



DW_asymfifo_s2_sf

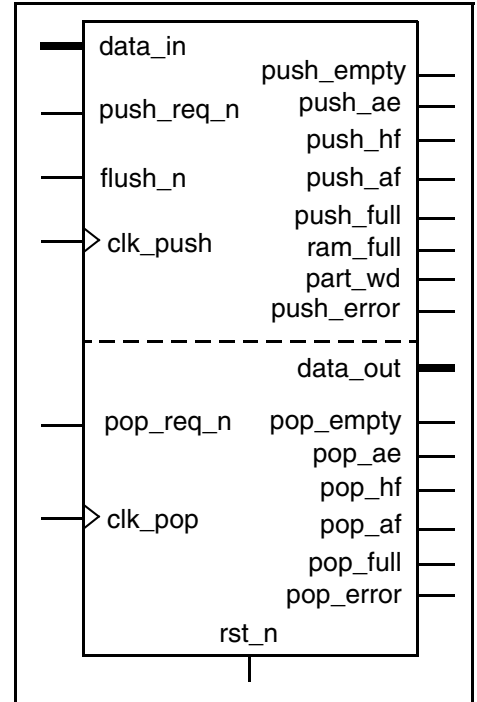
Asymmetric Synchronous (Dual-Clock) FIFO - Static Flags

Version, STAR and Download Information: [IP Directory](#)

Features and Benefits

- Parameterized asymmetric input and output bit widths (must be integer-multiple relationship)
- Fully registered synchronous flag output ports
- Separate status flags for each clock system
- FIFO empty, half full, and full flags
- Parameterized almost full and almost empty flags
- FIFO push error (overflow) and pop error (underflow) flags
- D flip-flop-based memory array for high testability
- Single clock cycle push and pop operations
- Word integrity flag for $data_in_width < data_out_width$
- Partial word flush for $data_in_width < data_out_width$
- Parameterized byte order within a word
- Parameterized reset mode (synchronous or asynchronous, memory array initialized or not)

Revision History



Description

DW_asymfifo_s2_sf is an asymmetric I/O dual independent clock FIFO. It combines the DW_asymfifocntl_s2_sf FIFO controller and the DW_ram_r_w_s_dff flip-flop-based RAM foundation.

The DW_asymfifo_s2_sf is recommended for relatively small memory configurations. For large FIFOs, use the asymmetric FIFO controller, DW_asymfifocntl_s2_sf, in conjunction with a compiled, full-custom RAM array.

Table 1-1 Pin Description

Pin Name	Width	Direction	Function
clk_push	1 bit	Input	Input clock for push interface
clk_pop	1 bit	Input	Input clock for pop interface
rst_n	1 bit	Input	Reset input, active low
push_req_n	1 bit	Input	FIFO push request, active low

Table 1-1 Pin Description (Continued)

Pin Name	Width	Direction	Function
flush_n	1 bit	Input	Flushes the partial word into memory (fills in 0's for empty bits) (for <i>data_in_width</i> < <i>data_out_width</i> only), active low
pop_req_n	1 bit	Input	FIFO pop request, active low
data_in	<i>data_in_width</i> bits	Input	FIFO data to push
push_empty	1 bit	Output	FIFO empty ^a output flag synchronous to <i>clk_push</i> , active high
push_ae	1 bit	Output	FIFO almost empty ^a output flag synchronous to <i>clk_push</i> (determined by <i>push_ae_lvl</i> parameter), active high
push_hf	1 bit	Output	FIFO half full ^a output flag synchronous to <i>clk_push</i> , active high
push_af	1 bit	Output	FIFO almost full ^a output flag synchronous to <i>clk_push</i> (determined by <i>push_af_lvl</i> parameter), active high
push_full	1 bit	Output	FIFO's RAM full ^a output flag (including the input buffer of FIFO for <i>data_in_width</i> < <i>data_out_width</i>) synchronous to <i>clk_push</i> , active high
ram_full	1 bit	Output	FIFO's RAM (excluding the input buffer of FIFO for <i>data_in_width</i> < <i>data_out_width</i>) full output flag synchronous to <i>clk_push</i> , active high
part_wd	1 bit	Output	Partial word accumulated in the input buffer synchronous to <i>clk_push</i> (for <i>data_in_width</i> < <i>data_out_width</i> only; otherwise, tied low), active high
push_error	1 bit	Output	FIFO push error (overrun) output flag synchronous to <i>clk_push</i> , active high
pop_empty	1 bit	Output	FIFO empty ^b output flag synchronous to <i>clk_pop</i> , active high
pop_ae	1 bit	Output	FIFO almost empty ^b output flag synchronous to <i>clk_pop</i> (determined by <i>pop_ae_lvl</i> parameter), active high
pop_hf	1 bit	Output	FIFO half full ^b output flag synchronous to <i>clk_pop</i> , active high
pop_af	1 bit	Output	FIFO almost full ^b output flag synchronous to <i>clk_pop</i> (determined by <i>pop_af_lvl</i> parameter), active high
pop_full	1 bit	Output	FIFO's RAM full ^b output flag (excluding the input buffer of FIFO for case <i>data_in_width</i> < <i>data_out_width</i>) synchronous to <i>clk_pop</i> , active high
pop_error	1 bit	Output	FIFO pop error (underrun) output flag synchronous to <i>clk_pop</i> , active high
data_out	<i>data_out_width</i> bits	Output	FIFO data to pop

a. As perceived by the push interface

b. As perceived by the pop interface

Table 1-2 Parameter Description

Parameter	Values	Description
data_in_width	1 to 2048	Width of the <code>data_in</code> bus; must in an integer-multiple of <code>data_out_width</code> . That is, either: <ul style="list-style-type: none"> ■ $data_in_width = K \times data_out_width$, or ■ $data_out_width = K \times data_in_width$ Where K is an integer.
data_out_width	1 to 2048	Width of the <code>data_out</code> bus; must be an integer-multiple of <code>data_in_width</code> . That is, either: <ul style="list-style-type: none"> ■ $data_in_width = K \times data_out_width$, or ■ $data_out_width = K \times data_in_width$ Where K is an integer.
depth	4 to 1024	Number of words that can be stored in FIFO
push_ae_lvl	1 to $depth-1$	Almost empty level for the <code>push_ae</code> output port (the number of words in the FIFO at or below which the <code>push_ae</code> flag is active)
push_af_lvl	1 to $depth-1$	Almost full level for the <code>push_af</code> output port (the number of empty memory locations in the FIFO at which the <code>push_af</code> flag is active)
pop_ae_lvl	1 to $depth-1$	Almost empty level for the <code>pop_ae</code> output port (the number of words in the FIFO at or below which the <code>pop_ae</code> flag is active)
pop_af_lvl	1 to $depth-1$	Almost full level for the <code>pop_af</code> output port (the number of empty memory locations in the FIFO at which the <code>pop_af</code> flag is active)
err_mode	0 or 1	Error mode: <ul style="list-style-type: none"> ■ 0: Stays active until reset (latched) ■ 1: Active only as long as error condition exists (unlatched)
push_sync	1 to 3	Push flag synchronization mode: <ul style="list-style-type: none"> ■ 1: Single register synchronization from pop pointer ■ 2: Double register ■ 3: Triple register
pop_sync	1 to 3	Pop flag synchronization mode: <ul style="list-style-type: none"> ■ 1: Single register synchronization from push pointer ■ 2: Double register ■ 3: Triple register
rst_mode	0 or 3 Default: 1	Reset mode: <ul style="list-style-type: none"> ■ 0: Asynchronous reset including memory ■ 1: Synchronous reset including memory ■ 2: Asynchronous reset excluding memory ■ 3: Synchronous reset excluding memory

Table 1-2 Parameter Description (Continued)

Parameter	Values	Description
byte_order	0 or 1 Default: 0	Order of bytes or subword within a word: <ul style="list-style-type: none"> 0: First byte is in most significant bits position 1: First byte is in the least significant bits position

Table 1-3 Synthesis Implementations

Implementation Name	Function	License Feature Required
str	Synthesis model	DesignWare

Table 1-4 Simulation Models

Model	Function
DW06.DW_ASYMFIFO_S2_SF_CFG_SIM ^a	Design unit name for VHDL simulation
dw/dw06/src/DW_asymfifo_s2_sf_sim.vhd	VHDL simulation model source code
dw/sim_ver/DW_asymfifo_s2_sf.v	Verilog simulation model source code

a. For reliable simulation in VHDL, always use a configuration in the design specifying the design unit name from the DesignWare Foundation (for example, DW06.DW_ASYMFIFO_S2_SF_CFG_SIM).

Table 1-5 Push Interface Function Table

push_req_n	push_full	Action	push_err
0	0	Push operation	No
0	1	Overflow; incoming data dropped (no action other than error generation)	Yes
1	X	No action	No

Table 1-6 Pop Interface Function Table

pop_req_n	pop_empty	Action	pop_err
0	0	Pop operation	No
0	1	Underrun (no action other than error generation)	Yes
1	X	No action	No

Table 1-7 Flush Interface Function Table (for data_in_width < data_out_width)

flush_n	part_wd	ram_full	Action	push_err
0	0	X	No action	No
0	1	0	Flush	No
0	1	1	No action other than error generation	Yes
1	X	X	No action	No

Block Diagram

The FIFO provides parameterized input data width, output data width, depth, and a full complement of flags (full, almost full, half full, almost empty, empty, and error) for both clock systems. [Figure 1-1](#) on page 6 shows an internal block diagram of the FIFO.

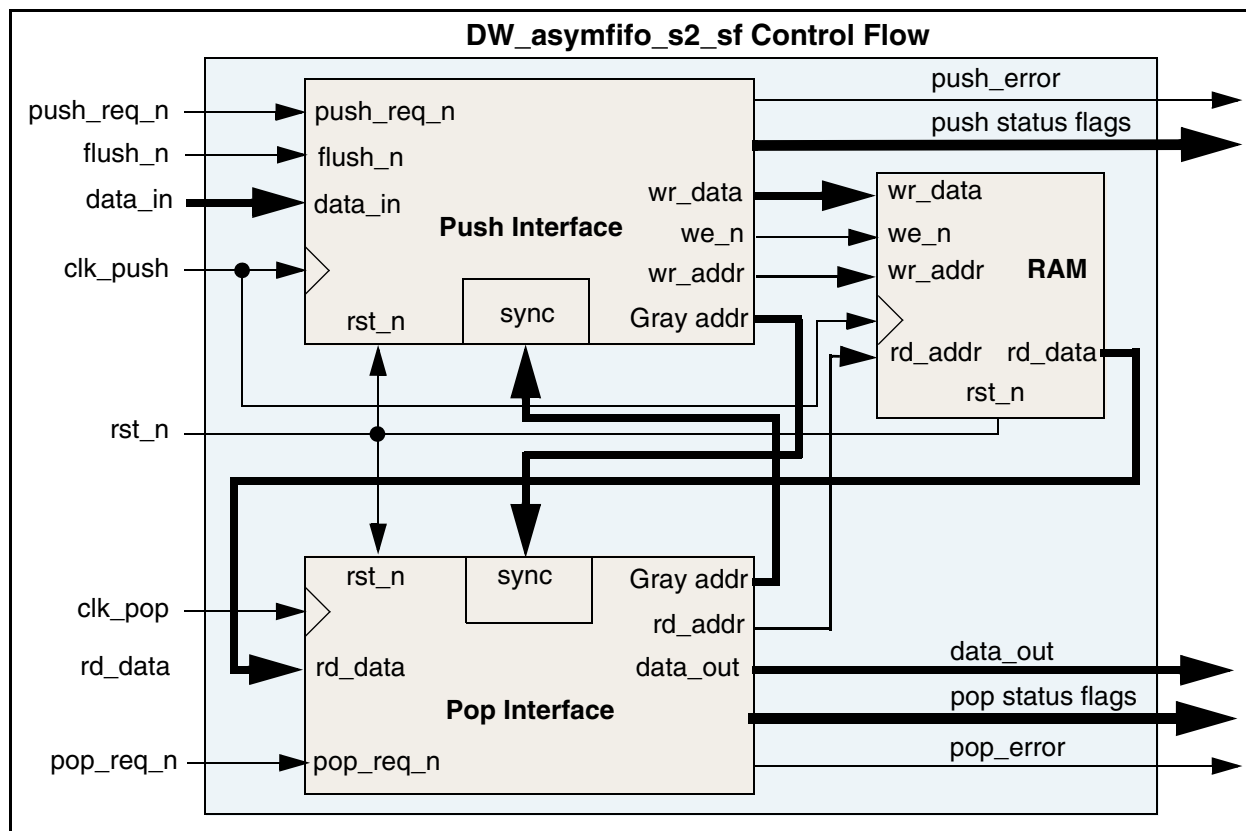
The input data bit width of DW_asymfifo_s2_sf can be different than its output data bit width, but the two must have an integer-multiple relationship (the input bit width being a multiple of the output bit width or vice versa). In other words, either of the following conditions must be true:

- The $data_in_width = K \times data_out_width$, or
- The $data_out_width = K \times data_in_width$ where K is a positive integer. For an example of $data_in_width \neq data_out_width$, see [Figure 1-2](#) on page 7

The asymmetric FIFO provides flag logic and operational error detection logic for each clock domain. Parameterizable features include FIFO depth (from 4 to 1024 locations), almost empty level, almost full level, level of error detection, type of reset (either asynchronous or synchronous, memory array initialized or not), and byte (subword) order within a word. The designer specifies these parameters when the FIFO is instantiated in the design.

Reset can be selected at instantiation to be either synchronous or asynchronous, and can either include or exclude the RAM array.

Figure 1-1 DW_asymfiffo_s2_sf Block Diagram



Input Bus > Output Bus ($\text{data_in_width} > \text{data_out_width}$)

Writing to the FIFO (Push)

For cases where $\text{data_in_width} > \text{data_out_width}$ (assuming that $\text{data_in_width} = K \times \text{data_out_width}$, where K is an integer larger than 1):

- The `flush_n` input pin is not used (at the system level, this pin should not be connected so that it is removed upon synthesis),
- The `part_wd` output pin is tied low.

For an example of $\text{data_in_width} > \text{data_out_width}$ case, refer to [Figure 1-2](#) on page 7.

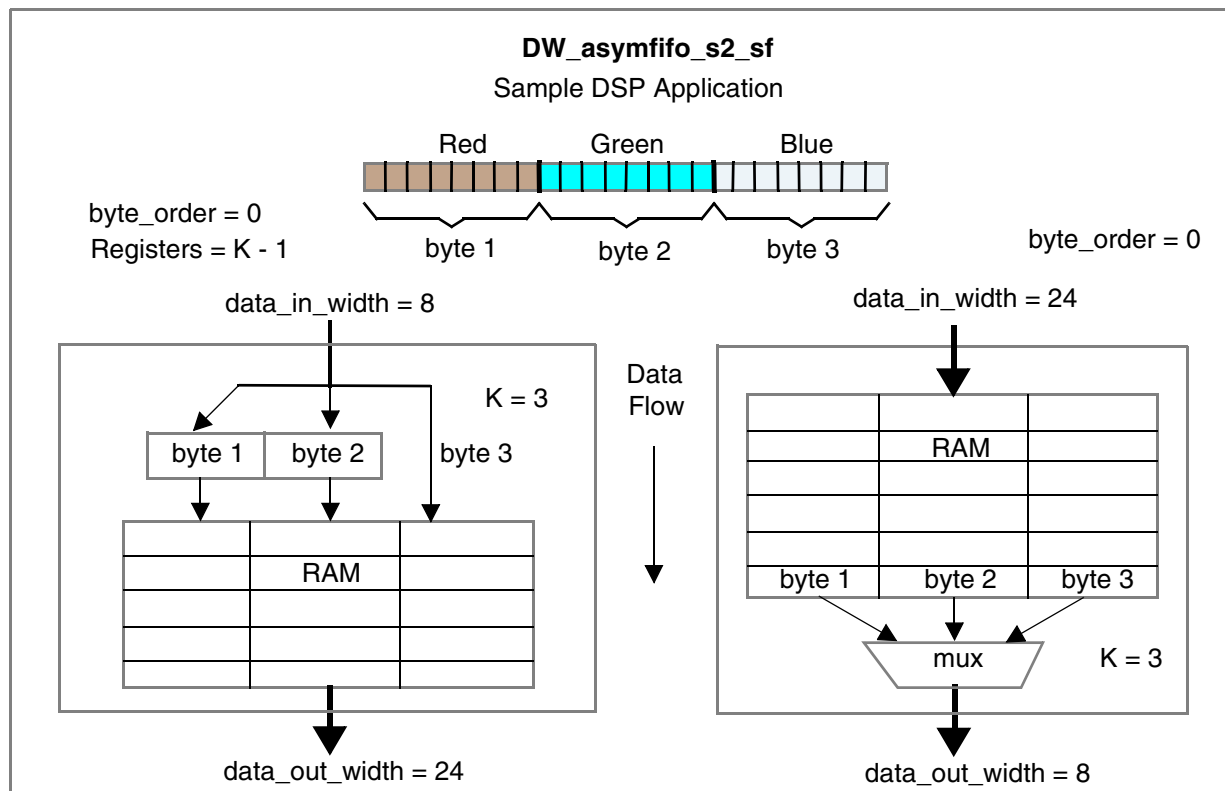
A push is executed when the `push_req_n` input is asserted (low) and the `full` flag is inactive (low) at the rising edge of `clk_push`.

Asserting `push_req_n` when the `full` flag is inactive causes:

- The FIFO to be written, and
- The write address register (which always points to the address of the next word to be pushed) to be incremented on the same rising edge of `clk_push` – the first clock after `push_req_n` is asserted.

This means that `push_req_n` must be asserted early enough to propagate to the FIFO before the next clock.

Figure 1-2 Example of Asymmetric FIFO Operation



Write Errors

An error occurs if a push operation is attempted while the FIFO is full (as perceived by the push interface). That is, the `push_error` output goes active if:

- The `push_req_n` input is asserted (low), and
- The `push_full` flag is active (high)

on the rising edge of `clk_push`.

Reading from the FIFO (Pop)

For cases where $data_in_width > data_out_width$ (assuming that $data_in_width = K \times data_out_width$, where K is an integer larger than 1), the number of bits in a word stored in memory is $data_in_width$. The bit width for each out-going byte (or subword) is $data_out_width$.

For every byte (or subword) to be read, `pop_req_n` must be active (low) at the positive edge of `clk_pop`. Each pop causes one byte (or subword) to be read. Popping K times results in one full word ($data_in_width$ bits) being read. The order of the output bytes within a word is determined by the `byte_order` parameter.

The `data_out` port of the FIFO is asynchronous. A pop operation occurs when `pop_req_n` is asserted (low) when the FIFO is not empty. Asserting `pop_req_n` when the output buffer is not empty causes the `data_out` output port to be switched to the next byte (or subword) on the next rising edge of `clk_pop`. Thus, memory read data must be captured on the `clk_pop` following the assertion of `pop_req_n`.

Refer to the timing diagrams for details of the pop operation.

Read Errors

An error occurs if a pop operation is attempted while the FIFO is empty (as perceived by the pop interface). That is, the `pop_error` output goes active on the rising edge of `clk_pop` if:

- The `pop_req_n` input is active (low), and
- The `pop_empty` flag is active (high).

Input Bus = Output Bus ($data_in_width = data_out_width$)

Writing to the FIFO (Push)

In this case, the FIFO is a symmetric I/O FIFO. Its function is the same as DW_fifo_s2_sf, except for the unused `part_wd`, `flush`, and `ram_full` pins.

A push is executed at the rising edge of `clk_push` when:

- The `push_req_n` input is asserted (low), and
- The `push_full` flag is inactive (low)

Asserting `push_req_n` when `push_full` is inactive causes

- The FIFO be written, and
- The write address register (which always points to the address of the next word to be pushed) to be incremented on the same rising edge of `clk_push` – the first clock after `push_req_n` is asserted.

This means that `push_req_n` must be asserted early enough to propagate to the FIFO before the next clock.

Write Errors

An error occurs if a push operation is attempted while the FIFO is full (as perceived by the push interface). That is, the `push_error` output goes active if:

- The `push_req_n` input is asserted (low), and
- The `push_full` flag is active (high)

on the rising edge of `clk_push`.

Reading from the FIFO (Pop)

In this case, the FIFO is a symmetric I/O FIFO. Its function is the same as the `DW_fifo_s2_sf`, except for the `part_wd`, `flush`, and `ram_full` pins, which are unused.

The `data_out` port of the FIFO is asynchronous.

A pop operation occurs when:

- The `pop_req_n` is asserted (low), and
- The `pop_empty` flag is not active (low) (the FIFO is not empty)

at the rising edge of `clk_pop`.

Asserting `pop_req_n` while `pop_empty` is not active causes the internal read pointer to be incremented on the next rising edge of `clk_pop`. Thus, the FIFO output data must be captured on the rising edge of `clk_pop` following the assertion of `pop_req_n`.

Read Errors

An error occurs if a pop operation is attempted while the FIFO is empty (as perceived by the pop interface). That is, the `pop_error` output goes active if:

- The `pop_req_n` input is active (low), and
- The `pop_empty` flag is active (high)

on the rising edge of `clk_pop`.

Input Bus < Output Bus ($\text{data_in_width} < \text{data_out_width}$)

Writing to the FIFO (Push)

For cases where $\text{data_in_width} < \text{data_out_width}$ (assuming that $\text{data_out_width} = K \times \text{data_in_width}$, where K is an integer larger than 1), every byte (or subword) written to the FIFO is first assembled into a full word with data_out_width bits. For an example of this case, refer to [Figure 1-2](#) on page 7.

A push of a partial word is executed when `push_req_n` is asserted (low) and the `full` flag is inactive (low) at the rising edge of `clk_push`. Thus, a push can occur even if the FIFO is full as long as a pop is executed in the same cycle.

The order of bytes within a word is determined by the `byte_order` parameter. For every byte (or subword) to be written, `push_req_n` must be active (low) at the positive edge of `push_clk`. Pushing K times in either of the cases that enables a push causes the word accumulated in the input buffer to be written to the next available location in the FIFO. This write occurs on the `clk_push` following the assertion of `push_req_n`.

The data at the `data_in` port must be stable for a setup time before the rising edge of `clk_push`, and `push_req_n` must be asserted early enough to propagate through the FIFO to the RAM before the next clock.

In this way, the FIFO is written, and write address register (which always points to the address of the next word to be pushed) is incremented on the same rising edge of `clk_push` – the first clock after `push_req_n` is asserted K times.

Partial Word

When a partial word is in the input buffer register, output flag `part_wd` is active (high). After K times pushing, K bytes (or subwords) are assembled into a full word ($K - 1$ bytes in the input buffer register and the last byte on the `data_in` bus) by a combinational circuit (see [Figure 1-2](#) on page 7.) This achieves single clock cycle operation for the asymmetric FIFO. The full word is then written into memory.

When a full word is sent from the input buffer into memory, `part_wd` goes inactive (low). The order of bytes within a word is determined by the `byte_order` parameter.

Flushing a Partial Word

A flush feature is provided for the $\text{data_in_width} < \text{data_out_width}$ case. The flush feature pushes a partial word into memory when there are less than K bytes accumulated in the input buffer. The input buffer is cleared after a flush, and `part_wd` goes inactive (low).

The sender device activates `flush_n` so that the N bytes of data are pushed into memory without waiting for a complete word to be assembled.

A flush is allowed for data byte word alignment, or when:

- N bytes have been read since the last complete word (where $0 < N < K$), and
- The sender device has no more bytes (or subwords) to assemble the last full word,

while

- The higher level system requires that the receiver device be able to read these N bytes of data (from memory) without waiting.

When the receiver reads the partial word from the memory, the “leftover” bytes of the partial word ($K - N$) are filled with 0s.

A flush is executed when the `flush_n` input is asserted (low) and the `ram_full` flag is inactive (low) at the rising edge of `clk_push`.

When flushing the FIFO, asserting `flush_n` causes the partial word accumulated in the input buffer to be written to the next available memory location. This write occurs on the `clk_push` following the assertion of `flush_n`.

Flushing the FIFO when the input buffer is empty (when the `part_wd` flag is inactive) is a “null” operation and does not cause an error.

Simultaneous Push and Flush

DW_asymfiffo_s2_sf supports simultaneous push and flush under the following conditions:

- The `ram_full` is inactive (low),
- The `part_wd` is active (high),
- The `push_req_n` is active (low), and
- The `flush_n` is active (low).

On the leading edge of `clk_push`, the partial word in the input buffer is flushed into the RAM of the FIFO, and the byte (or subword) at the `data_in` port is pushed into the first byte location of the input buffer.

If there is no partial word in the input buffer (`part_wd` inactive (low)), a simultaneous push and flush generates a normal push. Under this condition, flush is a null action.

Table 1-8 details a simultaneous push and flush operation for $data_in_width < data_out_width$, when `flush_n` is low, `push_req_n` is low.

Table 1-8 Simultaneous Flush and Push Function Table (for $data_in_width < data_out_width$)

<code>part_wd</code>	<code>ram_full</code>	<code>push_full</code>	Action	<code>push_err</code>
0	X	0	No flush, push only	No
1	0	0	Flush and Push	No
1	1	0	No flush, push only. Potential misaligned word.	Yes
1	1	1	No action other than error generation. Data loss and potential misaligned word.	Yes

Write Errors

An error occurs if a push operation is attempted while the FIFO is full (as perceived by the push interface). That is, the `push_error` output goes active if:

- The `push_req_n` input is asserted (low), and
- The `push_full` flag is active (high)

on the rising edge of `clk_push`.

For $data_in_width < data_out_width$, `push_error` also goes active if:

- The `ram_full` is active (high),
- The `part_wd` is active (high), and
- The `flush_n` input is asserted (low)

on the leading edge of `clk_push`.

Reading from the FIFO (Pop)

For cases where $data_in_width < data_out_width$ (assuming that $data_out_width = K \times data_in_width$, where K is an integer larger than 1), the number of bits in a word stored in memory is $data_out_width$.

The `data_out` port of the FIFO is asynchronous.

A pop operation occurs when `pop_req_n` is asserted (low), as long as the FIFO is not empty (as perceived by the pop interface) on the next rising edge of `clk_pop`. Thus, the FIFO output data must be captured on the `clk_pop` following the assertion of `pop_req_n`.

For details of the pop operation, see [“Timing Waveforms”](#) on page 20.

Read Errors

An error occurs if a pop operation is attempted while the FIFO is empty (as perceived by the pop interface). That is, the `pop_error` output goes active if:

- The `pop_req_n` input is active (low), and
- The `pop_empty` flag is active (high)

on the rising edge of `clk_pop`.

Reset

`rst_mode`

The parameter selects whether reset is asynchronous (`rst_mode = 0`) or synchronous (`rst_mode = 1`).

If asynchronous mode is selected, asserting `rst_n` (setting it low) immediately causes:

- The internal address pointers to be set to 0, and
- The flags and error outputs to be initialized.

If synchronous mode is selected, after the assertion of `rst_n`, at the rising edge of `clk_push` the:

- Write address pointer,
- Push flags, and
- The `push_error` output

are initialized. At the rising edge of `clk_pop`, the

- The read address pointer,
- The pop flags, and

- The `pop_error` output are initialized.

Metastability Issues Regarding Reset

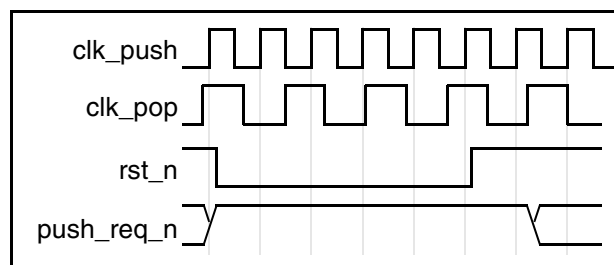
In order to avoid metastability upon reset, the assertion of `rst_n` (low) should be maintained for at least three cycles of the slower of the two clock inputs, `clk_push` and `clk_pop`. During the assertion of `rst_n` and for at least one cycle of `clk_push` after `rst_n` goes high, `push_req_n` must be inactive (high). For more, see [Figure 1-3](#) on page 13.



Note

Because the one input that is critical to proper reset sequencing (`push_req_n`) is in the domain of `clk_push`, it is recommended that the reset input, `rst_n`, should be synchronous to `clk_push`.

Figure 1-3 Avoiding Metastability Upon Reset



Error Outputs and Flag Status

The error outputs and flags are initialized as follows:

- The `push_empty`, `push_ae`, `pop_empty`, and `pop_ae` are initialized to 1 (high), and,
- All other flags and the error outputs are initialized to 0 (low).

Synchronization Between Clock Systems

Each interface (push and pop) operates synchronous to its own clock: `clk_push` and `clk_pop`. Each interface is independent, containing its own state machine and flag logic. The pop interface also has the primary read address counter and a synchronized copy of the write address counter. The push interface also has the primary write address counter and a synchronized copy of the read address counter. The two clocks may be asynchronous with respect to each other. The FIFO controller performs inter-clock synchronization in order for each interface to monitor the actions of the other. This enables the number of words in the FIFO at any given point in time to be determined independently by the two interfaces.

The only information that is synchronized across clock domain boundaries is the read or write address generated by the opposite interface. If an address is transitioning while being sampled by the opposite interface (for example, `wr_addr` sampled by `clk_pop`), sampling uncertainty can occur. By Gray coding the address values that are synchronized across clock domains, this sampling uncertainty is limited to a single bit. Single bit sampling uncertainty results in only one of two possible Gray coded addresses being sampled: the previous address or the new address. The uncertainty in the bit that is changing near a sampling clock

edge directly corresponds to an uncertainty in whether the new value will be captured by the sampling clock edge or whether the previous value will be captured (and the new value may be captured by a subsequent sampling clock edge). Thus there are no errors in sampling Gray coded pointers, just a matter of whether a change of pointer value occurs in time to be captured by a given sampling clock edge or whether it must wait for the next sampling clock edge to be registered

push_sync and pop_sync

The `push_sync` and `pop_sync` parameters determine the number of register stages (1, 2 or 3) used to synchronize the internal Gray code read pointer to `clk_push` (for `push_sync`) and internal Gray code write pointer to `clk_pop` (for `pop_sync`). A value of 1 indicates single-stage synchronization; a value of 2 indicates double-stage synchronization; a value of 3 indicates triple-stage synchronization.

Single-stage synchronization is only adequate when using very slow clock rates (with respect to the target technology). There must be enough timing slack to allow metastable synchronization events to stabilize and propagate to the pointer and flag registers.



Note

Since timing slack and selection of register types is very difficult to control and metastability characteristics of registers are extremely difficult to ascertain, single-stage synchronization is not recommended.

Double-stage synchronization is desirable when using relatively high clock rates. It allows an entire clock period for metastable events to settle at the first stage before being cleanly clocked into the second stage of the synchronizer. Double-stage synchronization increases the latency between the two interfaces, resulting in flags that are less up to date with respect to the true state of the FIFO.

Triple-stage synchronization is desirable when using very high clock rates. It allows an entire clock period for metastable events to settle at the first stage before being clocked into the second stage of the synchronizer. Then, in the unlikely event that a metastable event propagates into the second stage, the output of the second stage is allowed to settle for another entire clock period before being clocked into the third stage. Triple-stage synchronization increases the latency between the two interfaces, resulting in flags that are less up to date with respect to the true state of the FIFO.

Empty to Not Empty Transitional Operation

When the FIFO is empty, both `push_empty` and `pop_empty` are active (high). During the first push (`push_req_n` active (low)), the rising edge of `clk_push` writes the first word into the FIFO. The `push_empty` flag is driven low.

The `pop_empty` flag does not go low until one cycle (of `clk_pop`) after the new internal write pointer value has been synchronized to `clk_pop`. This could be as long as two or three cycles (depending on the value of the `pop_sync` parameter). For more information, see [“Timing Waveforms”](#) on page 20.

You should allow for this latency in the depth budgeting of the FIFO design.

Not Empty to Empty Transitional Operation

When the FIFO is almost empty, both `push_empty` and `pop_empty` are inactive (low) and `pop_ae` is active (high). During the final pop (`pop_req_n` active (low)), the rising edge of `clk_pop` reads the last word out of the FIFO. The `pop_empty` flag is driven high.

The `push_empty` flag is not asserted (high) until one cycle (of `clk_push`) after the new internal read pointer value has been synchronized to `clk_push`. This could be as long as 2 or 3 cycles (depending on the value of the `push_sync` parameter). For more information, see [“Timing Waveforms”](#) on page 20.

You should be aware of this latency when designing the system data flow protocol.

Full to Not Full Transitional Operation

When the FIFO is full, `part_wd`, `push_full`, and `pop_full` are active (high). During the first pop (`pop_req_n` active (low)), the rising edge of `clk_pop` reads the first word out of the FIFO. The `pop_full` flag is driven low.

The `push_full` flag does not go low until one cycle (of `clk_push`) after the new internal read pointer value has been synchronized to `clk_push`. This could be as long as two or three cycles (depending on the value of the `push_sync` parameter). For more information, see [“Timing Waveforms”](#) on page 20.

You should be aware of this latency when designing the system data flow protocol.

Not Full to Full Transitional Operation

When the FIFO is almost full, both `push_full` and `pop_full` are inactive (low), and `push_af` is active (high). During the final push (`push_req_n` active (low)), the rising edge of `clk_push` writes the last word into the FIFO. The `push_full` flag is driven high.

The `pop_full` flag is not asserted (high) until one cycle (of `clk_pop`) after the new internal write pointer value has been synchronized to `clk_pop`. This could be as long as two or three cycles (depending on the value of the `pop_sync` parameter). For more information, see [“Timing Waveforms”](#) on page 20.

You should allow for this latency in the depth budgeting of the FIFO design.

ram_full

The `ram_full` output is used for the $data_in_width < data_out_width$ case. This flag is synchronous to `clk_push`. The `ram_full` output indicates that the RAM (excluding the FIFO input buffer) is full and there is no space available for flushing a partial word into the RAM. However, if `part_wd` is inactive (low), there are still some spaces in the input buffer for incoming sub-words. The `ram_full` output is set low when `rst_n` is applied.

For $data_in_width \geq data_out_width$, `ram_full` is tied to the `push_full` output.

part_wd

The `part_wd` output is only used for the $data_in_width < data_out_width$ case. This flag is synchronous to `clk_push`. The `part_wd` output indicates that the FIFO has a partial word accumulated in the input buffer. The `part_wd` output is set low when `rst_n` is applied.

For $data_in_width > data_out_width$ and $data_in_width = data_out_width$ cases, `part_wd` is tied low, since the input data is always a full word.

Errors

err_mode

The `err_mode` parameter determines whether the `push_error` and `pop_error` outputs remain active until reset (persistent), or for only the clock cycle in which the error is detected (dynamic).

When the `err_mode` parameter is set to 0 at design time, persistent error flags are generated. When the `err_mode` parameter is set to 1 at design time, dynamic error flags are generated.

push_error

The `push_error` output signal indicates a push request was seen while the `push_full` output was active (high) (an overrun error). When an overrun condition occurs, the write address register cannot advance, and the RAM in the FIFO is not written.

Therefore, a push request that would overrun the FIFO is, in effect, rejected, and an error is generated. This guarantees that no data already in the FIFO is destroyed (overwritten). Other than the loss of the data accompanying the rejected push request, FIFO operation can continue without reset.

For `data_in_width < data_out_width` case, the `push_error` output signal may also indicate a flush request was seen while the `ram_full` output was active (high). This indicates a potential overrun error and/or word misalignment error. For details, see [Table 1-8](#) on page 11.

pop_error

The `pop_error` output signal indicates a pop request was seen while the `pop_empty` output signal was active (high) (an underrun error). When an underrun condition occurs, the read address register cannot decrement, as there is no data in the FIFO to retrieve.

The FIFO timing is such that the logic controlling the `pop_req_n` input would not see the error until 'nonexistent' data had already been accepted by the receiving logic. This problem is easily avoided if the logic controlling the `pop_req_n` input can pay close attention to the `pop_empty` output and thus avoid an underrun completely.

Controller Status Flag Outputs

The two halves of the FIFO controller each have their own set of status flags indicating their separate view of the state of the FIFO. It is important to note that both the push interface and the pop interface perceives the state of fullness of the FIFO independently based on information from the opposing interface that is delayed up to three clock cycles for proper synchronization between clock domains.

The push interface status flags respond immediately to changes in state caused by push operations but there is delay between pop operations and corresponding changes of state of the push status flags. This delay is due to the latency introduced by the registers used to synchronize the internal Gray coded read pointer to `clk_push`. The pop interface status flags respond immediately to changes in state caused by pop operations but there is delay between push operations and corresponding changes of state of the pop status flags. This delay is due to the latency introduced by the registers used to synchronize the internal Gray coded write pointer to `clk_pop`.

Most status flags have a property which is potentially useful to the designed operation of the FIFO controller. These properties are described in the following explanations of the flag behaviors.

push_empty

The `push_empty` output, active high, is synchronous to the `clk_push` input. `push_empty` indicates to the

push interface that the FIFO is empty. During the first push, the rising edge of `clk_push` causes the first word to be written into the FIFO, and `push_empty` is driven low.

The action of the last word being popped from a nearly empty FIFO is controlled by the pop interface. Thus, the `push_empty` output is asserted only after the new internal Gray code read pointer (from the pop interface) is synchronized to `clk_push` and processed by the status flag logic.

Property of `push_empty`

If `push_empty` is active (high) then the FIFO is truly empty. This property does not apply to `pop_empty`.

`push_ae`

The `push_ae` output, active high, is synchronous to the `clk_push` input. The `push_ae` output indicates to the push interface that the FIFO is almost empty when there are no more than `push_ae_lvl` words currently in the FIFO to be popped as perceived at the push interface.

The `push_ae_lvl` parameter defines the almost empty threshold of the push interface independent of that of the pop interface. The `push_ae` output is useful when it is desirable to push data into the FIFO in bursts (without allowing the FIFO to become empty).

Property of `push_ae`

If `push_ae` is active (high) then the FIFO has at least $(depth - push_ae_lvl)$ available locations. Thus such status indicates that the push interface can safely and unconditionally push $(depth - push_ae_lvl)$ words into the FIFO. This property guarantees that such a 'blind push' operation will not overrun the FIFO.

`push_hf`

The `push_hf` output, active high, is synchronous to the `clk_push` input, and indicates to the push interface that the FIFO has at least half of its memory locations occupied as perceived by the push interface.

Property of `push_hf`

If `push_hf` is inactive (low) then the FIFO has at least half of its locations available. Thus such status indicates that the push interface can safely and unconditionally push $(INT(depth/2) + 1)$ words into the FIFO. This property guarantees that such a 'blind push' operation will not overrun the FIFO.

`push_af`

The `push_af` output, active high, is synchronous to the `clk_push` input. `push_af` indicates to the push interface that the FIFO is almost full when there are no more than `push_af_lvl` empty locations in the FIFO as perceived by the push interface.

The `push_af_lvl` parameter defines the almost full threshold of the push interface independent of the pop interface. The `push_af` output is useful when more than one cycle of advance warning is needed to stop the flow of data into the FIFO before it becomes full (to avoid a FIFO overrun).

Property of `push_af`

If `push_af` is inactive (low) then the FIFO has at least $(push_af_lvl + 1)$ available locations. Thus such status indicates that the push interface can safely and unconditionally push $(push_af_lvl + 1)$ words into the FIFO. This property guarantees that such a 'blind push' operation will not overrun the FIFO.

`push_full`

The `push_full` output, active high, is synchronous to the `clk_push` input. `push_full` indicates to the push interface that the FIFO is full. During the final push, the rising edge of `clk_push` causes the last word to be pushed, and `push_full` is asserted.

The action of the first word being popped from a full FIFO is controlled by the pop interface. Thus, the `push_full` output goes low only after the new internal Gray code read pointer from the pop interface is synchronized to `clk_pop` and processed by the status flag state logic.

pop_empty

The `pop_empty` output, active high, is synchronous to the `clk_pop` input. `pop_empty` indicates to the pop interface that the FIFO is empty as perceived by the pop interface. The action of the last word being popped from a nearly empty FIFO is controlled by the pop interface. Thus, the `pop_empty` output is asserted at the rising edge of `clk_pop` that causes the last word to be popped from the FIFO.

The action of pushing the first word into an empty FIFO is controlled by the push interface. That means `pop_empty` goes low only after the new internal Gray code write pointer from the push interface is synchronized to `clk_pop` and processed by the status flag state logic.

pop_ae

The `pop_ae` output, active high, is synchronous to the `clk_pop` input. `pop_ae` indicates to the pop interface that the FIFO is almost empty when there are no more than `pop_ae_lvl` words currently in the FIFO to be popped as perceived by the pop interface.

The `pop_ae_lvl` parameter defines the almost empty threshold of the pop interface independent of the push interface. The `pop_ae` output is useful when more than one cycle of advance warning is needed to stop the popping of data from the FIFO before it becomes empty (to avoid a FIFO underrun).

Property of pop_ae

If `pop_ae` is inactive (low) then there are at least $(pop_ae_lvl + 1)$ words in the FIFO. This status indicates that the pop interface can safely and unconditionally pop $(pop_ae_lvl + 1)$ words out of the FIFO. This property guarantees that such a 'blind pop' operation will not underrun the FIFO.

pop_hf

The `pop_hf` output, active high, is synchronous to the `clk_pop` input. `pop_hf` indicates to the pop interface that the FIFO has at least half of its memory locations occupied as perceived by the pop interface.

Property of pop_hf

If `pop_hf` is active (high) then at least half of the words in the FIFO are occupied. This status indicates that the pop interface can safely and unconditionally pop $\text{INT}((depth + 1)/2)$ words out of the FIFO. This property guarantees that such a 'blind pop' operation will not underrun the FIFO.

pop_af

The `pop_af` output, active high, is synchronous to the `clk_pop` input. `pop_af` indicates to the pop interface that the FIFO is almost full when there are no more than `pop_af_lvl` empty locations in the FIFO as perceived by the pop interface.

The `pop_af_lvl` parameter defines the almost full threshold of the pop interface independent of that of the push interface. The `pop_af` output is useful when it is desirable to pop data out of the FIFO in bursts (without allowing the FIFO to become empty).

Property of pop_af

If `pop_af` is active (high) then there are at least $(depth - pop_af_lvl)$ words in the FIFO. This status indicates that the pop interface can safely and unconditionally pop $(depth - pop_af_lvl)$ words out of the FIFO. This property guarantees that such a 'blind pop' operation will not underrun the FIFO.

pop_full

The `pop_full` output, active high, is synchronous to the `clk_pop` input. The `pop_full` output indicates to the pop interface that the FIFO is full as perceived by the pop interface. The action of popping the first word out of a full FIFO is controlled by the pop interface. Thus, the `pop_full` output goes low at the rising edge of the `clk_pop` that causes the first word to be popped.

The action of the last word being pushed into a nearly full FIFO is controlled by the push interface. This means the `pop_full` output is asserted only after the new write pointer from the pop interface is synchronized to `clk_pop` and processed by the status flag state logic.

Property of pop_full

If `pop_full` is active (high) then the FIFO is truly full. This property does not apply to `push_full`.

Simulation Methodology

DW_asymfifo_s2_sf contains synchronization of Gray coded pointers between clock domains for which there are two methods for simulation.

- The first method is to use the simulation models, which emulate the RTL model, with no modeling of metastable behavior. Using this method requires no extra action.
- The second method (only available for Verilog simulation models) is to enable modeling of random skew between bits of the Gray coded pointers that traverse to and from each domain.

To use the second method, a Verilog preprocessing macro named `DW_MODEL_MISSAMPLES` must be defined in one of the following ways:

- Specify the Verilog preprocessing macro in Verilog code:


```
`define DW_MODEL_MISSAMPLES
```
- Or, include a command line option to the simulator, such as


```
+define+DW_MODEL_MISSAMPLES
```

 (which is used for the Synopsys VCS simulator)

Suppressing Warning Messages During Verilog Simulation

The Verilog simulation model includes macros that allow you to suppress warning messages during simulation.

To suppress all warning messages for all DWBB components, define the `DW_SUPPRESS_WARN` macro in either of the following ways:

- Specify the Verilog preprocessing macro in Verilog code:


```
`define DW_SUPPRESS_WARN
```
- Or, include a command line option to the simulator, such as:


```
+define+DW_SUPPRESS_WARN
```

 (which is used for the Synopsys VCS simulator)

The warning messages for this model include the following:

- If values other than 1 or 0 are present on a clock port, the following message is displayed:

```
WARNING: <instance_path>.<clock_name>_monitor:  
        at time = <timestamp>, Detected unknown value, x, on <clock_name> input.
```

To suppress only this warning message for all DWBB components, use the following macro:

- Define the DW_DISABLE_CLK_MONITOR macro. You can define this macro in the following ways:
 - Specify the Verilog preprocessing macro in Verilog code:
`define DW_DISABLE_CLK_MONITOR
 - Or, include a command line option to the simulator, such as:
+define+DW_DISABLE_CLK_MONITOR (which is used for the Synopsys VCS simulator)

This message is also suppressed using the DW_SUPPRESS_WARN macro explained earlier.

Timing Waveforms

The figures in this section show the waveforms for various conditions of DW_asymfifo_s2_sf.

Also refer to the DW_fifo_s2_sf datasheet for the following timing diagrams where Input = Output:

- [push](#) and [pop](#) timing waveforms
- [single word push and pop](#) timing waveforms
- FIFO *depth* $\neq 2^n$ [push](#) and [pop](#) timing waveforms
- FIFO *depth* $\neq 2^n$ [single word](#) timing waveforms

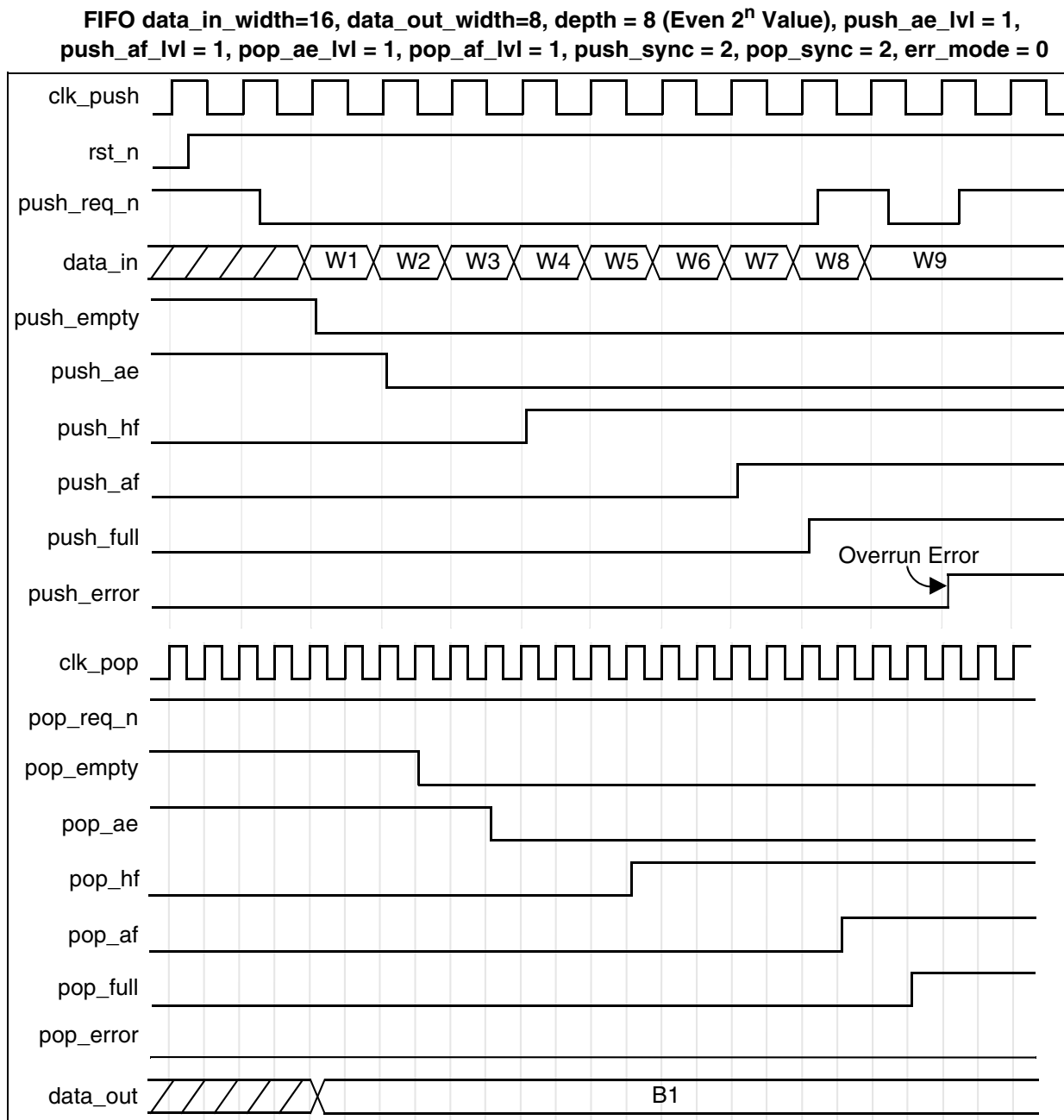
Figure 1-4 Push Timing Waveforms for Input > Output

Figure 1-5 Push Timing Waveforms for Input < Output

FIFO data_in_width=8, data_out_width=16, depth = 8 (Even 2ⁿ Value), push_ae_lvl = 1,
push_af_lvl = 1, pop_ae_lvl = 1, pop_af_lvl = 1, push_sync = 2, pop_sync = 2, err_mode = 0

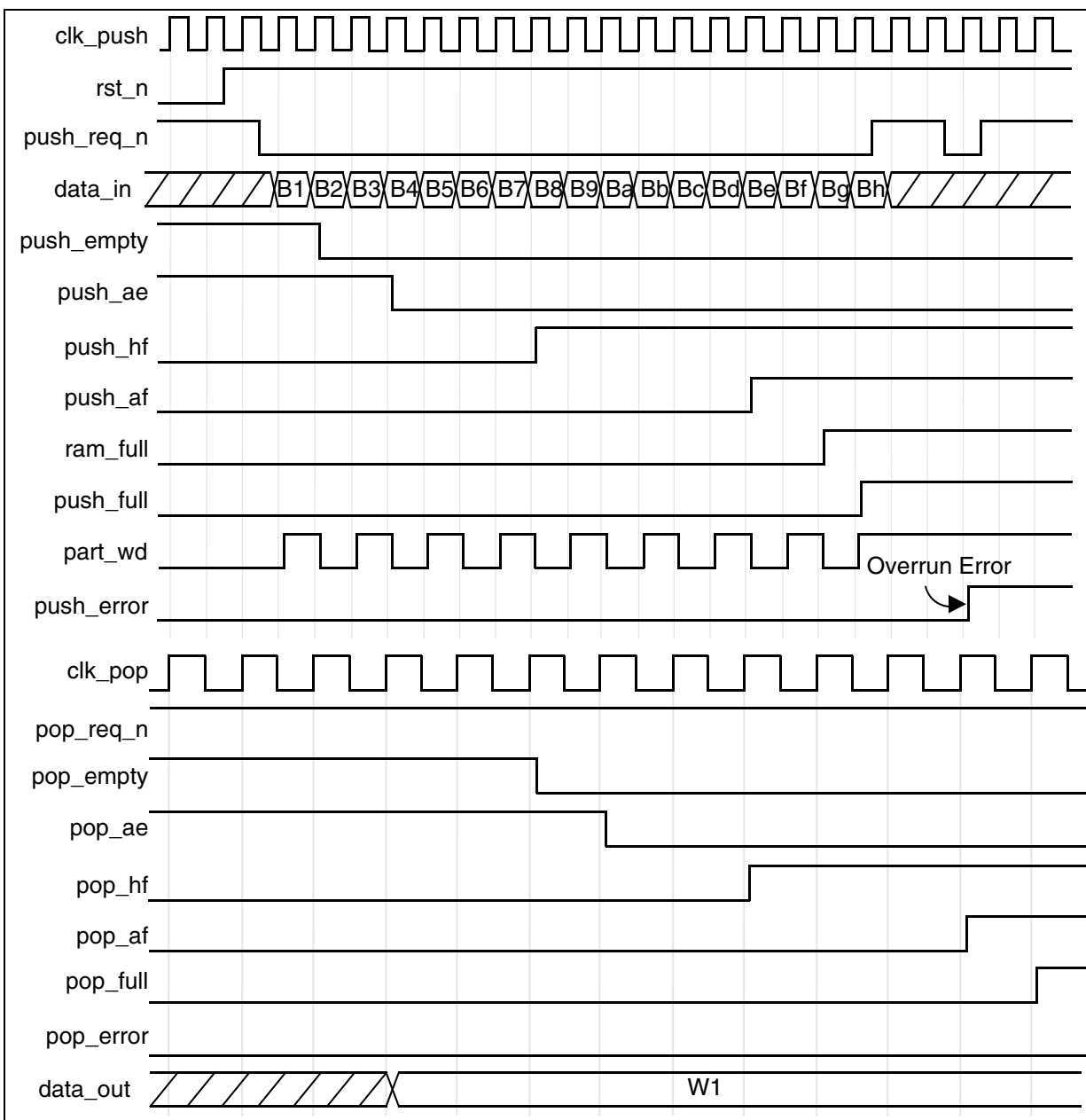


Figure 1-6 Pop Timing Waveforms for Input > Output

**FIFO data_in_width=16, data_out_width=8, depth = 8 (Even 2ⁿ Value), push_ae_lvl = 1,
push_af_lvl = 1, pop_ae_lvl = 1, pop_af_lvl = 1, push_sync = 2, pop_sync = 2**

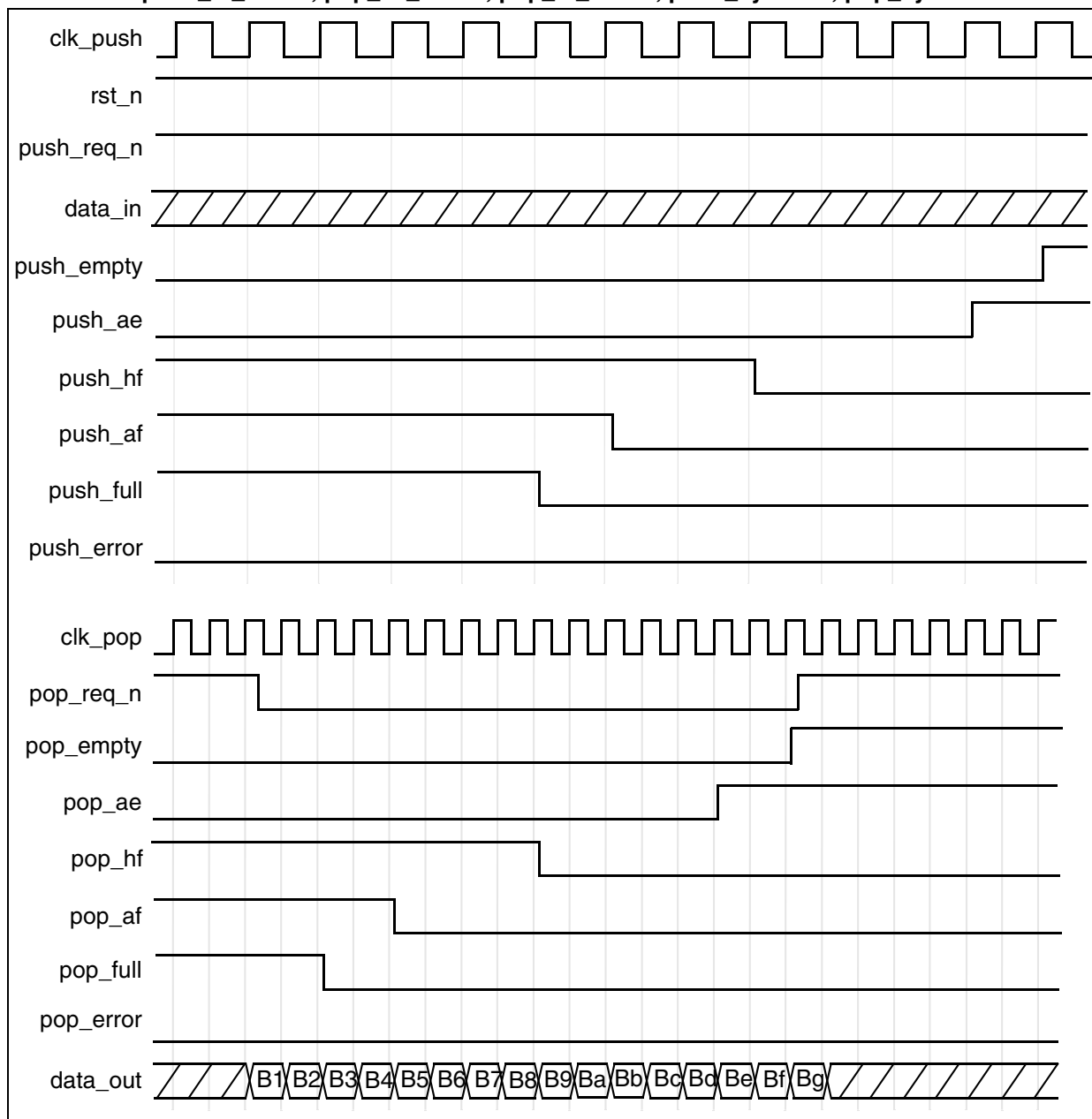


Figure 1-7 Pop Timing Waveforms for Input < Output

**FIFO data_in_width=8, data_out_width=16, depth = 8 (Even 2ⁿ Value), push_ae_lvl = 1,
push_af_lvl = 1, pop_ae_lvl = 1, pop_af_lvl = 1, push_sync = 2, pop_sync = 2**

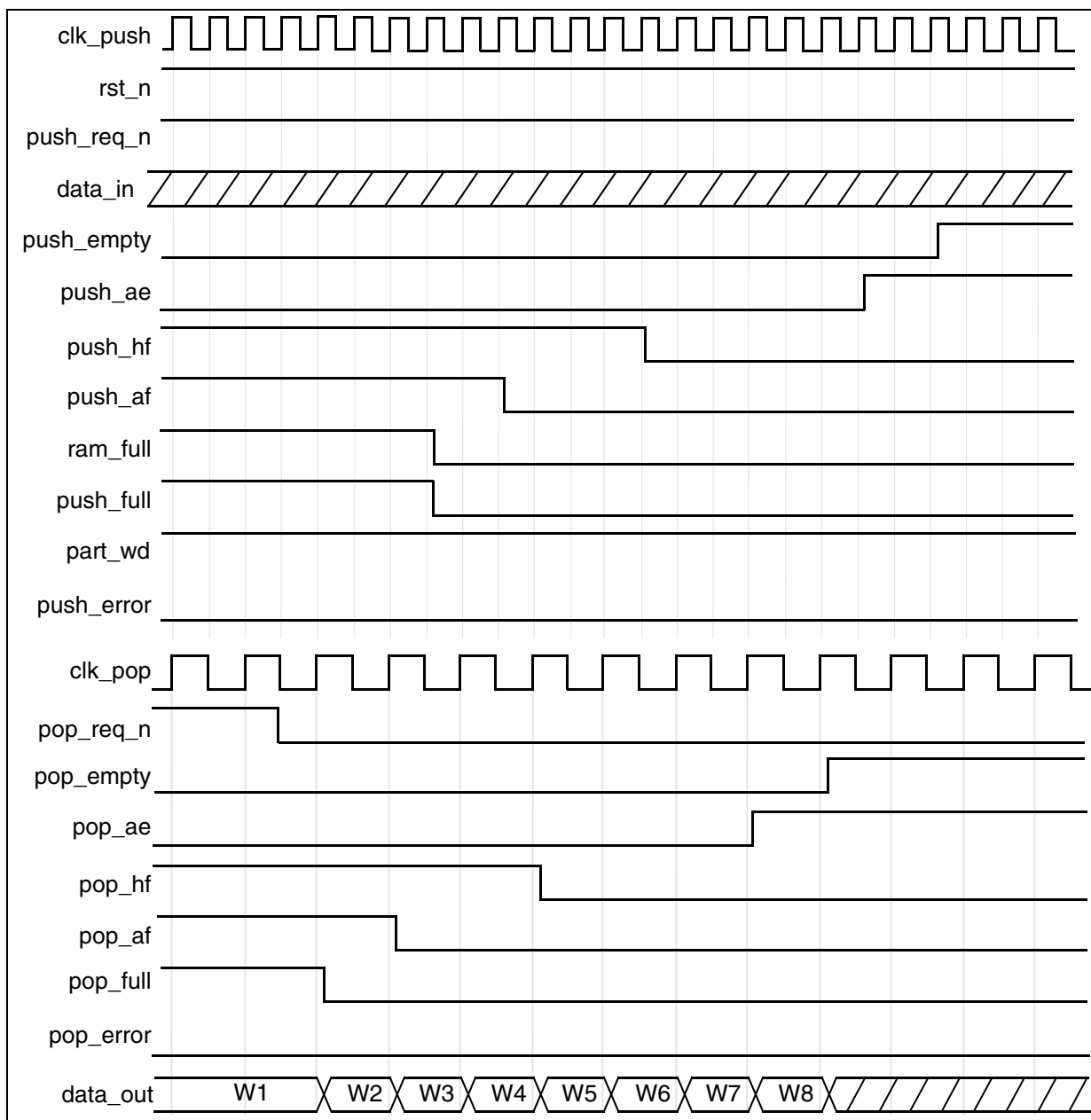


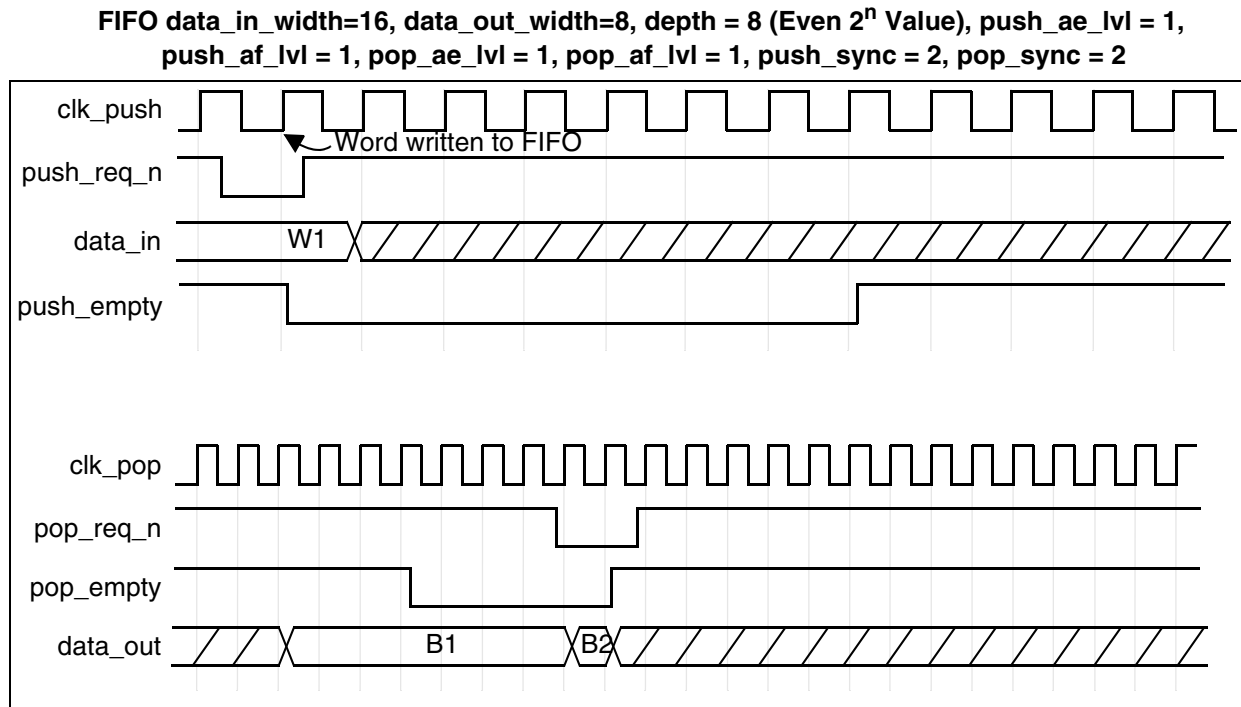
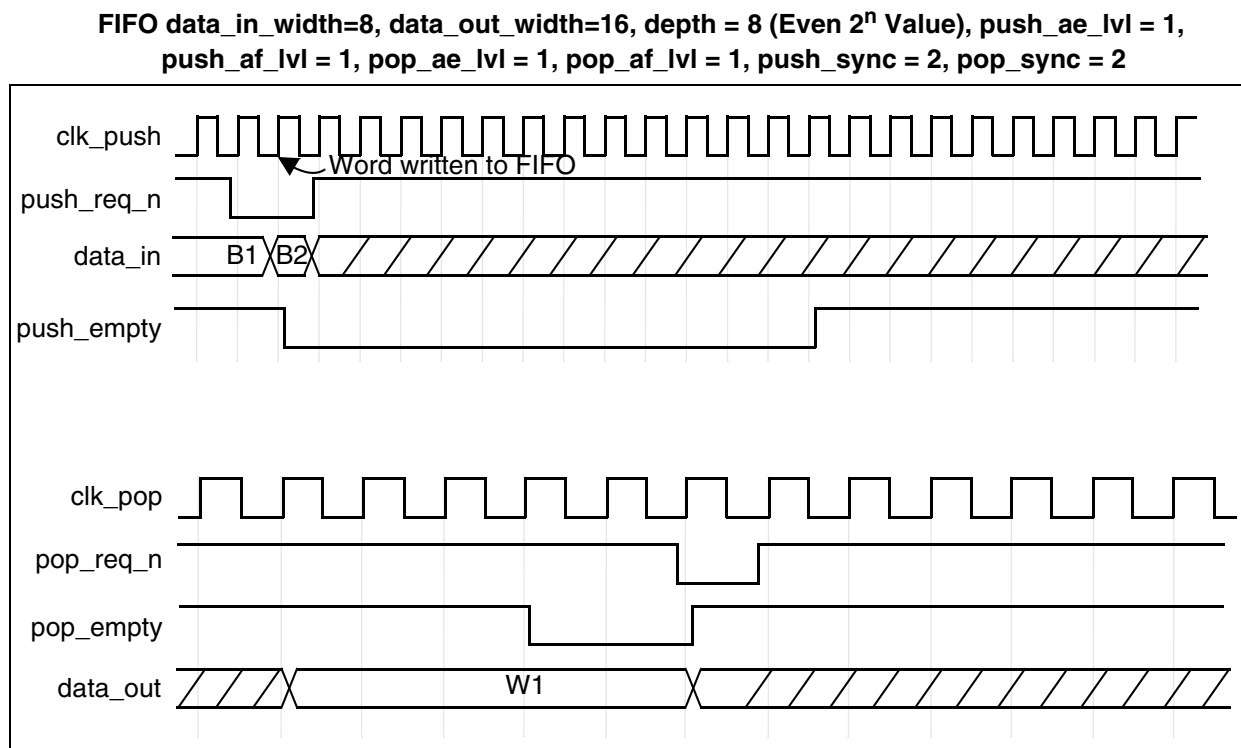
Figure 1-8 Single Word Push and Pop Timing Waveforms for Input > Output**Figure 1-9 Single Word Push and Pop Timing Waveforms for Input < Output**

Figure 1-10 FIFO Depth $\neq 2^n$ Push Timing Waveforms for Input > Output

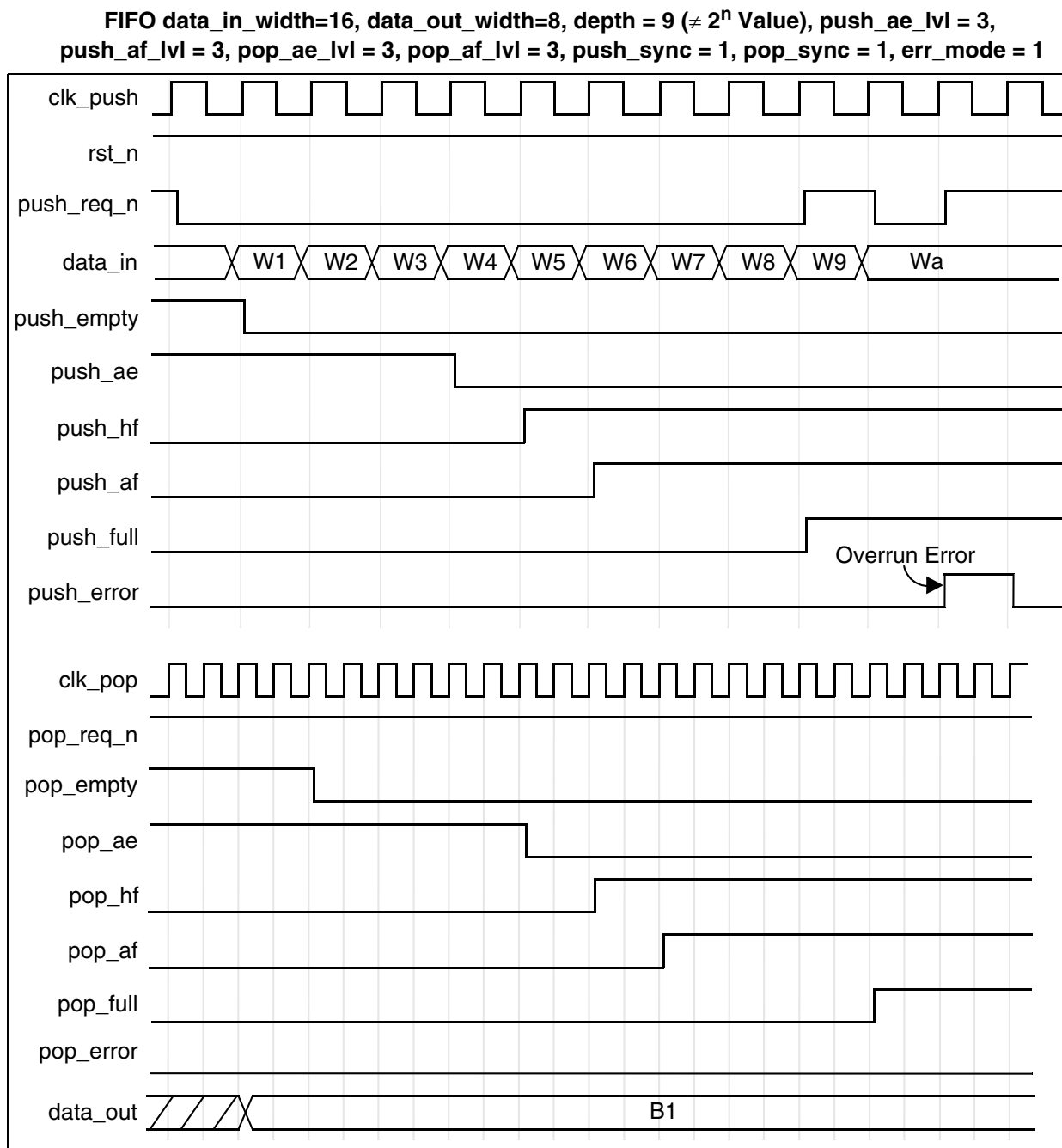


Figure 1-11 FIFO Depth $\neq 2^n$ Push Timing Waveforms for Input < Output

FIFO data_in_width=8, data_out_width=16, depth = 9 ($\neq 2^n$ Value), push_ae_lvl = 3,
push_af_lvl = 3, pop_ae_lvl = 3, pop_af_lvl = 3, push_sync = 1, pop_sync = 1, err_mode = 1

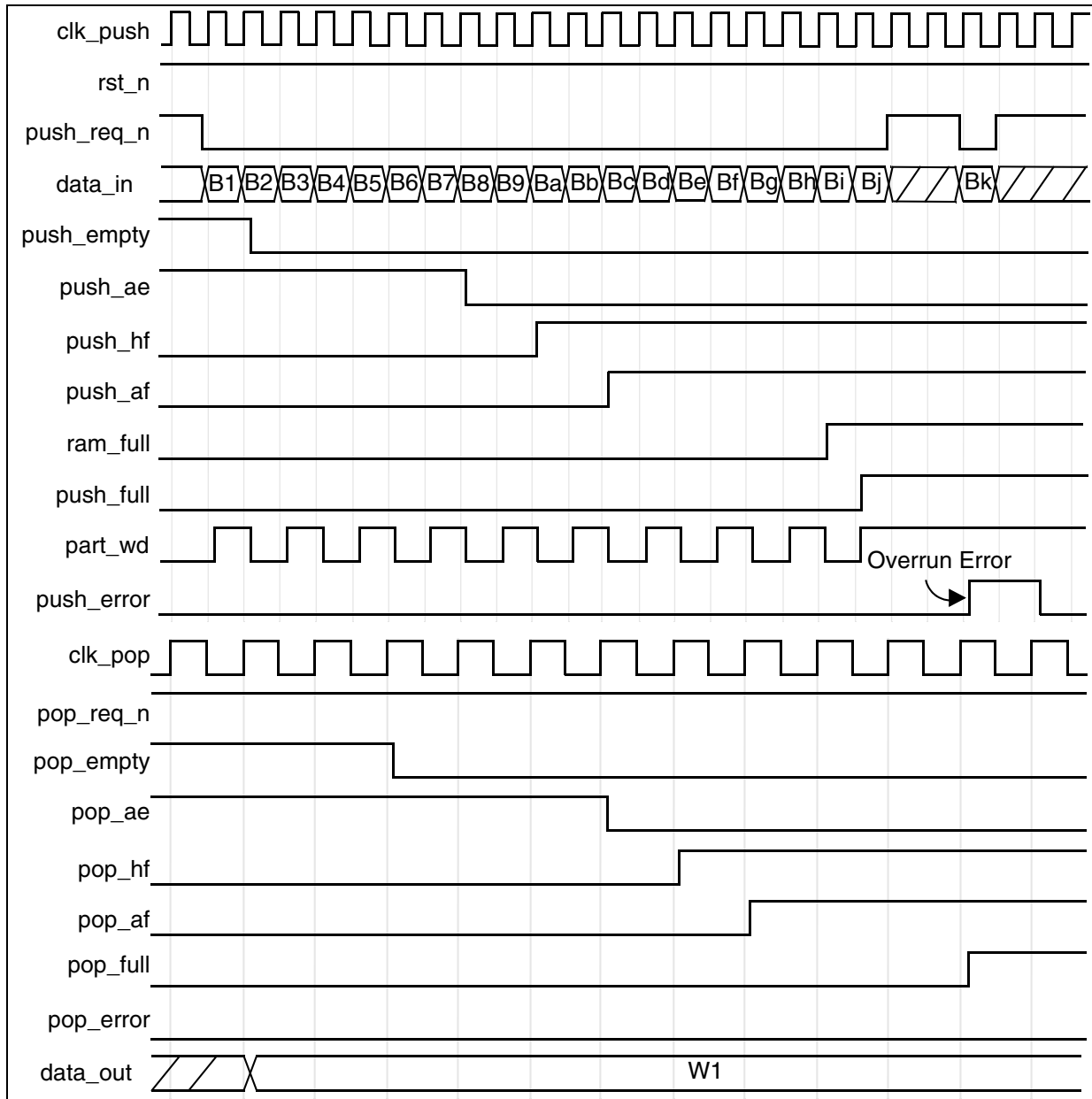


Figure 1-12 FIFO Depth $\neq 2^n$ Pop Timing Waveforms for Input > Output

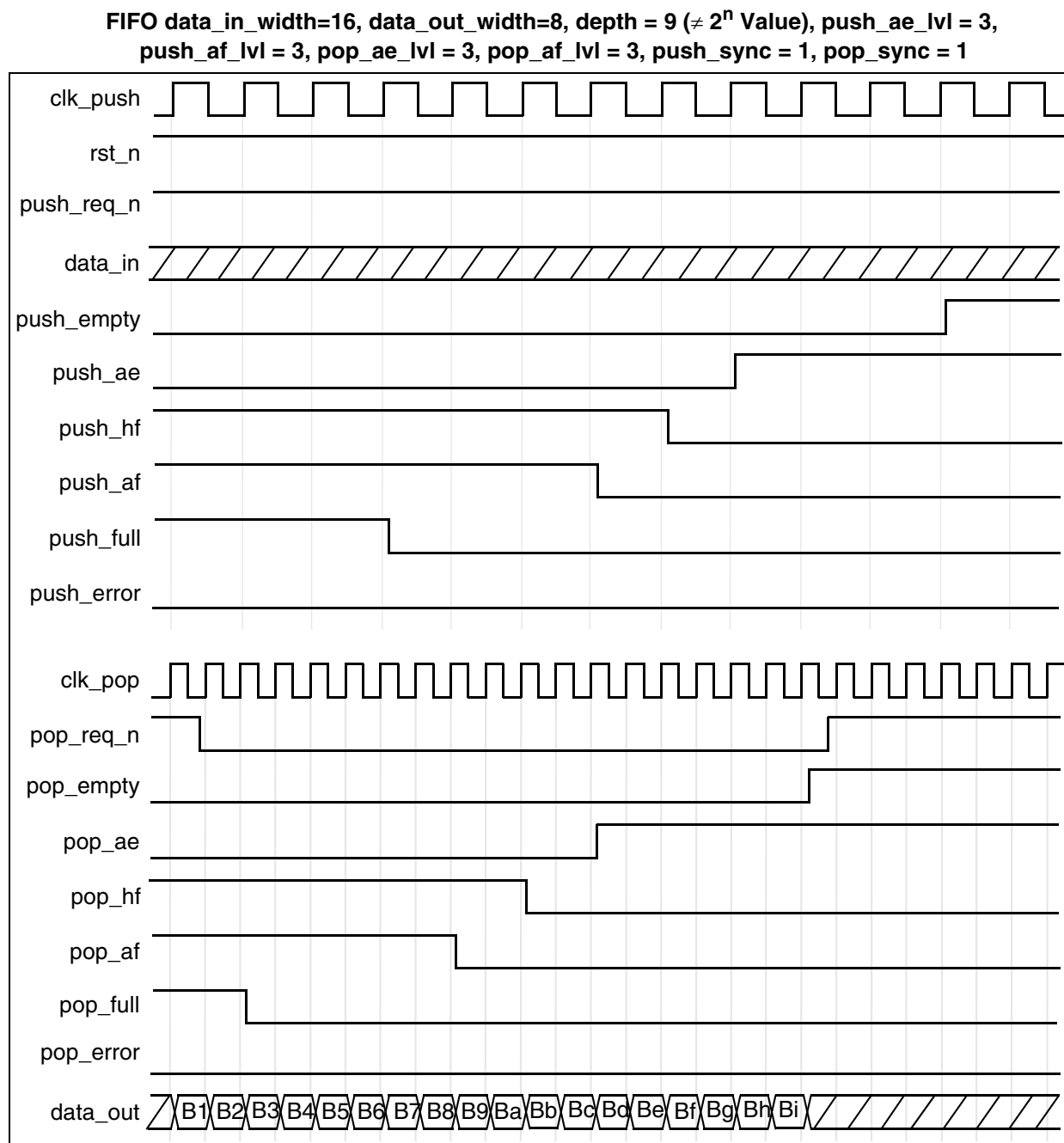


Figure 1-13 FIFO Depth $\neq 2^n$ Pop Timing Waveforms for Input < Output

FIFO data_in_width=8, data_out_width=16, depth = 9 ($\neq 2^n$ Value), push_ae_lvl = 3,
push_af_lvl = 3, pop_ae_lvl = 3, pop_af_lvl = 3, push_sync = 1, pop_sync = 1

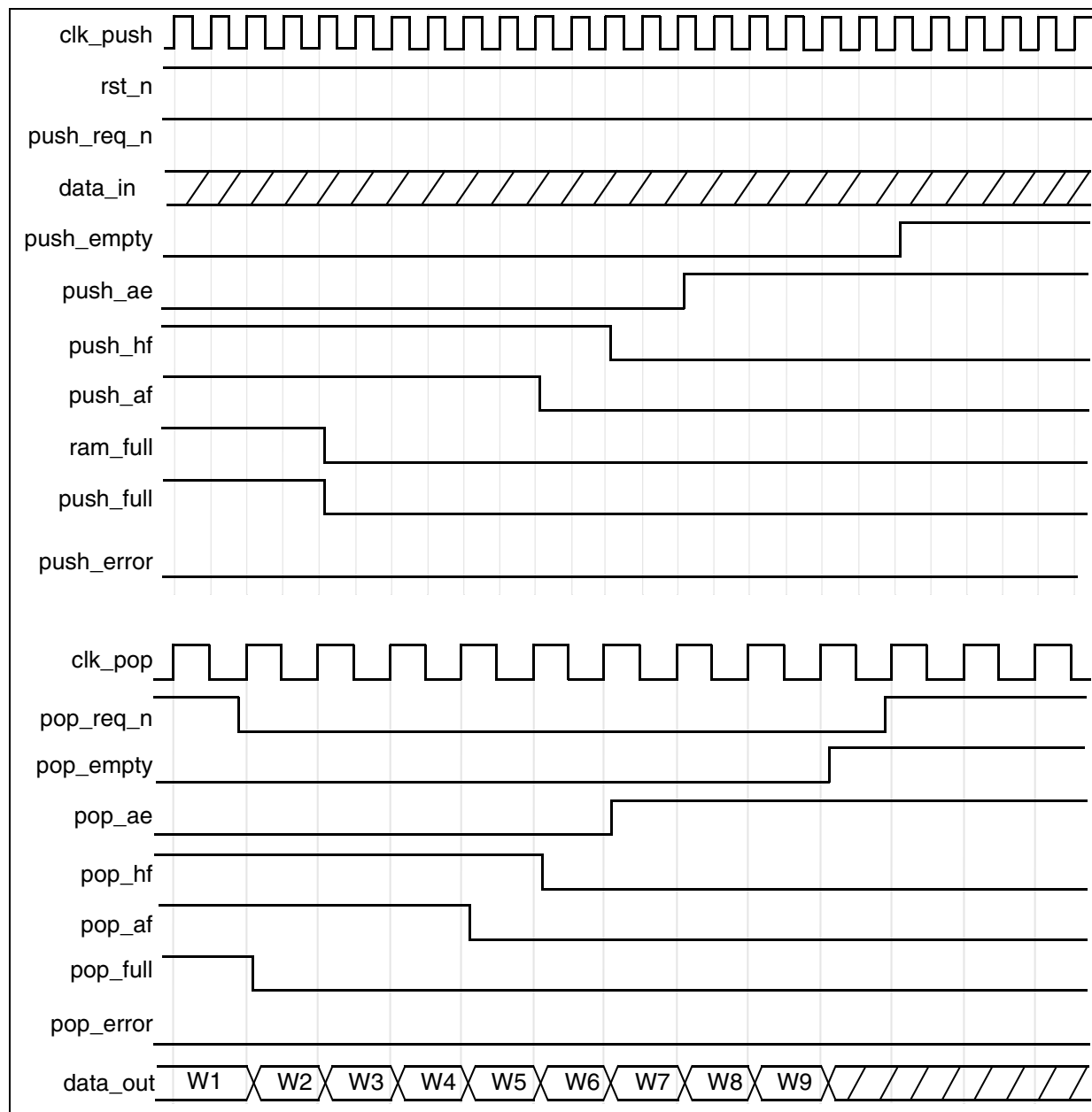


Figure 1-14 FIFO Depth $\neq 2^n$ Single Word Timing Waveform for Input > Output

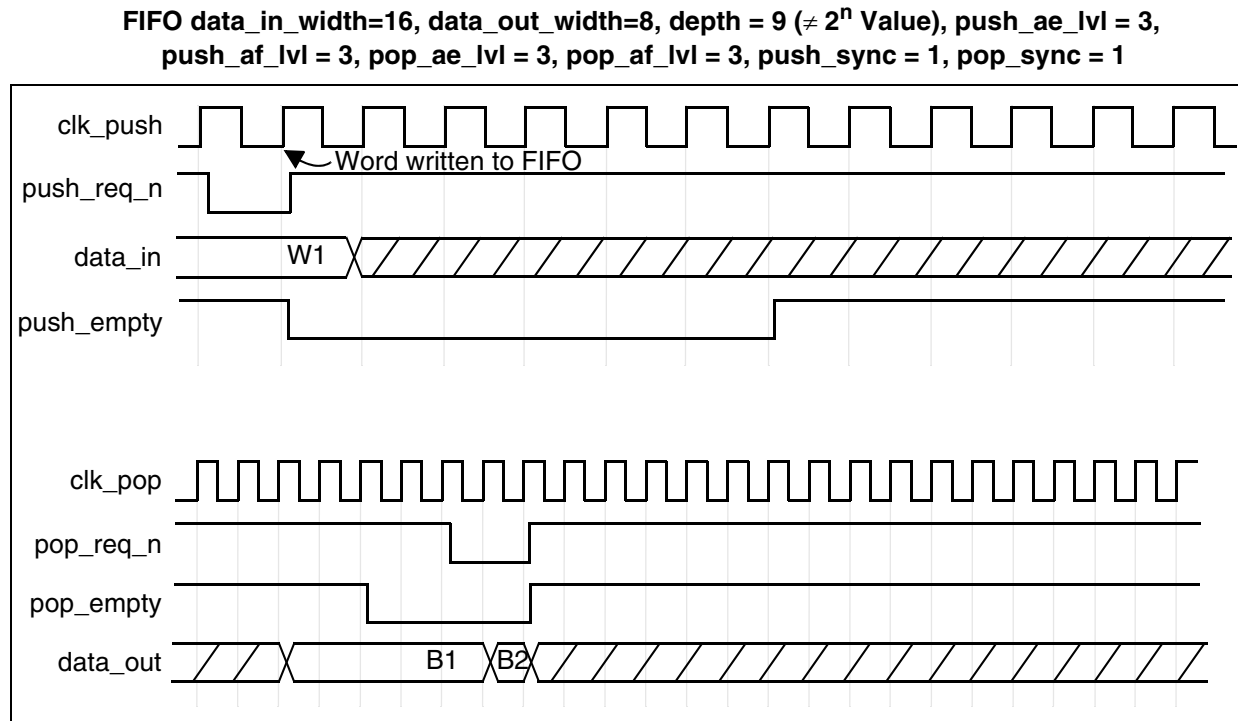


Figure 1-15 FIFO Depth = 2^n Single Word Timing Waveform for Input < Output

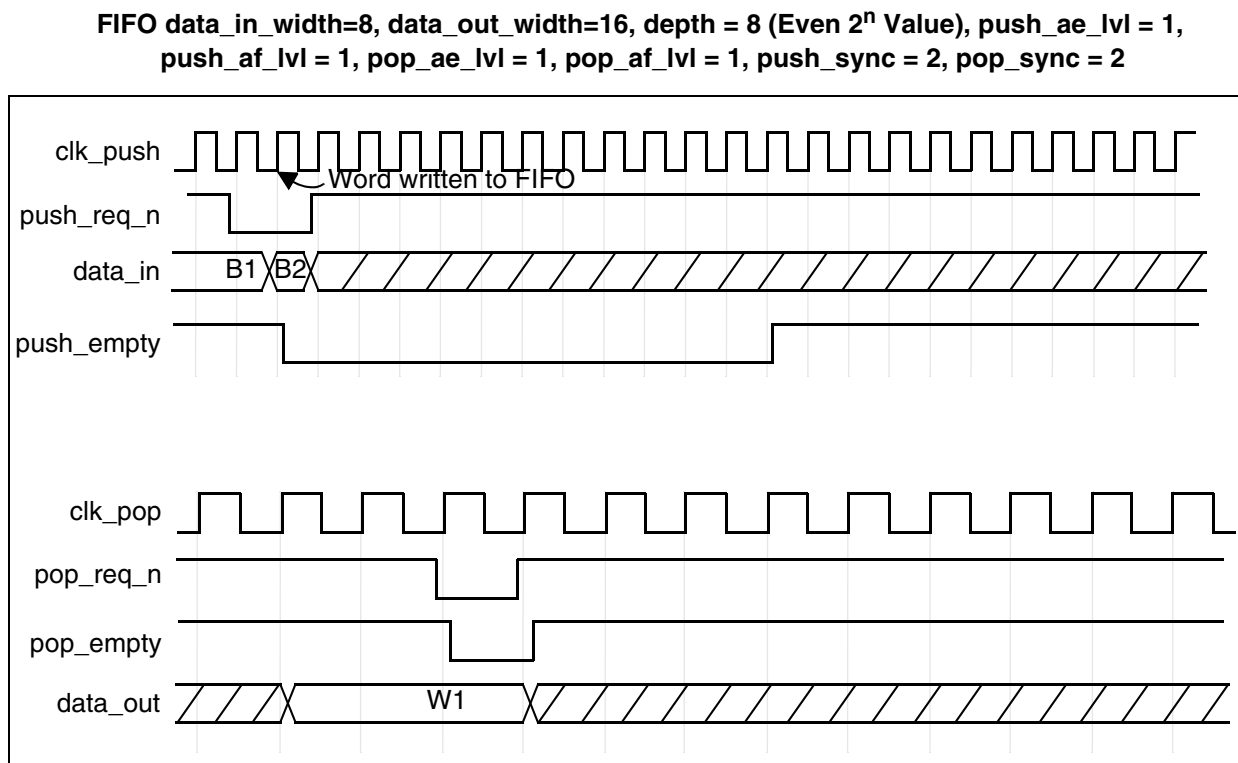
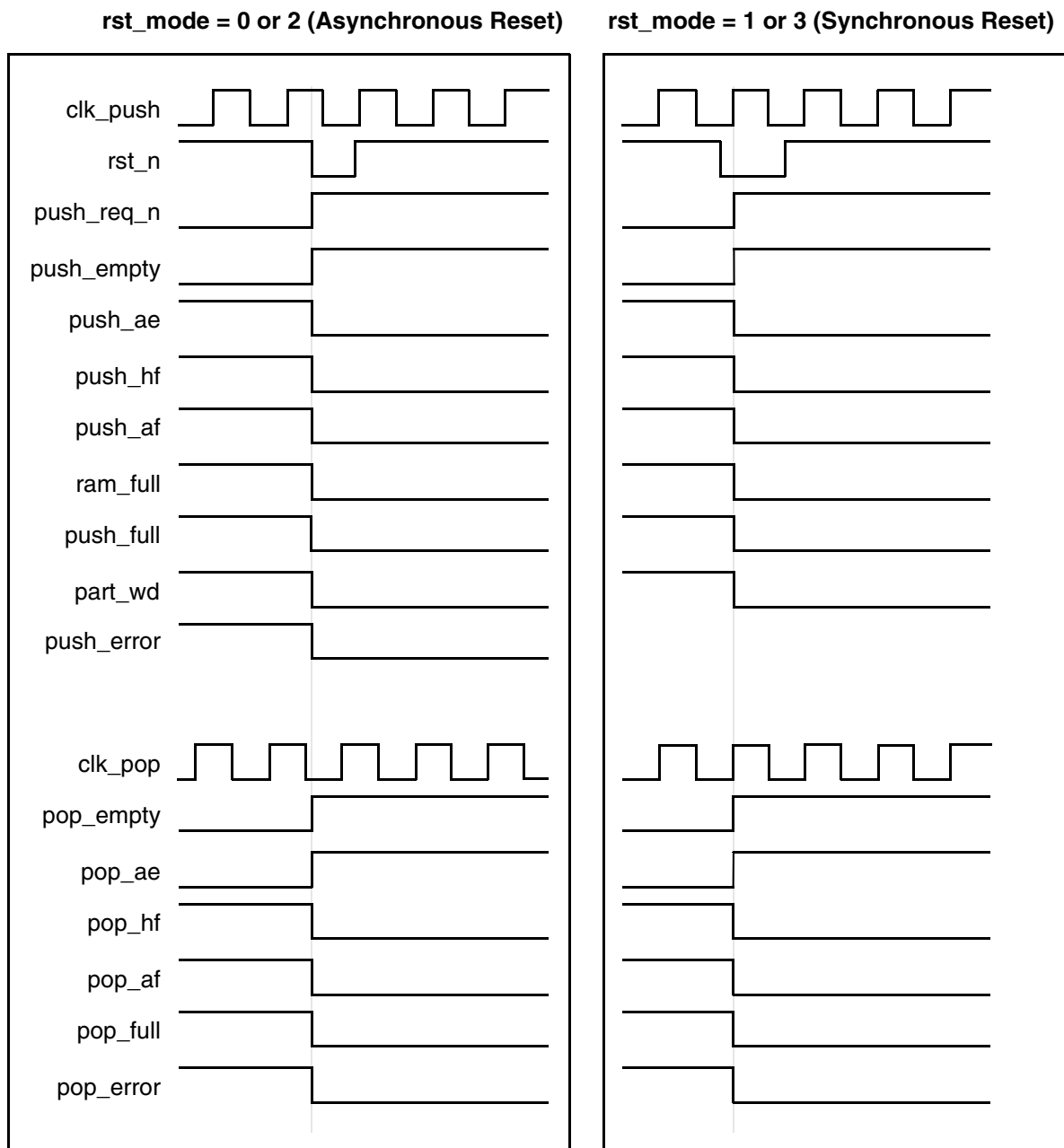


Figure 1-16 Reset Timing Waveforms

Related Topics

- [Memory - FIFO Overview](#)
- [DesignWare Building Block IP User Guide](#)

HDL Usage Through Component Instantiation - VHDL

```
library IEEE,DWARE,DW06;
use IEEE.std_logic_1164.all;
use DWARE.DW_Foundation_arith.all;
use DWARE.DW_Foundation_comp_arith.all;

entity DW_asymfifo_s2_sf_inst is
  generic (inst_data_in_width  : INTEGER := 24;
           inst_data_out_width : INTEGER := 8;
           inst_depth           : INTEGER := 8;
           inst_push_ae_lvl     : INTEGER := 2;
           inst_push_af_lvl     : INTEGER := 2;
           inst_pop_ae_lvl      : INTEGER := 2;
           inst_pop_af_lvl      : INTEGER := 2;
           inst_err_mode        : INTEGER := 0;
           inst_push_sync       : INTEGER := 2;
           inst_pop_sync        : INTEGER := 2;
           inst_rst_mode        : INTEGER := 1;
           inst_byte_order      : INTEGER := 0 );
  port (inst_clk_push   : in std_logic;
        inst_clk_pop    : in std_logic;
        inst_rst_n      : in std_logic;
        inst_push_req_n : in std_logic;
        inst_flush_n    : in std_logic;
        inst_pop_req_n  : in std_logic;
        inst_data_in    : in std_logic_vector(inst_data_in_width-1 downto 0);
        push_empty_inst : out std_logic;
        push_ae_inst    : out std_logic;
        push_hf_inst    : out std_logic;
        push_af_inst    : out std_logic;
        push_full_inst  : out std_logic;
        ram_full_inst   : out std_logic;
        part_wd_inst    : out std_logic;
        push_error_inst : out std_logic;
        pop_empty_inst  : out std_logic;
        pop_ae_inst     : out std_logic;
        pop_hf_inst     : out std_logic;
        pop_af_inst     : out std_logic;
        pop_full_inst   : out std_logic;
        pop_error_inst  : out std_logic;
        data_out_inst   : out std_logic_vector(inst_data_out_width-1 downto 0)
        );
end DW_asymfifo_s2_sf_inst;
```


architecture inst of DW_asymfifo_s2_sf_inst is
begin

-- Instance of DW_asymfifo_s2_sf

U1 : **DW_asymfifo_s2_sf**

```
generic map (data_in_width => inst_data_in_width,
             data_out_width => inst_data_out_width,
             depth => inst_depth,
             push_ae_lvl => inst_push_ae_lvl,
             push_af_lvl => inst_push_af_lvl,
             pop_ae_lvl => inst_pop_ae_lvl,
             pop_af_lvl => inst_pop_af_lvl,   err_mode => inst_err_mode,
             push_sync => inst_push_sync,   pop_sync => inst_pop_sync,
             rst_mode => inst_rst_mode,   byte_order => inst_byte_order )
```

```
port map (clk_push => inst_clk_push,   clk_pop => inst_clk_pop,
          rst_n => inst_rst_n,   push_req_n => inst_push_req_n,
          flush_n => inst_flush_n,   pop_req_n => inst_pop_req_n,
          data_in => inst_data_in,   push_empty => push_empty_inst,
          push_ae => push_ae_inst,   push_hf => push_hf_inst,
          push_af => push_af_inst,   push_full => push_full_inst,
          ram_full => ram_full_inst,   part_wd => part_wd_inst,
          push_error => push_error_inst,   pop_empty => pop_empty_inst,
          pop_ae => pop_ae_inst,   pop_hf => pop_hf_inst,
          pop_af => pop_af_inst,   pop_full => pop_full_inst,
          pop_error => pop_error_inst,   data_out => data_out_inst );
```

end inst;

-- pragma translate_off

configuration DW_asymfifo_s2_sf_inst_cfg_inst of DW_asymfifo_s2_sf_inst is
for inst

end for; -- inst

end DW_asymfifo_s2_sf_inst_cfg_inst;

-- pragma translate_on

HDL Usage Through Component Instantiation - Verilog

```

module DW_asymfifo_s2_sf_inst(inst_clk_push, inst_clk_pop, inst_rst_n,
                              inst_push_req_n, inst_flush_n, inst_pop_req_n,
                              inst_data_in, push_empty_inst, push_ae_inst,
                              push_hf_inst, push_af_inst, push_full_inst,
                              ram_full_inst, part_wd_inst, push_error_inst,
                              pop_empty_inst, pop_ae_inst, pop_hf_inst,
                              pop_af_inst, pop_full_inst, pop_error_inst,
                              data_out_inst );

parameter data_in_width = 24;
parameter data_out_width = 8;
parameter depth = 8;
parameter push_ae_lvl = 2;
parameter push_af_lvl = 2;
parameter pop_ae_lvl = 2;
parameter pop_af_lvl = 2;
parameter err_mode = 0;
parameter push_sync = 2;
parameter pop_sync = 2;
parameter rst_mode = 1;
parameter byte_order = 0;
`define bit_width_depth 3 // ceil(log2(depth))

input inst_clk_push;
input inst_clk_pop;
input inst_rst_n;
input inst_push_req_n;
input inst_flush_n;
input inst_pop_req_n;
input [data_in_width-1 : 0] inst_data_in;
output push_empty_inst;
output push_ae_inst;
output push_hf_inst;
output push_af_inst;
output push_full_inst;
output ram_full_inst;
output part_wd_inst;
output push_error_inst;
output pop_empty_inst;
output pop_ae_inst;
output pop_hf_inst;
output pop_af_inst;
output pop_full_inst;
output pop_error_inst;
output [data_out_width-1 : 0] data_out_inst;

// Instance of DW_asymfifo_s2_sf

```

```
DW_asymfifo_s2_sf #(data_in_width, data_out_width, depth,
                    push_ae_lvl, push_af_lvl, pop_ae_lvl,
                    pop_af_lvl, err_mode, push_sync,
                    pop_sync, rst_mode, byte_order)
U1 (.clk_push(inst_clk_push), .clk_pop(inst_clk_pop),
    .rst_n(inst_rst_n), .push_req_n(inst_push_req_n),
    .flush_n(inst_flush_n), .pop_req_n(inst_pop_req_n),
    .data_in(inst_data_in), .push_empty(push_empty_inst),
    .push_ae(push_ae_inst), .push_hf(push_hf_inst),
    .push_af(push_af_inst), .push_full(push_full_inst),
    .ram_full(ram_full_inst), .part_wd(part_wd_inst),
    .push_error(push_error_inst), .pop_empty(pop_empty_inst),
    .pop_ae(pop_ae_inst), .pop_hf(pop_hf_inst),
    .pop_af(pop_af_inst), .pop_full(pop_full_inst),
    .pop_error(pop_error_inst), .data_out(data_out_inst) );
endmodule
```

Revision History

For notes about this release, see the [DesignWare Building Block IP Release Notes](#).

For lists of both known and fixed issues for this component, refer to the [STAR report](#).

For a version of this datasheet with visible change bars, click [here](#).

Date	Release	Updates
July 2020	DWBB_201912.5	■ Adjusted content and title of “ Suppressing Warning Messages During Verilog Simulation ” on page 19 and added the DW_SUPPRESS_WARN macro
October 2019	DWBB_201903.5	■ Added the “Disabling Clock Monitor Messages” section
April 2018	DWBB_201709.5	■ For STAR 9001317257, updated the maximum value for the <i>data_in_width</i> , <i>data_out_width</i> , and <i>depth</i> parameters in Table 1-2 on page 3 and elsewhere
December 2017	DWBB_201709.2	■ Added “ Simulation Methodology ” on page 19 to explain how to simulate synchronization of Gray coded pointers between clock domains
October 2017	DWBB_201709.1	■ Replaced the synthesis implementations in Table 1-3 on page 4 with the str implementation ■ Added this Revision History table and the document links on this page

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