

Exam Review

D-don't Act Dumb

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1 God Please Don't Forget These

- Define Algorithm
- Summations, Logs, Ceiling, Floor
- Counting Operations
- Complexities (O, Ω, Θ)
- Divide & Conquer
- Greedy Algorithms
- Dynamic Programming

2 Logs

When it comes to CS, log always has a base of 2 unless specified otherwise.
Let's take a look at $\Theta(n \log n)$.

2.1 Log Identities

$$\log(x^y) = y \log x$$

$$\log(xy) = \log x + \log y$$

$$\log_b a = \frac{\log_x a}{\log_x b}$$

2.2 $\log_{2742} n$ wtf why

$$\begin{aligned}\log_{2742} n &\leq c \log n \quad n \geq k \\ \frac{\log n}{\log 2742} &\leq c \log n \\ \frac{1}{\log 2742} \log n &\leq c \log n \quad c = \frac{1}{\log 2742}\end{aligned}$$

Doing it the limit way:

$$\begin{aligned}\lim_{n \rightarrow \infty} \frac{\log_{2742} n}{\log n} \\&= \lim_{n \rightarrow \infty} \frac{\log n}{\log 2742} \frac{1}{\log n} \\&= \lim_{n \rightarrow \infty} \frac{\log n}{\log 2742 * \log n} \\&= \frac{1}{\log 2742}\end{aligned}$$

3 Complexities

$$\begin{aligned}\Theta(n) + \Theta(1) &= \Theta(n+1) = \Theta(n) \\ \Theta(n) + \Theta(n) &= \Theta(n) \\ \Theta(n) * \Theta(n) * \Theta(n) &= \Theta(n^3)\end{aligned}$$

4 Dynamic Programming: Optimal Substructure

Optimal answers to smaller problems are still applicable. For example, in the coin changing case, it's still better to use the two 6 cent pieces and not the bigger 10 cent piece.