# Ayudantia 3

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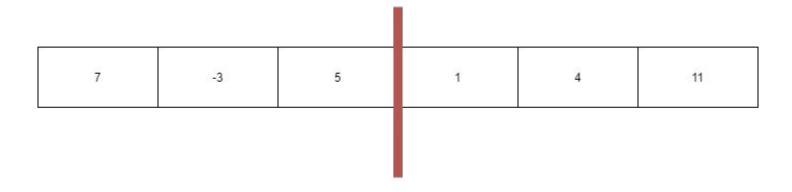
Sergio Gutiérrez: sergio.gutierrez@uc.cl

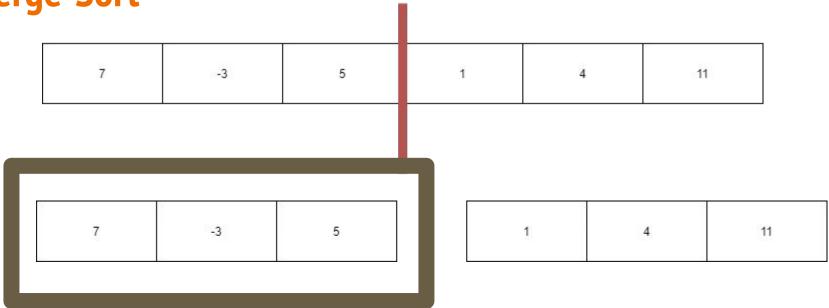
### **Divide and Conquer**

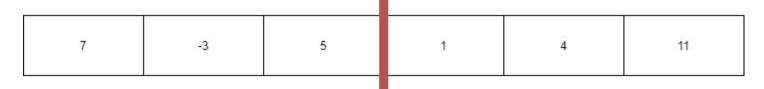
- Merge-sort
- Quick-sort

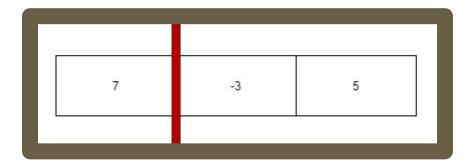
7 -3 5 1 4 11
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- Dividir nuestro array en 2
- Ordenar cada mitad de forma recursiva
- Unir las mitades

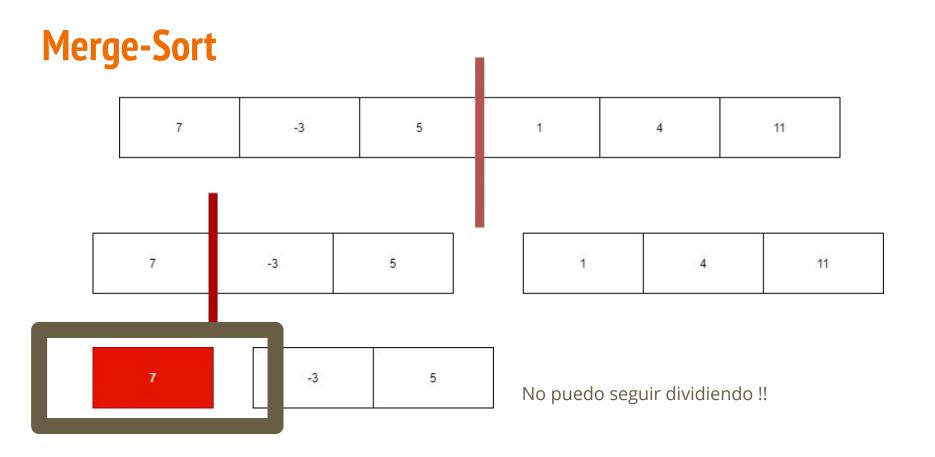


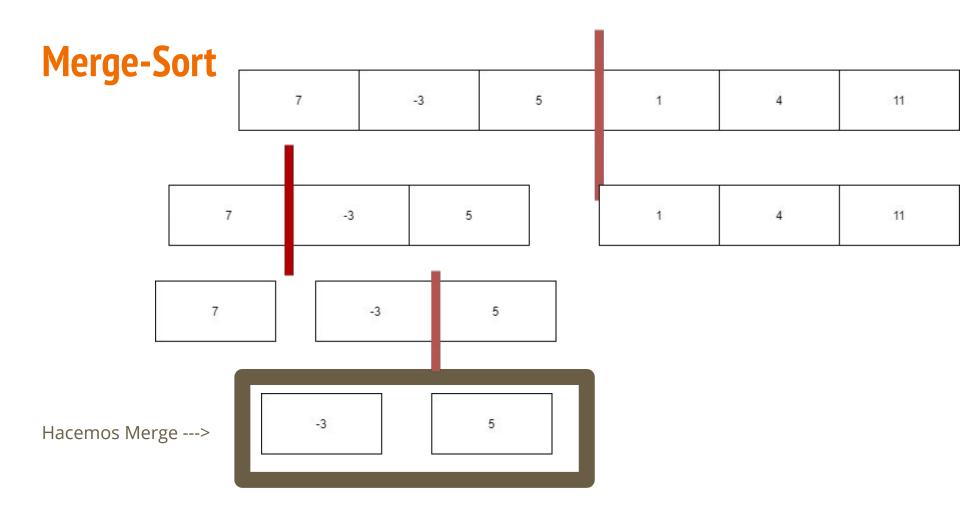


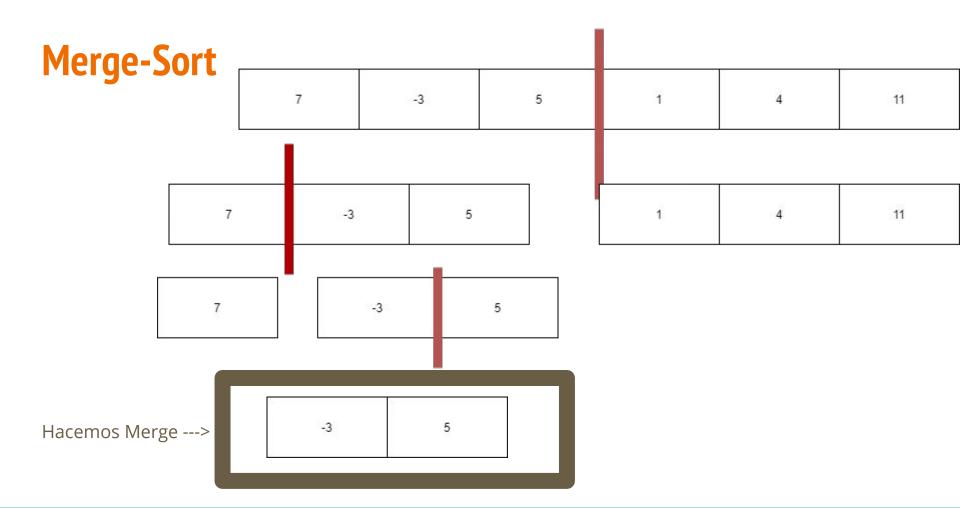


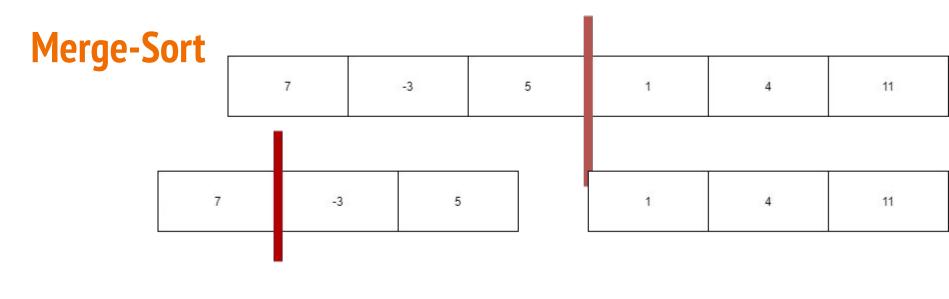


1	4	11

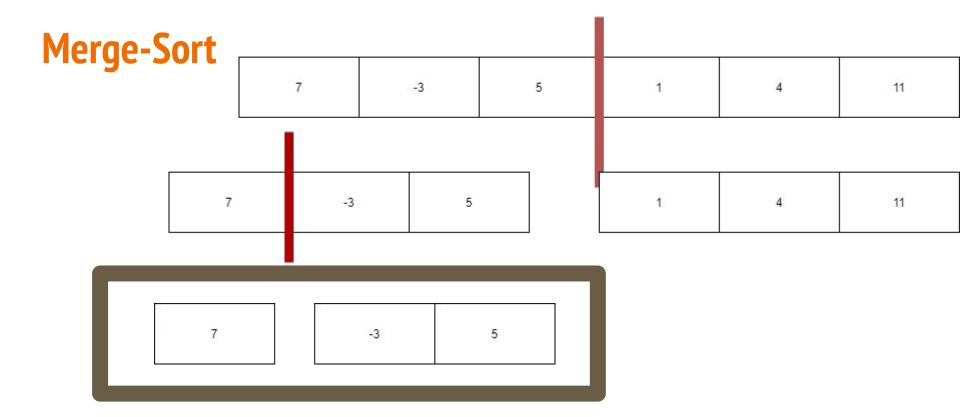


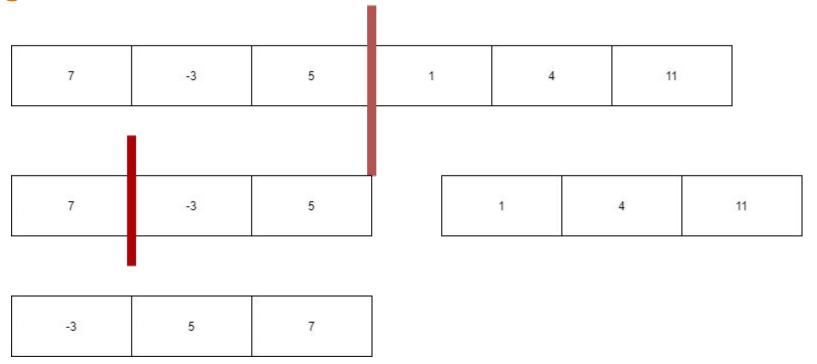


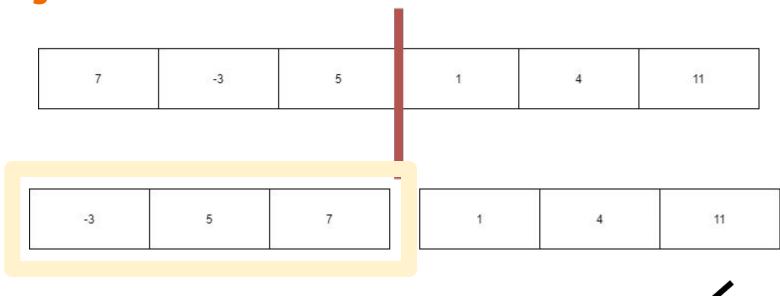




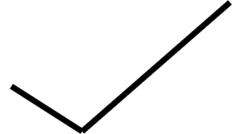
7 -3 5

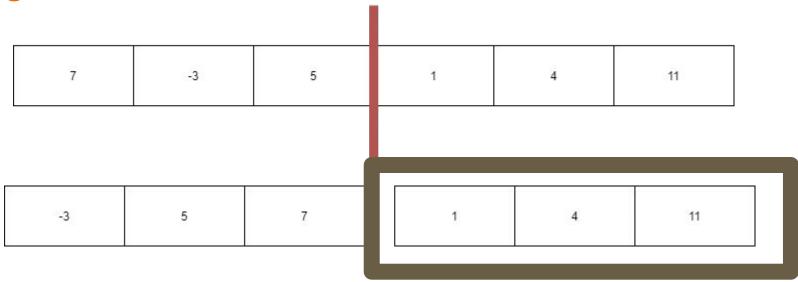




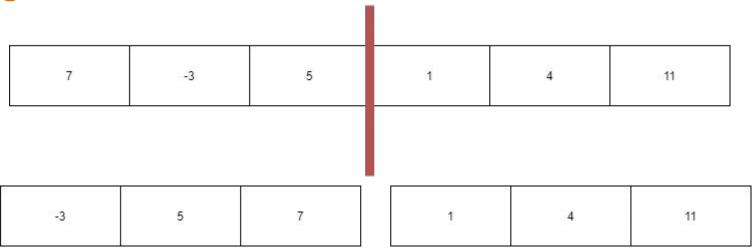


• Quedamos listos con el lado Izquierdo

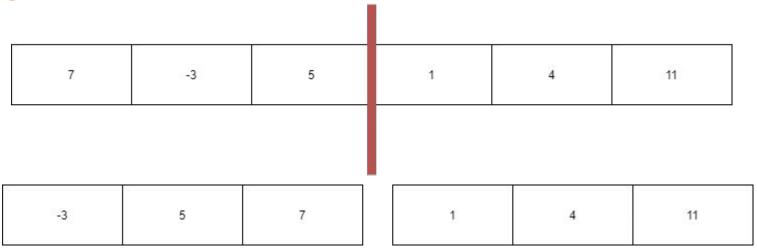




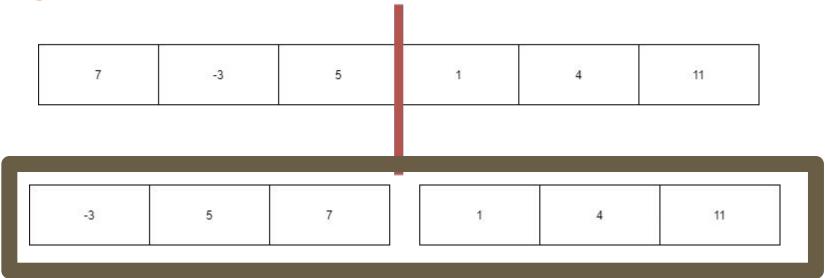
Aplicamos Merge al lado derecho



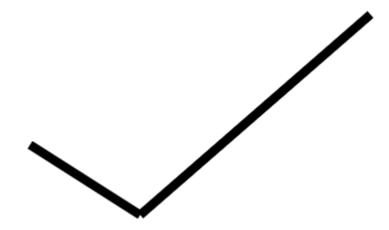
Hacemos FF



• Está Ordenado!



-3 1 4 5 7 11



```
void merge(int arr[], int 1, int m, int r)
   int i, j, k;
   int n1 = m - 1 + 1;
   int n2 = r - m;
   /* create temp arrays */
   int L[n1], R[n2];
   for (i = 0; i < n1; i++)
       L[i] = arr[l + i];
   for (j = 0; j < n2; j++)
       R[j] = arr[m + 1 + j];
   /* Merge the temp arrays back into arr[1..r]*/
   i = 0; // Initial index of first subarray
   j = 0; // Initial index of second subarray
   k = 1; // Initial index of merged subarray
   while (i < n1 && j < n2) {
       if (L[i] \leftarrow R[j]) {
           arr[k] = L[i];
           i++;
       else {
           arr[k] = R[j];
           j++;
       k++;
```

### Implementación C

```
/* Copy the remaining elements of L[], if there
are any */
while (i < n1) {
    arr[k] = L[i];
    i++;
    k++;
/* Copy the remaining elements of R[], if there
are any */
while (j < n2) {
    arr[k] = R[j];
   j++;
    k++;
```

#### Llamado recursivo

```
void mergeSort(int arr[], int 1, int r)
   if (1 < r) {
       // Same as (l+r)/2, but avoids overflow for
       // large 1 and h
        int m = 1 + (r - 1) / 2;
       // Sort first and second halves
       mergeSort(arr, 1, m);
       mergeSort(arr, m + 1, r);
       merge(arr, 1, m, r);
   UTILITY FUNCTIONS */
```

### QuickSort (breve repaso)

- Usar el algoritmo **partition** para dividir
- Ordenar cada partición de forma recursiva

#### **Partition**

Para entender QuickSort, primero hay que entender partition

```
partition(A, i, f):
         x \leftarrow un \ indice \ aleatorio \ en \ [i, f], \quad p \leftarrow A[x]
         A[x] \rightleftarrows A[f]
         j \leftarrow i
         for k \in [i, f-1]:
                    if A[k] < p:
                               A[j] \rightleftarrows A[k]
                               j \leftarrow j + 1
         A[j] \rightleftarrows A[f]
         return j
```

## Partition (ejemplo)

Partition(A, 0, 5)

i = 0

x = 4

f = 5

15

4

7

9

13

1

$$p = 7$$

### Partition (ejemplo)

Cambiamos A[x] por A[f]

i = 0

x = 4

f = 5

15

4

1

9

13

7

$$p = 7$$

#### **Entramos al For**

$$p = 7$$

15 > 7, no pasa nada, avanzamos

$$p = 7$$

$$p = 7$$

$$k = 2$$

# 1 < 7, se cambian 1 con 15 y se incrementa j, avanza k

$$p = 7$$

j = 2

k = 3

$$p = 7$$

j = 2

k = 4

#### Terminó el loop, se cambian A[j] con A[f] y finaliza, retorna j

$$p = 7$$

1

7

9

13

15

k = 4

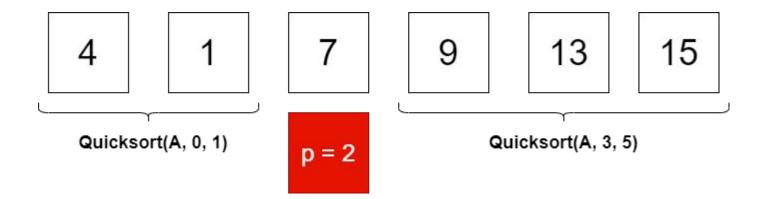
### QuickSort

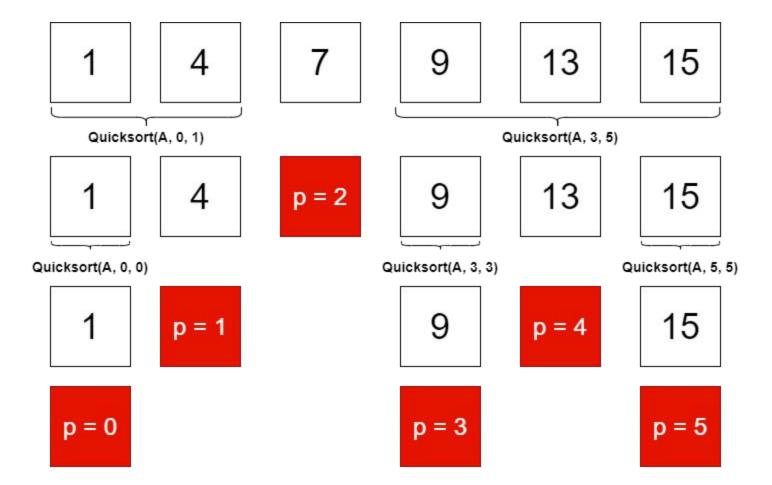
Ahora que tenemos claro partition, vemos quicksort

```
quicksort(A, i, f):
if \ i \leq f:
p \leftarrow partition(A, i, f)
quicksort(A, i, p - 1)
quicksort(A, p + 1, f)
```

### QuickSort (ejemplo)

Seguimos con el ejemplo de partition





### QuickSort (ejemplo)

Seguimos con el ejemplo de partition

