

Exámenes

Self-Assessment Test Theme 2

[Volver a la Lista de Exámenes](#)

Parte 1 de 3 - Second

2.0/ 2.0 Puntos

Preguntas 1 de 10

1.0/ 1.0 Puntos. Puntos descontados por fallo: 0.33

Given a *wrong* program sentence like:

```
if (x>0) else x:=x+1 then x:=x-1
```

which of the compiling phases detects its incorrectness?

- ☐ A. Lexical analysis.
- ☒ B. Syntactic analysis.
- ☐ C. Semantic analysis.
- ☐ D. Static analysis.

Preguntas 2 de 10

1.0/ 1.0 Puntos. Puntos descontados por fallo: 0.33

Which of the following statements is TRUE?

- ☐ A. The static semantics analyzes the most stable part of the code.
- ☒ B. Errors due to type incompatibilities are detected during the semantic analysis.
- ☐ C. Syntactic errors in programs are detected during the linking phase.
- ☐ D.

The static semantics detects all errors in compilation time; hence the dynamic semantics is executed with no error.

Parte 2 de 3 - First

1.84/ 3.5 Puntos

Preguntas 3 de 10

-0.33/ 0.5 Puntos. Puntos descontados por fallo: 0.33

Given a Hoare triple $\{P\} S \{Q\}$:

- ☐ A. P and Q are programs and S is a machine state.
- ☐ B. P is called the precondition, Q the postcondition, and $P \Rightarrow Q$ always holds.
- ☐ C. The correctness of the triple is guaranteed if $P \Rightarrow \text{pmd}(S, Q)$ holds.
- ☒ D. The correctness of the triple is guaranteed if $\text{pmd}(S, Q) \Rightarrow P$ holds.

Preguntas 4 de 10

0.5/ 0.5 Puntos. Puntos descontados por fallo: 0.33

Consider the following axiomatic definitions for the assignment and conditional instructions

$$\text{wp}(X := \text{exp}, Q) = Q[X \mapsto \text{exp}]$$
$$\text{wp}(\text{if } B \text{ then } i1 \text{ else } i2, Q) = (B \wedge \text{wp}(i1, Q)) \vee (\neg(B) \wedge \text{wp}(i2, Q))$$

Which of the following expressions corresponds to the weakest precondition P of the following program with respect to the postcondition Q

 $\{P\} = \{?\}$ if $x < 0$ then $y := -x$ else $y := x + y$ $\{Q\} = \{y = 0\}$

- ☒ A. $P = (x \geq 0 \wedge x + y = 0)$
- ☐ B. $P = (x \geq 0 \wedge y \geq 0)$
- ☐ C. $P = (x \geq 0 \wedge x = y)$
- ☐ D. $P = (x = 0 \wedge x = y)$

Preguntas 5 de 10

-0.33/ 0.5 Puntos. Puntos descontados por fallo: 0.33

Which of the following semantic descriptions of a programming language is the most useful in compiler design?

- ☐ A. Axiomatic semantics.
 - ☐ B. Hoare's triples.
 - ☒ C. Static semantics.
 - ☐ D. Operational semantics.
-

Preguntas 6 de 10

0.5/ 0.5 Puntos. Puntos descontados por fallo: 0.33

Mark the kind of semantics definition to which next rule belongs:

$$\frac{\langle i_0, e \rangle \Downarrow e'' \wedge \langle i_1, e'' \rangle \Downarrow e'}{\langle i_0; i_1, e \rangle \Downarrow e'}$$

$$\frac{\langle i_0, e \rangle \Downarrow e'' \wedge \langle i_1, e'' \rangle \Downarrow e'}{\langle i_0; i_1, e \rangle \Downarrow e'}$$

- ☐ A. Axiomatic semantics.
- ☒ B. Big-step operational semantics.
- ☐ C. Small-step operational semantics.
- ☐ D. Its own semantics.

Preguntas 7 de 10

0.5/ 0.5 Puntos. Puntos descontados por fallo: 0.33

Which is the function of the following operation $X \vdash a$ defined by the following semantic rule:

$$\frac{\langle X, e \rangle \Rightarrow n_0 \quad \langle a, e \rangle \Rightarrow n_1}{\langle X \vdash a, e \rangle \rightarrow \langle \text{skip}, e[X \mapsto n_0 + n_1] \rangle}$$

- ☐ A. The value of expression a is assigned two times to the variable X .
- ☒ B. The value of variable X is incremented by the value of expression a .
- ☐ C. Variable X receives the double of the value of expression a .
- ☐ D. Variable X receives the value of expression a .

Preguntas 8 de 10

0.5/ 0.5 Puntos. Puntos descontados por fallo: 0.33

Which configuration is required (in *) to make the following evaluation complete by using the small-step operational semantics?

$\langle \text{if } X > Y \text{ then } Y := Y + X \text{ else } Y := 0, \{X \mapsto 42, Y \mapsto 0\} \rangle$
 $\langle X > Y, \{X \mapsto 42, Y \mapsto 0\} \rangle \Rightarrow \text{true}$
 $\langle X, \{X \mapsto 42, Y \mapsto 0\} \rangle \Rightarrow 42$
 $\langle Y, \{X \mapsto 42, Y \mapsto 0\} \rangle \Rightarrow 0$
 $\rightarrow \langle Y := Y + X, \{X \mapsto 42, Y \mapsto 0\} \rangle$
 $\langle Y + X, \{X \mapsto 42, Y \mapsto 0\} \rangle \Rightarrow 42$
 $\langle Y, \{X \mapsto 42, Y \mapsto 0\} \rangle \Rightarrow 0$
 $\langle X, \{X \mapsto 42, Y \mapsto 0\} \rangle \Rightarrow 42$
 $\rightarrow (*)$

- ☐ A. $\langle Y := 0, \{X \mapsto 42, Y \mapsto 0\} \rangle$
- ☐ B. $\langle \text{skip}, \{X \mapsto 42, Y \mapsto 0\} \rangle$
- ☐ C. $\langle \text{if } X > Y \text{ then } Y := Y + X \text{ else } Y := 0, \{X \mapsto 42, Y \mapsto 0\} \rangle$
- ☒ D. $\langle \text{skip}, \{X \mapsto 42, Y \mapsto 42\} \rangle$

Preguntas 9 de 10

0.5/ 0.5 Puntos. Puntos descontados por fallo: 0.33

Given the following transition rules for the small-step semantics associated to the conditional instruction

$\frac{}{\langle b, e \rangle \Rightarrow \text{true}} \quad \frac{}{\langle b, e \rangle \Rightarrow \text{false}}$

$\langle \text{if } b \text{ then } i0 \text{ else } i1, e \rangle \rightarrow \langle i0, e \rangle \quad \langle \text{if } b \text{ then } i0 \text{ else } i1, e \rangle \rightarrow \langle i1, e \rangle$

determine which is the next configuration for $\langle \text{if } X < 4 \text{ then } X := X - 1 \text{ else } X := X + 1, \{X \mapsto 3\} \rangle$

- ☐ A. $\langle \text{skip}, \{X \mapsto 2\} \rangle$
- ☐ B. $\langle \text{skip}, \{X \mapsto 4\} \rangle$
- ☒ C. $\langle X := X - 1, \{X \mapsto 3\} \rangle$
- ☐ D. $\langle X := X + 1, \{X \mapsto 3\} \rangle$

Preguntas 10 de 10

-0.33/ 1.0 Puntos. Puntos descontados por fallo: 0.33

Indicate which of the following sentences is **FALSE**:

- ☐ A. The mix implementation of a language combines the advantages of compilation and interpretation.
- ☐ B. Interpreters are mainly used in script languages such as Perl or Postscript.
- ☒ C. Intermediate code is generated by a compiler (e.g. P-code) providing portability to any platform.
- ☐ D. Usually, an interpreter generates intermediate code that it is compiled into machine code.

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