



THE GRAINGER COLLEGE OF ENGINEERING  
SIEBEL SCHOOL OF COMPUTING AND DATA SCIENCE

# CS 521

## Technological Foundations of Blockchain and Cryptocurrency

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Topic 5 – Consensus

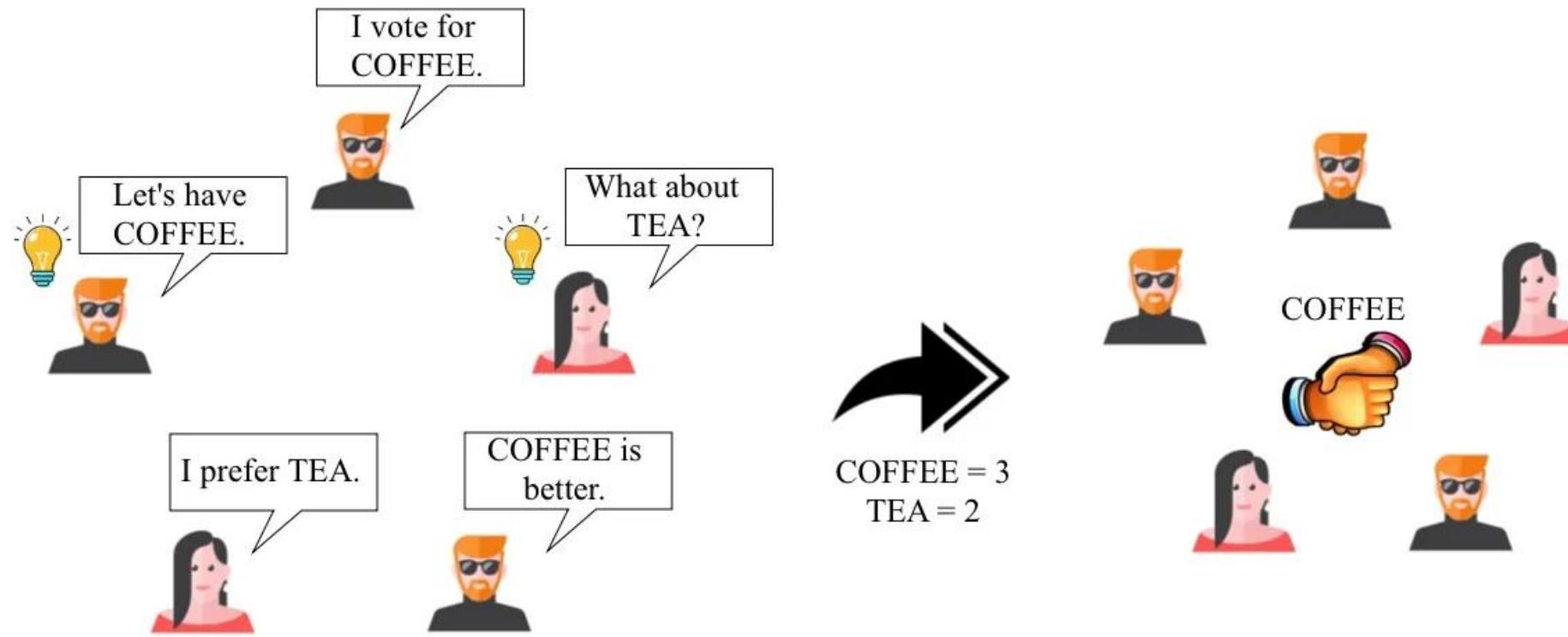
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# Thanks

- Musab Al Turki

# The Consensus Problem

- Distributed processes agreeing on a single, correct value, despite failures
- Much harder when decentralized



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## Key Properties

- **Agreement:** All non-faulty nodes must agree on the same value
- **Validity:** The agreed value must be a correct one (proposed by a non-faulty node)
- **Termination:** All non-faulty nodes must eventually decide on a value

## Main Challenges

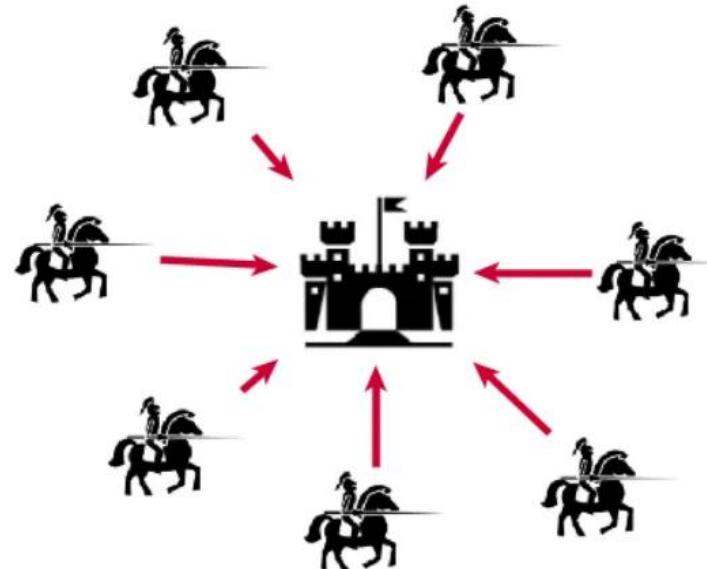
- **Byzantine Faults :** Some nodes may act arbitrarily or maliciously or may crash
- **Asynchrony:** Nodes operate at different speeds with unpredictable message delivery delays

# The Byzantine Generals Problem

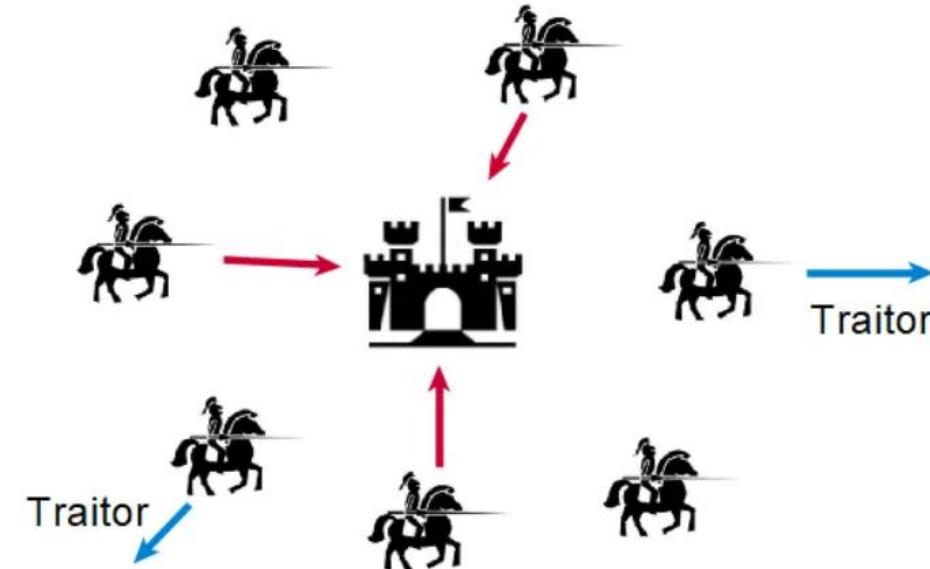
1982: Lamport, Shostak, and Pease

## Formalization of BFT Consensus

[source](#)



Coordinated attack  
leading to victory



Uncoordinated attack  
leading to defeat

Can only tolerate up to  $f$  faulty generals if there are  $3f + 1$  generals

# The FLP Impossibility Result

1985: Fischer, Lynch, and Paterson

No **deterministic** consensus algorithm can simultaneously guarantee **safety** and **liveness** (correctness) in an **asynchronous** distributed system where even one process may **crash**

Real-world protocols must introduce additional assumptions, such as:

- Crash failure recovery mechanisms (e.g., Paxos, Raft)
- Synchrony or partial synchrony (e.g., PBFT, Tendermint, HotStuff).
- Randomization (e.g., Bitcoin, Algorand, Bullshark)

Or they may sacrifice safety for liveness or vice versa (e.g. Paxos, PBFT, Bitcoin, Ethereum)

# Blockchain Consensus

A state-machine-replication problem: Agreement on a sequence of states

- Clients submit transactions to nodes
- Each node locally maintains an ordered sequence of txs (in blocks)
- Nodes need to agree on a canonical, totally ordered sequence of txs (defining the canonical sequence of states)
  - Assume an initial state (the genesis state)
  - In each round, the proposed state is valid if the ordered list of transactions in the proposed block applied to the last consensus state is valid and yields proposed state
- A blockchain protocol guarantees a *total order* on transactions
  - Original motivation in Bitcoin: Simulate centralized ledgers
  - Long believed to be a practical necessity for real-world decentralized applications
  - Enforced globally on all applications built on top