01204211 Discrete Mathematics Lecture 18: Primality testing (3)

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The first algorithm

This lecture presents the primality testing algorithm based on the Fermat test.

To perform the test if q is a prime, we pick an integer a from the range $\{2,3,\ldots,q-1\}$ and check that the condition specified by the Fermat's Little Theorem is satisfied.

PROCEDURE FermatTest(q,a)

- 1. if $GCD(q, a) \neq 1$ then return "COMPOSITE" endif
- 2. Find $y = a^{q-1} \mod q$
- 3. if $y \neq 1$ then
- 4. return "COMPOSITE"
- 5. **else**
- 6. return "PRIME"
- 7. endif

Guarantees

Let's consider all input/output possibilities. Let z = FermatTest(q, a).

	z=PRIME	z=COMPOSITE
q is prime	Correct	Incorrect
q is composite	Incorrect	Correct

However, from the Fermat's Little Theorem, we know that if q is prime, the test always return "PRIME"; thus the test is always correct.

Guarantees

Let's consider all input/output possibilities. Let z = FermatTest(q, a).

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q is prime	Correct	
q is composite	Incorrect	Correct

However, from the Fermat's Little Theorem, we know that if q is prime, the test always return "PRIME"; thus the test is always correct.

We then need to consider the case when q is composite. In this case, the Fermat's Little Theorem does not provide any guarantee. However, we hope that if q is not a prime, we may be able to find a such that FermatTest(q,a) reveals the truth, i.e., it returns COMPOSITE.

Witness

- ▶ Let's try to be precise. For a composite q, if a is an integer such that FermatTest(q,a) returns COMPOSITE, we say that a is a witness for q.
- From the code of FermatTest, an integer a can be a witness for q in two cases:
 - (1) when $GCD(q, a) \neq 1$, or
 - (2) when $a^{q-1} \mod q \neq 1$.
- Note that if we pick a randomly from the set $\{2,3,\ldots,q-1\}$ and composite q is a product of very large primes, it is very unlikely that we pick a that satisfies (1) because a must be a multiple of some of q's prime factors.

Carmichael numbers

It would be great if, for any composite q, we can quickly find its witness. But unfortunately, there exists a composite q with very few witnesses. I.e., every witness a for q shares its factor (i.e., it is such that $gcd(a,q) \neq 1$). They are called **Carmichael numbers**.

A **Carmichael** number is a composite number q such that

$$b^{q-1} \mod q = 1,$$

for all integers 1 < b < n which are relatively prime to q.

► The first Carmichael number is 561. The next ones are 1105, 1729, and 2465.

Witnesses for non-Carmichael numbers (1)

Let's focus on the bright side. Suppose that q is not Carmichael, can we say anything about its witnesses? Can we say that there are plenty of them?

Witnesses for non-Carmichael numbers (2)

Let a be q's witness such that gcd(a,q)=1, i.e., we also have that

$$a^{q-1} \mod q \neq 1$$
.

Let's consider a non-witness b, i.e., b is an integer such that gcd(b,q)=1 and $b^{q-1} \mod q=1.$

Now, consider $ab \mod q$. We have that

$$(ab \mod q)^{q-1} \mod q = (ab)^{q-1} \mod q$$

= $(a^{q-1}b^{q-1}) \mod q$
= $(a^{q-1} \mod q)(b^{q-1} \mod q) \mod q$
= $a^{q-1} \mod q$
 $\neq 1$

Can we use this to say that there are a lot of witnesses?

A lot of witnesses

Since we can show that given a witness a and a non-witness b, ab is also witness, we might be able to construct a lot of witnesses from non-witnesses.

Let $A=\{2,3,\ldots,q-1\}$ be the set of all candidates. Let B be the set of non-witnesses which are relatively prime to q, i.e.,

$$B = \{ b \in A : (gcd(b, q) = 1) \land (b^{q-1} \mod q = 1) \}.$$

Let $C = \{ab \mod q : b \in A\}$. We know that every $c \in C$ is a witness.

If we know that |C|=|B|, can you show that the probability of choosing a witness from the set A is large (i.e, at least 1/2). Now, is it obvious that |C|=|B|? What is missing?

The missing argument

To show that |C|=|B|, we need to argue that when we multiply every element of B, we do not get duplicate elements. I.e., we need to prove that for $x\in B$ and $y\in B$ such that $x\neq y$,

$$ax \mod q \neq ay \mod q$$
,

when gcd(a,q) = 1.

Quick check: prove this statement.

(Note: From the definitions of a and B , you may assume that $\gcd(q,a)=\gcd(q,x)=\gcd(q,y)=1.$)

Conclusions

From the previous discussion, we know that for non-Carmichael numbers, the Fermat test succeeds with probability at least 1/2. Further developments:

- In 1976, Miller and Rabin show that one can deal with Carmichael numbers, providing the first randomized algorithm for testing primes.
- ▶ In 2002, Agrawal, Kayal, and Saxena devise an $O(m^{12})$ -time deterministic algorithm for primality testing.