Introduction to Formal Methods
Chapter 02: Modeling Transition Systems

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Outline

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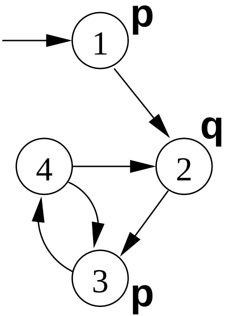
1. Transition Systems as Kripke Models
2. Languages for Transition Systems
3. Properties of Transition Systems

Modeling the system: Kripke models

- Kripke models are used to describe reactive systems:
  - nonterminating systems with infinite behaviors
    (e.g. communication protocols, hardware circuits);
  - represent the dynamic evolution of modeled systems;
  - a state includes values to state variables, program counters, content of communication channels.
  - can be animated and validated before their actual implementation

Kripke model: formal definition

- A Kripke model $\langle S, I, R, AP, L \rangle$ consists of



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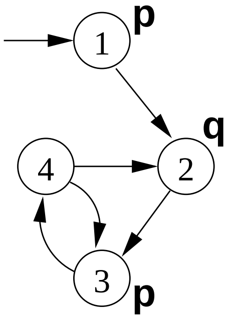

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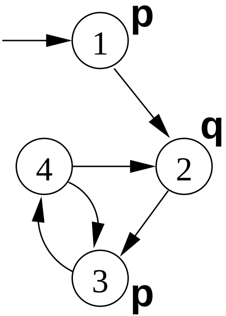

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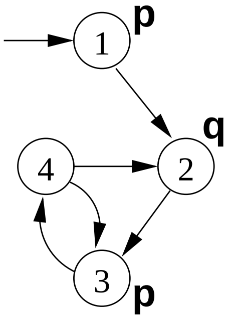
  - a set of transitions $R \subseteq S \times S$;



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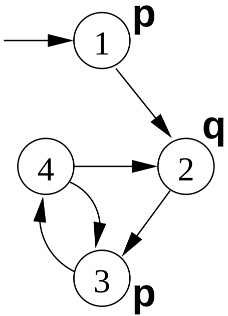

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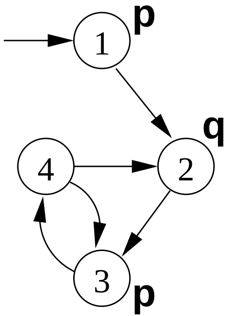

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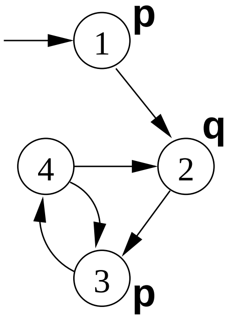

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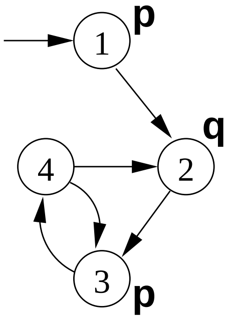

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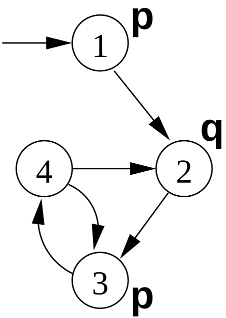
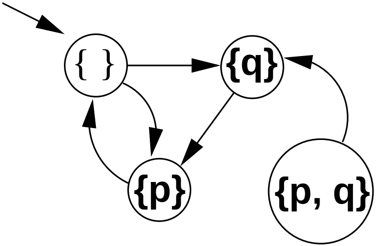


Remark

Unlike with other types of Automata (e.g., Buechi), in Kripke structures the value of every variable is always assigned in each state.

Kripke Structures: two alternative representations:

- each state identifies univocally the values of the atomic propositions which hold there



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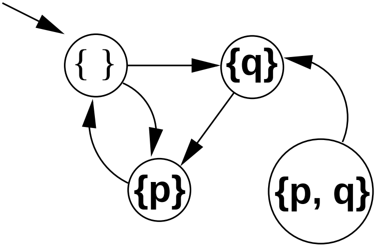
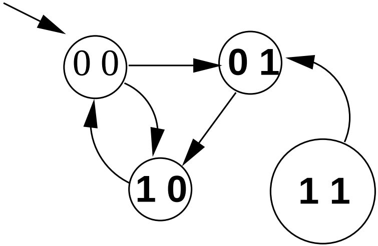
- each state identifies univocally the values of the atomic propositions which hold there
- each state is labeled by a bit vector

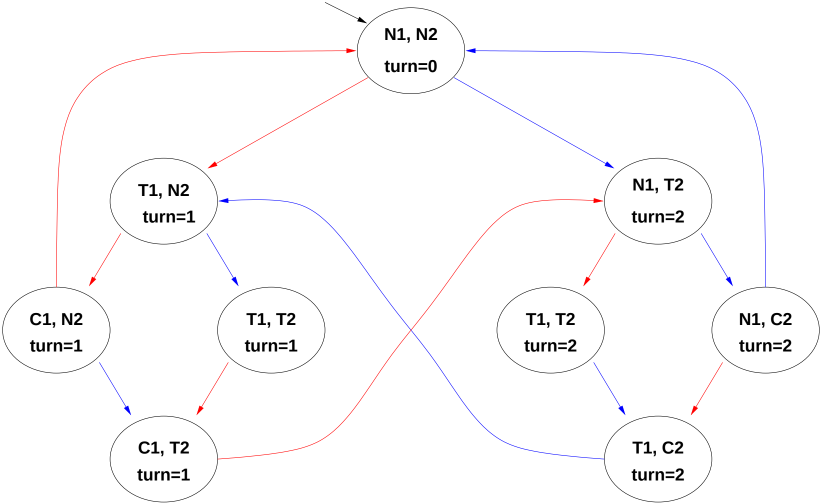




Other representations of finite state machines

- Moore machines
- Mealy machines
- Finite automata
- Büchi automata
- ...

Example: a Kripke model for mutual exclusion



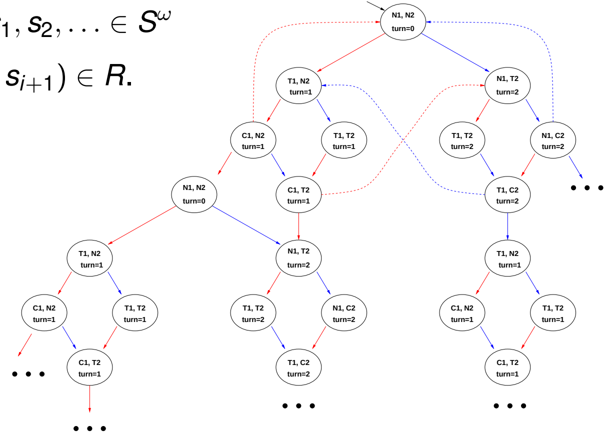
N = noncritical, T = trying, C = critical **User 1** **User 2**

Path in a Kripke Model

A path in a Kripke model M is an infinite sequence of states

$$\pi = s_0, s_1, s_2, \ldots \in S^\omega$$

such that $s_0 \in I$ and $(s_i, s_{i+1}) \in R$.



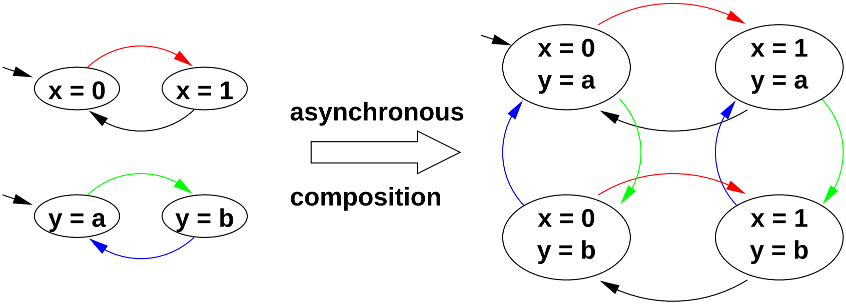
A state s is reachable in M if there is a path from the initial states to s.

Composing Kripke Models

- Complex Kripke Models are tipically obtained by composition of smaller ones
- Components can be combined via
  - asynchronous composition.
  - synchronous composition,

Asynchronous Composition

- Interleaving of evolution of components.
- At each time instant, one component is selected to perform a transition.



- Typical example: communication protocols.

Asynchronous Composition/Product: formal definition

Asynchronous product of Kripke models

Let $M_1 \stackrel{\text{def}}{=} \langle S_1, I_1, R_1, AP_1, L_1 \rangle$, $M_2 \stackrel{\text{def}}{=} \langle S_2, I_2, R_2, AP_2, L_2 \rangle$. Then the asynchronous product $M \stackrel{\text{def}}{=} M_1 || M_2$ is $M \stackrel{\text{def}}{=} \langle S, I, R, AP, L \rangle$, where

- $S \subseteq S_1 \times S_2$ s.t.,
  $\forall \langle s_1, s_2 \rangle \in S,\ \forall l \in AP_1 \cap AP_2, l \in L_1(s_1)$ *iff* $l \in L_2(s_2)$
- $I \subseteq I_1 \times I_2$ s.t. $I \subseteq S$
- $$R(\langle s_1, s_2 \rangle, \langle t_1, t_2 \rangle) \textit{ iff } \begin{array}{l} (R_1(s_1, t_1) \textbf{ and } s_2 = t_2) \textbf{ or} \\ (s_1 = t_1 \textbf{ and } R_2(s_2, t_2)) \end{array}$$
- $AP = AP_1 \cup AP_2$
- $L : S \longmapsto 2^{AP}$ s.t. $L(\langle s_1, s_2 \rangle) \stackrel{\text{def}}{=} L_1(s_1) \cup L_2(s_2)$.

Note: combined states must agree on the values of Boolean variables.

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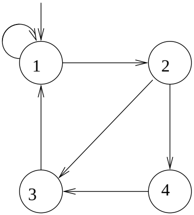
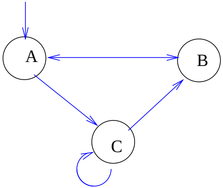
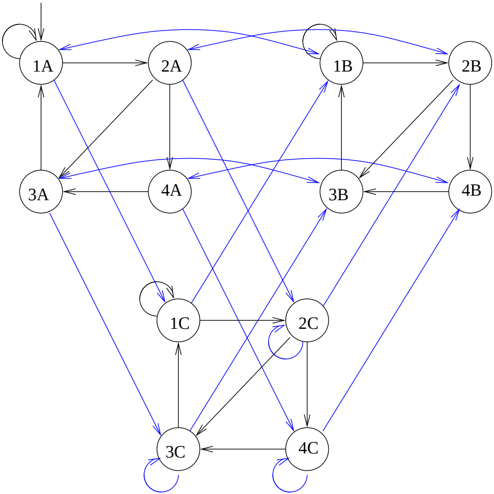
Asynchronous composition is associative:
$(...(M_1||M_2)||...)||M_n) = (M1||(M_2||(...||M_n)...) = M_1||M_2||...||M_n$

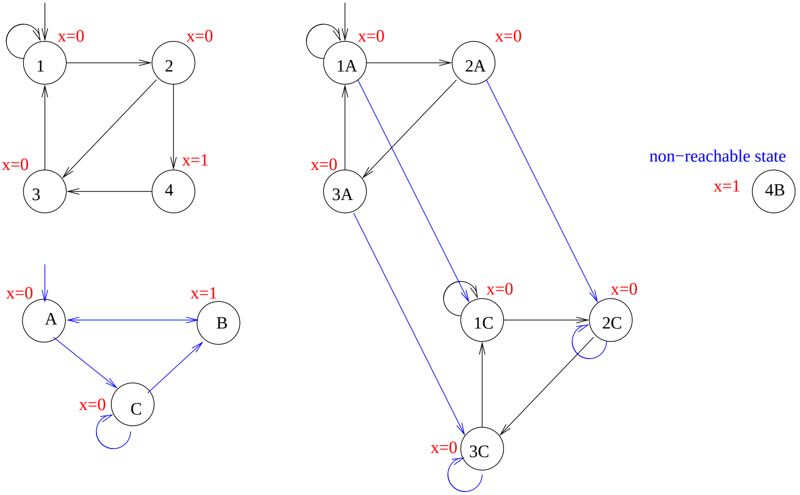
Asynchronous Composition: Example 1







Asynchronous Composition: Example 2



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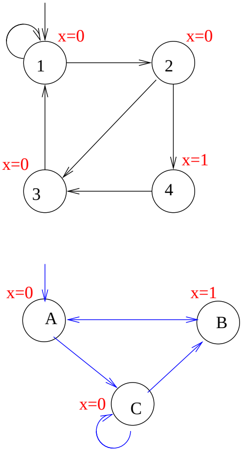

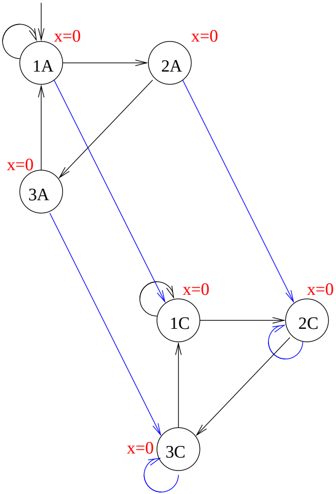

Synchronous Composition

- Components evolve in parallel.
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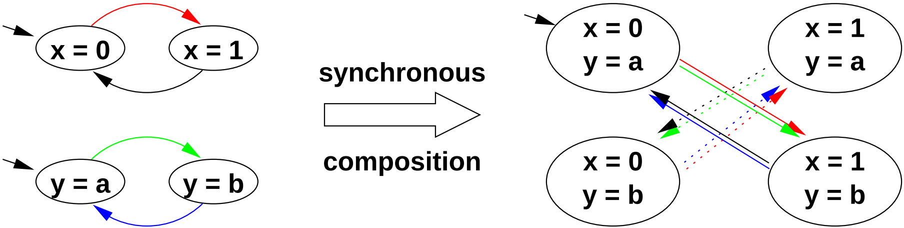

- Typical example: sequential hardware circuits.

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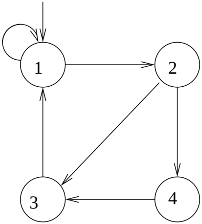
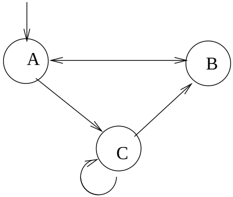
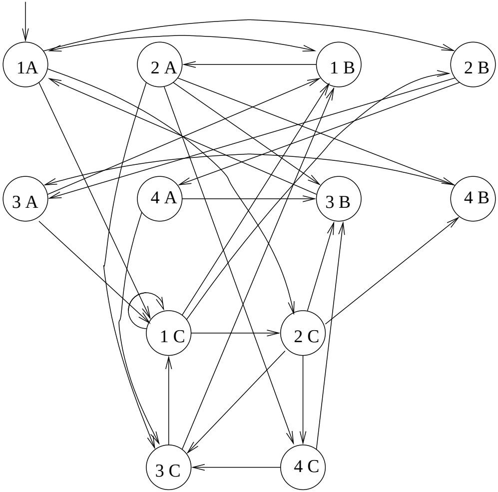
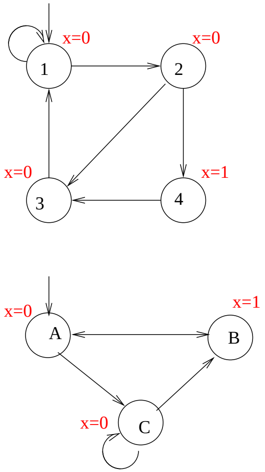
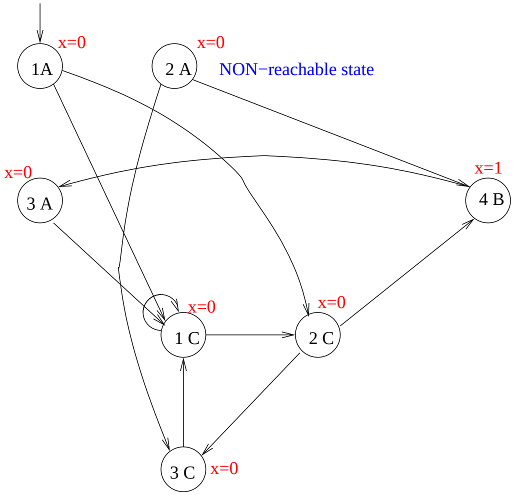




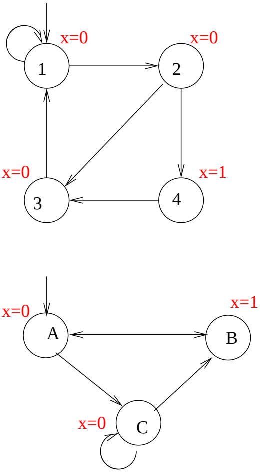


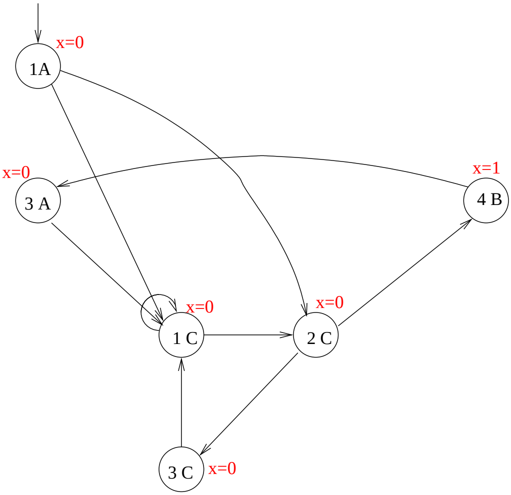
Synchronous Composition: Example 2





Synchronous Composition: Example 2 (cont.)





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Description languages for Kripke Model

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Remark

tipically these description are much more compact (and intuitive) than the explicit representation of the Kripke model.

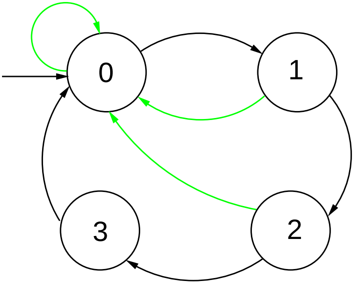
The SMV language

- The input language of the SMV M.C. (and NuSMV)
- Booleans, enumerative and bounded integers as data types
- now enriched with other constructs, e.g. in NuXMV language
- An SMV program consists of:
  - Declarations of the state variables (e.g., `b0`);
  - Assignments that define the valid initial states (e.g., `init(b0) := 0`).
  - Assignments that define the transition relation (e.g., `next(b0) := !b0`).
- Allows for both synchronous and asyncronous composition of modules (though synchronous interaction more natural)

The SMV language: example

Example: The modulo 4 counter with reset

```
MODULE main
 VAR
   b0     : boolean;
   b1     : boolean;
   reset  : boolean;
   out    : 0..3;

 ASSIGN
   init(b0) := 0;
   next(b0) := case
                 reset = 1 : 0;
                 reset = 0 : !b0;
               esac;

   init(b1) := 0;
   next(b1) := case
                 reset = 1 : 0;
                 reset = 0 : (b0 xor b1);
               esac;
   out := toint(b0) + 2*toint(b1);
```

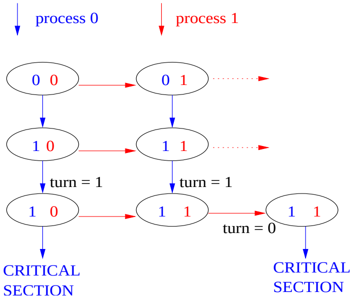


The PROMELA language

- PROMELA (Process Meta Language) is the modeling language of the M.C. SPIN
- The syntax is C-like
- A system in PROMELA consists of a set of *processes* that interact by means of:
  - shared variables
  - communication channels
    - rendez-vous communications
    - buffered communications
- Processes can be created dynamically
- Allows for both synchronous and asyncronous composition of processes (though asynchronous interaction more natural)

The PROMELA language: example

Example: A Mutual Exclusion Algorithm



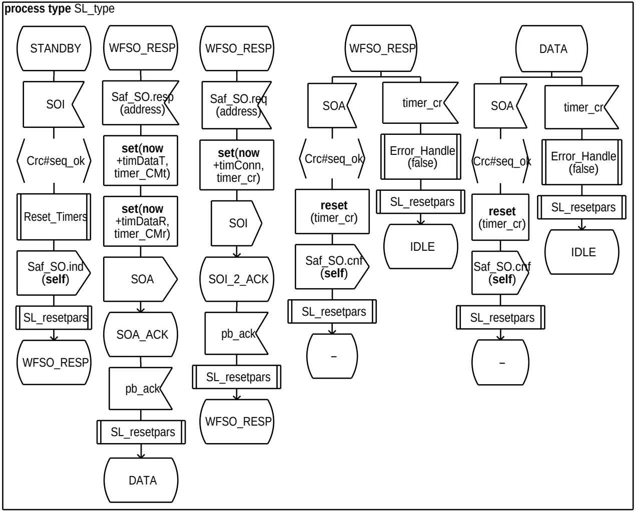

```
bool turn;
bool flag[2];

proctype User(bool pid) {
  flag[pid] = 1;
  turn = 1-pid;
  (flag[1-pid] == 0 || turn == pid);
  /* Begin of critical section */
  ...
  /* End of critical section   */
  flag[pid] = 0;
}
init { run User(0); run User(1) }
```

The SDL language

- An ITU standard
- Allows for booleans, enumerative and bounded integers as data types
- Allows for representing TIME (time elapse, clocks, ...)
- represents states, message I/O, conditions, clock operations, subroutines
- Allows for both synchronous and asyncronous composition of processes (though asynchronous interaction more natural)

The SDL Language: example

Example: the Safety Layer protocol



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Safety properties

- bad events never happen
  - deadlock: two processes waiting for input from each other, the system is unable to perform a transition.
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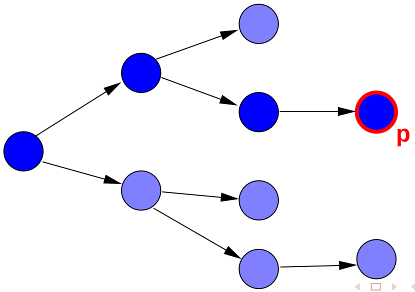
- can be refuted by a finite behaviour
- Ex.: it is never the case that p.



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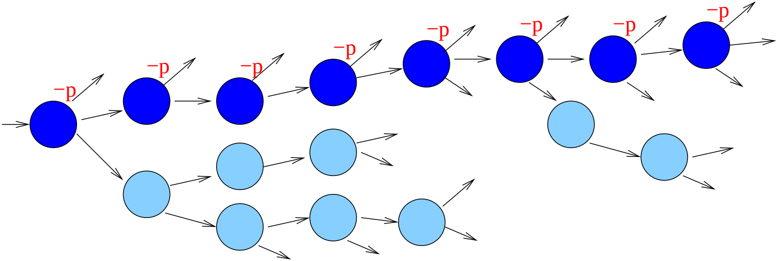

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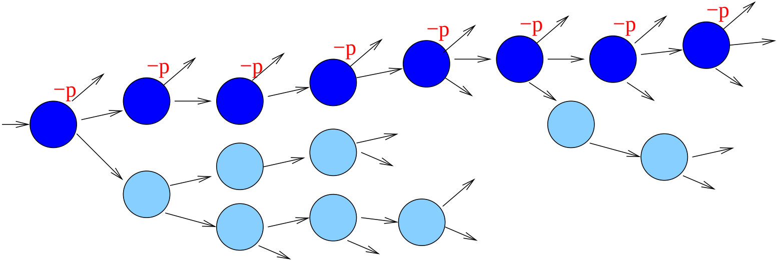

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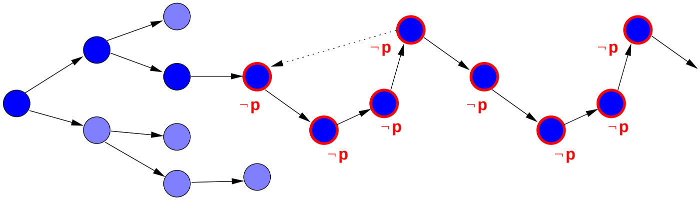

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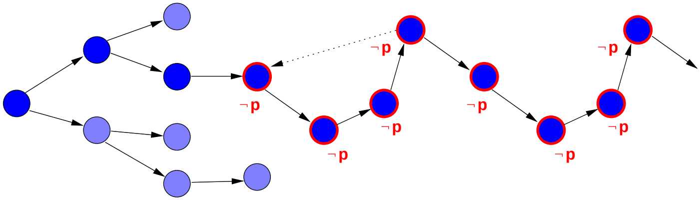

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