Behaviour-Oriented Concurrency

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Jan 5, 2025

https://github.com/Foreverhighness/boc-talk

Links

Behaviour-Oriented Concurrency Paper

Basic C# implementation

Core C++ implementation

Presentation video

Supplementary

KAIST CS431: Concurrent Programming

Concurrency

Parallelism

- Thread
- Task
- Coroutines
- Async
- ...

Coordination

- Promises
- Locks
- Condition variables
- Transactions
- ...

Concurrency

Parallelism

- Thread
- Task
- Coroutines
- Async

 $\operatorname{spawn}(A \xrightarrow{\operatorname{transfer } 100} B) \qquad \operatorname{spawn}(B \xrightarrow{\operatorname{transfer } 100} C)$

$$\operatorname{spawn}(C \xrightarrow{\operatorname{transfer } 100} D) \qquad \operatorname{spawn}(D \xrightarrow{\operatorname{transfer } 100} A)$$

Coordination

- Promises
- Locks
- Condition variables
- Transactions

$$\operatorname{spawn}(B \xrightarrow{\operatorname{transfer } 100} C)$$

$$\operatorname{spawn}(D \xrightarrow{\operatorname{transfer } 100} A)$$

- Isolation
- Parallelism
- Deadlock Freedom
- Causal Ordering

$$\operatorname{spawn}(A \xrightarrow{\operatorname{transfer 100}} B) \qquad \operatorname{spawn}(B \xrightarrow{\operatorname{transfer 100}} C)$$

$$\operatorname{spawn}(C \xrightarrow{\operatorname{transfer 100}} D) \qquad \operatorname{spawn}(D \xrightarrow{\operatorname{transfer 100}} A)$$

- Isolation → exclusive access (Mutex Transaction)
- Parallelism
- Deadlock Freedom
- Causal Ordering

$$spawn(\&mut A \xrightarrow{transfer 100} \&mut B) \qquad spawn(\&mut B \xrightarrow{transfer 100} \&mut C)$$

$$spawn(\&mut C \xrightarrow{transfer 100} \&mut D) \qquad spawn(\&mut D \xrightarrow{transfer 100} \&mut A)$$

- Isolation → exclusive access (Mutex Transaction)
- Parallelism
- Deadlock Freedom
- Causal Ordering

$$spawn(\&mut A \xrightarrow{transfer 100} \&mut B) \qquad spawn(\&mut B \xrightarrow{transfer 100} \&mut C)$$

- Isolation \rightarrow exclusive access (Mutex Transaction)
- Parallelism
- Deadlock Freedom
- Causal Ordering

spawn(&mut A
$$\xrightarrow{\text{transfer 100}}$$
 &mut B)

$$spawn(\&mut C \xrightarrow{transfer 100} \&mut D)$$

- Isolation → exclusive access (Mutex Transaction)
- Parallelism
- Deadlock Freedom → Deadlock avoidance (Sort)
- Causal Ordering

$$spawn(\&mut A \xrightarrow{transfer 100} \&mut B) \qquad spawn(\&mut B \xrightarrow{transfer 100} \&mut C)$$

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- Isolation → exclusive access (Mutex Transaction)
- Parallelism
- Deadlock Freedom → Deadlock avoidance (Sort)
- Ordering → DAG (Dependency Graph)

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- Deadlock Freedom → Deadlock avoidance (Sort)
- Causal Ordering → DAG (Dependency Graph)

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- Cown: protects a piece of separated data \rightarrow Mutex
- Behaviour: unit of concurrent execution \rightarrow Thread
- When: spawns a behaviour with a set of required cowns \rightarrow Spawn

when(Cown, Cown**; &mut A
\$\$\xrightarrow{\text{transfer 100}}\$\$
 &mut B\) when\(Cown**, Cown; &mut B \$\xrightarrow{\text{transfer 100}}\$ &mut C\) when\(Cown, Cown; &mut C \$\xrightarrow{\text{transfer 100}}\$ &mut D\) when\(Cown, Cown; &mut D \\$\xrightarrow{\text{transfer 100}}\\$ &mut A\\)****

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Abstraction

Example:

```
1 when (c1) { /* b1 */ }
2 when (c3) { /* b2 */ }
3 when (c1, c2) { /* b3 */ }
4 when (c1) { /* b4 */ }
5 when (c2, c3) { /* b5 */ }
6 when (c3) { /* b6 */ }
```

c3

```
when (c1) { /* b1 */ }
when (c3) { /* b2 */ }
when (c1, c2) { /* b3 */ }
when (c1) { /* b4 */ }
when (c2, c3) { /* b5 */ }
when (c3) { /* b6 */ }
```

c2

c1

```
when (c1) { /* b1 */ }
```

```
when (c3) { /* b2 */ }
```

```
when (c1,c2)
{ /* b3 */ }
```

when (c2,c3) { /* b5 */ }

when (c3) { /* *b6* */ }

Request

when (c1)

{ /* b1 */ }

c1

when (c3) { /* b2 */ }

when (c1,c2) { /* b3 */ } when (c1) { /* *b4* */ } when (c2,c3) { /* b5 */ }

when (c3) { /* b6 */ }

Request

c1

```
when (c1) { /* b1 */ } when (c3) { /* b2 */ }
```

```
when (c1,c2)
{ /* b3 */ }
```

when (c2,c3) { /* b5 */ }

when (c3) { /* b6 */ }

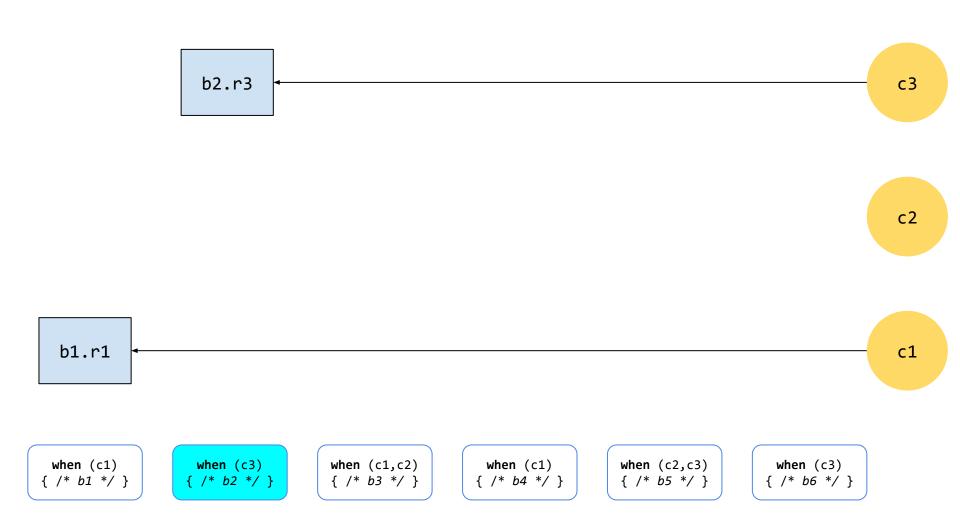
when (c1) { /* b1 */ }

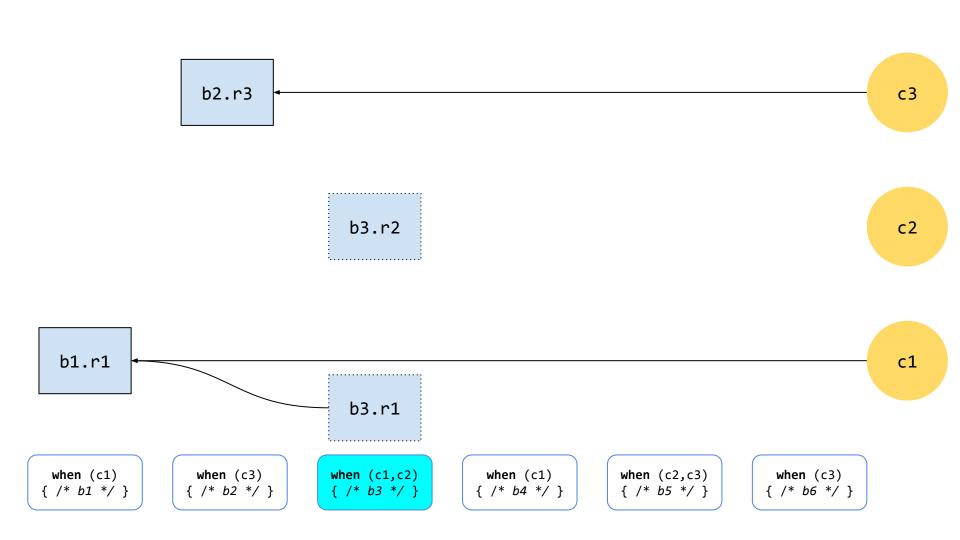
when (c3) { /* b2 */ }

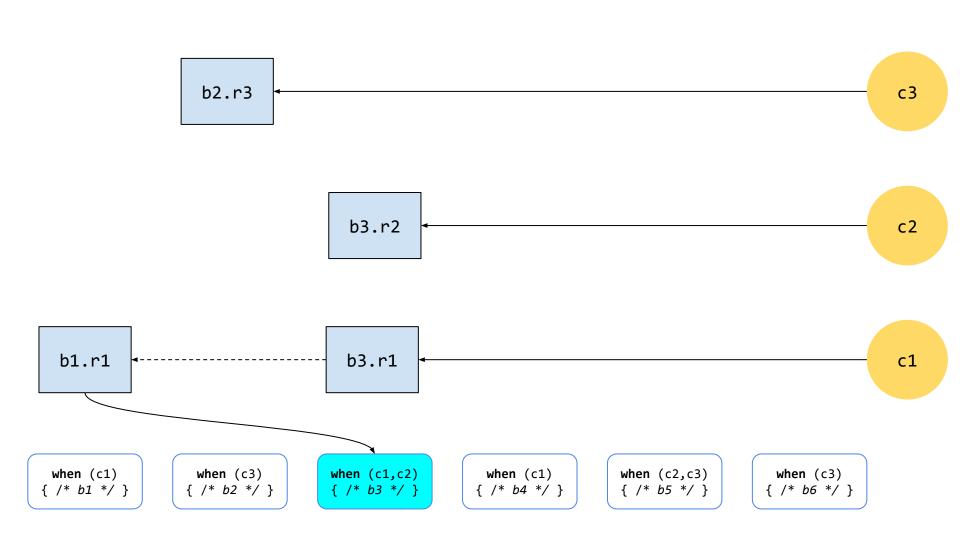
when (c1,c2) { /* b3 */ }

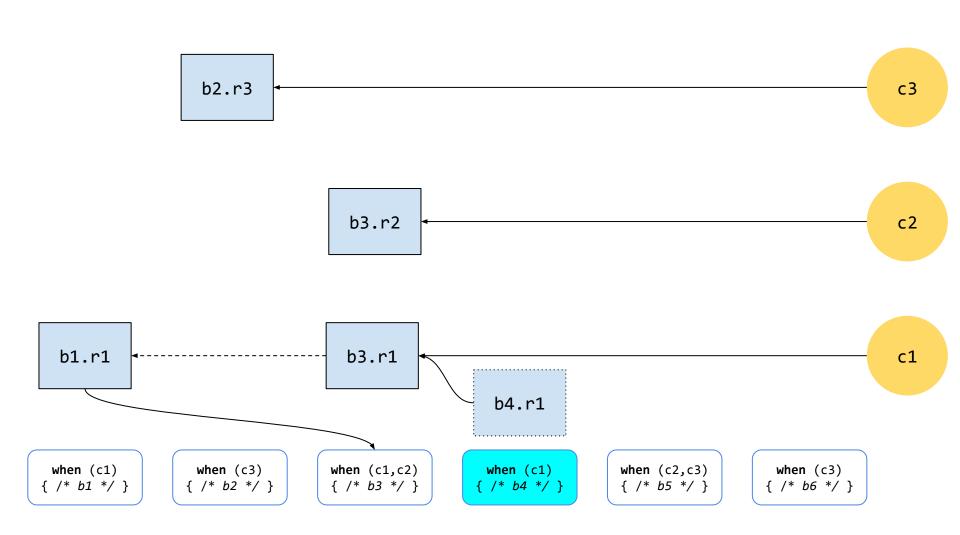
when (c1) { /* b4 */ } when (c2,c3) { /* *b5* */ }

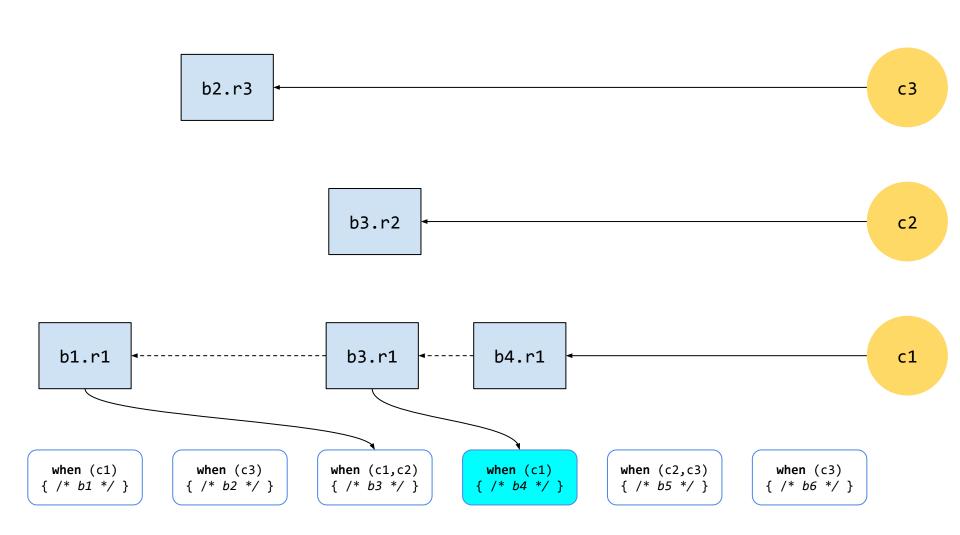
when (c3) { /* b6 */ }

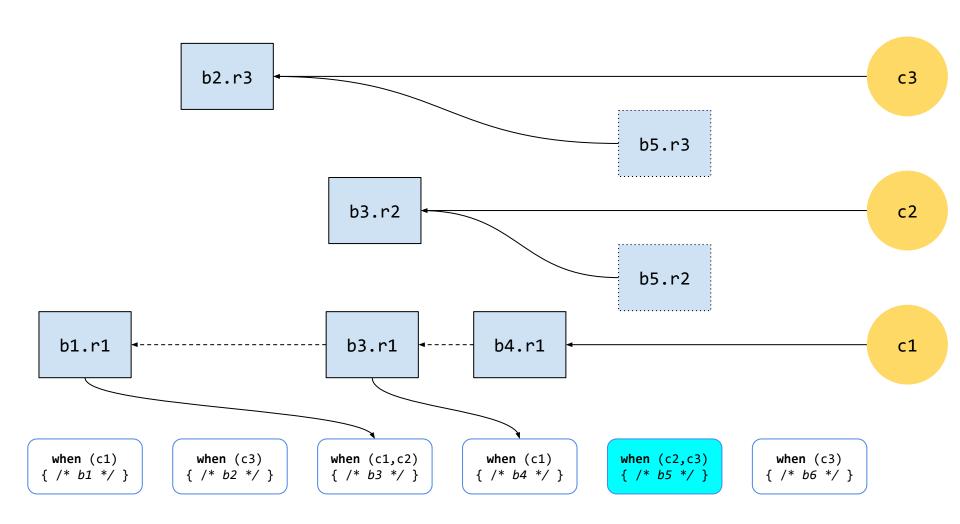


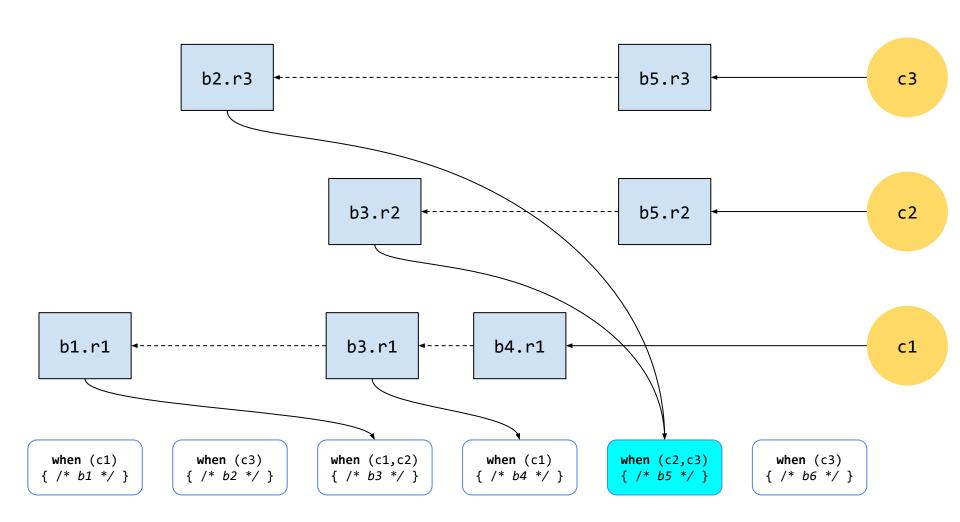


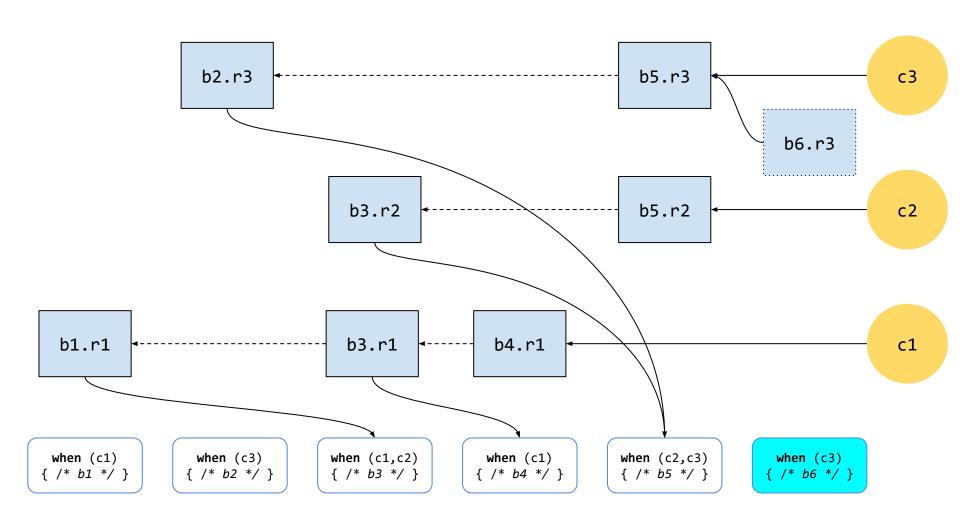


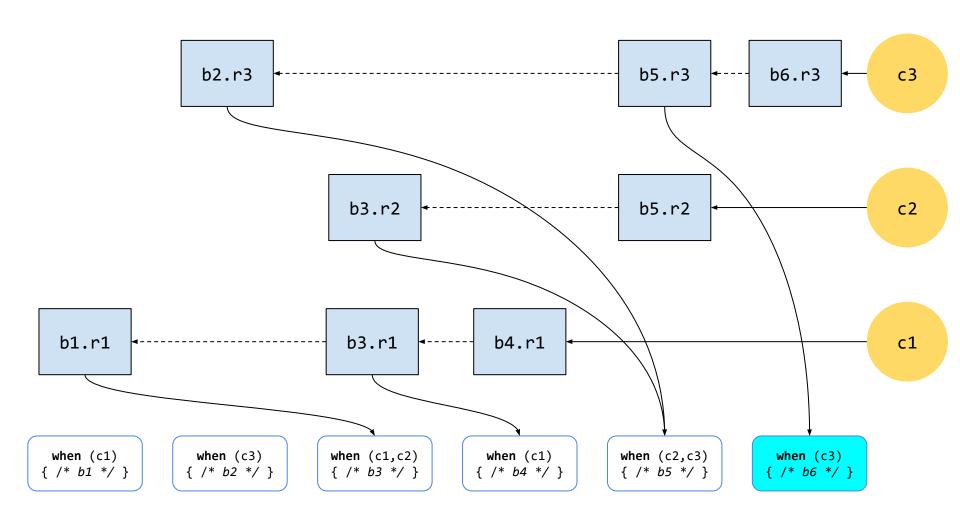


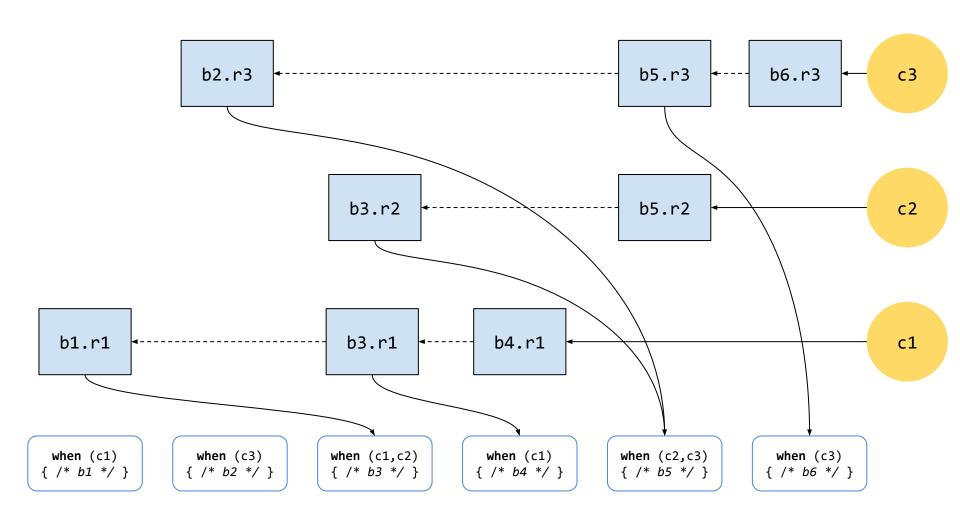












Implementation with lock

- Additional count
- Scheduled flag

b1.r2

b2.r2

c1

b1.r1

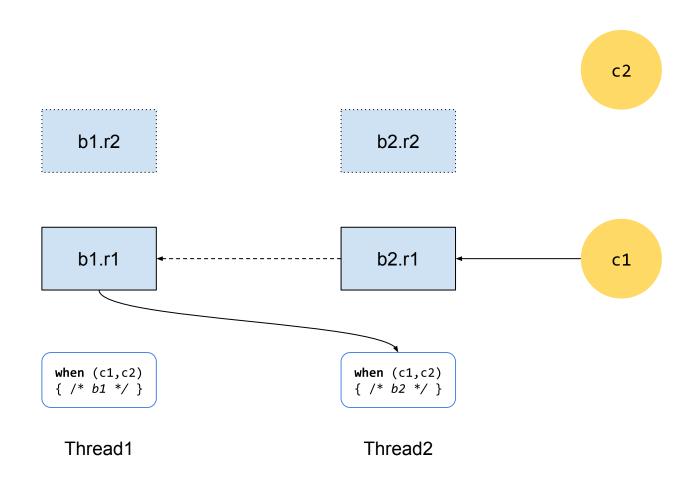
when (c1,c2) { /* b1 */ }

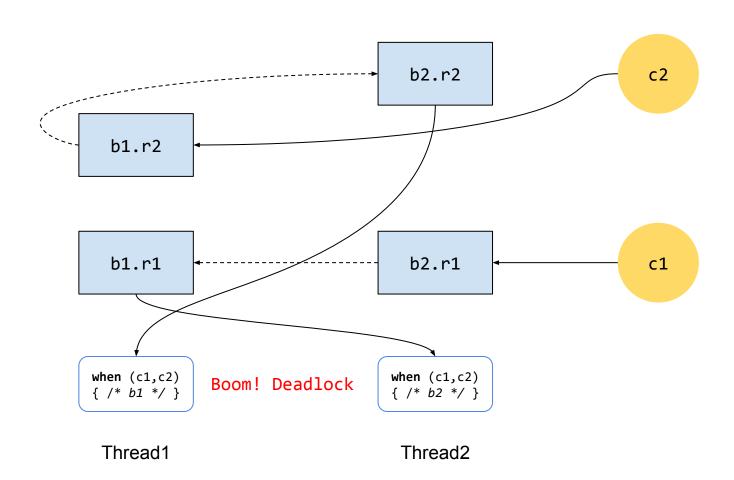
Thread1

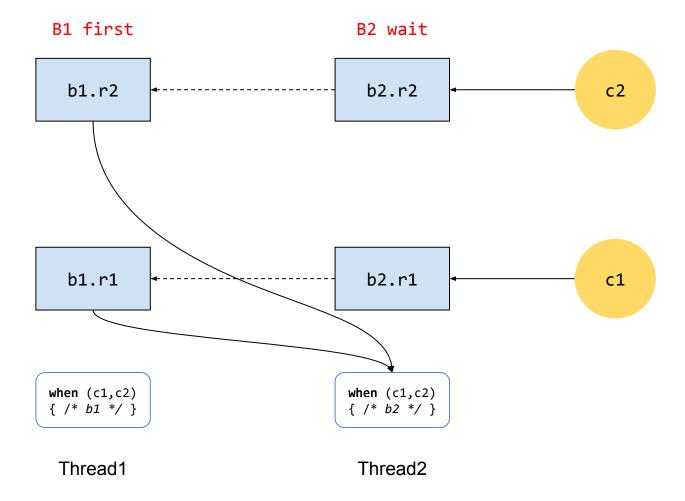
b2.r1

when (c1,c2) { /* b2 */ }

Thread2







Implementation without lock

- Behaviour, Request, and Cown all on heap
- Pin semantics

Related topic

- Actor
- Transaction
- Distribute Programming

Thanks for watching!