

CSE331: Automata and Computability
Practice Sheet
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Question 1: Designing Context-Free Grammars

Give a Context-free Grammar that generates the language

- a) $L = \{w \in \{0,1\}^* : w \text{ starts with '0' and the length of } w \text{ is even.}\}$
- b) $L = \{w \in \{a,b\}^* : \text{every second letter in } w \text{ is a 'b'.}\}$
- c) $L = \{w \in \{a,b\}^* : \text{the length of } w \text{ is divisible by three}\}$
- d) $L = \{w \in \{0,1\}^* : w = 0^n 1^{2n+3}, n \geq 0\}$
- e) $L = \{w \in \{0,1\}^* : w = 0^n 1^n, \text{ where } n \text{ is odd.}\}$
- f) $L = \{w \in \{a,b\}^* : \text{the number of 'a' is at least the number of 'b' in } w.\}$
 [Hint: <https://www.youtube.com/watch?v=QF3lTaM56lo&list=PLBENQsMXh3gz85EJ3ZCSa9l9hnUiOer-H&index=35>]
- g) $L = \{w \in \{0,1,2\}^* : w = 0^i 1^j 2^k \text{ where } j \geq 2i + 3k, \text{ and } i, k \geq 0\}$
- h) $L = \{w \in \{0,1,2,3\}^* : w = 0^i 1^j 2^k 3^m \text{ where } i=m, j \geq 3k+2, \text{ and } m, k \geq 0\}$
- i) $L = \{w \in \{0,1,2\}^* : w = 0^i 1^j 2^k \text{ where } i > 2j + 3k, \text{ and } j, k \geq 0\}$
- j) If $A = \{w \in \{0, 1\}^* : w \text{ contains at least two 0s}\}$, then construct $L = \{w \in \{0, 1\}^* : w = 0^{3i} v 1^{2i} \text{ where } v \in A \text{ and } i \geq 0\}$
- k) Recall that for a string w , $|w|$ denotes the length of w . $\Sigma = \{0,1\}$
 $L1 = \{w \in \Sigma^* : w \text{ contains } 11\}$
 $L2 = \{x\#y : y \in L1, x \in \Sigma^*, |x| = |y|\}$
 Construct a CFG for $L2$.
- l) $L = \{w \in \{0,1\}^* : w_1\#w_2 : \text{number of 0s } w_1 \text{ is equal to number of 1s in } w_2.\}$
- m) $L = \{w \in \{0,1\}^* : w_1\#w_2 : \text{length of } w_2 \text{ is double of length of } w_1.\}$
- n) Recall that for a string w , $|w|$ denotes the length of w . $\Sigma = \{0,1\}$
 $L1 = \{w \in \Sigma^* : w \text{ contains exactly two 1s}\}$
 $L2 = \{x\#y : x \in \Sigma^*, y \in L1, |x| = |y|\}$
 Construct a CFG for $L2$.
- o) $L = \{1^i 0 2^j 1^k \mid i, j, k \geq 0, 3i \geq 4k + 2, j \text{ is not divisible by three}\}$

Construct a Context-free Grammar for the language expressed by the following regular expression:

- a) $(ab)^* + (a+b^3)cb^*$
- b) $(a^*b^* + (ac+b^*c)b)^*$

Question 2: Derivation, Parse tree, Ambiguity

I. $S \rightarrow ASB \mid SS \mid SAS \mid A$

$$A \rightarrow ASS|BS|B$$

$$B \rightarrow 00|11|01|1$$

Given the above Context-Free Grammar, answer the following questions:

- Show the Left Most Derivation of 00010111
- Sketch the parse tree corresponding to the derivation you gave in (a).
- Show the Right Most Derivation of 00010111
- Sketch the parse tree corresponding to the derivation you gave in (c).

II. $A \rightarrow A1 | 0A1 | 01$

Given the above Context-Free Grammar, answer the following questions:

- Give a leftmost derivation for the string 001111.
- Sketch the parse tree corresponding to the derivation you gave in (a).
- Demonstrate that there are two more parse trees (apart from the one you already found in (b)) for the same string.
- Find a string w of length six such that w has exactly one parse tree in the grammar above.

III. $S \rightarrow 0S1 | 1S0 | SS | 01 | 10$

Given the above Context-Free Grammar, answer the following questions:

- Show that the grammar above is ambiguous by demonstrating two different parse trees for 011010.
- Find a string w of length six such that w has exactly one parse tree in the grammar above.

Question 3: Chomsky Normal Form - CNF [Each Question Carries 10 marks]

- Convert the following grammar into Chomsky Normal Form. You must show work.

$$S \rightarrow bBS | \epsilon$$

$$B \rightarrow bYb | bY$$

$$C \rightarrow Ccb | c | A$$

$$Y \rightarrow bcY | \epsilon | YBF$$

Here b , and c are terminals and the rest are variables.

- Convert the following grammar into Chomsky Normal Form. You must show work.

$$S \rightarrow YaSb$$

$$X \rightarrow aXY | bX | Y$$

$$Y \rightarrow X | b | \epsilon$$

Here a , and b are terminals and the rest are variables.

- Convert the following grammar into Chomsky Normal Form. You must show work.

$$S \rightarrow bXaY | ZXb$$

$$X \rightarrow aY | bY | Y$$

$$Y \rightarrow X | c | \epsilon$$

$$Z \rightarrow ZaX$$

Here a , b , c are terminals and the rest are variables.

d)

Problem 3 (CO2): Chomsky Normal Form (10 points)

- (a) List the rules that violate the conditions of Chomsky Normal form in the following grammar. Here a , b , and c are terminals and the rest are variables.

$$\begin{aligned} A &\rightarrow BC \mid bB \mid a \\ B &\rightarrow bb \mid Cb \mid b \mid C \\ C &\rightarrow c \end{aligned}$$

- (b) Write down the additional rules that need to be added to the following grammar if the production $B \rightarrow c$ is removed. Here 0 and 1 are terminals and the rest are variables.

$$\begin{aligned} S &\rightarrow AB \mid 1 \\ A &\rightarrow BAB \mid ABA \mid B \mid 1 \\ B &\rightarrow 00 \mid \epsilon \end{aligned}$$

- (c) Write down the additional rules that need to be added to the following grammar if the unit productions are removed. Here 0 and 1 are terminals and the rest are variables.

$$\begin{aligned} S &\rightarrow XYX \mid YX \mid X \mid Y \\ Y &\rightarrow XY \mid X0 \mid 0 \\ X &\rightarrow 1 \mid Y \end{aligned}$$

Question 4: CYK Algorithm [Each Question Carries 10 marks]

- a) Consider the following grammar in CNF form.

$$\begin{aligned} S &\rightarrow BC \mid CD \\ B &\rightarrow CB \mid c \\ C &\rightarrow DD \mid c \\ D &\rightarrow BC \mid b \end{aligned}$$

Show if string $w=ccbb$ can be derived from the grammar above.

- b) Apply the CYK algorithm to determine whether the string “abcaa” can be derived in the following grammar. You must show the entire CYK table. Here a , b , c are terminals and the rest are variables.

$$\begin{aligned} S &\rightarrow CA \\ A &\rightarrow AA \mid AD \mid a \\ B &\rightarrow AB \mid BC \mid b \\ C &\rightarrow CA \mid BC \mid c \\ D &\rightarrow a \end{aligned}$$

Question 5: Pushdown Automata [Each Question Carries 5 marks]

- $L = \{ w \in \{0,1\}^* : w = 0^n 1^n, \text{ where } n \text{ is odd.} \}$
- $L = \{ w \in \{0,1\}^* : w = 0^m 1^n, \text{ where } m, n \geq 1 \text{ and } m \geq n. \}$
- $L = \{ w \in \{0,1\}^* : w = 0^{n+2} 1^n, \text{ where } n \geq 0. \}$
- $L = \{ w \in \{0,1\}^* : w = 0^n 1^{2n}, \text{ where } n \geq 0. \}$
- $L = \{ w \in \{0,1\}^* : w \text{ is a palindrome.} \}$
[what if it says even length palindrome or odd length palindrome specifically]
- $L = \{ w \in \{0,1\}^* : w \mid \text{the length of } w \text{ is divisible by four} \}$
- $L_1 = \{ w \in \{0,1\}^* : \text{the number of 1s in } w \text{ is multiple of 3.} \}$
 $L_2 = \{ w \in \{0,1\}^* : w \text{ contains even numbers of 0.} \}$

Construct a PDA for $L = \{ w \in \{0,1\}^* : w = uv, \text{ where } u \in L_1, v \in L_2 \text{ and } |u| = |v| \}$

h) $L = \{ w \in \{0,1\}^* : w = 0^i 1^j 0^k, \text{ where } j = i+k \text{ and } i, k \geq 0. \}$

i) $L = \{ w \in \{0,1\}^* : w_1 \# w_2 : \text{number of } 00 \text{ substrings in } w_1 \text{ is equal to number of } 11 \text{ in } w_2. \}$

i)

Construct a pushdown automaton for the following language.

$$L = \{ u \bar{u}^R : u \in \{0,1\}^* \}$$

Here, \bar{u}^R denotes the reverse complement of the string u .

For example, $01001101 \in L$ because the second half of the string, 1101 , is the reverse complement of the first half, 0100 , i.e. $1101 = \overline{0100}^R$.