01204211 Discrete Mathematics Lecture 2: Quantifiers and proofs

Jittat Fakcharoenphol

August 18, 2015

This lecture covers:

- More on quantifiers
- ► How to prove a proposition
- Basic proof techniques

Review: Quantifiers

- ▶ A *predicate* is a statement with variables, which can be either true or false, after all its variables are specified.
- ▶ If we quantify a predicate completely, the quantified expression now has a truth value, and it is called a quantified proposition.
- Two ways to quantify:
 - ▶ **Using universal quantifiers** (∀). This quantifier states that the quantified proposition is true when the predicate is true for every value of the variable in the specified set.
 - ▶ **Using existential quantifiers (**∃**).** This quantifier states that the quantified proposition is true when the predicate is true for at least one value of the variable in the specified set.
- Quantifiers can be nested. E.g.,
 - $\forall x \forall y P(x,y) \equiv \forall x (\forall y (P(x,y)))$
 - $\forall x \exists y P(x,y) \equiv \exists x (\forall y (P(x,y)))$

Quick check 1

Negations (1)

Let consider a set of positive integers \mathbb{Z}^+ as our universe. Let predicate $P(x)\equiv$ "x is a prime number." Consider this proposition

$$(\forall x \in \mathbb{Z}^+)P(x).$$

How can we show that this is false?

Negations (1)

Let consider a set of positive integers \mathbb{Z}^+ as our universe. Let predicate $P(x)\equiv$ "x is a prime number." Consider this proposition

$$(\forall x \in \mathbb{Z}^+)P(x).$$

How can we show that this is false?

When showing that a universally quantified proposition is false, we need to show "one" counter example. In this case, since P(4) is false, $\forall x P(x)$ is false.

Negations (1)

Let consider a set of positive integers \mathbb{Z}^+ as our universe. Let predicate $P(x)\equiv$ "x is a prime number." Consider this proposition

$$(\forall x \in \mathbb{Z}^+)P(x).$$

How can we show that this is false?

When showing that a universally quantified proposition is false, we need to show "one" counter example. In this case, since P(4) is false, $\forall x P(x)$ is false.

This way of disproving a statement is equivalent to showing that

$$(\exists x)(\neg P(x)).$$

Negations of quantified propositions

Let consider a set of positive integers \mathbb{Z}^+ as our universe. Let predicate $Q(x)\equiv$ "if x>2, then $x^2\leq 2x$." Consider this proposition

$$(\exists x \in \mathbb{Z}^+)Q(x).$$

How can we show that this is false?

Negations of quantified propositions

Let consider a set of positive integers \mathbb{Z}^+ as our universe. Let predicate $Q(x)\equiv$ "if x>2, then $x^2\leq 2x$." Consider this proposition

$$(\exists x \in \mathbb{Z}^+)Q(x).$$

How can we show that this is false?

When showing that an existential quantified proposition is false, we need to show that Q(x) is false for every possible values of x. In this case, since $x^2 = x \cdot x > 2 \cdot x$ for every x > 2, we have that $(\exists x)Q(x)$ is false.

Negations of quantified propositions

Let consider a set of positive integers \mathbb{Z}^+ as our universe. Let predicate $Q(x)\equiv$ "if x>2, then $x^2\leq 2x$." Consider this proposition

$$(\exists x \in \mathbb{Z}^+)Q(x).$$

How can we show that this is false?

When showing that an existential quantified proposition is false, we need to show that Q(x) is false for every possible values of x. In this case, since $x^2 = x \cdot x > 2 \cdot x$ for every x > 2, we have that $(\exists x)Q(x)$ is false.

This way of disproving a statement is equivalent to showing that

$$(\forall x)(\neg Q(x)).$$



Negations (3)

Thus, the following equivalences:

Quick check 2

How to prove a mathematical statement

Given propositions P and Q, these are a very useful logical equivalences (referred to as the De Morgan's Laws).

$$\neg (P \lor Q) \equiv \neg P \land \neg Q$$

$$\neg (P \land Q) \equiv \neg P \lor \neg Q$$

(Note that \neg takes precedence over \lor or \land .)

How can we prove that the first statement is true?

Proof by exhaustion

For any proposition
$$P$$
 and Q , $\neg(P \lor Q) \equiv \neg P \land \neg Q$.

Proof.

We will prove by exhaustion. There are 4 cases as in the truth table below.

P	Q	$P \lor Q$	$\neg (P \lor Q)$	$\neg Q \wedge \neg P$
T	T	T	F	F
T	F	T	F	F
F	T	T	F	F
F	F	F	T	T

Note that for all possible truth values of P and Q, $\neg(P \lor Q)$ equals $\neg P \land \neg Q$. Thus, the statement is true.

