

$$T(n) = \begin{cases} 1, & n=1 \\ T(n-1) + n, & n \geq 2 \end{cases}$$

Supongamos que $n \geq 2$.

$$P_1 = T(n-1) + n$$

$$P_2 = [T(n-2) + n-1] + n = T(n-2) + n-1 + n$$

$$P_3 = [T(n-3) + n-2] + n-1 + n = T(n-3) + n-2 + n-1 + n$$

$$P_4 = [T(n-4) + n-3] + n-2 + n-1 + n$$

$$P_i = T(n-i) + \sum_{k=0}^{i-1} (n-k) = T(n-i) + \sum_{k=0}^{i-1} n - \sum_{k=0}^{i-1} k = T(n-i) + (i-1-0+1)n - \frac{i(i-1)}{2}$$

$$T(n-i) + i n - \frac{i^2-i}{2} = T(n-i) + i n - \frac{i^2}{2} + \frac{i}{2}$$

- Ahora debe encontrar el valor de i que haga $T(n)$:

$$n-i=1$$

$$n-1=i$$

- Reemplazamos el caso general:

$$T(n-(n-1)) + (n-1)n - \frac{(n-1)^2}{2} + \frac{(n-1)}{2} = T(1) + n^2 - n - \frac{(n^2-2n+1)}{2} + \frac{(n-1)}{2} =$$

$$T(1) + n^2 - n - \frac{n^2}{2} + n - \frac{1}{2} + \frac{n}{2} - \frac{1}{2} = T(1) + n^2 - \frac{n^2}{2} - 1 + \frac{n}{2} = 2 + \frac{n^2}{2} + \frac{n}{2} - 1$$

$$\boxed{\frac{1+n^2+n}{2}} \quad O(n^2/2)$$

- Existen constantes C y n_0 tales que: $1+n^2/2+n/2 \leq C \cdot n^2/2$

$$1 \leq C \cdot n^2/2 \quad \text{Si } n_0=1 \Rightarrow C=2 \checkmark$$

$$1 \leq 1 \cdot 1/2$$

$$1 \leq 1$$

$$n^2/2 \leq C \cdot n^2/2 \quad \text{Si } n_0=1 \Rightarrow C=1 \checkmark$$

$$1/2 \leq 1/2$$

$$n/2 \leq C \cdot n^2/2 \quad \text{Si } n_0=1 \Rightarrow C=1 \checkmark$$

$$1/2 \leq 1/2$$

$$\frac{1+n^2+n}{2} \leq (2+1+1) \cdot \frac{n^2}{2}$$

$$\frac{1+n^2+n}{2} \leq 4 \cdot \frac{n^2}{2} = O(n^2/2) \quad \text{¿Esta bien o es } O(n^2)?$$

$$b_2 \quad T(n) = \begin{cases} 2, & n=1 \\ T(n-1) + n/2, & n \geq 2 \end{cases}$$

1- Encontrar caso general

$$P_1 = T(n-1) + n/2$$

$$P_2 = [T(n-2) + (n-1)/2] + n/2 = T(n-2) + (n-1)/2 + n/2$$

$$P_3 = [T(n-3) + (n-2)/2] + (n-1)/2 + n/2 = T(n-3) + (n-2)/2 + (n-1)/2 + n/2$$

$$P_i = T(n-i) + \sum_{k=0}^{i-1} n-k = T(n-i) + \sum_{k=0}^{i-1} n - \sum_{k=1}^{i-1} k = T(n-i) + (i-1+1)n - \frac{(i-1)i}{2}$$

$$T(n-i) + i n - \left(\frac{i^2}{2} - \frac{i}{2} \right) = T(n-i) + i n - \frac{i^2}{2} + \frac{i}{2}$$

$$n-i=1$$

$$n-1=i$$

$$T(n-(n-1)) + (n-1)n - \frac{(n-1)^2}{2} + \frac{(n-1)}{2} = T(1) + n^2 - n - \frac{(n^2 - 2n + 1)}{2} + \frac{(n-1)}{2}$$

$$2 + n^2 - n - \frac{n^2}{2} + n - \frac{1}{2} + \frac{n}{2} - \frac{1}{2} = 2 - 1 + \frac{n^2}{2} + n - n + \frac{n}{2} = 1 + \frac{n^2}{2} + \frac{n}{2}$$

$$C = \begin{cases} 1 & n=1 \\ 2T(n/4) + \sqrt{n} & n \geq 2 \end{cases}$$

$$P_1 = 2T(n/4) + \sqrt{n}$$

$$P_2 = 2[2T(n/4/4) + \sqrt{n/4}] + \sqrt{n} = 4T(n/16) + 2\sqrt{n/4} + \sqrt{n}$$

$$P_3 = 4[2T(n/16/4) + \sqrt{n/16}] + 2\sqrt{n/4} + \sqrt{n} = 8T(n/64) + 4\sqrt{n/16} + 2\sqrt{n/4} + \sqrt{n}$$

$$P_i = 2^i T(n/4^i) + \sum_{k=0}^{i-1} 2^k \sqrt{n/4^k} = 2^i T(n/4^i) + \sum_{k=0}^{i-1} 2^k \cdot \sqrt{n}/(2^k)$$

$$P_i = 2^i T(n/4^i) + \sum_{k=0}^{i-1} 2^k \cdot \frac{1}{2^k} \cdot \sqrt{n} = 2^i T(n/4^i) + \sum_{k=0}^{i-1} \sqrt{n} = 2^i T(n/4^i) + i \cdot \sqrt{n}$$

$$n/4^i = 1$$

$$n = 4^i$$

$$\log_4(n) = \log_4(4^i)$$

$$\log_4(n) = i \cdot \log_4(4)$$

$$\log_4(n) = i$$

$$2^{\log_4(n)} T(n/4^{\log_4(n)}) + \log_4(n) \cdot \sqrt{n} = \sqrt{n} \cdot T(n/n) + \log_4(n) \cdot \sqrt{n}$$

$$\sqrt{n} \cdot T(1) + \log_4(n) \cdot \sqrt{n} = \sqrt{n} \cdot 1 + \log_4(n) \cdot \sqrt{n} = \sqrt{n} + \log_4(n) \cdot \sqrt{n}$$

El orden de la función $\sqrt{n} + \log_4(n) \cdot \sqrt{n}$ es $\log_4(n) \cdot \sqrt{n}$

Existen constantes $C \geq 1$ y n_0 tales que:

$$\sqrt{n} \leq C \cdot \log_4(n) \cdot \sqrt{n} \quad \text{si } n \geq n_0$$

$$\frac{\sqrt{n}}{\sqrt{n}} \leq C \cdot \log_4(n)$$

$$1 \leq 1 \cdot \log_4(4) \quad \text{con } n_0 = 4, C = 1$$

$$\log_4(n) \cdot \sqrt{n} \leq C \cdot \log_4(n) \cdot \sqrt{n} \quad \text{con } n_0 = 4, C = 1$$

$$\log_4(4) \cdot \sqrt{1} \leq 1 \cdot \log_4(4) \cdot \sqrt{1}$$

$$1 \leq 1$$

Por lo tanto:

$$\sqrt{n} + \log_4(n) \cdot \sqrt{n} \leq (1+1) \cdot \log_4(n) \cdot \sqrt{n}$$

$$n \leq 2 \cdot \log_4(n) \cdot \sqrt{n} \quad \text{Para todo } C \geq 2 \text{ y } n_0 \geq 4$$

$$d- \quad T(n) = \begin{cases} 1 & n=1 \\ 4T(n/2) + n^2 & n \geq 2 \end{cases}$$

$$P_1 = 4T(n/2) + n^2$$

$$P_2 = 4[4T(n/4) + (n/2)^2] + n^2 = 16T(n/4) + 4\left(\frac{n}{2}\right)^2 + n^2 = 16T\left(\frac{n}{4}\right) + 2n^2$$

$$P_3 = 16[4T(n/8) + (n/4)^2] + 2n^2 = 64T(n/8) + 16\left(\frac{n}{4}\right)^2 + 2n^2 = 64T\left(\frac{n}{8}\right) + 3n^2$$

$$P_i = 4^i T(n/2^i) + i n^2$$

$$n/2^i = 1$$

$$n = 2^i$$

$$\log_2(n) = \log_2(2^i)$$

$$\log_2(n) = i$$

$$4^{\log_2(n)} T(n/2^{\log_2(n)}) + \log_2(n) \cdot n^2 = (n^2) T(1) + \log_2(n) \cdot n^2 = n^2 + \log_2 n$$

$$O(n^2)$$

Existen constantes $c > 0$ tales que:

$$n^2 \leq c n^2 \quad n_0 = 1 \quad c = 1$$

$$1 \leq 1$$

$$\log_2(n) \leq c \cdot n^2 \quad n_0 = 1 \quad c = 1$$

$$0 \leq 1$$

Entonces

$$n^2 + \log_2(n) \leq c(n)^2$$

$$n^2 \leq (1+1)n^2$$

$$n^2 \leq 2 \cdot n^2 \quad \text{Para todo } n_0 \geq 1 \Rightarrow c \geq 2$$