

01204211 Discrete Mathematics

Lecture 1: Introduction

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The goals of this course

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- ▶ To learn how to make mathematical arguments.

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- ▶ To learn how to make mathematical arguments.
- ▶ To learn various fundamental mathematical concepts that are very useful in computer science.

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IMHO, mathematics is a mean to communicate *precise* ideas.

It's like learning a new language

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if a > b:  
    return a  
else:  
    return b
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The author claims that this program takes two variables a and b and returns the larger one.

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Finding the maximum value. (1)

Now look at this program.

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if a > b:
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    else:
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Finding the maximum value. (2)

Finally, look at this program.

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// Input: array A with n elements: A[0], ..., A[n-1]
m = 0
for i = 0, 1, ..., n-1:
    if A[i] > m:
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Can we try to test the code with all possible inputs?

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Let's try again.

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// Input: an integer n
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Note that if n is a prime number, the for-loop repeats for $n - 2$ times. Thus, the running time is approximately proportional to n . Can we do better?

Another example: testing primes (2)

Consider the following code.

```
// Input: an integer n
if n <= 1:
    return False
let s = square root of n
i = 2
while i <= s:
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Ok, it should be faster. **But is it correct?**

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- ▶ Now, if $2 \leq a \leq \sqrt{n}$, at some point during the execution of the algorithm, $i = a$ and i should divide n ; thus the algorithm correctly returns False.
- ▶ Are we done?

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- ▶ Note that if we can argue that $2 \leq b \leq \sqrt{n}$, we are done. (why?)
- ▶ How can we do that?

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- ▶ **Original goal:** To show that the faster algorithm is correct.
- ▶ **Current (sub) goal:** Consider a positive composite n and its positive divisor a , where $a > \sqrt{n}$. Let $b = n/a$. We want to show that $2 \leq b \leq \sqrt{n}$.