Real World Algorithms: A Beginners Guide Errata to the First Printing

Last updated 8 February 2018

This document lists the changes that should be made to *Real World Algorithms* to correct mistakes that made their way to printing, to improve infelicities that the author spotted too late, or update the material with something that the author did not know at the time of writing the book.

There are three different kinds of changes noted here. In all of them the date that they became known to the author is given at the first line of each item. The name of the person who suggested the change is also given at the end of each change.

Page 1, line 1 1 Ja	ın 1
These are technical or typographical errors.	
Page 1, line 1 1 Ja	ın 1
These as changes that improve the book, even if they do not correct an er. They include small rewordings, or material that became known to the autafter the book was published.	
Page 1, line 1 1 Ja These are minor fixes that although they do not make a big difference they do hurt the aut	
Some of them might strain the reader's eye to see where the improvement is exactly	

► Page xii, line 2	24 Apr 2017
they can proved	(S. Subramanya)
Page 5, line -17 and its last element is the $(n-1)$ th \rightsquigarrow and so its last element is the $(n-1)$ th	08 Feb 2018
Page 5, line – 10 Move "y;" to previous line.	08 Feb 2018
Page 8, line -8 and -2 big-Oh $\uparrow \downarrow \rightarrow$ big O	12 Aug 2017
Page 9, line 4 big-Ohs	12 Aug 2017
Page 9, line −11 In terms of big-Oh notation, we have by definition that $\ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \$	
Page 10, line –15 Move "of" to the next line.	08 Feb 2018
► Page 11, line -2	01 Apr 2017
$f(n) = e^x \rightsquigarrow f(n) = e^n$	(P. Tsanakas)
Page 13, line −11 big-Oh ∕\→ big O	12 Aug 2017
► Page 13, line -8	12 Aug 2017
This is called "big-Omega," or $\Omega(n)$, and the precise definition called "big Omega," $\Omega(f(n))$; the precise definition	on ∕्→ This is
Page 13, line −6 Having defined big-Oh and big-Omega 🆴 Having defined big O and big Ome	_
Page 13, line −5 big-Theta Ŋ→ big Theta	12 Aug 2017
► Page 20, line -4 line 3 \ → line 4	30 Mar 2017
Page 20, line −3 line 11 \rightarrow line 12	30 Mar 2017
Page 20, line −1 line 6 _→ line 7	30 Mar 2017
·	12 Aug 2017

► Page 41, lines -4 to -3	30 Jan 2018
Room 6 still has one unvisited room ✓→ Room 5 still h (Yi-Ming Lai)	as one unvisited room
▶ Page 57, line 4	24 Apr 2017
When you insert an item in the queue, you increase similarly, when you remove an item from the queue, of the tail. \(\shi \) When you insert an item in the queue, of the tail; similarly, when you remove an item from the index of the head.	you increase the index you increase the index
▶ Page 65, line 2	06 Mar 2017
011110 ♦ 011011	
► Page 71, algorithm 3.1, line 1	26 Mar 2017
Size ∕√→ SizePQ	
▶ Page 73, line −11	24 Apr 2017
root of the three	(S. Subramanya
▶ Page 80, line −6	25 May 2017
Joyces's ∕ _{√→} Joyce's	
▶ Page 80, line −5	29 Jun 2017
41% ∕ → 53%	
▶ Page 84, line 6	30 Jan 2018
by assigning it to wc in line 13 \rightsquigarrow by assigning to it w	c in line 13 (Yi-Ming Lai
Page 91, line –17	14 Dec 2017
"1110" \ → "1110"	
Page 95, figure 4.1, caption	21 Apr 2017
An encryption $\wedge \rightarrow$ A decryption	
▶ Page 140, lines −2 to −1	
SHA-2 (Secure Hash Standard-2) ✓ SHA-2 (Secure H	_
Page 144, line 2 command packet √→ command packet	21 Apr 2017
▶ Page 145, line −14	01 Jun 2017
$OR_3 \longrightarrow OR_2$	

► Page 145, line −12	01 Jun 2017
Alice $\bigwedge \rightarrow OR_1$	
Page 147, line –13	17 Jul 2017
SHA-224. \ → SHA-224,	
Page 157, figure 6.6, caption	21 Mar 2017
weighted $ \searrow $ weighted	
Page 162, line -1	30 Jan 2018
$prev$, that is, $prev[i] \longrightarrow pred$, that is, $pred[i]$	(Yi-Ming Lai)
Page 165, lines −2 to −1 move line break before "then"	01 Feb 2018
Page 166 , figure 6.13, second panel, label under t 13 $\$ \rightarrow 13/ $-\infty$	21 Apr 2017
Page 166 , figure 6.13, fourth panel, label under t 13 $\$ $\$ $\$ $\$ $\$ $\$ $\$ $\$ $\$ $\$	21 Apr 2017
Page 166, figure 6.13, fifth panel, label under t	21 Apr 2017
Page 170, figure 7.1, caption	30 Jan 2018
Breaking lines into paragraphs $\wedge \rightarrow$ Breaking paragraphs Lai)	into lines (Yi-Ming
Page 178, algorithm 7.1, line 12	23 Apr 2017
Page 179, line 10	24 Apr 2017
line 11 ∕√→ line 14	(S. Subramanya)
Page 179, line 12	24 Jul 2017
line 11 ∕√→ line 14	
Page 180, line 13	26 Mar 2017
lines $1-7 \longrightarrow lines 1-10$	
Page 181, line −4re-weighting	23 Jul 2017
• Page 182, figure 7.11	22 Jul 2017
$\lim_{N \to \infty} 0 \xrightarrow{0} 2 \xrightarrow{N} 0 \xrightarrow{0} 2 \text{ and } \lim_{N \to \infty} 0 \xrightarrow{0} 3 \xrightarrow{N} 0 \xrightarrow{7} 3$	

Page 182, figure 7.11, caption re-weighted ∕√→ reweighted	23 Jul 2017
▶ Page 184, line -12 , exercise 1 a better path goes through u , we can check whether $u $	
► Page 196, line 10	30 Jan 2018
We underline edges \→ We underline nodes	(Yi-Ming Lai)
Page 206, line 1 Euros ^→ euros	23 Apr 2017
► Page 214, line 8 $P_{B_i} \land \rightarrow B_{P_i}$	04 Apr 2017
▶ Page 217, line -3 page 3 \ → page 6	04 Apr 2017
▶ Page 217, line −2 page 4 √→ page 5	04 Apr 2017
► Page 219, line 10	30 Jan 2018
from node 4 to nodes 3 and 2 $\uparrow \rightarrow$ from node 4 to nodes 2 and 1	(Yi-Ming Lai)
Page 222, figure 9.6	28 Apr 2017
► Page 229, line −16 support √→ supported	_ 04 May 2017
► Page 230, line -3	23 Apr 2017
If there are <i>n</i> voters, then candidate <i>A</i> gets $(60 \times 2)n = 120n$ points are $100m$ voters, candidate <i>A</i> gets $(60 \times 2)m = 120m$ points	s √→ If there
► Page 230, line -2	23 Apr 2017
► Page 230, line -2	23 Apr 2017
► Page 231, heading 10.2 Shulze ∧→ Schulze	23 Apr 2017

▶ Page 233, algorithm 10.1, line 4	23 Apr 2017
$P[i][j] \searrow P[i,j]$	
▶ Page 234, line -8	04 May 2017
$P[i,j] \longrightarrow P[c_i,c_j]$	
▶ Page 234, line -7	04 May 2017
$P[j,i] \longrightarrow P[c_j,c_i]$	
▶ Page 234, line −6	04 May 2017
$P[i,j] - P[j,i] \longrightarrow P[c_i, c_j] - P[c_j, c_i]$	
Page 236, line -4	28 Apr 2017
$(k+1) \stackrel{\wedge}{\searrow} k+1$	
▶ Page 238, algorithm 10.2, line 6	23 Apr 2017
$S[i][j] \longrightarrow S[i,j]$	
▶ Page 238, algorithm 10.2, line 9	23 Apr 2017
$S[i][j] \longrightarrow S[i,j]$	
► Page 241, algorithm 10.3, second line of output	23 Apr 2017
$s[i,j_k] > s[j_k,i] \land S[i,j_k] > S[j_k,i]$	
▶ Page 242, line 6	30 Jan 2018
D would beat B, C, and D, while A would beat C, B would beat	$D \longrightarrow D$ would
beat both <i>B</i> and <i>C</i> , while <i>A</i> would beat <i>C</i> , <i>B</i> would beat <i>C</i>	(Yi-Ming Lai)
	23 Apr 2017
all $pred$ and $dist \land \rightarrow pred$ and $dist$	
▶ Page 249, algorithm 11.1	24 Apr 2017
a array of items $\wedge \rightarrow$ an array of items	(S. Subramanya)
▶ Page 249, algorithm 11.1	24 Apr 2017
a element we are searching for	ing for (S. Subra-
Page 249, figure 11.1	28 Apr 2017
Change the array to: 114 480 149 903 777 65 680 437 4 181 613 551 10 3	31 782 507
We need not use sequential search in a sorted array.	

► Page 250, line -3	30 Jan 2018
real and complex parts $ \searrow$ real and imaginary parts	(Yi-Ming Lai)
► Page 254, line -5	24 Apr 2017
figure 11.3 ∕√→ figure 11.6	
► Page 259, line -8	30 Jan 2018
whether the match is in the head of the list \rightsquigarrow whether the m the head of the list	atch is not in (Yi-Ming Lai)
► Page 260, algorithm 11.2	24 Apr 2017
a element we are searching for	g for (S. Subra-
► Page 260, algorithm 11.2, line 10	24 Apr 2017
$NULL; \longrightarrow NULL$	
► Page 261, algorithm 11.3	28 Jul 2017
${\sf TranspositionSearch}(A,s) \not \searrow {\sf TranspositionSearch}(L,s)$	
Page 261, algorithm 11.3 a list of items	24 Apr 2017
► Page 261, algorithm 11.3	24 Apr 2017
a element we are searching for	g for (S. Subra-
► Page 261, algorithm 11.3, line 12	25 Apr 2017
$NULL; \longrightarrow NULL$	
► Page 262, algorithm 11.4	24 Apr 2017
a array of items	(S. Subramanya)
► Page 262, algorithm 11.4	24 Apr 2017
a element we are searching for	g for (S. Subra-
► Page 262, line 1	30 Jan 2018
the same search as in figure 11.11 $\uparrow \searrow \rightarrow$ the same search as in figure Ming Lai)	ıre 11.10 (Yi-
► Page 264, algorirthm 11.5	25 Apr 2017
SecretarySearch(A , s) \fiver SecretarySearch(A)	

► Page 264, algorithm 11.5	24 Apr 2017
a array of items ∕√→ an array of items	(S. Subramanya)
► Page 264, algorirthm 11.5, line 4	24 Apr 2017
$Compare(A[i],A[b]) \ \\ \\ \\ \\ \\ \\ \\ \\ \\ \\ \\ \\ \\ \\ \\ \\ \\ \\ $	(S. Subramanya)
► Page 264, algorirthm 11.5, line 6	25 Apr 2017
$i \leftarrow m + 1 \nearrow \rightarrow i \leftarrow m$	
► Page 267, line 18	6 May 2017
Unless you are not psychic	
► Page 268, algorithm 11.6	24 Apr 2017
a element we are searching for	ng for (S. Subra-
► Page 270, figure 11.14b, last row	31 May 2017
$ \begin{array}{ccc} l = 7 \\ m = 7 \end{array} \qquad \begin{array}{c} l = 8 \\ m = 8 \end{array} $	(I. Kafetzaki)
► Page 275, line -2	02 May 2017
one's complement	
► Page 278, algorithm 11.7	24 Apr 2017
a element we are searching for $ \searrow $ an element we are searching manya)	ng for (S. Subra-
► Page 287, algorithm 12.1	24 Apr 2017
a array of items $\wedge \rightarrow$ an array of items	(S. Subramanya)
► Page 289, algorithm 12.2	24 Apr 2017
a array of items ∕√→ an array of items	(S. Subramanya)
► Page 291, algorithm 12.3	24 Apr 2017
a array of items ∕√→ an array of items	(S. Subramanya)
► Page 297, line -5	30 Jan 2018
we want to have $A[i] \ge A[i] \longrightarrow$ we want to have $A[0] \ge A[i]$	(Yi-Ming Lai)
► Page 298, caption of figure 12.6b	28 Apr 2017
$1 \longrightarrow one$	

Page 299, algorithm 12.4	24 Apr 2017
a array of items $\wedge \rightarrow$ an array of items	(S. Subramanya)
Page 310, figure 12.12, third panel	08 May 2017
$i \to 5 \nearrow i \to 37$	
Page 327, line −16, exercise 2 characters like " ", "_", and "+" _→ characters like " ", "-", and "+"	
Page 327, line −15, exercise 3	20 Dec 2017
The in-place array merge, algorithm 12.7 \fine The inrithm 12.7,	place array merge, algo-
Page 333, line −11 minimal perfect mapping \(\rightarrow \) minimal perfect mapping	09 May 2017
Page 340, line −3 456, 976 \square 456, 976	09 May 2017
Page 343, figure 13.5 4, 847 ∖¬→ 4,847	09 May 2017
Page 343, figure 13.5	09 May 2017
Page 343, figure 13.5	09 May 2017
Page 343, line 8	30 Jan 2018
in line $4 \rightsquigarrow$ in line 3	(Yi-Ming Lai)
Page 346, line 3 binary fractional number	09 May 2017
Page 353, line −12	23 Jul 2017
An successful search cannot take longer than a succe ful search cannot take longer than an unsuccessful o	·
Page 359, line −9	13 May 2017
Page 359, line −9	13 May 2017
=	31 May 2017
the number of frequency peaks in the song, and there is even a number of frequency peaks in the song, and there is even a notati	on for it:
Page 361, line 16	31 May 2017

► Page 362, line -1	31 May 2017
the data are not the	
Page 367, line 7	13 May 2017
$(1-1/m)^{m(\frac{k}{m})} \rightsquigarrow (1-1/m)^{m(\frac{k}{m})}$	
► Page 370, figure 13.20, third panel	13 May 2017
The solid arrows should emanate from "this".	
► Page 371, line 2	30 Jan 2018
Our hash algorithms take a specific and produce a spec hash algorithms take a specific input and produce a spec Lai)	
Page 383, table 14.1, caption	14 May 2017
letter	14 May 2017
► Page 386, line 9, 12, 19	25 May 2017
Gibb's ∕ → Gibbs's	
Page 387, line −14 "ineligible" ∕\→ "ineligible."	16 May 2017
► Page 390, line 3	16 May 2017
six ∕√→ five	
► Page 395, line -15	30 Jan 2018
we get the values shown in figure 14.7 \rightsquigarrow we get the ure 14.8 Lai)	values shown in fig- (Yi-Ming
► Page 396, figure 14.8, fourth panel	17 May 2017
$H = 0.40 \checkmark H = 0.940$	
▶ Page 397, line −9	16 May 2017
tox ∕√→ to	
► Page 400, figure 14.10	08 Jun 2017
$\{1, 2,, 14\}$: outlook $\land \rightarrow \{1, 2,, 15\}$: outlook	(V. Malandrakis)
► Page 400, line 5	30 Jan 2018
happens in the normal branch $\wedge \rightarrow$ happens in the high	branch (Yi-Ming Lai)

$r \leftarrow \text{CreateMap}() \land \rightarrow dt \leftarrow \text{CreateMap}()$ (Y	
$T \leftarrow Creaternap() \cdot \sqrt{\pi} \ ur$ $\leftarrow Creaternap()$	i-Ming Lai)
Page 413, figure 14.12 22 add label "high" on the first, left, edge emanating from the root node	2 Dec 2017
Page 414, line 3 12 because in terms of the big-Oh notation it is $\$ because in terms of the big O notation	_
Page 417, line -3	6 Feb 2017
Witten, Frank, and Hall $\ \ \ \ \ \ \ \ \ \ \ \ \ $	
Page 426, figure 15.10 Change the gray letters from 40% gray to gray.	3 Feb 2018
Page 427, graphics	3 Feb 2018
Page 428, second and fourth graphics0 Change the gray letters from 40% gray to gray.	3 Feb 2018
Page 430, line -17	May 2017
Page 430, line -16	May 2017
Page 430, line -41	4 Sep 2017
all A, AB, and ABA are ✓→ substrings A and ABA are	P. Mpellos)
Page 431, fourth graphic	
\	
	<i>7</i> 72
Page 431, line −1023 of the pattern → of the matched pattern	May 2017
Page 431, fifth graphic23	
\$33333333 \$33333333 \$33333333 \$33333333	<u>2/2</u> 1
\$333333333 \$3333333 \$444444444444444	<i>77</i> 2

Page 431, line –9	_ 24 May 2017
So we get:	
Page 431, line –1	_ 24 May 2017
$longer\ shifts\ \searrow \ longer\ shifts$	
Page 432, second graphic	
^ →	
· · · · · · · · · · · · · · · · · · ·	//////X
	<i></i>
Page 432, line 7	_ 24 May 2017
AABAABAA \ → AABAABAAAA	
Page 432, third graphic	_ 24 May 2017
A A B A A B A A B A A A	
A A B A A A	
Page 432, fifth graphic	03 Feb 2018
Change the gray letters from 40% gray to gray.	
Page 432, line -4	_ 24 May 2017
define its length to be zero $ \searrow $ define its border length as zero	
Page 433, line 13	_ 25 May 2017
borders array ↑ border array	
Page 434, algorithm 15.2, line 9	02 Jun 2017
$p[i] \longrightarrow p[j]$	(A. Tsalapatis)
Page 434, line 4	22 Dec 2017
to a queue $q ightharpoonup to the queue q$	
Page 435, figure 15.5 caption	_ 24 May 2017
Another trace the Knuth-Morris-Pratt algorithm; the borders ar	ray is at the
bottom.	; the border
Page 437, line 3	_ 25 May 2017
borders array ↑ border array	
8 7 7 8	03 Feb 2018
Change the gray letters from 40% gray to gray.	

>	Page 440, line 12	. 30 May 2017
	mattern ∕→ pattern	
	Page 441, figure 15.9b	_ 02 Feb 2018
	E M B E R	
	r = 1	
	· · · · · · · · · · · · · · · · · · ·	
	r = 1	
•	Page 443, algorithm 15.4	_ 23 Dec 2017
	$\texttt{CreateRtOccurrencesTable}(p,t,s) \ \nwarrow \\ \texttt{TreateRtOccurrencesTable}(p,t,s) \ \\ \texttt{CreateRtOccurrencesTable}(p,t,s) \ \\ \texttt{CreateRtOccurrencesTable}(p,t,$	able(p, s)
•	Page 448, line 7	_ 23 Dec 2017
	Try using a different data structure, like a hash table or a set, inste then using a different data structure, like a hash table, instead.	
	Page 449, line 16 50-50 √→ 50-50	. 23 May 2017
•	Page 462, line 10	. 20 May 2017
	line 6 $\uparrow \rightarrow$ line 7	
•	Page 463, line 4	. 20 May 2017
	change ∕√→ maybe fix	
>	Page 466, lines 18, 21, 23	. 20 May 2017
	ECC √→ EEC	
•	Page 466, line -17	_ 30 Jan 2018
	Counting of Ministers ✓→ Council of Ministers	(Yi-Ming Lai)
•	Page 467, lines 12, 19, 23	. 20 May 2017
	ECC √→ EEC	·
>	Page 467, paragraph –2	. 22 May 2017
	Rewrite the paragraph as follows: To tackle this kind of question, we must adopt a systematic We have a set of voters, $V = \{v_1, v_2, \dots, v_n\}$, and a set of we $\{w_1, w_2, \dots, w_m\}$. A voter v_i has a weight w_j given by a mapping For a decision to be taken, it needs to meet a <i>quota Q</i> . In the exa EEC, we have $Q = 12$. The setup of V, W, f , and Q is called a <i>voti</i>	eights, $W = f: V \to W$. mple of the

▶ Page 468, line 3 such as ^→ such that	_ 21 May 2017
► Page 468, line 4in obtaining a losing coalition	_ 21 May 2017
► Page 468, line 14 ECC \rightarrow EEC	_ 21 May 2017
► Page 468, line -7 then then \rightarrow then the	_ 21 May 2017
▶ Page 468, lines -3 to -1 As an example, take four voters $V = \{A, B, C, D\}$ with correspond $W = \{4, 2, 1, 3\}$ and quota $Q = 6$. The critical coalitions are (we uncritical voters) $\{\underline{A}, \underline{B}\}, \{\underline{A}, \underline{D}\}, \{\underline{A}, \underline{B}, C\}, \{\underline{A}, B, D\}, \{\underline{A}, C, \underline{D}\}, \{\underline{B}, C\}, \{\underline{A}, B, D\}, \{\underline{A}, C, \underline{D}\}, \{\underline{B}, C\}, \{\underline{A}, B, D\}, \{\underline{A}, C, \underline{D}\}, \{\underline{A}, C, $	ling weights nderline the
As an example, let us take four voters A , B , C , D with corresponded equal to 4, 2, 1, 3, and quota $Q = 6$. The critical coalitions then are, the critical voters: $\{\underline{A},\underline{B}\}$, $\{\underline{A},\underline{D}\}$, $\{\underline{A},\underline{B},C\}$, $\{\underline{A},B,D\}$, $\{\underline{A},C,\underline{D}\}$, and	ling weights underlining
▶ Page 469, lines 6–7	
► Page 472, line −1 zero \rightarrow one	
▶ Page 473, line 1 one ∕√→ zero	05 Sep 2017 (N. Batsal)

Page 476, table 16.3 _______ 05 Feb 2017

Table 16.3 was built with data from 2008. To update it for 2016, it should be as follows:

Table 16.3 2016 U.S. electoral college number of electors and Banzhaf measure.

CA	55	0.471	MN	10	0.076	NM	5	0.038
TX	38	0.298	MO	10	0.075	WV	5	0.038
FL	29	0.223	WI	10	0.076	HI	4	0.03
NY	29	0.224	AL	9	0.068	ID	4	0.03
IL	20	0.153	CO	9	0.068	ME	4	0.03
PA	20	0.153	SC	9	0.068	NH	4	0.03
OH	18	0.136	KY	8	0.06	RI	4	0.03
GA	16	0.121	LA	8	0.061	AK	3	0.023
MI	16	0.121	CT	7	0.053	DC	3	0.023
NC	15	0.114	OK	7	0.052	DE	3	0.023
NJ	14	0.106	OR	7	0.053	MT	3	0.023
VA	13	0.098	AR	6	0.045	ND	3	0.023
WA	12	0.091	IA	6	0.045	SD	3	0.023
ΑZ	11	0.083	KS	6	0.045	VT	3	0.023
IN	11	0.083	MS	6	0.045	WY	3	0.023
MA	11	0.083	NV	6	0.045			
TN	11	0.083	UT	6	0.046			
MD	10	0.076	NE	5	0.038			

	Page 476, lines -6 to -5	_ 05 Feb 2017
1	Page 476, lines -3 to -2 California's Banzhaf measure is about 20.65 times that of Vermont. $\ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \$	
	Page 479, line −4 primes ♦ composites	. 21 May 2017
	Page 479, lines -4 to -3	. 21 May 2017
	Page 479, line -3 (1/2 + 1/3 + 1/5 \cdots + 1/k) \sqrt{1/2 + 1/3 + 1/5 + \cdots + 1/k)}	. 21 May 2017
	Page 485, algorithm 16.11Output: (r, q) , such that $n = 2^r q \land \!$	•

odd

Page 498, reference 219	26 Mar 2017
Ian H. Witten, Eibe Frank, and Mark A. Hall. Data Mining Learning Tools and Techniques. Morgan Kaufmann Publis cisco, CA, 3rd edition, 2011.	hers Inc., San Fran-
Ian H. Witten, Eibe Frank, Mark A. Hall, and Christopher <i>Practical Machine Learning Tools and Techniques</i> . Elsevie 4th edition, 2016.	_
► Page 502, first column	12 2017
big-Oh $(O(f(n)) \rightsquigarrow \text{big O}(O(f(n)))$ big-Omega $(\Omega(f(n))) \rightsquigarrow \text{big Omega}(\Omega(f(n)))$ add big Theta $(\Theta(f(n)))$, 13	
Page 502, first columnadded binary fractional number	09 May 2017
► Page 503, second column	20 May 2017
European Economic Community (ECC) ✓→ European Eco (EEC)	onomic Community
Page 504, first column graph re-weighting ∕√→ graph reweighting	23 Jul 2017
Page 504, first columnremove length (move to path, length)	03 Feb 2018
▶ Page 505, first column	30 Jan 2018
Lember-Ziv-Welch $\bigwedge \rightarrow$ Lempel-Ziv-Welch	(Yi-Ming Lai
Page 505, second columnadded mapping, minimal perfect	09 May 2017
Page 506, first columnadd path, length	03 Feb 2018