



INSTITUT
POLYTECHNIQUE
DE PARIS



Quantum Communication PAF Day 5

24 Juin 2022

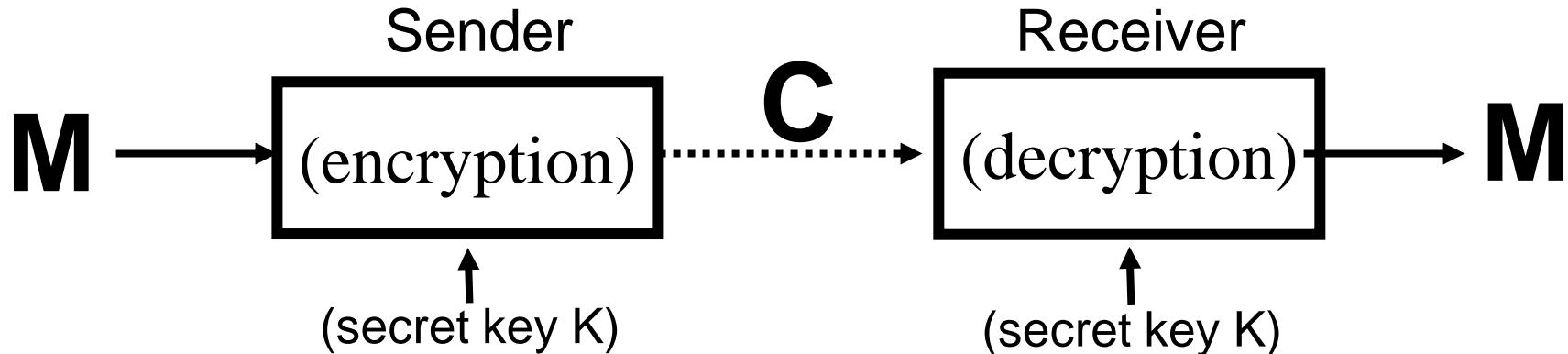
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Plan du cours

- Cryptographie quantique et cryptographie classique
- Principe de la Distribution Quantique de Clé (QKD)
- Real-World QKD

Cryptographie quantique et cryptographie classique

Symmetric-Key Cryptography



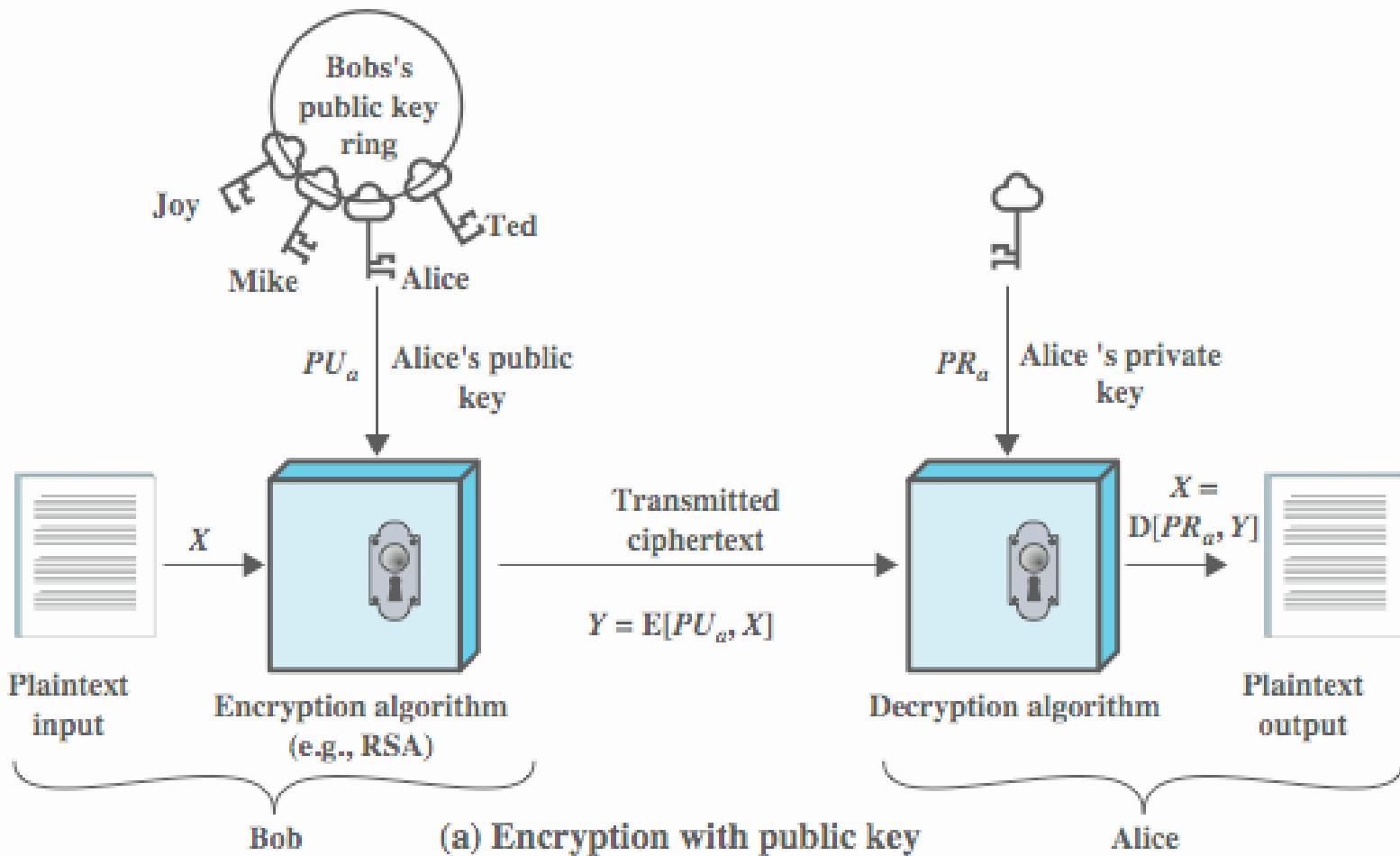
- Traditional (before 1970's) **private/secret/symmetric key** cryptography uses **one key**
- Same key shared by both sender and receiver => **Symmetric**

- *If this key is disclosed communications are compromised*
- *Does not protect sender from receiver forging a message & claiming is sent by sender (encryption ≠ authentication)*

Public-Key Cryptography

- Probably most significant advance in the 3000 year history of cryptography
- **Uses two keys – a public & a private key**
- **Asymmetric** since parties are **not** equal
- Uses clever application of number theoretic concepts to function
- **Complements rather than replaces private key crypto**
- Developed to address mainly two key issues:
 - **key distribution** – how to have secure communications in general without having to trust a KDC with your key
 - **digital signatures** – how to verify a message comes intact from the claimed sender

Example: RSA Public-Key encryption



Any person (here Bob), in posession of Alice public key can send confidential message to Alice

Asymmetric : $B \rightarrow A$

Cryptographie: Symétrique vs. asymétrique

Chiffrement symétrique (à clé secrète)

- Pros
 - Clé courte (~100 bits)
 - Chiffrement / déchiffrement rapide
- Cons
 - Distribution sécurisée de la clé
- Utilisation
 - Chiffrement de grand volume de données
- Exemples d'algorithmes
 - AES
 - DES

Chiffrement asymétrique (à clé publique)

- Pros
 - Pas nécessaire d'échanger la clé secrète (seule la clé publique est publiée)
- Cons
 - Clé longue (~ 1000 bits)
 - Calcul intensif
- Utilisation
 - Distribution de clés secrètes
 - Signature numérique
- Exemples d'algorithmes
 - RSA
 - Diffie-Helman

Modern cryptography : computational assumptions

Example1: Hardness of breaking AES128 encryption

Assumptions: AES (block cipher) is a secure one-way function

→ Best attack is exhaustive search, requires 2^{128} operations

Example2: Hardness of factoring

Assumption: Best known factoring algorithm (General Field Number Sieve) is **sup-exponential**

Factoring large number N, requires $\text{Exp} [O((\ln n)^{1/3})]$ operations

Remark: what about practical computing power ?

(individu, ~ 10 k\$)

1 GHz * 100 (parallélisation) * 1 an $\sim 2^{52}$

(grande organisation type NSA ~ 1000 M\$)

10 Petaflops * 1 an $\sim 2^{78}$

One Time Pad OTP – Masque Jetable (Vernam 1917)

$$M=C=K=\{0,1\}^n$$

Chiffrement

$$C = E(k, m) = k \oplus m$$

Déchiffrement

$$D(k, c) = k \oplus c$$

Msg M: 0 1 1 0 1 1 1 
Key K: 1 0 1 1 0 1 0

Ciphertext

C: 1 1 0 1 1 0 1

Shannon (1949): OTP vérifie la propriété de sécurité inconditionnelle

Key distribution problem and public-key crypto

Shannon positive result (1949):

- One-Time-Pad verifies the perfect secrecy condition

Shannon negative result (1949):

- Perfect secrecy condition requires $|K| \geq |M|$
- Secret-key distribution problem
- Information-theoretic security considered non-practical

Public-key cryptography

- Diffie-Hellman 1976
- RSA 1978

Current solution to the key distribution problem

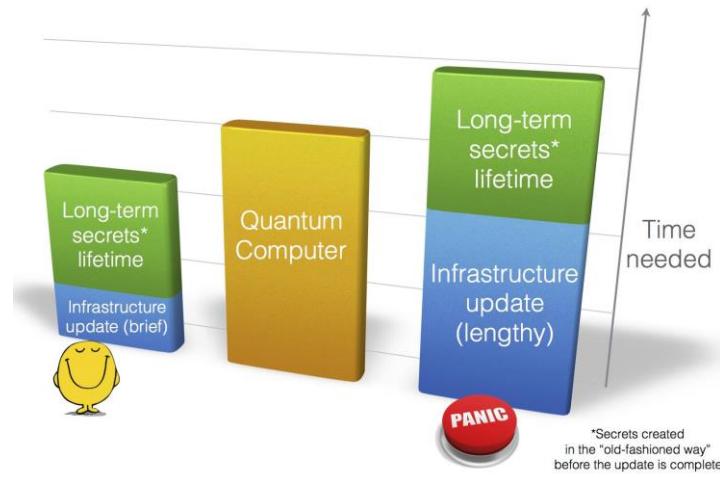


Need for Quantum Resistant (Post-Quantum) Crypto

Threats on existing public-key cryptography

~~RSA, Elliptic Curve, Discrete Log (DH,...)~~

If $X+Y > Z \rightarrow \text{PANIC}$ (Mosca Th.)



NIST Call for Quantum Resistant

First call in 2017

→ first standards in 2022/ 2024

Public-key cryptosystem	Example	Year
Code-based cryptography	McEliece encryption scheme	1978
Hash-based cryptography	Merkle's hash-tree signature system	1979
Lattice-based cryptography	NTRU encryption scheme	1996
Multivariate-quadratic-equations	HFE signature scheme	1996

Table 1.1: List of post-quantum cryptography

Principe de la QKD (distribution quantique de clé)

Protocole BB84

Precursor of quantum crypto: Quantum Money

Uncloseable Quantum Banknotes

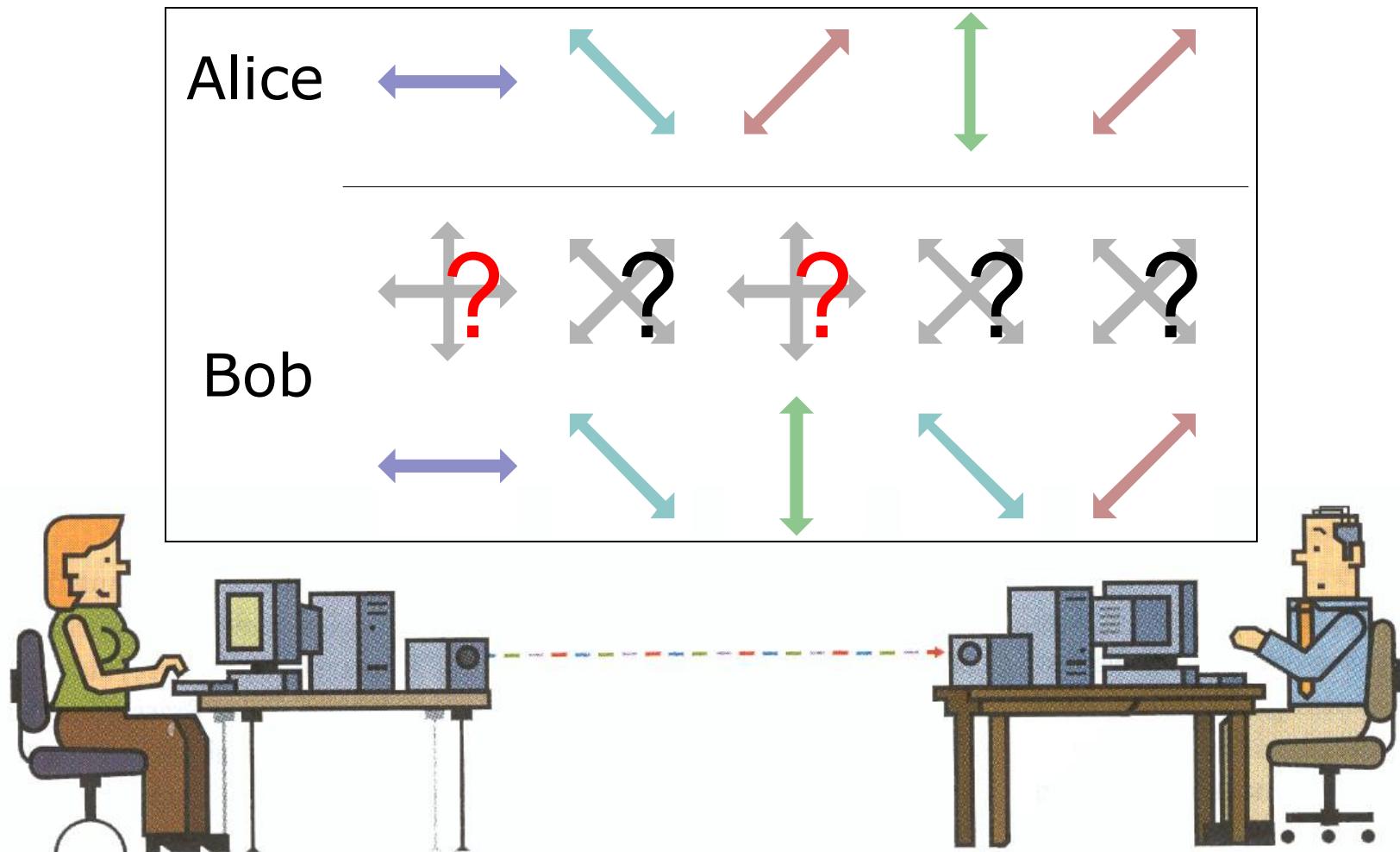
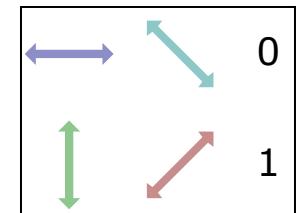
Wiesner, 1969 .. Published 1983



Crucial ideal: Conjugate Coding

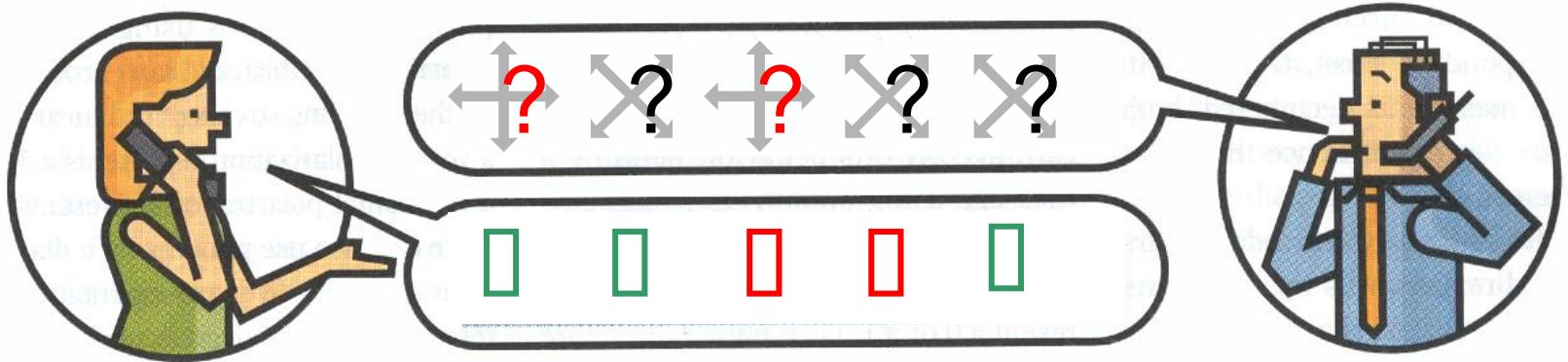
Encode 1 bit into 1 qubit using 2 complementary basis

BB84: Transmission

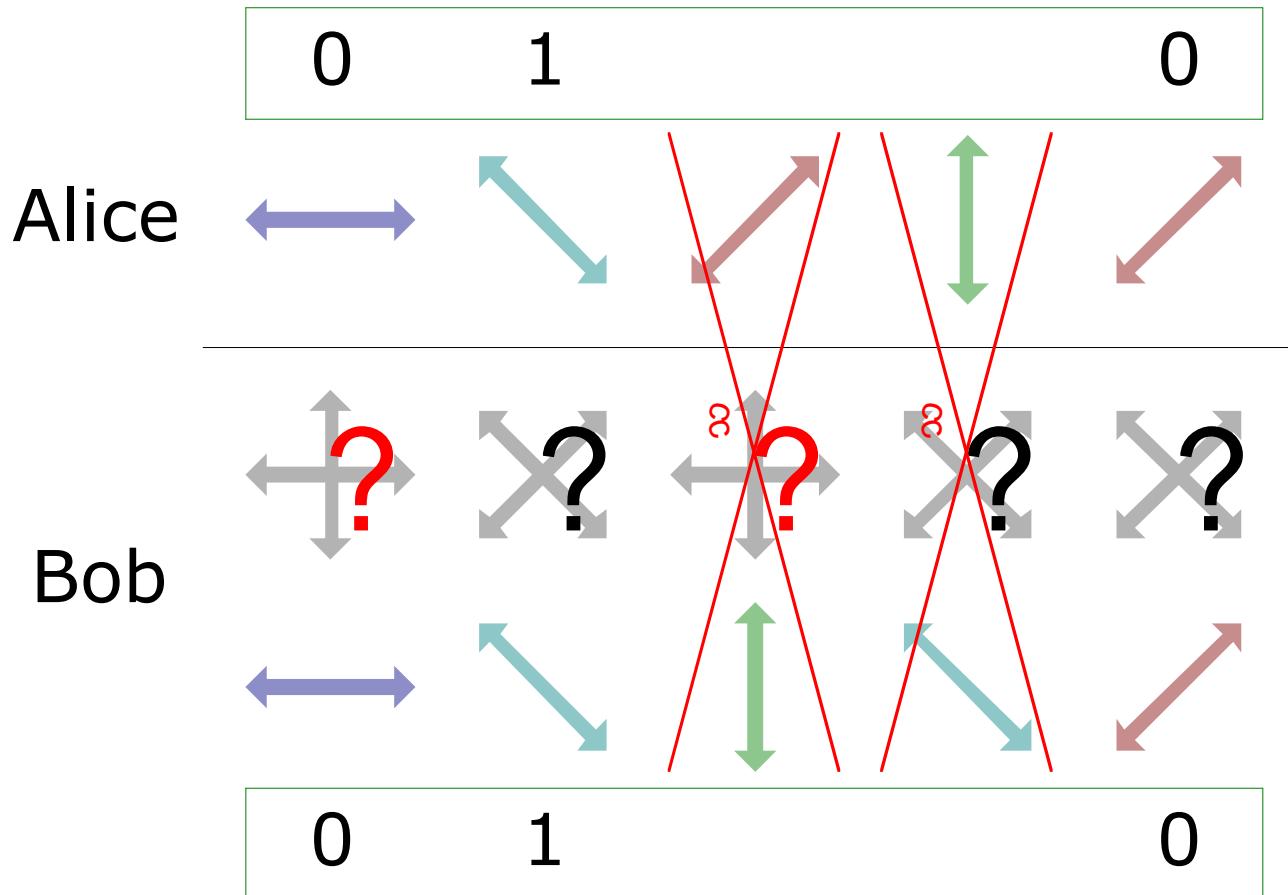
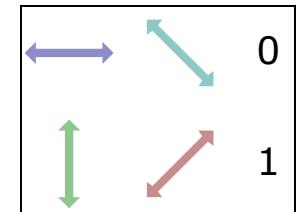


BB84: Tamisage/ Sifting

- Eliminer les mesures qui ne correspondent pas
 - Utilisation d'un canal **public**



BB84: Partage de la Clé



Nécessite d'étapes supplémentaire pour “fabriquer une clé”

Données initiales (clé brute)

- contiennent des erreurs
- Peuvent donc aussi être correlées à l'espion Eve

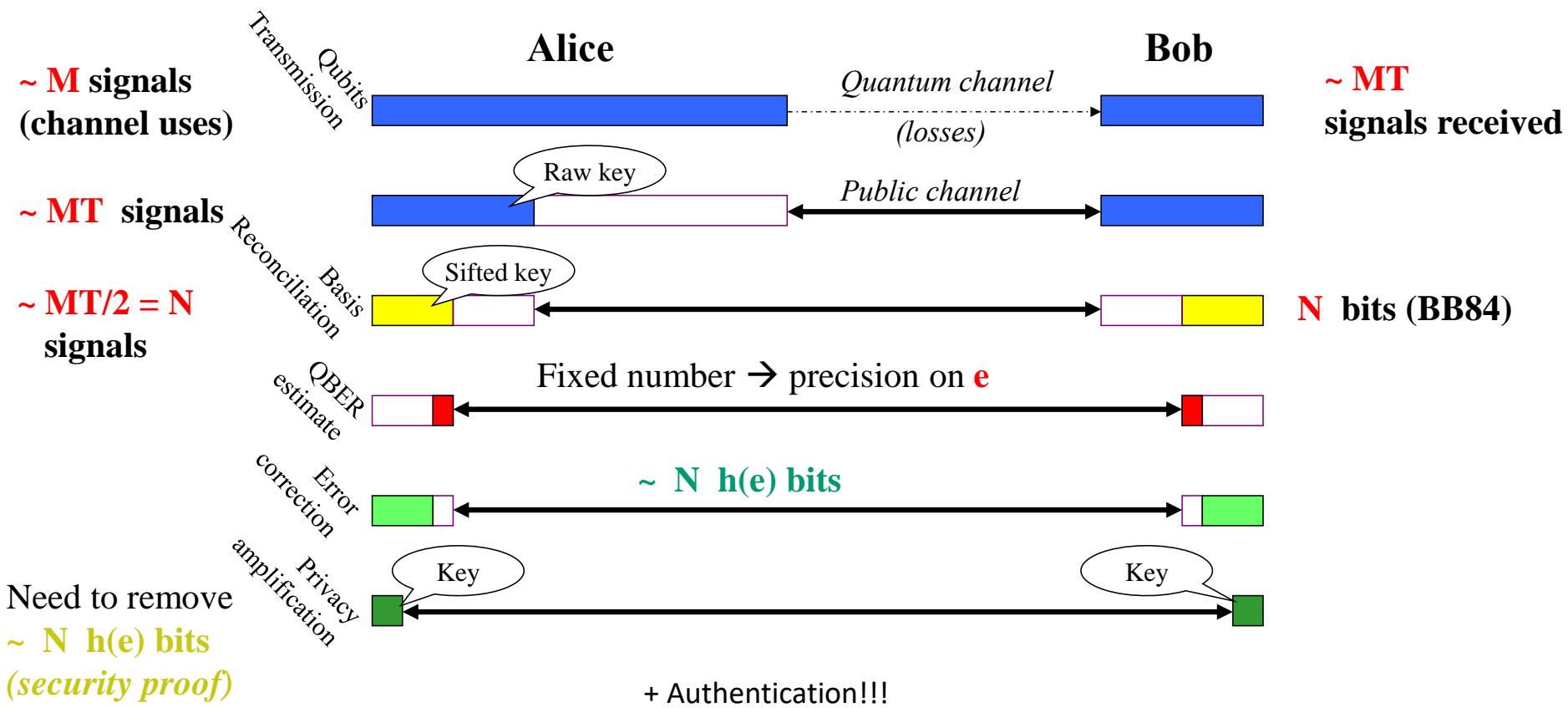


Il faut corriger les erreurs

Il faut distiller une clé totalement inconnue d'Eve

➔ *Etapes de reconciliation (canal Classique)*

The steps to a secret key



Key Rate _{BB84} per ch.use $\sim T (1-2 h(e))$

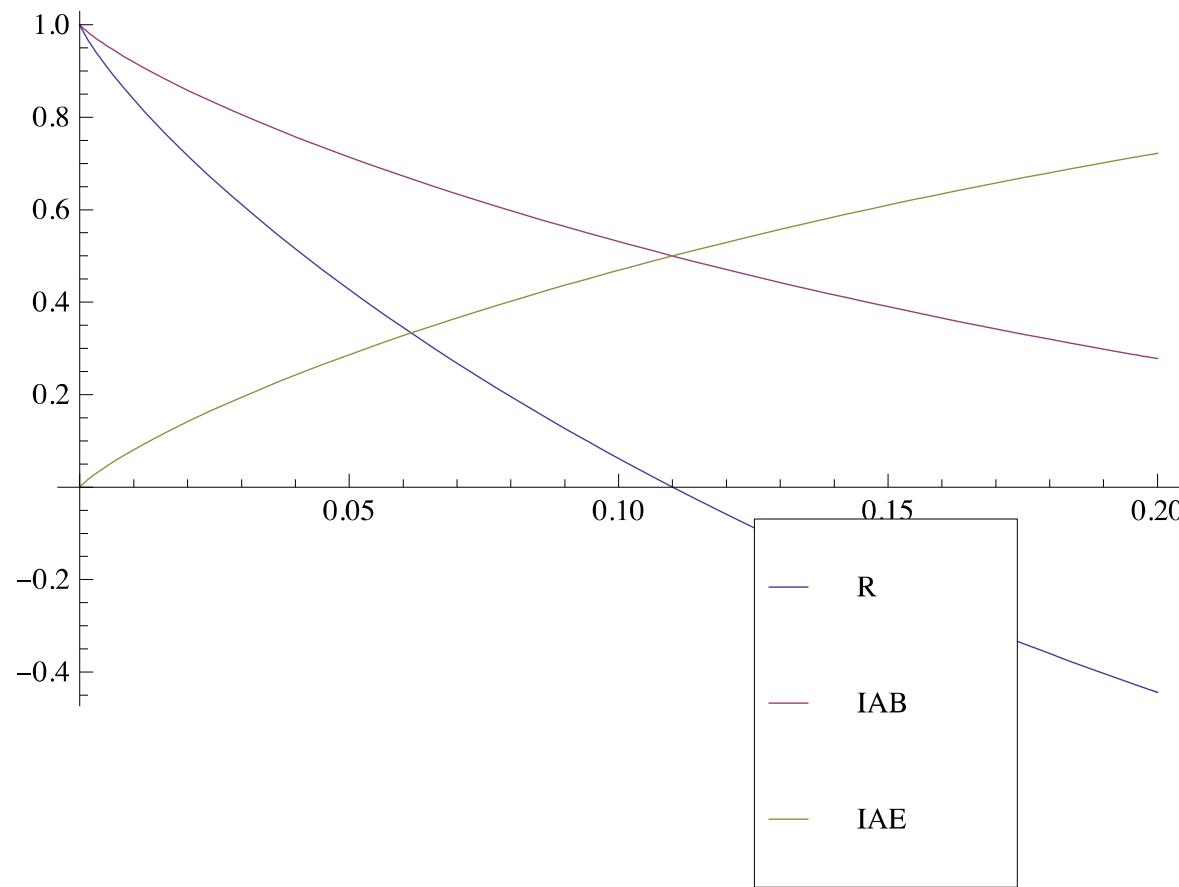
Secure key rate: (depends on Eve attack power)

Main idea: Csiszar Körner, (for classical correlated data X_A, X_B, X_E),

$$R(XA ; XB \mid\mid X_E) \geq \max[I(X_A ; X_B) - I(X_A ; X_E), \\ I(X_B ; X_A) - I(X_B ; X_E)]$$

- **Direct reconciliation:** $R \geq I(X_A ; X_B) - I(X_A ; X_E)$
- Reverse reconciliation: $R \geq I(X_B ; X_A) - I(X_B ; X_E)$

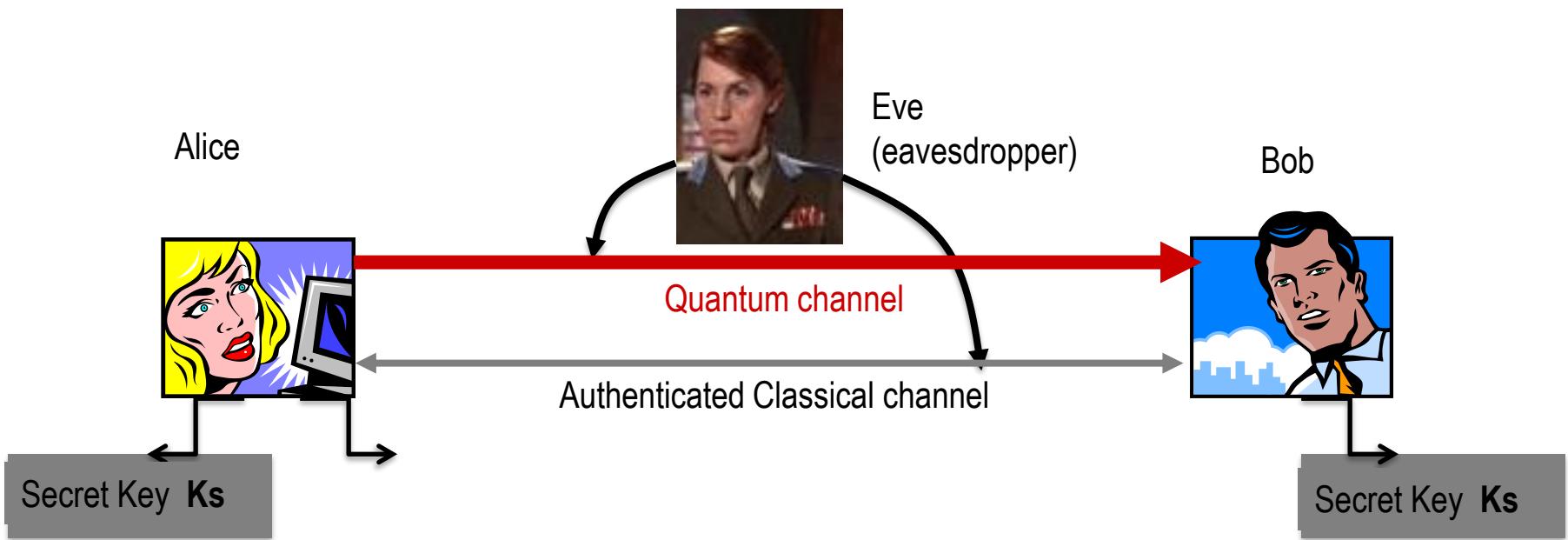
BB84 Secure key rate against collective and coherent attacks



BB84 : $R_{(\text{secure key rate})} \sim n (1 - 2 h(QBER))$

Taux d'erreur tolérable maximal $\sim 11\%$

Quantum Key Distribution (QKD): general setting



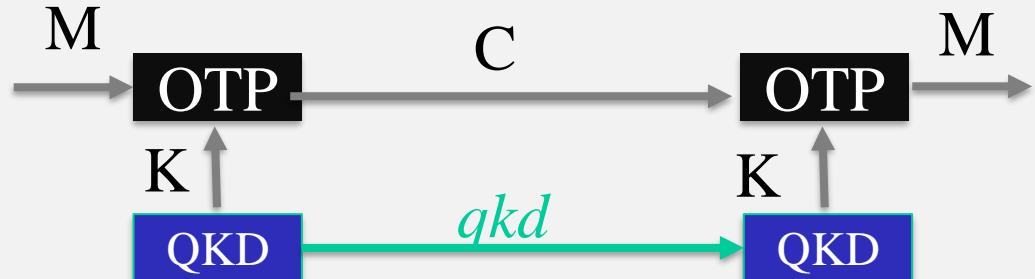
- Intuition for security:
Any measurement by Eve leads to detectable perturbation by Alice/Bob
- **Specificity: Information-Theoretic Security (ITS)** [Unconditional Security]
 - No assumption about Eve computational power
 - « Future-proof »

QKD combined with encryption: Secure Communication

One-Time Pad rekeying

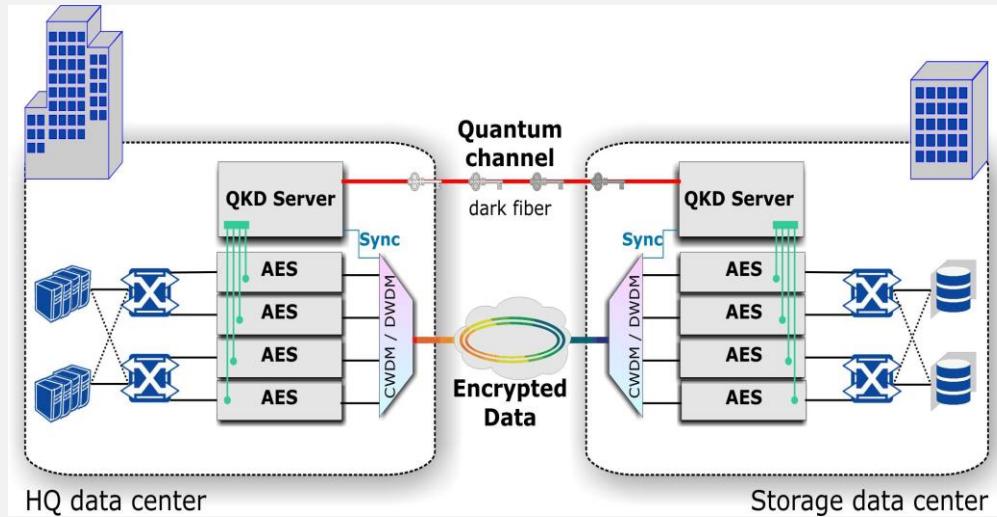
- ☺ ITS
- ➔ Perfect Secrecy

☹ Rate = QKD Rate \leq 1Mb/s



AES Rekeying

- ☺ High Rate \sim 10 Gb/s
- ☹ Less Security Gain



QKD mostly useful when long-term security is needed

Long-term secrets

- Industry, IP
- Military
- Governmental



Personal Data

- Medical record
- Genomic
- Private



Computational Cryptography

Based on hardness of mathematical problems

Generic Vulnerability (incl. PQC)

Harvesting Attack

"Intercept now, decrypt later "



NSA Bullrun program



Long-Term Secure Storage

Need:

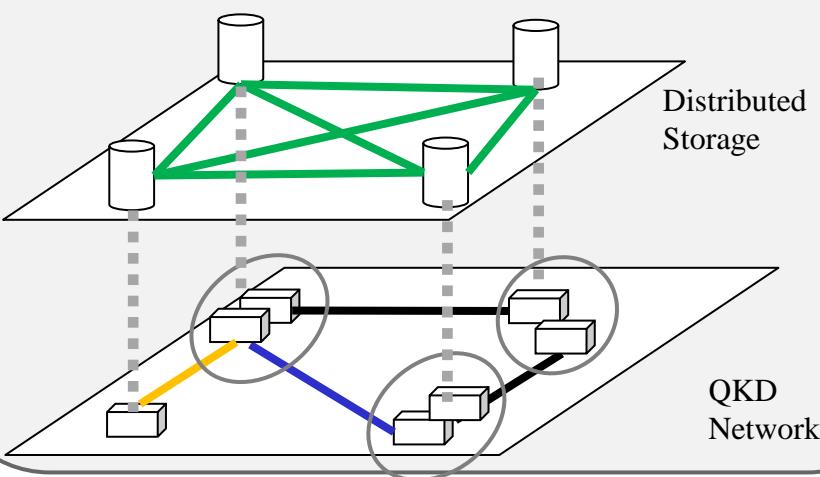
Confidentiality (integrity, availability) over 30+ years

Principle: Proactive Secret Sharing

→ Protects against storage node corruption

Requirements:

- Distributed Storage Infrastructure
- Secure Communication with ITS
 - *Impossible Classically*
 - → **QKD + OTP**



Initial demonstration

Braun, Johannes, J. Buchmann et al. "LINCOS: A storage system providing long-term integrity, authenticity, and confidentiality." *Proceedings of the 2017 ACM on Asia Conference on Computer and Communications Security*. 2017.

+ 2 OpenQKD Use-Cases
& Pilot implementations in Japan

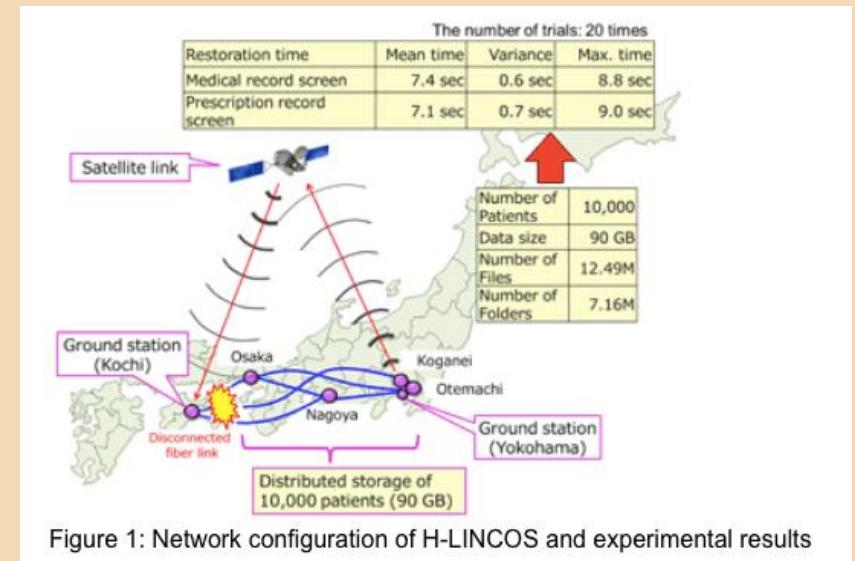
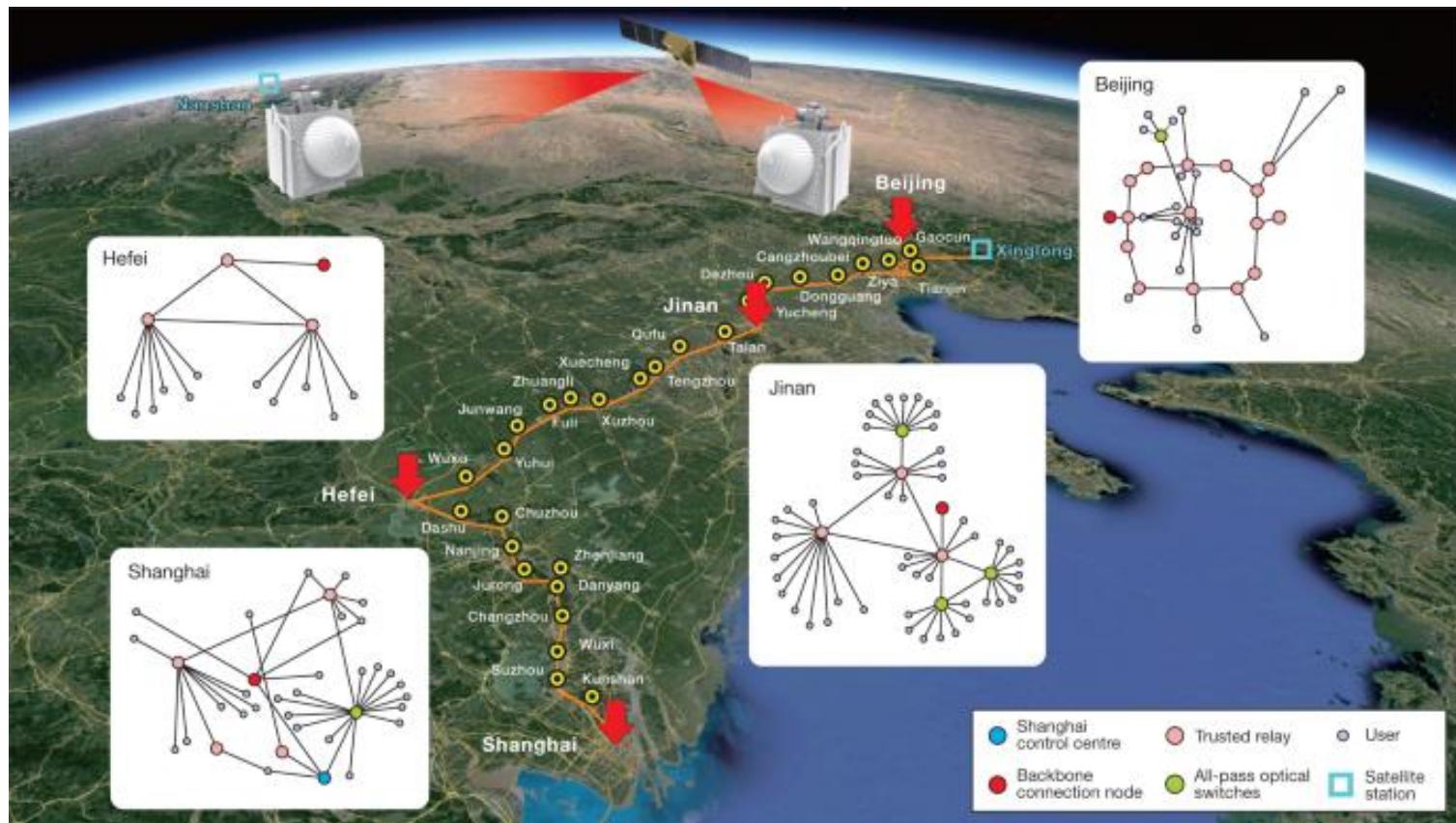


Figure 1: Network configuration of H-LINCOS and experimental results

Real-World QKD

QKD Networks – overview of major prototypes



- 4 metropolitan areas
- 32 trusted relays
- 150 users
- 700 fibre links

Span $\sim 4600 \text{ km}$

QKD Networks – overview of major prototypes

DECLARATION ON A
**QUANTUM COMMUNICATION
INFRASTRUCTURE**
FOR THE EU



Initiated by the European Commission in 2019
Currently under negotiation

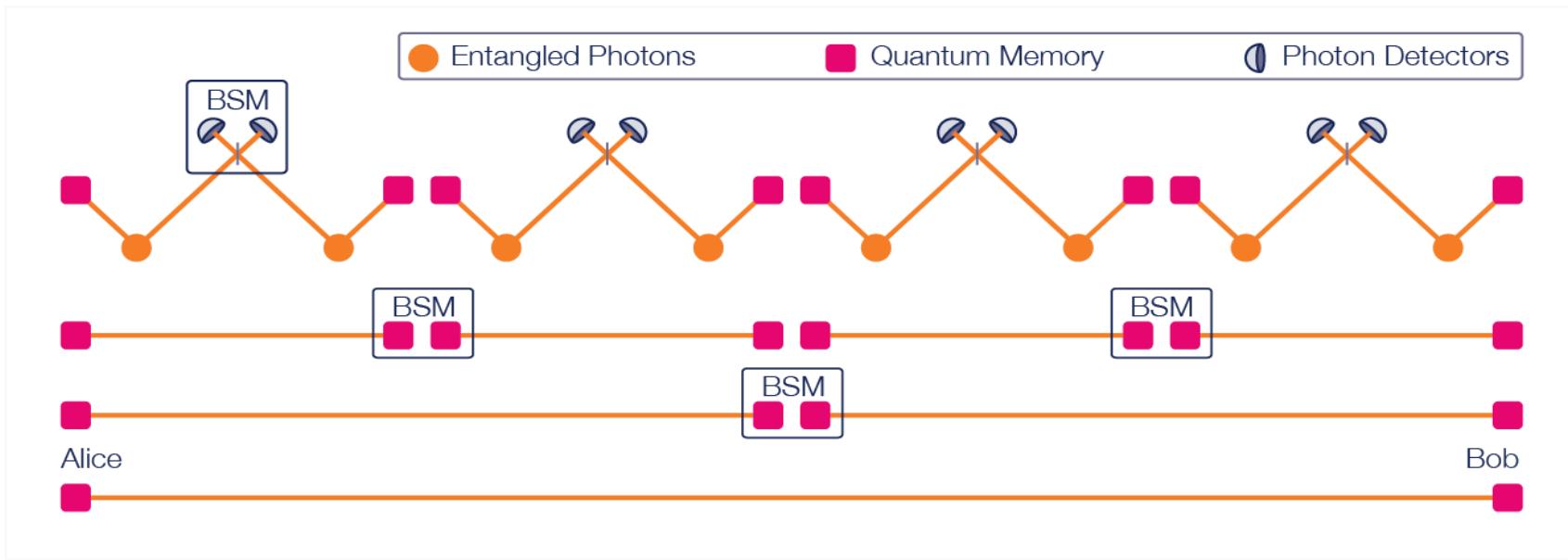
Aims at the **deployment**, by **2030**, of a pan-European **publicly controlled secure quantum communication infrastructure** linking selected EU strategic sites by using **both terrestrial and space links**.

Current signature of 24 countries (incl. France)

Large Industry Involvement – Important foreseen investment (1B€)

→ Opportunities for Quantum Engineers and Researchers

Long-term vision for Q networks: Quantum Internet



Quantum Internet Alliance
Flagship Project