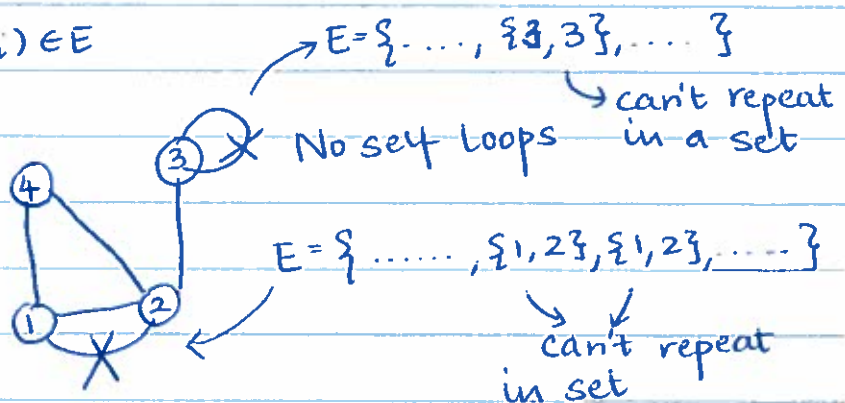
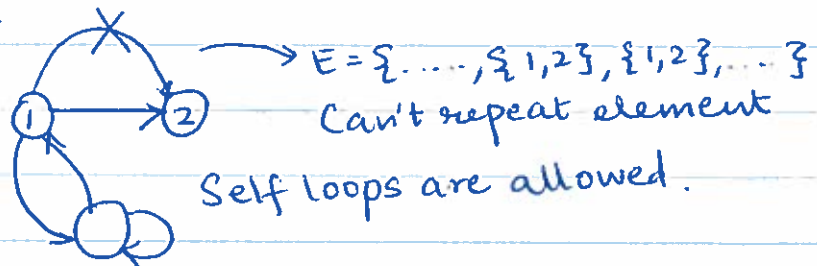


- Recall: Path from u to v is a sequence of vertices (v_0, v_1, \dots, v_k) such that $v_0 = u$, $v_k = v$ and $\forall i = 1, \dots, k$
- v is reachable from u if \exists a path from u to v
 - Note: Vertex is always ~~for~~ reachable from itself

→ Undirected $(v_{i-1}, v_i) \in E$

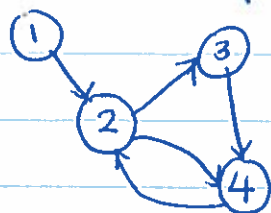


Directed $\{v_{i-1}, v_i\} \in E$



- Think in terms of how a set can be formulated.
- If you can't think of the choices for a graph, think about the underlying representation (sets)

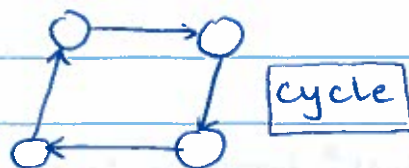
→ Path is simple if all of its vertices are different



$(1, 2, 4) \rightarrow$ SIMPLE!

$(1, 2, 3, 4, 2, 3, 4) \rightarrow$ NOT SIMPLE!

→ Directed Case :



Starts at a vertex and ends back there.

Smallest cycle →  self-loop

→ Cycle: A path forms a cycle if $v_0 = v_k$ and $k \geq 1$

at least 1 edge

and all vertices other than v_0, v_k are distinct

For a
directed
graph

→ simple cycle

→ What is a cycle in an undirected case?

- Same as that of directed? What is the problem?

• Smallest cycle is



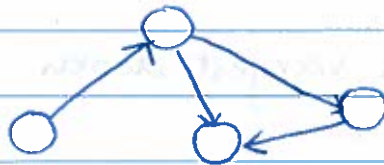
→ cycle for an undirected graph: Same but $k \geq 3$

at least 3² edges

→ Graph is acyclic if \nexists cycle.

⇒ SPECIAL TYPES OF GRAPHS:

→ DAG : Directed Acyclic Graph



Application:

- Course Pre-requisites

(A) → (B) A is required for B

If there is a cycle what does it mean?

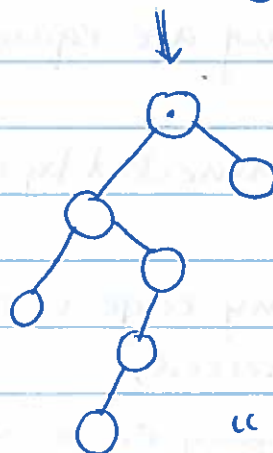
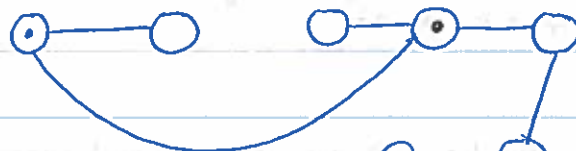
- can't resolve dependence

Another application is loading of packages.

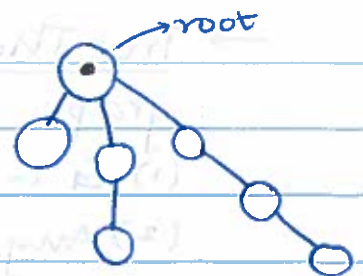
? Given a graph, how can you tell if it is acyclic or not? → will see this later on.

→ Tree : Undirected, acyclic, connected graph

every vertex is reachable from every other



Same graph rooted at different nodes



Outside info → which node is the root
"Tree rooted at"

• Other terms associated with a tree:

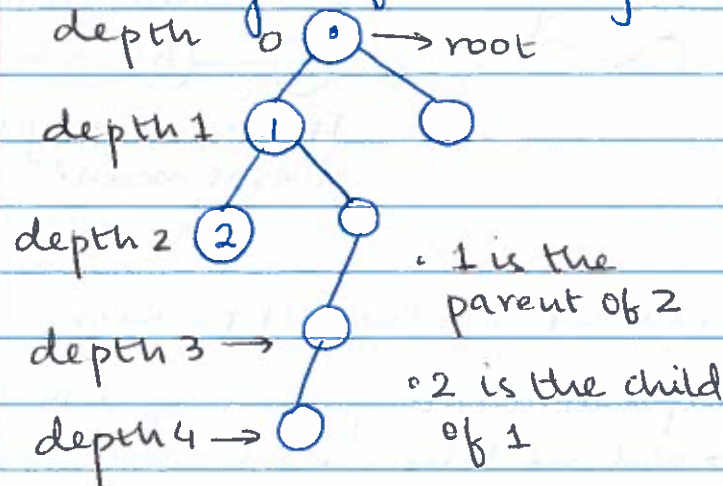
- root

- child/parent

- leaf: vertex without children

- depth of u :

- height of tree: largest depth



Application:

- Any hierarchical structure.
- File systems

(Q) which node is the root?

(A) It's upto you! or according to definition from outside \rightarrow "Tree rooted at"

\rightarrow Tree Theorem: Let $G=(V,E)$ be an undirected graph. Then the following are equivalent:

(1) G is a tree

(2) Any two vertices are connected by a unique simple path

(3) G is connected, but if any edge is removed, it's not. (minimally connected)

(4) G is ~~connected~~^{acyclic}, but if any edge is added, the graph will contain a cycle.

Notation for
size of

$|E| \rightarrow$ No. of edges
 $|V| \rightarrow$ No. of vertices
Size of set of vertices

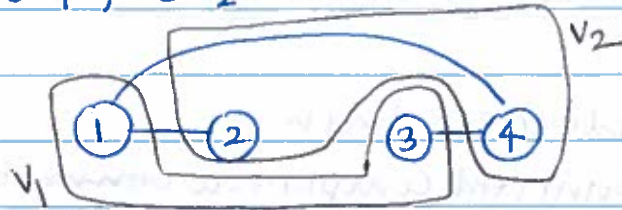
(5) G is connected and $|E| = |V| - 1$

(6) G is a cyclic and $|E| = |V| - 1$

→ Bipartite Graph: Undirected graph in which V can be partitioned into V_1 and V_2 st. every edge $\{u, v\}$ has

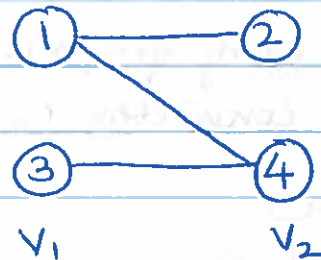
$$u \in V_1, v \in V_2$$

every vertex is in exactly one of V_1 or V_2



1 & 4, 2 & 3 → a partition? → No; edge between 1 & 4

Can be redrawn as

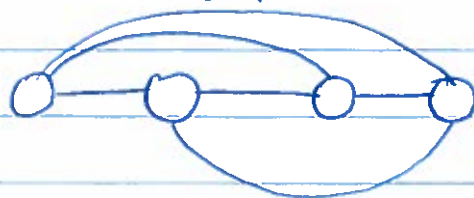


- No edge within the same partition.

- Application:

- Connections between routers and computers where only routers and computers can be connected
- Will be seen later in stable marriage problem.

→ Complete Graph: Undirected graph with an edge between every pair of nodes.



In a sense opposite to minimally connected tree as this has maximum number of edges

? How many edges are there with n vertices?

(A) $nC_2 = \frac{n(n-1)}{2}$ edges for n vertices.

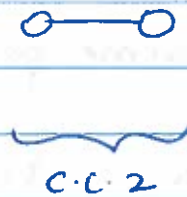
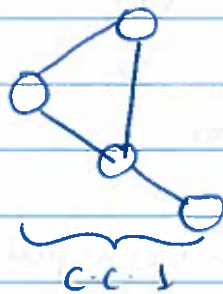
⇒ BASICS OF GRAPH CONNECTIVITY:

1. Undirected Graph:

Connected Component ~~connect~~ containing vertex u : set of all vertices reachable from u

◦ Note: Every graph can be partitioned into connected components

every vertex belongs to exactly one connected component



- Tree: exactly one connected component.