

Fundamentos de los Sistemas Operativos (FSO)

Departamento de Informática de Sistemas y Computadoras (DISCA)

Universitat Politècnica de València

Part 2: Process Management

Unit 3

Process concept and implementation

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- **Goals**

- To define **process** concept
- To analyze the differences between **sequential** and **concurrent** execution
- To define the **process states** as well as the causes of the transitions between them
- To study the **basic structures to support processes** in the OS
- To analyze the **context switch** mechanism

- **Bibliography**

- “Operating System Concepts” Silberschatz 8ª Ed
- “Sistemas operativos: una visión aplicada” Carretero 2º Ed

- **Previous concepts**
- Executable files
- Process concept
- Process states
- Process implementation: PCB
- System calls table for processes and signals
- Exercises

- In order to understand the process concept some terms have to be clarified, particularly the differences between the following pairs of concepts:
 - Program / Process
 - Sequential execution / Concurrent execution

Program vs. Process

- **Program:**

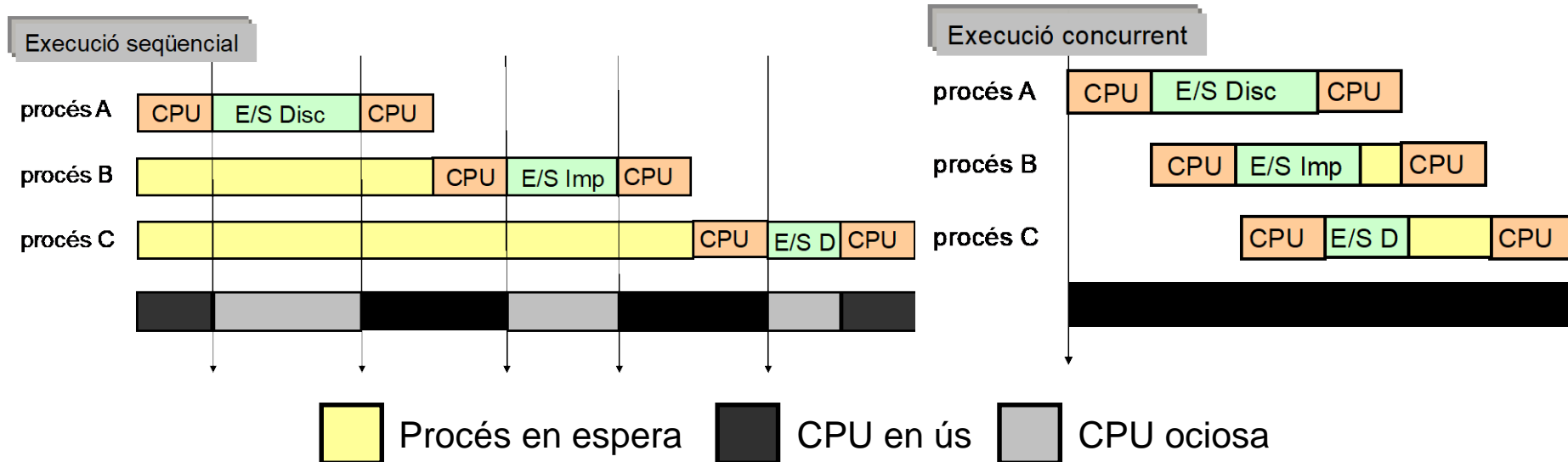
- **File** in executable format generated by a compilation of source code followed by a linking operation.
- **Passive entity** it doesn't change along time.

- **Process:** a certain program running:

- Working unit for the OS
- Resource consumer
- Every activity happening on a computer is related to a process
- **Abstract entity belonging to the OS** that allows modeling computer activity

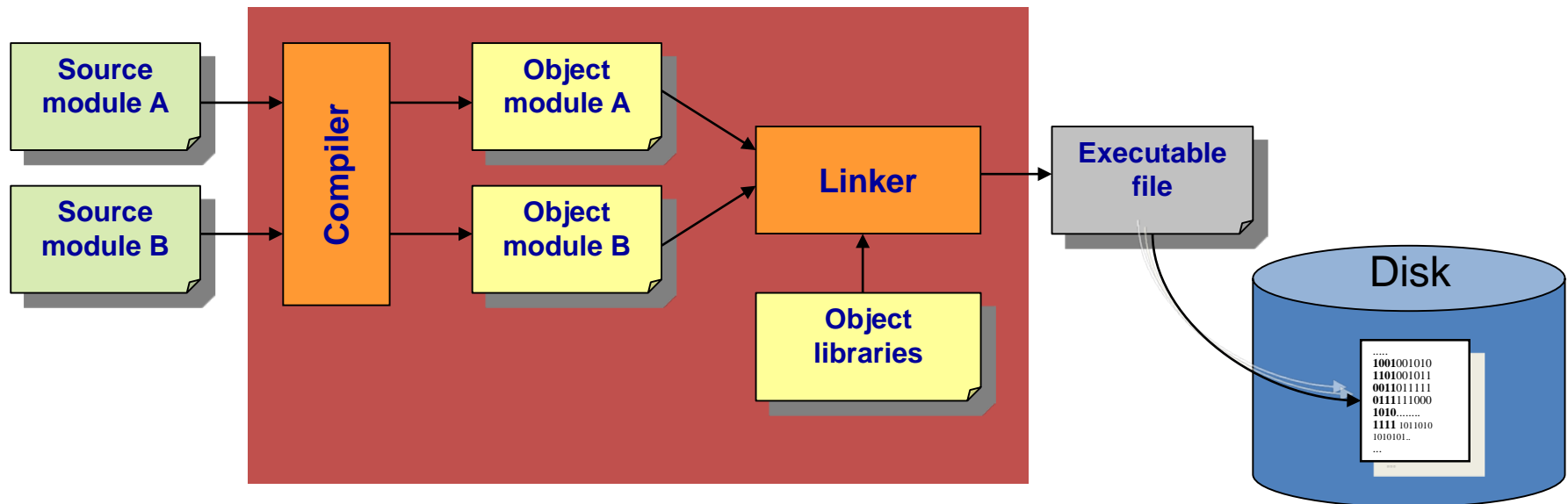
A process is an alive entity that experiments changes while it exists

- **Sequential execution:** when the CPU is used by a single process from the beginning till the end of its execution
- **Concurrent execution:** when the CPU is used alternatively by several processes during their lifetime
 - The most direct benefit of concurrent execution is the increase in CPU utilization, since processes are not constantly demanding CPU because they alternate CPU with I/O bursts

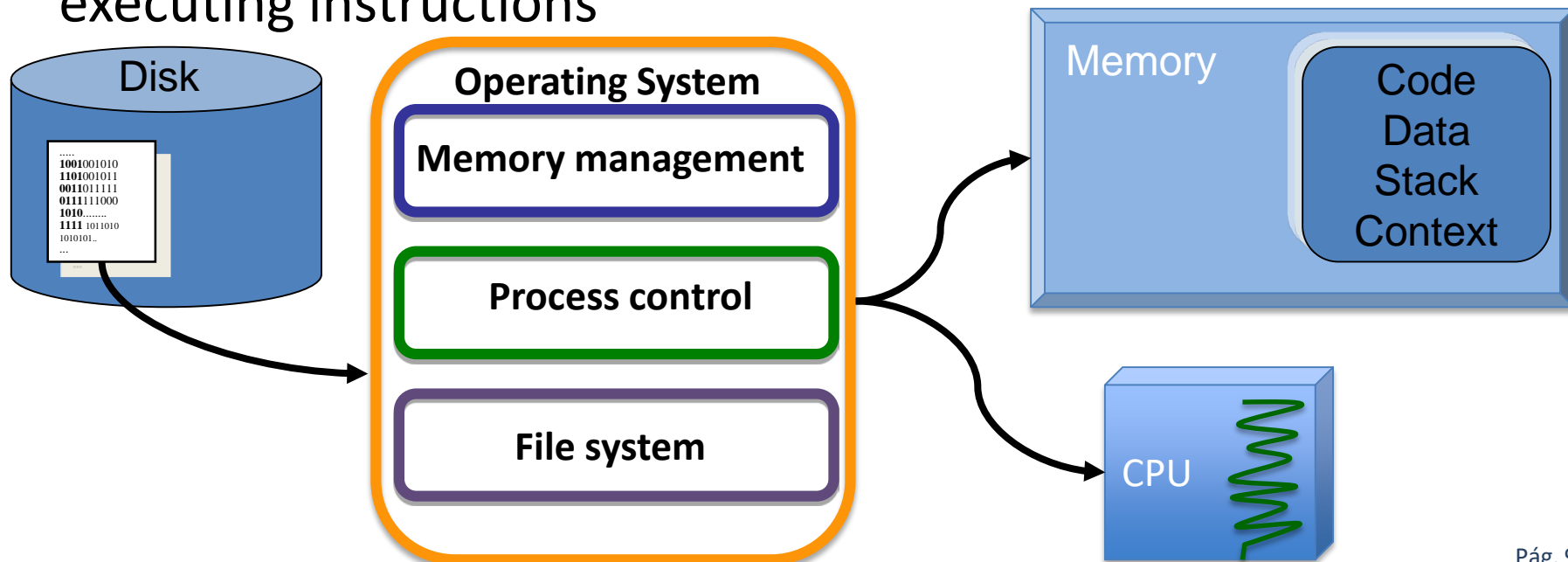


- Previous concepts
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- Steps to get an executable file:
 - The first step is **to compile** the **source code** to get **object code**
 - Later it is **linked** to system or other user **library code** -> this allows incorporating **external code**
 - The final result is an **executable file**

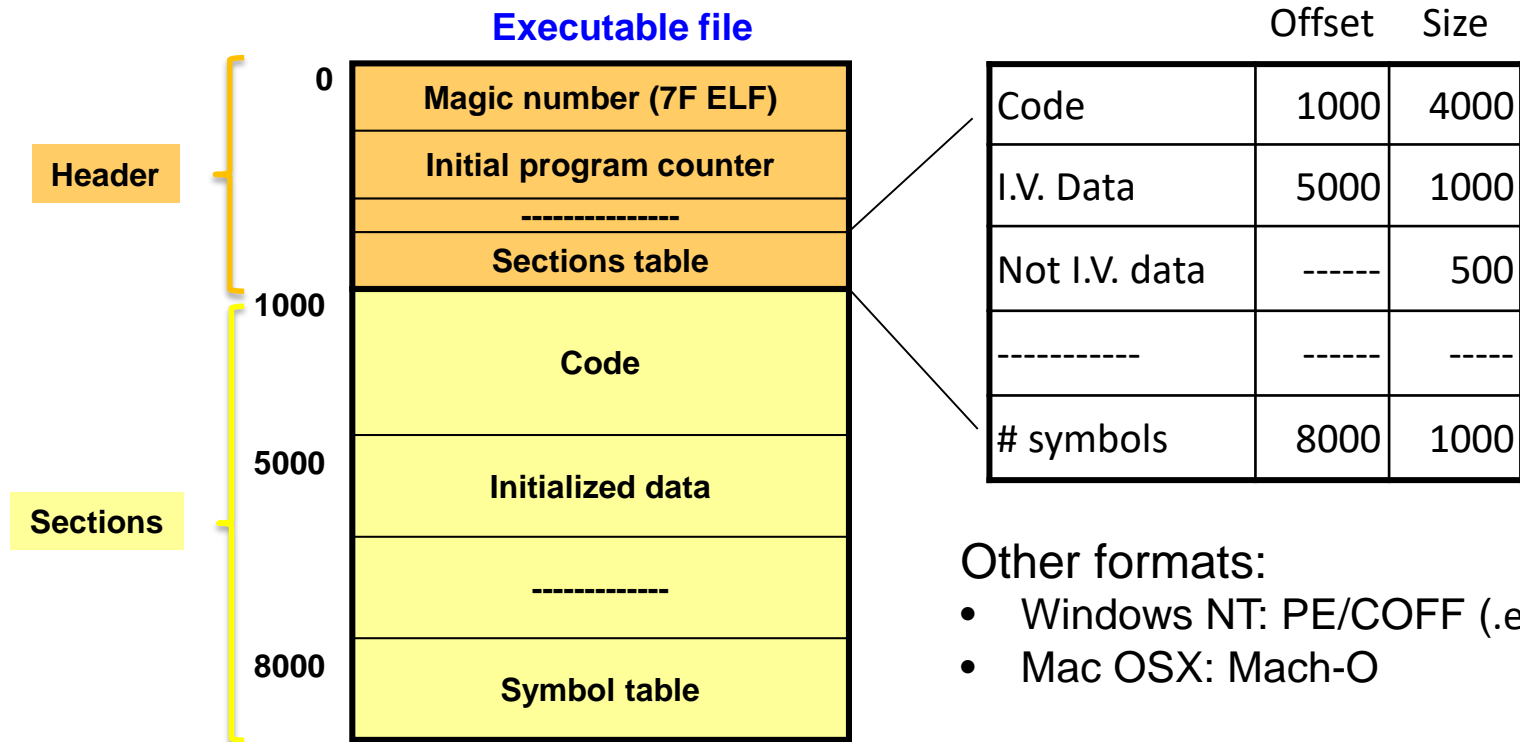


- An executable is a file whose structure is well known by the operating system because it contains:
 - The code to be executed
 - Initialized data
 - Library functions (static linking)
- With this information the OS can allocate the necessary resources for execution. Load data, code etc. and start executing instructions



- **Executable and Linking Format**

- File format for executables, object code, libraries and memory dumps
 - Widely used in Unix, Linux, Solaris and BSD
 - Extensions: .o, .so, .elf



Other formats:

- Windows NT: PE/COFF (.exe, .dll)
- Mac OSX: Mach-O

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- In order to define a process there are three aspects to be considered:
 - **Attributes** or features that define a process and allow to manage them
 - **Behaviour**, as being an active entity it will move between states so these states and their transitions have to be defined
 - **Operations** that can be done with processes

- **Process attributes**

Resources and features owned by a process and kept by the OS

- **Identity**

- Process identifier (PID)

- **Runtime environment**

- Working directory
- Opened file descriptors

- **State**

- Process state
- Machine context (program counter, stack pointer, general purpose register values)
- Etc.

- **Memory**

- Memory addresses for code, data and stack

- **Scheduling**

- CPU consumption time
- Priority

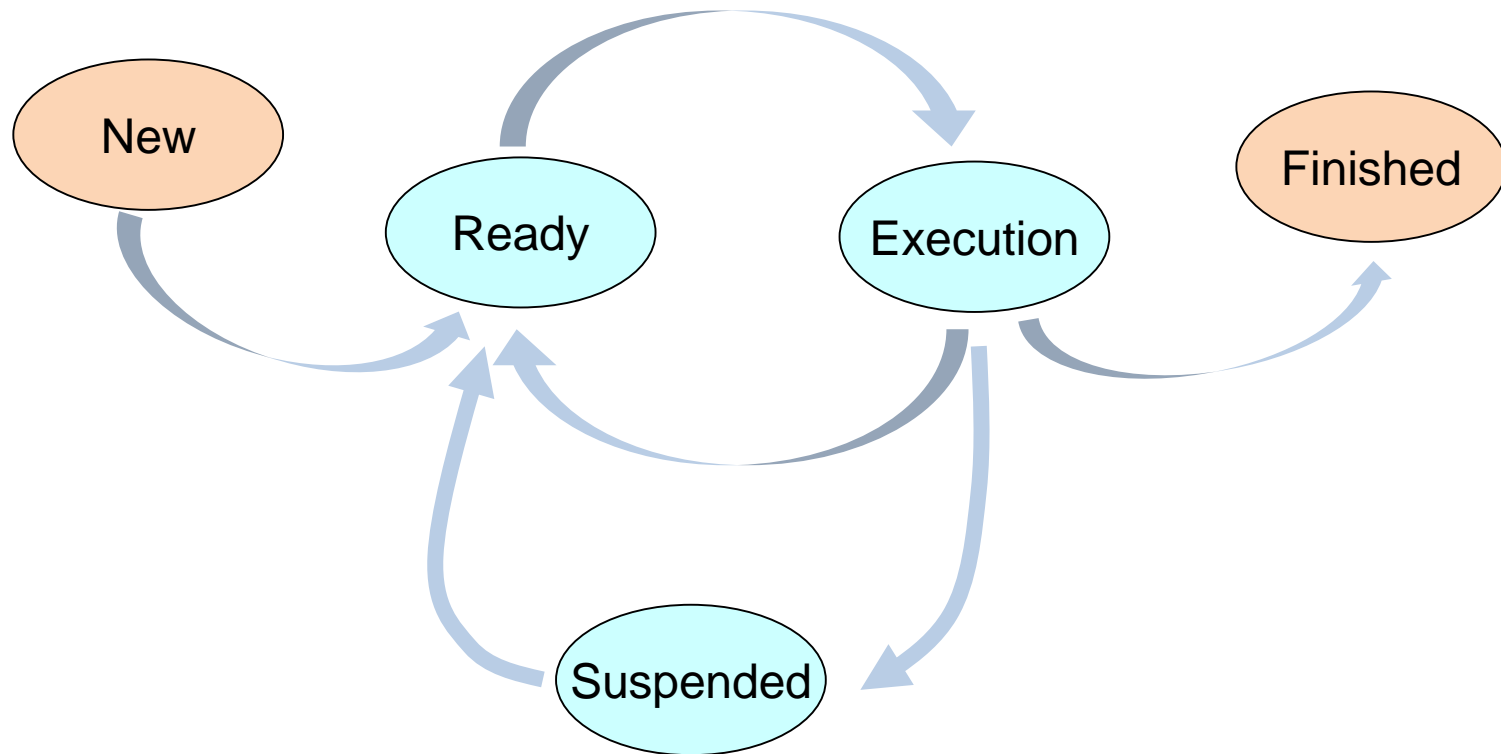
- **Monitoring**

- **Operations on processes**

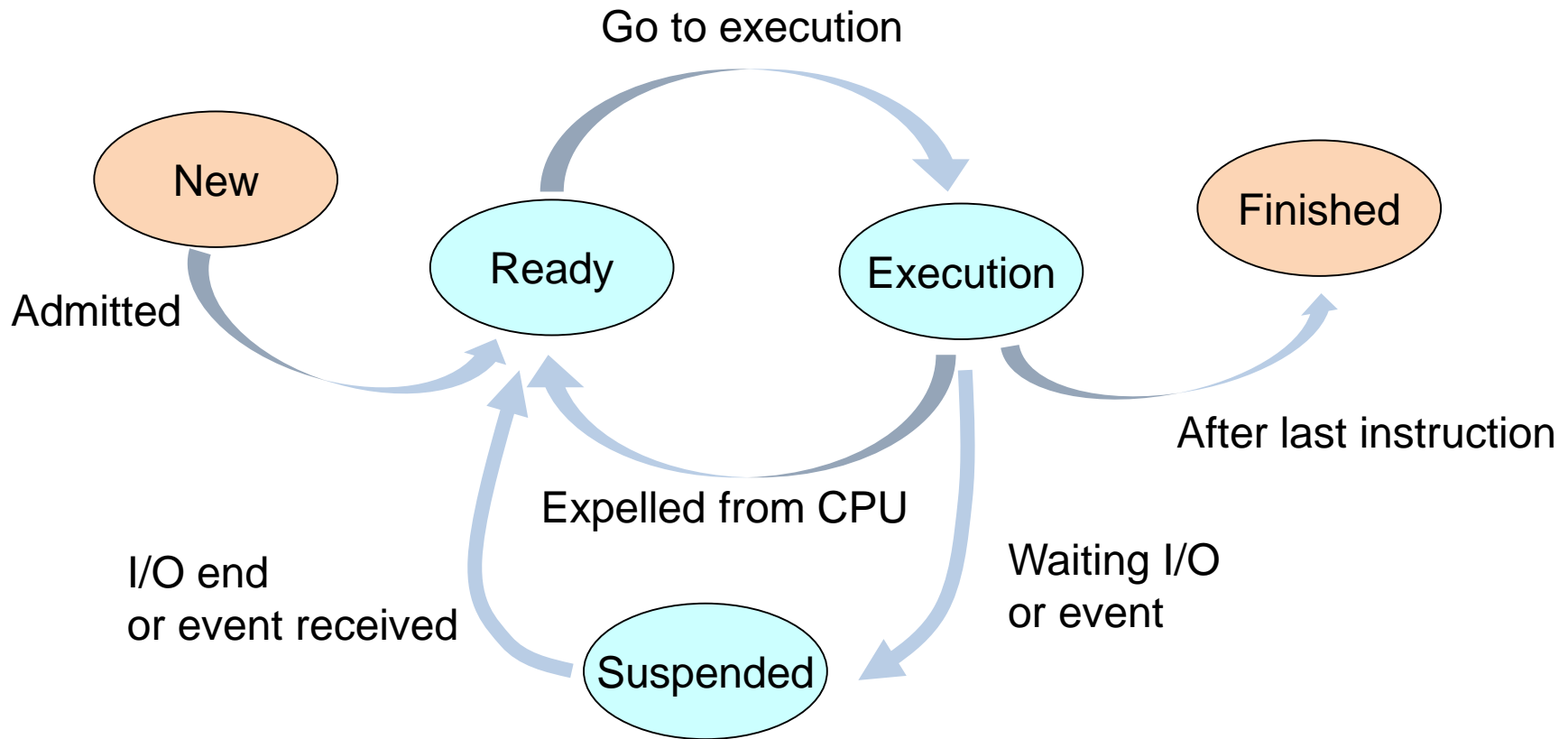
- As happens with the attributes, the number and the type of operations that can be done on processes **depend on the OS considered**
- Anyway, the following operations are always supported:
 - Creation
 - Communication
 - Waiting
 - Resource access
 - Ending

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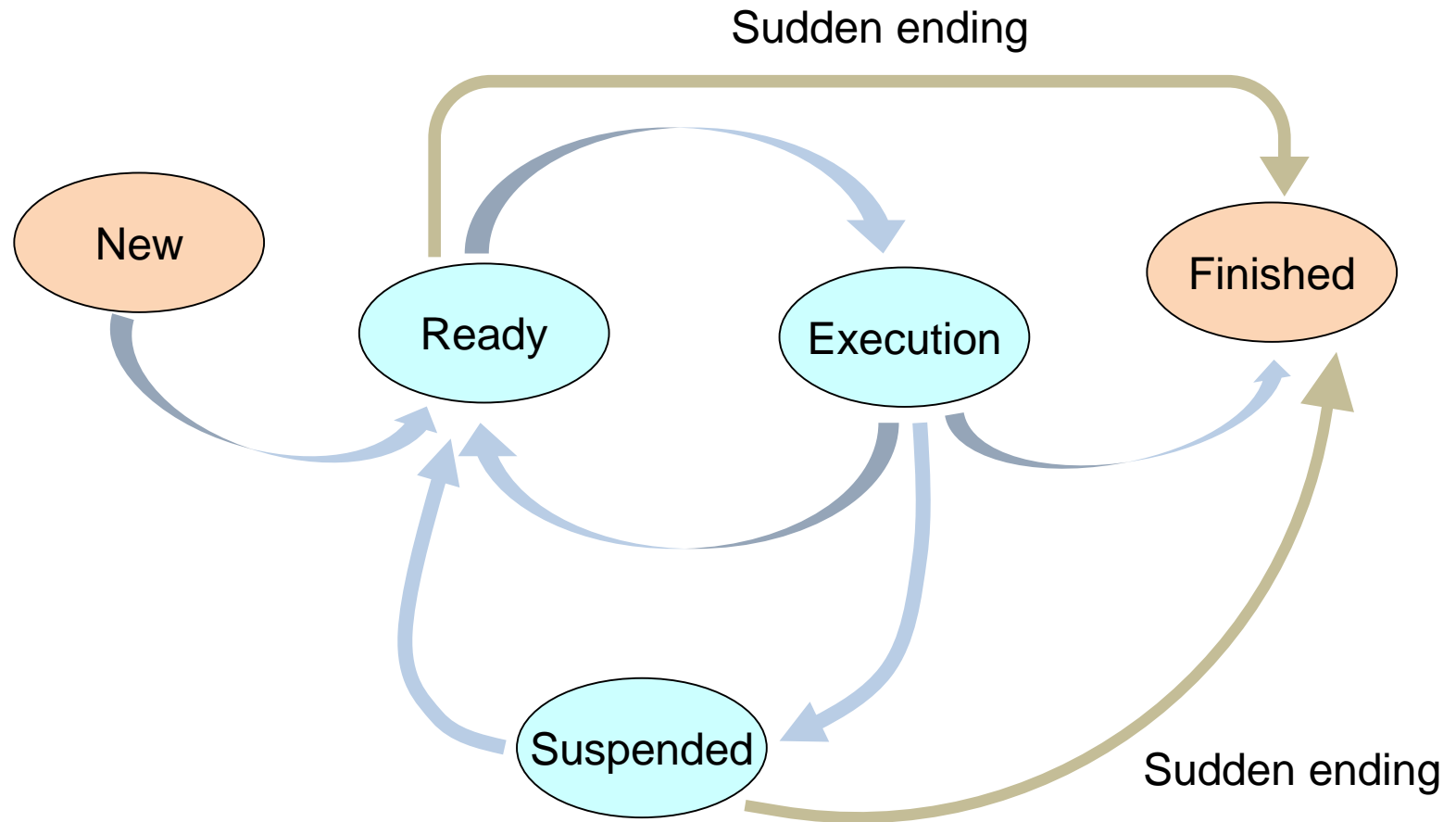
- **Processes** are **active entities** so they evolve during his lifetime through a **set of states**

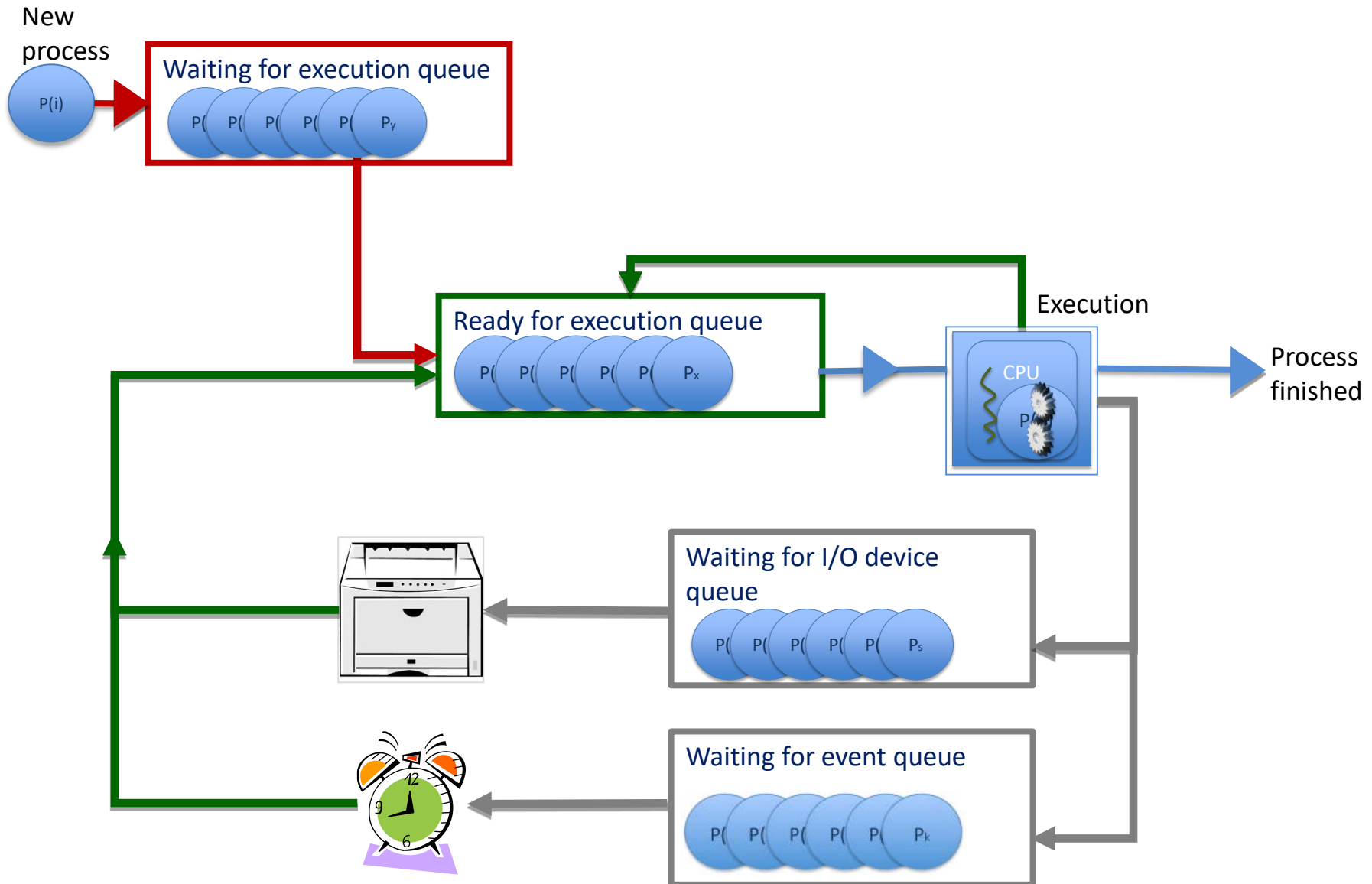


- Common process state transitions



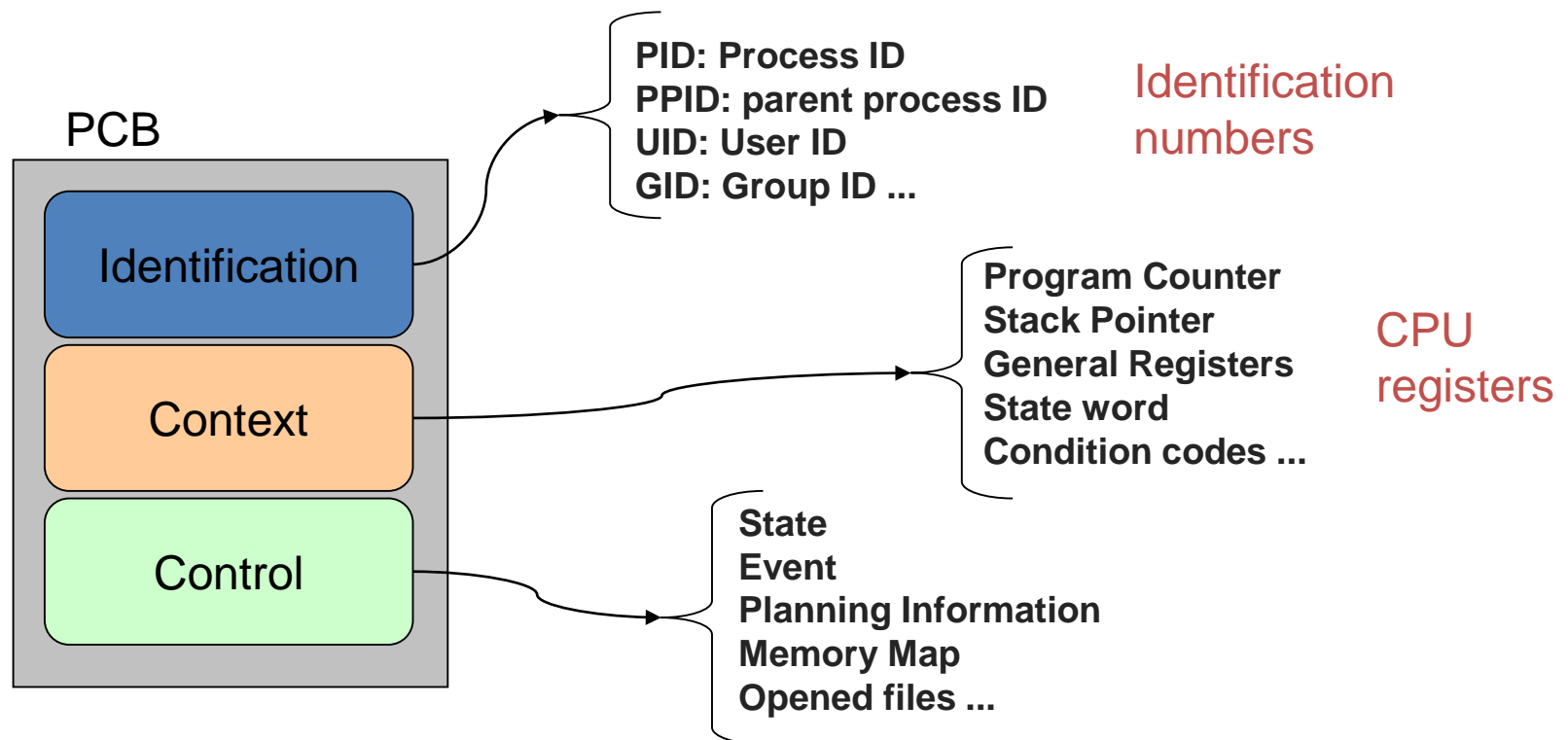
- **Transitions due to anomalous behavior**

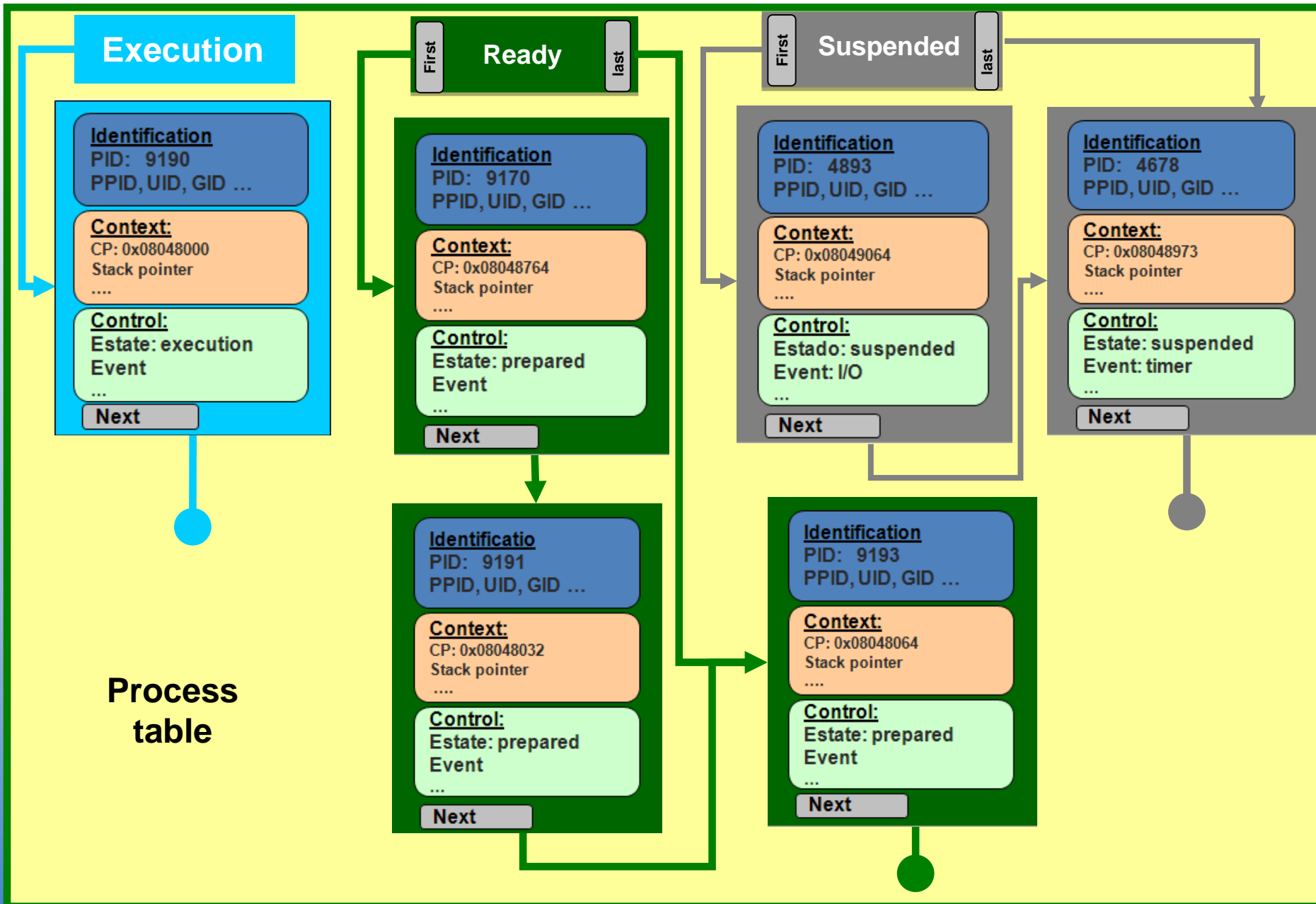




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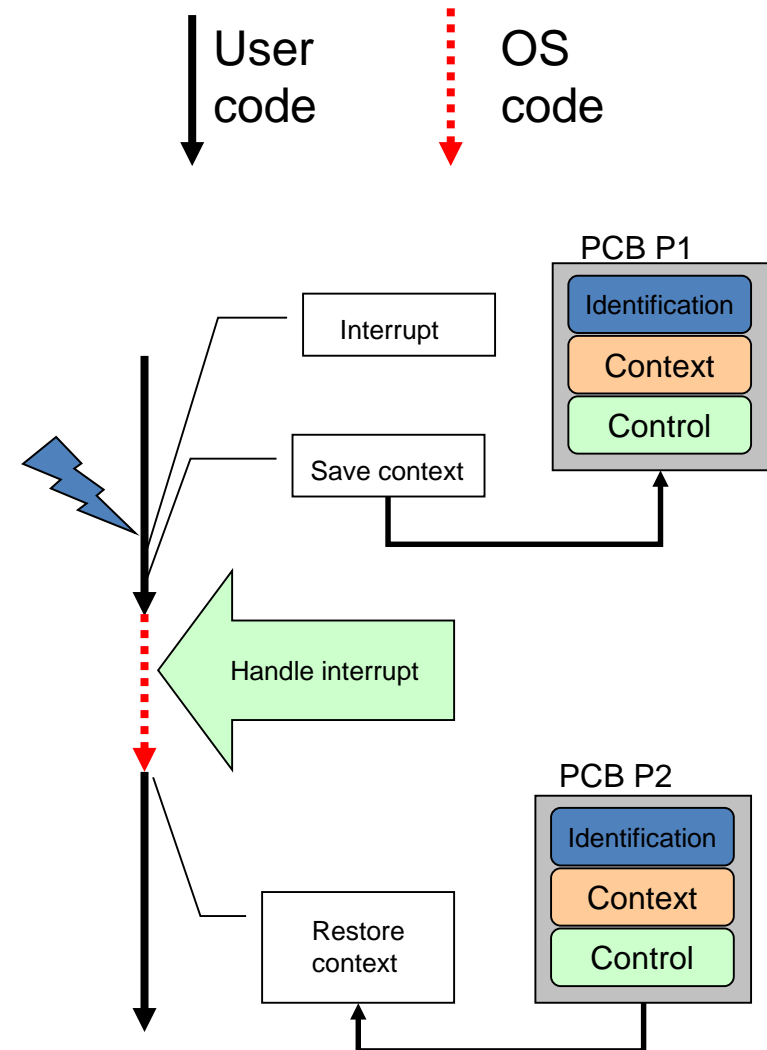
- A PCB (Process Context Block) is the **data structure** which supports the concept (abstraction) process
 - An OS is also a program, based on using **algorithms** and **data structures**
- PCB keeps process relevant information which changes during process lifetime -> each operating system has its own structure





- **Context switch**

- Mechanism that allows an OS to suspend the execution of the current process in order to start or resume the execution of another process.
 - This process is activated by an interrupt (i.e. clock interrupt)
- What is done?
 1. The state information relevant to the running process (context) is saved
 2. The scheduler updates PCBs and queues
 3. The new process takes the CPU



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System calls table for processes and signals fso

	Processes
fork	Create a child process
exit	End process
wait	Wait for a process ending
exec	Execute a program
getpid	Get process attributes

	Signals
kill	Send signals
alarm	Generate an alarm (clock signal)
sigemptyset	Init a mask with no signals set
sigfillset	Init a mask with signals set
sigaddset	Append a signal to a signal set
sigdelset	Delete a signals in a signal set
sigismember	Check if a signals belongs to a signal set
sigprocmask	Check/Modify/Set a signal mask
sigaction	Capture/Manage a signal
sigsuspend	Wait signals capture

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- In the following actions list, indicate from each one if it is shell code or OS code the one that performs it. Mark one, none or both options with a cross:

OS	Shell	
		Reading the command line and parse it
		Programming a device controller
		Providing a system calls interface
		Selecccting a process to assign it the CPU
		Performing a system call
		Providing a confortable user interface

- Which state (new, ready, execution, suspended, finished) corresponds to every one of the following processes:

Process		State
P1	The CPU is executing instructions belonging to P1	
P2	P2 has asked for a disk access, but the disk is busy serving to process P3	
P3	P3 is doing disk access	
P4	P4 belongs to a user that has finished all his/her jobs in a terminal and is logging out the session	
P5	P5 has a process identifier assigned and only its control tables are already created	
P6	P6 has all its SO control structures and its memory image stored in main memory	

- In the following command line:

```
$ cat f1 f2 f3 | grep start | wc -l >tracd
```

Indicate:

- How many processes will be created during its execution in a UNIX system
- What files has associated every command

- Given the following code:

```
#include <stdio.h>
#include <sys/types.h>

int main(void)
{
    pid_t pid;
    int i;

    for (i=0; i<2; i++)
        pid=fork();
    return 0;
}
```

How many processes will be created along its execution?