between each pair of cities, find the shortest possible route that visits each city exactly once and returns to the origin city.
- (c) *Job Sequencing*. Given a set of unit-time jobs, each of which has an integer deadline and a nonnegative penalty for missing the deadline. Does there exist a job sequence that has a total penalty $w \leq k$?

Solution. (a) This is a decision problem.

The search problem is: Find a subset M' of M of size n where no two elements of M' agree in any coordinate.

Certificate: A subset M' of M of size n .

Certifier: Check that in any coordinate there are no two elements of M' agree.

(b) This is a search problem.

The decision problem is: Is there a route that visits each city exactly once and returns to the origin city and its length $l \leq k$.

Certificate: A route r_0 that visits each city exactly once and returns to the origin city.

Certifier: Assume that the length of r_0 is l_0 . Check that $l_0 \leq k$.

(c) This is a decision problem.

The search problem is: Find a job sequence with a smallest total penalty.

Certificate: A job sequence s .

Certifier: Check that the total penalty w of s is less or equal to k .

□

Remark: Please include your .pdf, .tex files for uploading with standard file names.