# **Data Privacy Homework 1**

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## 1. (15') Laplace Mechanism

- (a) (5') Given the function  $f(x) = \frac{1}{6} \sum_{i=1}^{6} x_i$ , where  $x_i \in \{1, 2, ..., 10\}$  for  $i \in \{1, 2, ..., 6\}$ , compute the global sensitivity and local sensitivity when  $x = \{3, 5, 4, 5, 6, 7\}$ .
- **(b) (10')** Given a database x where each element is in  $\{1, 2, 3, 4, 5, 6\}$ , design  $\varepsilon$ -differentially private Laplace mechanisms corresponding to the following queries, where  $\varepsilon = 0.1$ :
  - 1.  $q_1(x) = \sum_{i=1}^6 x_i$
  - 2.  $q_2(x) = \max_{i \in \{1, 2, \dots, 6\}} x_i$

# 2. (15') Exponential mechanism

ID	sex	Chinese	Mathematics	English	Physics	Chemistry	Biology
1	Male	96	58	80	53	56	100
2	Male	60	63	77	50	59	75
3	Female	83	86	98	69	80	100
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4000	Female	86	83	98	87	82	92

Table 1: Scores of students in School A

Table 1 records the scores of students in School A in the final exam. We need to help the teacher query the database while protecting the privacy of students' scores. The domain of this database is  $\{\text{Male}, \text{Female}\} \times \{0, 1, 2, ..., 100\}^6$ . Answer the following questions.

- (a) (5') What is the sensitivity of the following queries:
  - 1.  $q_1(x) = \frac{1}{4000} \sum_{\text{ID}=1}^{4000} \text{Physics}_{\text{ID}}$
  - 2.  $q_2(x) = \max_{\text{ID} \in \{1, 2, \dots, 4000\}} \text{Biology}_{\text{ID}}$
- **(b) (10')** Design  $\varepsilon$ -differentially privacy mechanisms corresponding to the two queries in (a), where  $\varepsilon = 0.1$ . (Using Laplace mechanism for  $q_1$  and exponential mechanism for  $q_2$ )

### 3. (20') Composition

**Theorem 3.16.** Let  $\mathcal{M}_i : \mathbb{N}^{|\mathcal{X}|} \to \mathcal{R}_i$  be an  $(\varepsilon_i, \delta_i)$ -differentially private algorithm for  $i \in [k]$ . Then if  $\mathcal{M}_{[k]}(x) = (\mathcal{M}_1(x), ..., \mathcal{M}_{k(x)})$ , then  $\mathcal{M}_k$  is  $(\sum_{i=1}^k \varepsilon_i, \sum_{i=1}^k \delta_i)$ -differentially private.

**Theorem 3.20 (Advanced Composition)**. For all  $\varepsilon$ ,  $\delta$ ,  $\delta$   $\geq$  0, the class of  $(\varepsilon, \delta)$ -differentially private mechanisms satisfies  $(\varepsilon, k\delta + \delta)$ -differential privacy under k-fold adaptive composition for:

$$\varepsilon' = \varepsilon \sqrt{2k \ln\left(\frac{1}{\delta'}\right)} + k\varepsilon(\varepsilon^{\varepsilon} - 1)$$

- (a) (10') Given a database  $x = \{x_1, x_2, ..., x_{2000}\}$  where  $x_i \in \{0, 1, 2, ..., 100\}$  for each i and privacy parameters  $(\varepsilon, \delta) = (1.25, 10^{-5})$ , apply the Gaussian mechanism to protect 100 calls to the query  $q_1(x) = \frac{1}{2000} \sum_{i=1}^{2000} x_i$ . Determine the noise variances  $\sigma^2$  of the Gaussian mechanism to ensure  $(\varepsilon, \delta)$ -DP based on the composition and advanced composition theorems, respectively.
- **(b) (10')** Determine the noise variances  $\sigma^2$  of the Gaussian mechanism to protect 100 calls to the query  $q_2(x) = \max_{i \in \{1,2,\dots,2000\}} x_i$  to ensure (1.25,  $10^{-5}$ )-DP based on the composition and advanced composition theorems, respectively, where x is the database in (a).

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#### 4. (25') Randomized Response for Local DP

Consider a population of n users, where the true proportion of males is denoted as  $\pi$ . Our objective is to gather statistics on the proportion of males, prompting a sensitive question: "Are you male?" Each user responds with either a yes or no, but due to privacy concerns, they refrain from directly disclosing their true gender. Instead, they employ a biased coin with a probability of landing heads denoted as p, and tails as 1-p. When the coin is tossed, a truthful response is given if heads appear, while the opposite response is provided if tails come up.

- (a) (10') Demonstrate that the aforementioned randomized response adheres to local differential privacy and determine the corresponding privacy parameter,  $\varepsilon$ .
- (b) (15') Employing the perturbation method outlined above to aggregate responses from the n users yields a statistical estimate for the number of males. Assuming the count of "yes" responses is  $n_1$ , construct an unbiased estimate  $\pi$  for based on n,  $n_1$ , p. Calculate the variance associated with this estimate.

#### 5. (10') Accuracy Guarantee of DP

Consider the application of an  $(\varepsilon, \delta)$ -differentially private Gaussian mechanism denoted by  $\mathcal{M}$  to protect the mean estimator  $\bar{x} = \frac{1}{n} \sum_{i=1}^n x_i$  of a d-dimensional input database x, where  $x_i \in \{0, 1, ..., 100\}^d$  for each i. Let  $\mathcal{M}(x)$  represent the output of this Gaussian mechanism. Utilize both the tail bound and the union bound to derive the  $L_{\infty}$ -norm error bound of  $\mathcal{M}$ , denoted by  $\|\mathcal{M}(x) - \bar{x}\|_{\infty}$ , ensuring a probability of at least  $1 - \beta$ . Specifically, solve for the bound  $\mathcal{B}$  such that

$$\Pr[\|\mathcal{M}(x) - \bar{x}\|_{\infty} \le \mathcal{B}] \ge 1 - \beta$$

Hint: Refer to Zhihu link for descriptions of statistical inequalities.

#### 6. (15') Personalized Differential Privacy

Consider an n-element dataset D where the i-th element is owned by a user  $i \in [n]$ , where  $[n] = \{1, 2, ..., n\}$  and the privacy requirement of user i is  $\varepsilon_i$ -DP. A randomized mechanism  $\mathcal{M}$  satisfies  $\{\varepsilon_i\}_{i\in[n]}$ -personalized differential privacy (or  $\{\varepsilon_i\}_{i\in[n]}$ -PDP) if, for every pair of neighboring datasets D, D differing at the j-th element for an arbitrary  $j \in [n]$ , and for all sets S of possible outputs,

$$\Pr[\mathcal{M}(D) \in S] \le \exp(\varepsilon_j) \Pr[\mathcal{M}(D') \in S].$$

- (a) (5') Prove the composition theorem of PDP: if a mechanism is  $\left\{\varepsilon_{i}^{(1)}\right\}_{i\in[n]}$ -PDP and another is  $\left\{\varepsilon_{i}^{(2)}\right\}_{i\in[n]}$ -PDP, then publishing the result of both is  $\left\{\varepsilon_{i}^{(1)}+\varepsilon_{i}^{(2)}\right\}_{i\in[n]}$ -PDP.
- **(b) (10')** Given a dataset D and a privacy requirement set  $\{\varepsilon_i\}_{i\in[n]}$ , the *Sample mechanism* works as follows:
  - 1) We pick an arbitrary threshold value t > 0;
  - 2) We sample a subset  $D_S \subset D$  where the probability that the *i*-th element of D is contained in  $D_S$  equals  $\frac{\exp(\varepsilon_i)-1}{\exp(t)-1}$  if  $\varepsilon_i < t$  and 1 otherwise;
  - 3) We output  $\mathcal{M}(D_S)$ , where  $\mathcal{M}$  is a *t*-differentially private mechanism.

Prove that the Sample mechanism with any t > 0 is  $\{\varepsilon_i\}(i \in [n])$ -PDP.

**Hint**: Use the Bayes formula.