## THE GALE-SHAPLEY ALGORITHM

```
GALE-SHAPLEY()
    Initially all m \in M and w \in W are free
    while \exists m who is free and hasn't proposed to every w \in W
 3
         do Choose such a man m
            Let w be the highest ranked in m's preference list
              to whom m has not yet proposed
            if w is free
 6
               then (m, w) become engaged
               else w is currently engaged to m'
                    if w prefers m' to m
                      then m remains free
10
                      else w prefers m to m'
11
                            (m, w) become engaged
                           m' becomes free
12
13
    return the set S of engaged pairs
```

## UNDERSTANDING THE SOLUTION

For a given problem instance, there may be several stable matchings. Do all executions of Gale-Shapley yield the same stable matching? If so, which one?

An instance with two stable matchings:

A-X, B-Y, C-Z

A-Y, B-X, C-Z

|        | 1 <sup>s†</sup> | 2 <sup>nd</sup> | 3 <sup>rd</sup> |
|--------|-----------------|-----------------|-----------------|
| Xavier | Α               | В               | С               |
| Yancey | В               | Α               | С               |
| Zeus   | Α               | В               | С               |

|        | 1 <sup>st</sup> | 2 <sup>nd</sup> | 3 <sup>rd</sup> |
|--------|-----------------|-----------------|-----------------|
| Amy    | У               | X               | Z               |
| Bertha | X               | У               | Z               |
| Clare  | X               | У               | Z               |