

# Assignment 11

516030910259 **Xinpeng Liu**

## 11.1

- a. It's easy for the implementation of the allocation, but it may result in more internal fragmentation,
- b. It's complicated for allocation, and there might be more external fragmentation.
- c. It achieves intermediate complexity and flexibility compared with the former two schemes.

## 11.3 .

- a. Yes. We can search the entire directory structure, find the empty spaces and link them into a new free-space list.
- b. 4.
- c. Store the pointer on the disk.

## 11.6

- a. Assume  $X$  is the start address of the file,  $L$  is the logical address,  $Y$  is the physical block number translated from logical address, and  $Z$  is the block offset.
  - (a) Contiguous.  $Y=L/512$ ,  $Z=L\%512$ . Load the block with address  $X + Y$  into memory.  $Z$  is the offset.
  - (b) Linked.  $Y=L/511$ ,  $Z=L\%511$ . Load the block with address  $X$ , then chase down the list by  $Y + 1$  elements.  $Z$  is the offset.
  - (c) Indexed.  $Y=L/512$ ,  $Z=L\%512$ . Load the block with address  $X$ , then the  $Y$ th line is the target physical block number.  $Z$  is the offset.
- b.
  - (a) Contiguous. 1.
  - (b) Linked. 4.
  - (c) Indexed. 2.