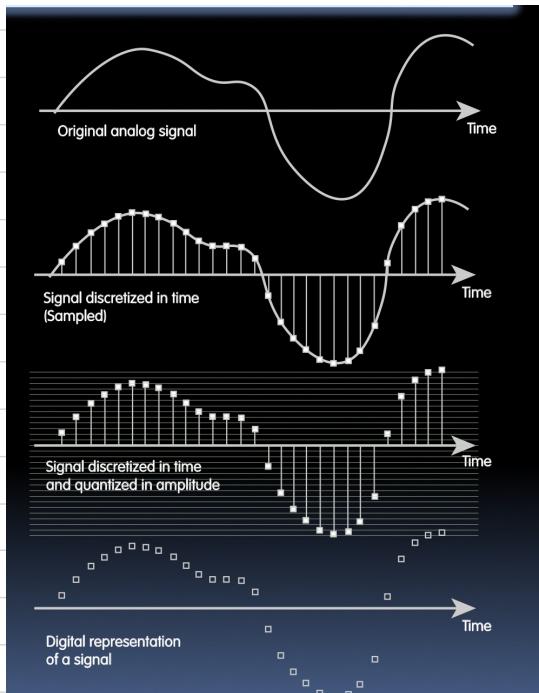


Number Representation

- Real world is analog!
- To import analog information, we must do two things
 - Sample
 - Quantize



Binary Decimal Hex

Numeral

A symbol or name that stands for a number
e.g., 4, four, quattro, IV, IIII, ...
...and **Digits** are symbols that make numerals

Number

The "idea" in our minds...there is only ONE of these
e.g., the concept of "4"

Base 10 (Ten) #s, Decimals

Digits: 0, 1, 2, 3, 4, 5, 6, 7, 8, 9

Example:

$$3271_{10} = 3 \times 10^3 + 2 \times 10^2 + 7 \times 10^1 + 1 \times 10^0$$

Base 2 (Two) #s, Binary (to Decimal)

Digits: 0, 1 [binary digits \rightarrow bits]

Example: "1101" in binary? ("0b1101")

$$1101_2 = (1 \times 2^3 + 1 \times 2^2 + 0 \times 2^1 + 1 \times 2^0)$$

$$= 8 + 4 + 0 + 1$$

$$= 13$$

Convert from Decimal to Binary

- E.g., 13 to binary
- Start with the column

13	2 ³ =8	2 ² =4	2 ¹ =2	2 ⁰ =1
5				
1	1	0	1	
0				

- Left to right, is (column) \leq number n?
 - If yes, put how many of that column fit in n, subtract col* that many from n, keeping going.
 - If not, put 0 and keep going. (and Stop at 0)

Convert from Decimal to Hexadecimal

- E.g., 165 to hexadecimal?
- Start with the columns

Start with the columns				
165	16 ³ =4096	16 ² =256	16 ¹ =16	16 ⁰ =1
5				
0	0	0	(10) A	5

- Left to right, is (column) \leq number n ?
 - If yes, put how many of that column fit in n , subtract col* that many from n , keeping going.
 - If not, put 0 and keep going. (and Stop at 0)

Convert Binary \longleftrightarrow Hexadecimal

- Binary \rightarrow Hex
- Always left-pad with 0s to make full 4-bit values, then look up!
- E.g., Db 11110 to Hex?

• 0b11110 \rightarrow 0b00011110
 • Then look up: 0x1E

- Hex \rightarrow Binary
- Just look up, drop leading 0s

• 0x1E \rightarrow 0b00011110 \rightarrow 0b11110

D	H	B
00	0	0000
01	1	0001
02	2	0010
03	3	0011
04	4	0100
05	5	0101
06	6	0110
07	7	0111
08	8	1000
09	9	1001
10	A	1010
11	B	1011
12	C	1100
13	D	1101
14	E	1110
15	F	1111

Decimal vs Hexadecimal vs Binary

- 4 Bits
 - 1 "Nibble"
 - 1 Hex Digit = 16 things
- 8 Bits
 - 1 "Byte"
 - 2 Hex Digits = 256 things
 - Color is usually

0-255 Red,
 0-255 Green,
 0-255 Blue.
 #D0367F = 

Number Representations

- add, subtract, multiply, divide, compare
- what if too big?

- Binary bit patterns are simply **representatives** of numbers. Abstraction!
- Strictly speaking they are called "**numerals**"
- Numerals really have an ∞ number of digits
 - with almost all being some 100...0 or 11...1) except for a few of the rightmost digits
- If result of add (or -, *, /) cannot be represented by these rightmost HW bits, we say **overflow** occurred

How to Represent Negative Numbers?

(C's `unsigned int`, C18's `uintN_t`)

- So far, **unsigned** numbers
- Obvious solution: define leftmost bit to be sign!
 - 0 \rightarrow + 1 \rightarrow - ...and rest of bits are numerical value
- Representation called **Sign and Magnitude**

Shortcoming of sign and Magnitude

- Arithmetic circuit complicated
 - Special steps depending on if signs are the same or not
- **two** zeros
- Incrementing "binary odometer", sometimes increases values, and sometimes decreases!
- Therefore Sign and magnitude used only in signal processors

Another try! complement the bits

- Example: $7_{10} = 00111_2$ $-7_{10} = 11000_2$
- Called One's Complement
- Note: positive numbers have leading 0s, negative numbers have leading 1s.



- What is -00000 ? Answer: 11111
- How many positive numbers in N bits?
- How many negative numbers?

D. 1.1

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Shortcomings of One's Complement

- Arithmetic still somewhat complicated
- Still two zeros
 - $0x00000000 = +0_{10}$
 - $0xFFFFFFFF = -0_{10}$
- Although used for a while on some computer products, one's complement was eventually abandoned because another solution was better.

Two's Complement & Bias Encoding

- Problem is the negative mappings “overlap” with the positive ones (the two 0s). Want to shift the negative mappings left by one.
 - Solution! For negative numbers, complement, then add 1 to the result
 - As with sign and magnitude, & one's compl. leading 0s → positive, leading 1s → negative
 - 00000...xxx is ≥ 0 , 11111...xxx is < 0
 - except 1...1111 is -1, not -0 (as in sign & mag.)
 - This representation is Two's Complement
 - This makes the hardware simple!
- (C's int, C18's intN_t, aka a “signed integer”)

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Two's Complement Formula

- Can represent positive and negative numbers in terms of the bit value times a power of 2:
 $d_{31} \times -(2^{31}) + d_{30} \times 2^{30} + \dots + d_2 \times 2^2 + d_1 \times 2^1 + d_0 \times 2^0$
- Example: 1101_{two} in a nibble?
$$\begin{aligned} &= 1 \times -(2^3) + 1 \times 2^2 + 0 \times 2^1 + 1 \times 2^0 \\ &= -2^3 + 2^2 + 0 + 2^0 \\ &= -8 + 4 + 0 + 1 \\ &= -8 + 5 \\ &= -3_{\text{ten}} \end{aligned}$$

Example: -3 to +3 to -3
(again, in a nibble):

x :	1101_{two}	-3
x' :	0010_{two}	3
+1 :	0011_{two}	3
()' :	1100_{two}	-3
+1 :	1101_{two}	-3

In Summary

- We represent “things” in computers as particular bit patterns: N bits $\Rightarrow 2^N$ things
- These 5 integer encodings have different benefits; 1s complement and sign/mag have most problems.

- **unsigned** (C18’s `uintN_t`) :

00000 00001 ... 01111 10000 ... 11111
↔ →

- **Two’s complement** (C18’s `intN_t`) universal, learn!

00000 00001 ... 01111
↔ →
10000 ... 11110 11111

- Overflow: numbers ∞ ; computers finite, errors!

Garcia,