

Lab 2

DTD validation

- Provide a DTD D such that the document t is valid with regard to D
- Write D as a regular grammar G
- Show that t is in L(G)
- Generate another non-trivial document that is in L(G)

D.dtd:

```
<?xml version="1.0" encoding="UTF-8"?>
<!ELEMENT a ((b|e)*)>
<!ELEMENT b ((c|f)*)>
<!ELEMENT e ((c|f|d)*)>
<!ELEMENT d (#PCDATA)>
<!ELEMENT f (#PCDATA)>
<!ELEMENT c (d?)>
```

Regular Grammar G

Reminder : $G = (N, T, S, P)$

where N = non-terminal, T = terminal, S = start, P = rules

- $N = (A, B, C, D, F)$
- $T = (a, b, e, c, f, d)$
- $S = (a)$
- $P =$
 - $A \rightarrow a[B, E]$
 - $B \rightarrow b[(C, F)^*]$
 - $C \rightarrow c[D?]$
 - $D \rightarrow d[\text{epsilon}]$
 - $E \rightarrow e[D, C^+]$
 - $F \rightarrow d[\text{epsilon}]$

In order to prove that t is in $L(G)$, we transform the grammar in automaton

- $a(q_b, q_e) \rightarrow q_a$
- $b((q_c, q_f)^*) \rightarrow q_b$
- $c(q_+) \rightarrow q_c$
- $d(\epsilon) \rightarrow q_D$
- $e(q_d, q_{c+}) \rightarrow q_e$
- $f(\epsilon) \rightarrow q_f$

$\Sigma = \{a, b, c, d, e, f\}$

$Q = \{q_a, q_b, q_c, q_d, q_e, q_f\}$

$F = \{q_a\}$

Another non-trivial example

```
<a>
  <b>
    <c>
      </c>
    </b>
  </a>
```

Properties of Regular Tree Grammars

1

- FAUX, student without address in G_1 cannot be in G_2

2

- TO G_1 no
- TO G_2 yes

```
<?xml version="1.0" encoding="UTF-8"?>
<!ELEMENT directory ((student|teacher)*)>
<!ELEMENT student (contact)>
<!ELEMENT teacher (contact)>
```

```
<!ELEMENT contact (Name,address)>
<!ELEMENT name (#PCDATA)>
<!ELEMENT address (#PCDATA)>
```

3

```
<directory>
  <student>
    <contact>
      <name> coucou </name>
      <address> 46 rue barrault </address>
    </contact>
  </student>
</directory>
```

XML Schema

regular tree Grammar

$R = (N, T, S, P)$

$N = (A, B, C, D, E, F, G, H, I, J, K, L, M)$

$T = (\text{shiporder}, \text{orderperson}, \text{shipto}, \text{name}, \text{address}, \text{country}, \text{item}, \text{title}, \text{note}, \text{quantity}, \text{price})$

$S = (A)$

$P =$

- $A \rightarrow \text{shiporder}[B, C, H, M]$
- $B \rightarrow \text{orderperson}[\text{epsilon}]$
- $C \rightarrow \text{shipto}[D, E, F, G]$
- $D \rightarrow \text{name}[\text{epsilon}]$
- $E \rightarrow \text{city}[\text{epsilon}]$
- $F \rightarrow \text{address}[\text{epsilon}]$
- $G \rightarrow \text{city}[\text{epsilon}]$
- $H \rightarrow \text{item}[I, J, K, L]$
- $I \rightarrow \text{title}[\text{epsilon}]$
- $J \rightarrow \text{note}[\text{epsilon}]$
- $K \rightarrow \text{qty}[\text{epsilon}]$
- $L \rightarrow \text{price}[\text{epsilon}]$

DTD and RTG

1

The grammar allows only patient* follow doctors * therefore there is a valid DTD D file that cannot be generated by the RTG G.

2

The opposite is true because of the flexibility of dtd syntax (correct even if there are more attributes)

3

G written in DTD

```
<!ELEMENT files ((file|patient|Doctor)*)>
<!ELEMENT Doctor (person)>
<!ELEMENT patient
<!ELEMENT person (FName,LName,Tel)>
<!ELEMENT file ((file|person|Doctor)*)>
```

THIS IS NOT POSSIBLE -> CANNOT DEFINE TWO PERSON FOR PATIENT AND DOCTOR

4

```
<Files>
  <file>
    <consultation>
      <symptom>
        data
      </symptom>
      <Prescription>
        <Medication>
          data
        </Medication>
      </Prescription>
    </consultation>
  </file>
</patient>
```

```
<person>
  <FName>
  </FName>
  <LName>
  </LName>
</person>
</patient>
<Doctor>
  <person>
    <FName>
    </FName>
    <LName>
    </LName>
    <Tel></Tel>
  </person>
</Doctor>
```

Properties of Regular Tree Grammars

Consider two regular tree grammars $G_1 = (N_1, T, S_1, P_1)$ and $G_2 = (N_2, T, S_2, P_2)$, having the same terminal symbols T .

Can you build the regular tree grammar $G_3 = (N_3, T, S_3, P_3)$ that captures the intersection of $L(G_1)$ and $L(G_2)$?

idea :

$$N_3 = N_1 * N_2$$

$$S_3 = S_1 * S_2$$

$P_3 : n \in N_3 \rightarrow (R_1, R_2) \text{ where}$

$R_1 \in P_1$

and

$R_2 \in P_2$