
LTP > ☒ Exámenes

Exámenes

Self-Assessment Test Theme 2

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Parte 1 de 3 - Second

0.67/ 2.0 Puntos

Preguntas 1 de 10

1.0/ 1.0 Puntos. Puntos descontados por fallo: 0.33

Consider the following definition of an *identifier* using BNF notation

```
<digito> ::= <par> | <impar>
<par>  ::= 0 | 2 | 4 | 6 | 8
<impar> ::= 1 | 3 | 5 | 7 | 9
<letra> ::= x <letra> | y <par> | z <impar>
<identificador> ::= <letra> <digito>
```

Which of the following identifiers is **NOT** legal in the language defined by the previous rules?

- ☐ A. xxy00
- ☐ B. xxy01
- ☒ C. z00
- ☐ D. z11

Preguntas 2 de 10

-0.33/ 1.0 Puntos. Puntos descontados por fallo: 0.33

The following declaration of variables of a program in C

```
itn a
```

contains an error produced by misspelling the name of the default type for integers. This error is detected by the

- ☐ A. scanner
- ☐ B. parser
- ☒ C. semantic analyzer
- ☐ D. linker

Parte 2 de 3 - First

3.5/ 3.5 Puntos

Preguntas 3 de 10

0.5/ 0.5 Puntos. Puntos descontados por fallo: 0.33

Two programs **P1** and **P2** are *equivalent* ($P1 \approx P2$) if they have the same semantics. Considering the *big step* operational semantics as a basis to define such an equivalence, which of the following equivalence statements is **TRUE**?

- ☒ A. $(x := 5; \text{while } x > 2 \text{ do } x := x - 1) \approx (x := 1; x := x + x)$
- ☐ B. $(P1 \not\approx P2)$, siendo

P1: $x := 0; y := 2;$
 if false then $x := y$ else $y := x;$
 $x := y;$

P2: $x := 0; y := 2;$
 if false then $x := y$ else $\{y := x; x := y\};$

- ☐ C. $(x := 1; \text{while } x < 2 \text{ do } x := x + 3) \approx (x := 1)$
- ☐ D. $(x := 1; y := 3; \text{if } x \leq y \text{ then } x := 0 \text{ else } x := 1) \approx (x := 1; y := 3)$

Preguntas 4 de 10

0.5/ 0.5 Puntos. Puntos descontados por fallo: 0.33

Complete the following sentence:

"The following axiomatic semantic rule:

$$\text{wp}(X := \text{exp}; Q) = Q[X \mapsto \text{exp}]$$

establishes that the weakest precondition of the assignment instruction $X := \text{exp}$ with respect to the postcondition is obtained by ..."

- ☐ A. replacing exp in Q by X .
- ☐ B. adding the expression $X \mapsto \text{exp}$ to Q .
- ☐ C. removing the expression $X \mapsto \text{exp}$ from Q .
- ☒ D. replacing every occurrence of X in Q by exp .

Preguntas 5 de 10

0.5/ 0.5 Puntos. Puntos descontados por fallo: 0.33

Which configuration is required (in $*$) to make the following evaluation complete by using the small-step operational semantics?

$\langle \text{if } X > Y \text{ then } Y := Y + X \text{ else } Y := 0, \{X \mapsto 42, Y \mapsto 0\} \rangle$

$\langle X > Y, \{X \mapsto 42, Y \mapsto 0\} \rangle \Rightarrow \text{true}$

$\langle X, \{X \mapsto 42, Y \mapsto 0\} \rangle \Rightarrow 42$

$\langle Y, \{X \mapsto 42, Y \mapsto 0\} \rangle \Rightarrow 0$

$\rightarrow \langle Y := Y + X, \{X \mapsto 42, Y \mapsto 0\} \rangle$

$\langle Y + X, \{X \mapsto 42, Y \mapsto 0\} \rangle \Rightarrow 42$

$\langle Y, \{X \mapsto 42, Y \mapsto 0\} \rangle \Rightarrow 0$

$\langle X, \{X \mapsto 42, Y \mapsto 0\} \rangle \Rightarrow 42$

$\rightarrow (*$

- ☐ A. $\langle Y := 0, \{X \mapsto 42, Y \mapsto 0\} \rangle$
- ☐ B. $\langle \text{skip}, \{X \mapsto 42, Y \mapsto 0\} \rangle$
- ☐ C. $\langle \text{if } X > Y \text{ then } Y := Y + X \text{ else } Y := 0, \{X \mapsto 42, Y \mapsto 0\} \rangle$
- ☒ D. $\langle \text{skip}, \{X \mapsto 42, Y \mapsto 42\} \rangle$

Preguntas 6 de 10

0.5/ 0.5 Puntos. Puntos descontados por fallo: 0.33

Consider the following transition rules (for assignment and sequencing) of the small- step operational semantics of the small imperative language IMP studied in the course

$$\frac{}{\langle a, e \rangle \Rightarrow n}$$

$$\langle X := a, e \rangle \rightarrow \langle \text{skip}, e[X \mapsto n] \rangle$$

$$\frac{\langle i0, e \rangle \rightarrow \langle i'0, e' \rangle}{\langle (i0; i1), e \rangle \rightarrow \langle (i'0; i1), e' \rangle} \quad \frac{}{\langle (\text{skip}; i), e \rangle \rightarrow \langle i, e' \rangle}$$

Fill the gap below for an appropriate continuation of the following execution trace:

$$\langle x := x + y; y := x + 1, \{x \mapsto 1, y \mapsto 1\} \rangle$$

$$\langle x := x + y, \{x \mapsto 1, y \mapsto 1\} \rangle$$

$$\langle x + y, \{x \mapsto 1, y \mapsto 1\} \rangle \Rightarrow 2$$

$$\rightarrow \langle \text{skip}, \{x \mapsto 2, y \mapsto 1\} \rangle$$

$$\rightarrow \langle \text{skip}; y := x + 1, \{x \mapsto 2, y \mapsto 1\} \rangle$$

$$\rightarrow \langle y := x + 1, \{x \mapsto 2, y \mapsto 1\} \rangle$$

?

- ☐ A. $\rightarrow \langle \text{skip}, \{x \mapsto 2, y \mapsto 1\} \rangle$
- ☒ B. $\langle x + 1, \{x \mapsto 2, y \mapsto 1\} \rangle \Rightarrow 3$
 $\rightarrow \langle \text{skip}, \{x \mapsto 2, y \mapsto 3\} \rangle$
- ☐ C. $\{x \mapsto 2, y \mapsto 1\} \Rightarrow 3$
 $\rightarrow \langle \{x \mapsto 2, y \mapsto 3\} \rangle$
- ☐ D. $\langle x + 1, \{x \mapsto 2, y \mapsto 1\} \rangle \Rightarrow 3$

Preguntas 7 de 10

0.5/ 0.5 Puntos. Puntos descontados por fallo: 0.33

The big-step operational semantics

- ☒

A.

establishes a relationship between the initial configuration $\langle P, e_0 \rangle$ of the program and the final state e_f after the execution.

- ☐

B.

yields the sequence of configurations (*trace*) which are obtained during the step-by-step execution of the program.

- ☐ C. can be used to establish the correctness of the program by means of the weakest precondition calculus.
- ☐ D. cannot be used to guide the implementation of programming languages.

Preguntas 8 de 10

0.5/ 0.5 Puntos. Puntos descontados por fallo: 0.33

Given a semantics S for a programming language, we say that two programs $i1$ and $i2$ are equivalent (written $i1 \approx i2$) if they have the same semantics. If we consider the *big step* operational semantics for SIMP, which of the following program equivalence statements is **WRONG**?

- ☐ A. $(y:=y) \approx \text{skip}$

(donde skip es la instrucción vacía)

- ☒ B. $(x:=y; y:=x) \approx (y:=x; x:=y)$
- ☐ C. $(x:=y; y:=x) \approx x:=y$
- ☐ D. $(\text{if } x \geq 0 \text{ then } x:=x-x \text{ else } x:=x*0) \approx x:=0$

Preguntas 9 de 10

0.5/ 0.5 Puntos. Puntos descontados por fallo: 0.33

Supposing that we extend the syntax of SIMP with the instruction $:=$ and add the following small-step semantic rule:

$\frac{\langle X, e \rangle \Rightarrow n_0 \quad \langle Y, e \rangle \Rightarrow n_1}{\langle X := Y, e \rangle \rightarrow \langle \text{skip}, e[X \mapsto n_1][Y \mapsto n_0] \rangle}$

$\langle X := Y, e \rangle \rightarrow \langle \text{skip}, e[X \mapsto n_1][Y \mapsto n_0] \rangle$

Which of the following sentences defines its meaning:

- ☐ A. It's a regular assignment but restricted to two variables in such a way that X is assigned the value of Y.
- ☐ B. Both variables X and Y are assigned an initial value.
- ☒ C. It exchanges the values of variables X and Y.
- ☐ D. Checks whether the values of variables X and Y are the same or not.

Parte 3 de 3 - Third

1.0/ 1.0 Puntos

Preguntas 10 de 10

1.0/ 1.0 Puntos. Puntos descontados por fallo: 0.33

Which of the following sentences is **TRUE**

- ☐ A.
In pure compiled languages each instruction is simultaneously analyzed and executed, so that programs are more efficient.
- ☐ B. Pure interpreted languages are high-level languages which are translated to intermediate code.
- ☒ C.
In programming languages with a mixed implementation scheme the original language is translated into an intermediate language which is then interpreted.
- ☐ D.
In pure programming languages the original language is translated into an object language which is then interpreted.

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