

CSci 4061

Introduction to Operating Systems

OS Concepts and Structure

Reading

Read Chapter 1 (R&R)

Opt:

Chapter 1 (MOS) or

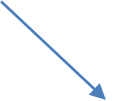
Chapters 1 and 2 (S&G)

The Kernel: core layer of the OS

- The kernel is a library of procedures shared by all user programs, **but the kernel is protected:**
 - User code cannot access internal kernel data structures (and associated code) directly
 - User code can invoke the kernel only at well-defined entry points, and **these are?** **system calls**
- Kernel code is like user code, but the kernel is privileged:
 - Kernel has direct access to all hardware, and handles interrupts and hardware exceptions
 - CPU is either executing OS code (**kernel-mode**) or your code (**user-mode**)

OS can be a mix of user-mode and kernel-mode

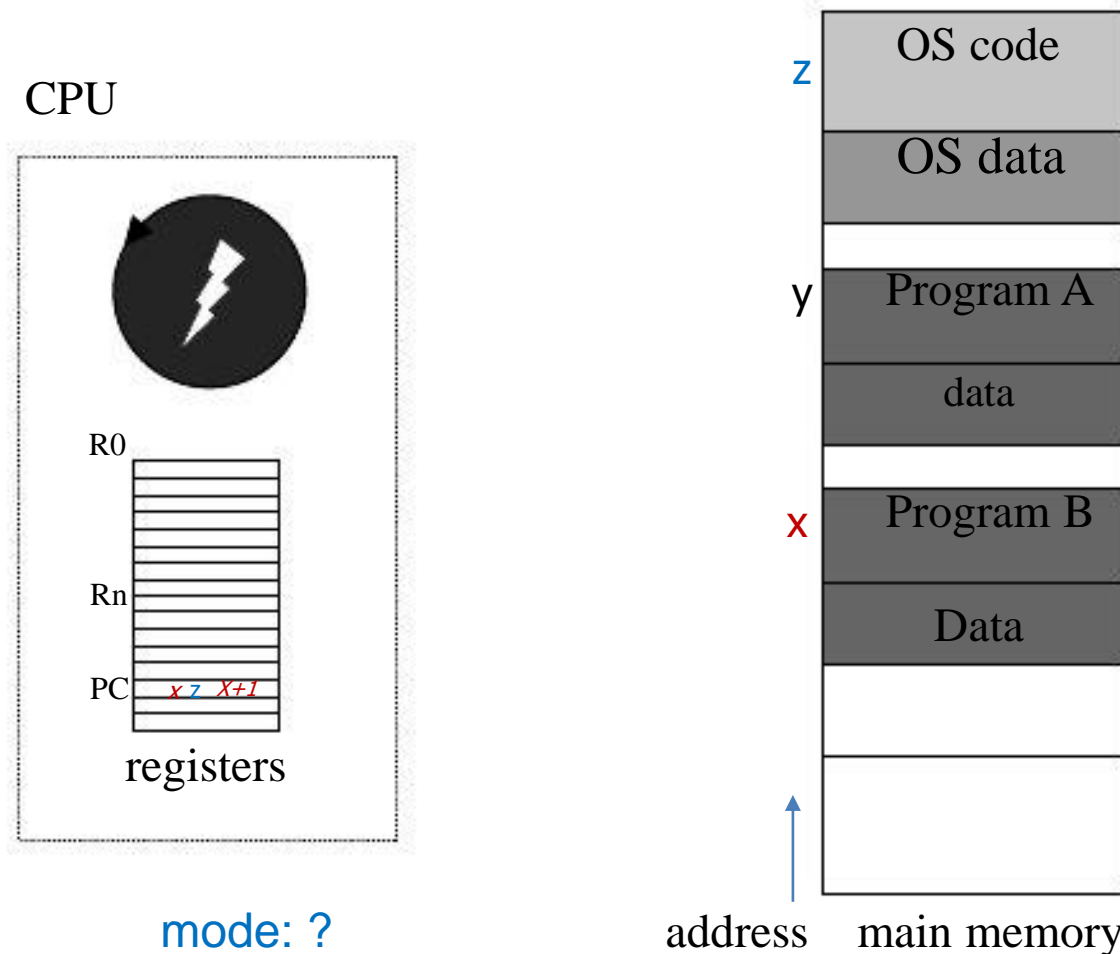
Systems Programmer Viewpoint

- Systems programmer can use system calls **directly** (in assembly)
 - executed by the OS (i.e. kernel mode)
 - when efficiency demands it
 - assembly code: x86 “`int`” instruction, e.g. `int 48`
- Alternatively, language-specific libraries can be used to access system calls
 - C programming language libraries (`libc.a`)
 - E.g. `read (...)`

Terminology Alert!

- I will often refer to low-level library calls as system calls
 - e.g. `read (...) ;`
 - becomes `int #`
- Library (or system calls) are not part of the C language

Running programs: memory and the CPU



Program B makes a system call
System call completes

Let's Look At

OS Concepts and Abstractions
Above the Hardware

Abstraction

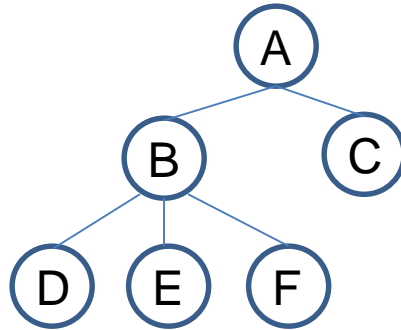
High-level construct

Useful, easy-to-use, understand

Hides lower-level details

PL: class or structure data-type

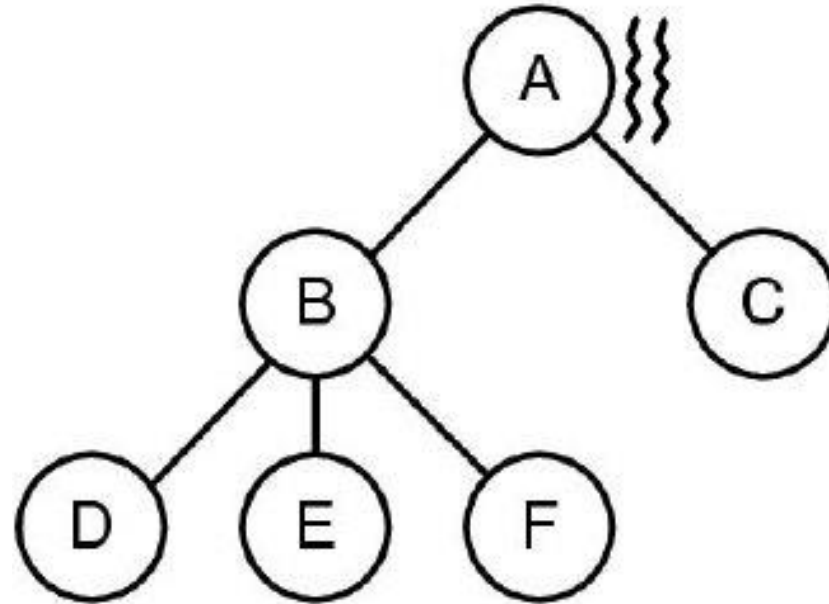
Operating System Concepts: Process



- Process is an executing program: **container** for computing resources (abstraction)
 - Process tree
 - A created two child processes, B and C
 - B created three child processes, D, E, and F

What resources?

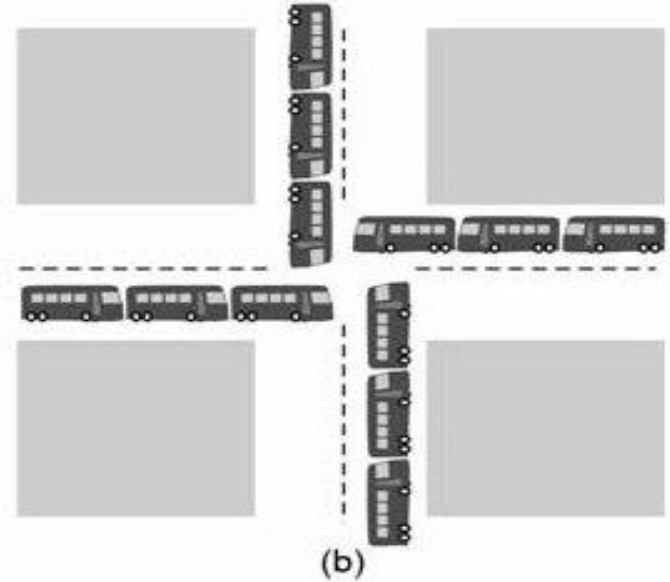
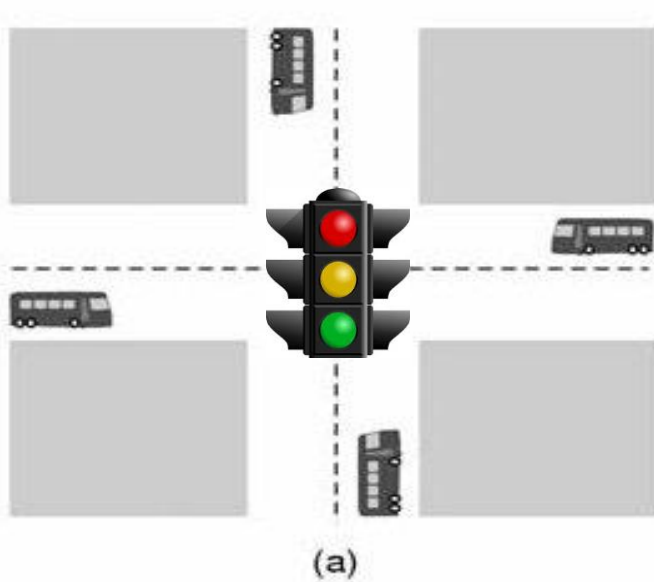
Operating System Concepts: Threads



- A thread is an executing stream of instructions normally within a process
 - A has two threads; share A's resources
 - Every process has at least one thread
 - Threads can also exist in the OS

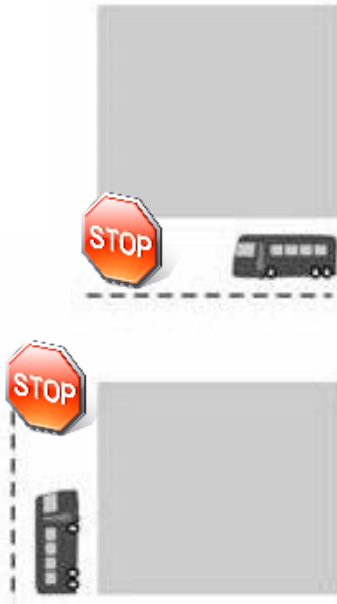
```
main () {  
    int i;  
    i=2;  
}
```

Operating System Concepts: Synchronization



- Concurrency (processes/threads run together) and shared resources can lead to problems:
 - (a) Race condition
 - (b) Deadlock
- Solution: Synchronization, e.g. case (a)?

Operating System Concepts: Synchronization Issues



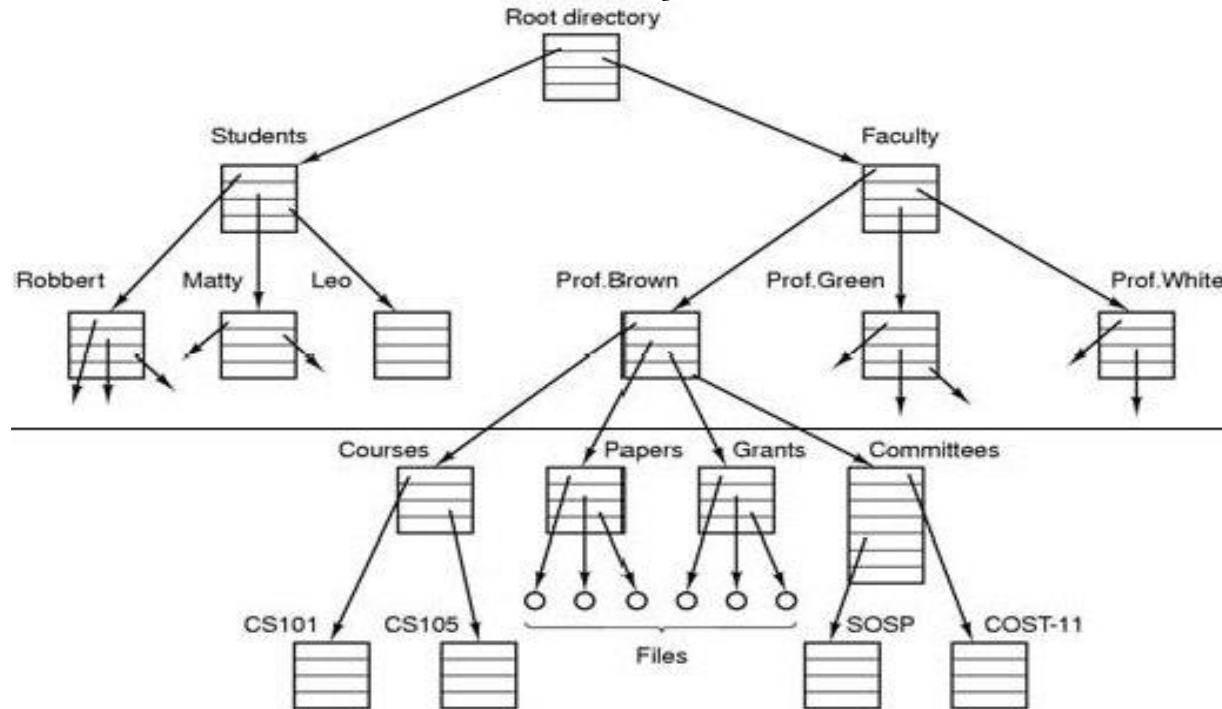
Livelock! (aka “Minnesota Nice”)

No one makes progress

Deadlock/Livelock is often caused by poor use of synchronization

Operating System Concepts:

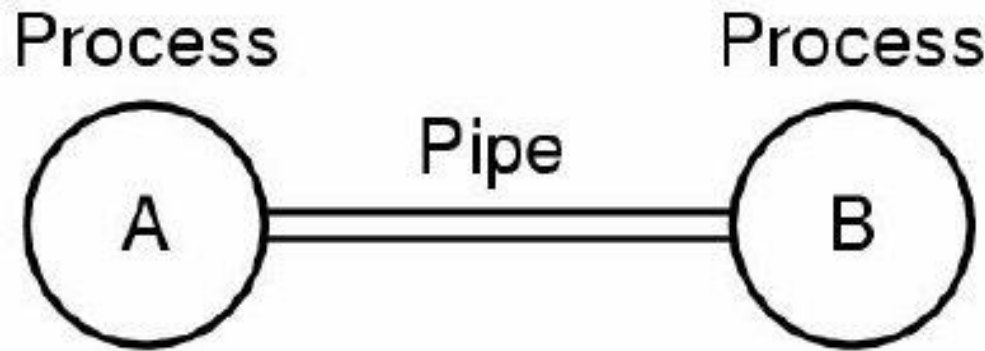
File system



Files/directories are an OS **abstraction** to make data storing and data sharing easier

What are they abstracting?

Operating System Concepts: Communication



- Two processes connected by a “pipe”, channel
- Processes need to communicate - why?

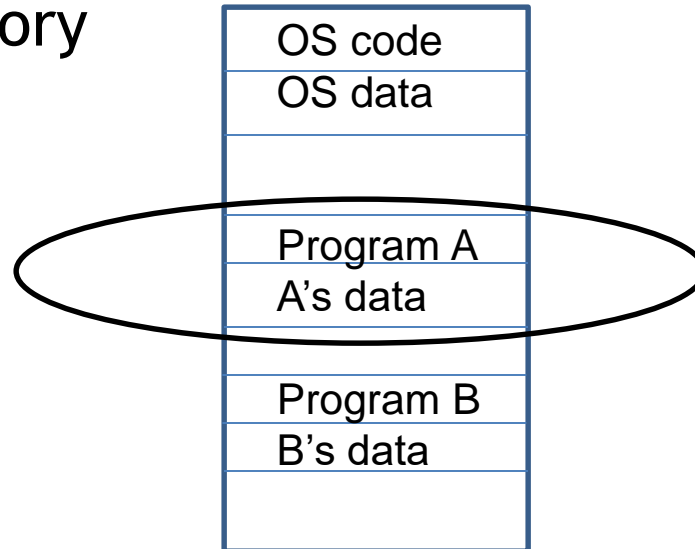
Decompose complex applications

Web browser- Web server

X windows/X11 applications

Operating System Concepts: Memory Management

- How is memory allocated to programs?
 - Largely an “inside” issue but
 - We will see how a program can make good use/bad use of memory



- **Abstraction = virtual memory**

Operating System Concepts: System Calls

- System calls are how user programs interact with the OS
 - Generally available as assembly-language instructions
 - C-Unix provides a library interface to system calls to avoid this messiness
 - e.g. `read (...)` gets compiled into the appropriate syscall linkage/assembly code
 - `a = read(b, c)` vs. `a = myfunc (d, e)`
 - Key differences between these two calls?
 - How parameters are “passed”, “address space” crossing
 - Performance

Example: Some “System Calls” For File Management

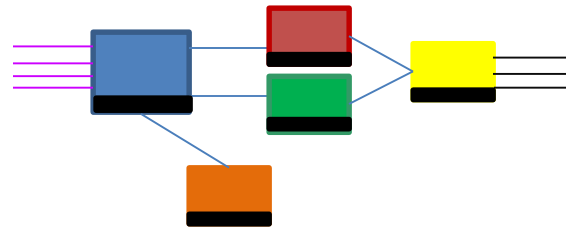
File management

Call	Description
<code>fd = open(file, how, ...)</code>	Open a file for reading, writing or both
<code>s = close(fd)</code>	Close an open file
<code>n = read(fd, buffer, nbytes)</code>	Read data from a file into a buffer
<code>n = write(fd, buffer, nbytes)</code>	Write data from a buffer into a file
<code>position = lseek(fd, offset, whence)</code>	Move the file pointer
<code>s = stat(name, &buf)</code>	Get a file's status information

In this course, we will use the term system call to refer to the C-Unix interface, e.g. **open**

Systems Concepts

“systems”: OS, Internet, ATC, ...



Granularity

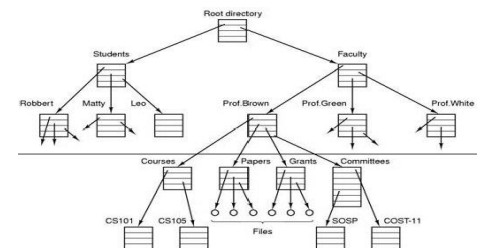
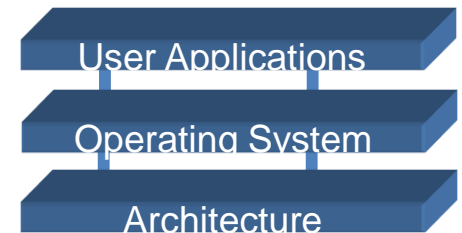
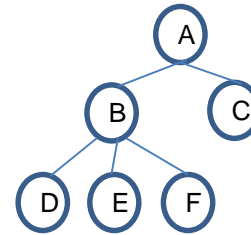
Modularity

Abstraction

Layering

Hierarchy

Complexity



Complexity?

- Different stakeholders => different metrics and requirements
 - Programmer => ease-of-problem-solving
 - End-user(s) => performance, ease-of-use
 - Owner (~ system) => fairness/priority, efficiency or utilization
 - Admin => security
 - OS Vendor => extensible, secure, reliable, ...
- Tradeoff and conflict lead to complexity

Projects and Groups

- Group work repository
 - [github.umn.edu](https://github.com/umn)
- IDE? Your favorite text editor, makefiles/gcc, gdb/ddd
- Standard 4061 environment
 - [posted soon](#)
- Group composition? Stay tuned.

This Weekend

- C/UNIX Refresh (or cram)
 1. Edit and write a simple C program
 2. Compile and run it
 3. Look at a debugger such as DDD, GDB

Next Time

Programs and Processes in C and UNIX

Read Chapter 2,3 (R&R), opt: Chapter 2 (MOS) or Chapter 3 (S&G)

Have a great weekend

Recitation on Monday: must attend