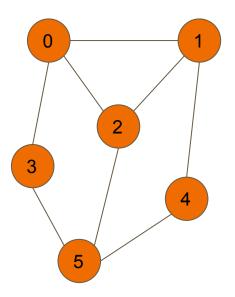
Ayudantía 10

Grafos y DFS

DFS (Iterative)

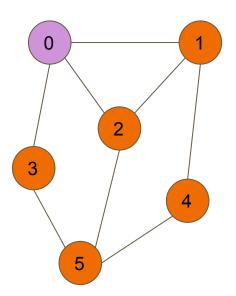
```
dfs(graph G, node start)
   stack s
   s.push(start)
   label start as discovered
   while not s.empty()
       node u = s.pop()
       for v in G.adjacent[u]
           if v is not discovered
               s.push(v)
               label v as discovered
```

DFS (Recursive)



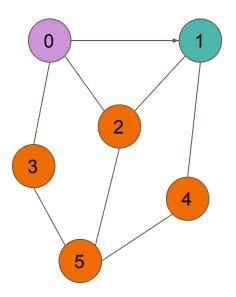
dfs(G, 0)

```
dfs(graph G, node u)
    label u as discovered
    for v in G.adjacent[u]
        if v not discovered
        dfs(G, v)
```



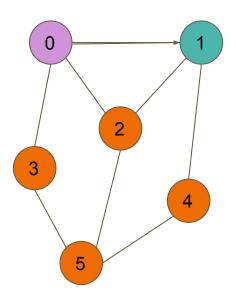
```
dfs(G, 0)

dfs(graph G, node u)
    label u as discovered
    for v in G.adjacent[u]
        if v not discovered
        dfs(G, v)
```



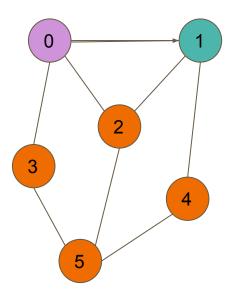
```
dfs(G, 0)

dfs(graph G, node u)
    label u as discovered
    for v in G.adjacent[u]
    if v not discovered
        dfs(G, v)
```



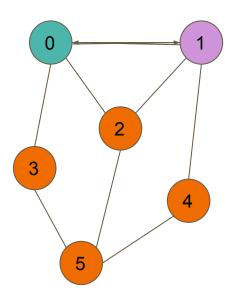
```
dfs(G, 0)

dfs(graph G, node u)
    label u as discovered
    for v in G.adjacent[u]
    if v not discovered
    dfs(G, v)
```



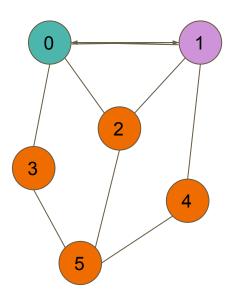
```
dfs(G, 0)

dfs(graph G, node u)
    label u as discovered
    for v in G.adjacent[u]
        if v not discovered
        dfs(G, v)
```



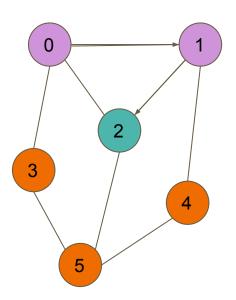
```
dfs(G, 0)

dfs(graph G, node u)
    label u as discovered
    for v in G.adjacent[u]
    if v not discovered
        dfs(G, v)
```



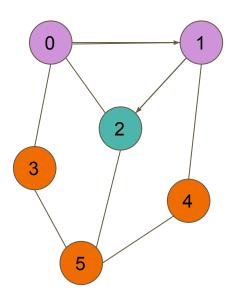
```
dfs(G, 0)

dfs(graph G, node u)
    label u as discovered
    for v in G.adjacent[u]
    if v not discovered
        dfs(G, v)
```



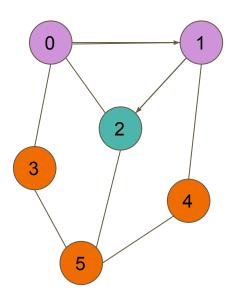
```
dfs(G, 0)

dfs(graph G, node u)
    label u as discovered
    for v in G.adjacent[u]
    if v not discovered
        dfs(G, v)
```



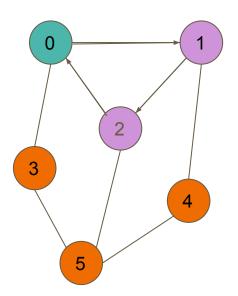
```
dfs(G, 0)

dfs(graph G, node u)
    label u as discovered
    for v in G.adjacent[u]
    if v not discovered
        dfs(G, v)
```



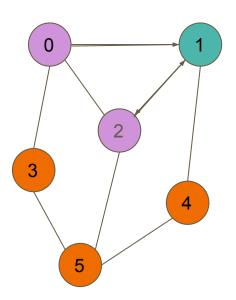
```
dfs(G, 0)

dfs(graph G, node u)
    label u as discovered
    for v in G.adjacent[u]
    if v not discovered
        dfs(G, v)
```



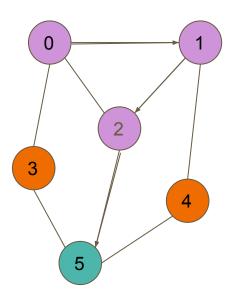
```
dfs(G, 0)
```

```
dfs(graph G, node u)
    label u as discovered
    for v in G.adjacent[u]
        if v not discovered
        dfs(G, v)
```



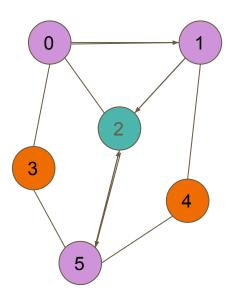
```
dfs(G, 0)
```

```
dfs(graph G, node u)
    label u as discovered
    for v in G.adjacent[u]
        if v not discovered
        dfs(G, v)
```



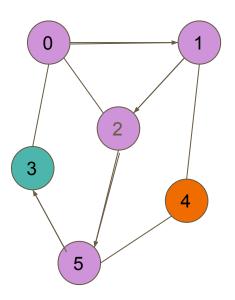
```
dfs(G, 0)
```

```
dfs(graph G, node u)
    label u as discovered
    for v in G.adjacent[u]
        if v not discovered
        dfs(G, v)
```



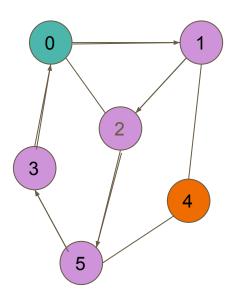
```
dfs(G, 0)
```

```
dfs(graph G, node u)
    label u as discovered
    for v in G.adjacent[u]
        if v not discovered
        dfs(G, v)
```



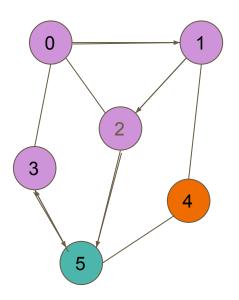
```
dfs(G, 0)
```

```
dfs(graph G, node u)
    label u as discovered
    for v in G.adjacent[u]
        if v not discovered
        dfs(G, v)
```



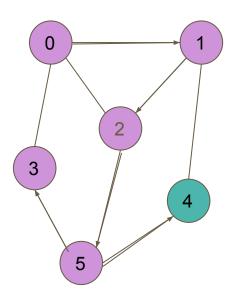
```
dfs(G, 0)
```

```
dfs(graph G, node u)
    label u as discovered
    for v in G.adjacent[u]
        if v not discovered
        dfs(G, v)
```



```
dfs(G, 0)
```

```
dfs(graph G, node u)
    label u as discovered
    for v in G.adjacent[u]
        if v not discovered
        dfs(G, v)
```

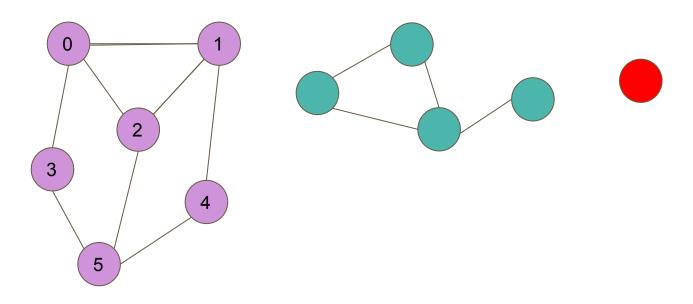


```
dfs(G, 0)
```

```
dfs(graph G, node u)
    label u as discovered
    for v in G.adjacent[u]
        if v not discovered
        dfs(G, v)
```

Nota: El algoritmo continúa por un más, pero no se mostrará el resto

Contar componentes conexas en un grafo no dirigido G



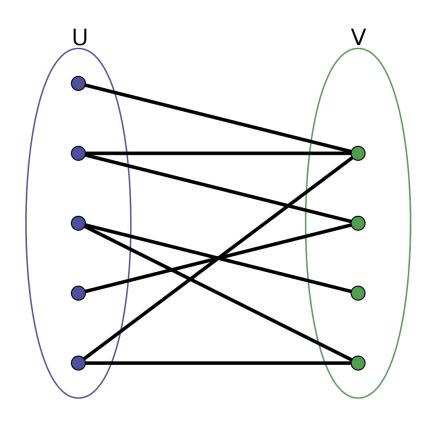
Notar: DFS recorre toda una componente conexa.

Idea: Ir por cada nodo y ver si ha sido visitado, si no correr DFS desde el nodo y sumar 1

Dado un grafo G verificar si es bipartito

Definición: Un grafo G es bipartito, si es que se puede particionar en dos conjuntos U y V, tal que ningún nodo en U es vecino de otro nodo en U, y análogamente para V.

Bipartito



Definición Bi-coloración de un Grafo: Sea un grafo G, se llama a un función f de los nodos del grafo a {0,1} una bi-coloración de G, si cada nodo tiene un color distinto a sus vecinos.

Nota: Un grafo es bi-coloreable si y sólo si es bipartito.

Idea: Usar DFS para intentar bi-colorear un grafo.

Componente Fuertemente Conexa

Dado un grafo dirigido G, un subconjunto de nodos A es una componente fuertemente conexa (SCC (Strongly Connected Component)) si para cada par de nodos u, v en A existe un camino de u a v y un camino de v a u, y si es maximal (no está contenido en una SCC más grande).

Nota 1: Si G es no dirigido, se tiene que una componente es conexa si y sólo si es fuertemente conexa.

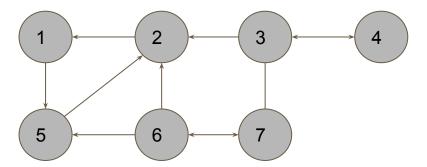
Nota 2: Dado dos SCC distintas A y B de un grafo dirigido G, la intersección de A con B es vacía.

Dado un grafo dirigido G, encontrar todas las SCC.

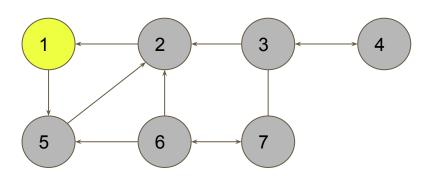
Solución: Algoritmo de Tarjan para SCC.

Idea general: Usar un stack e índices con un DFS para marcar las

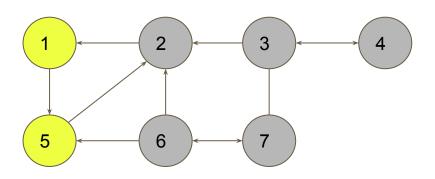
componentes



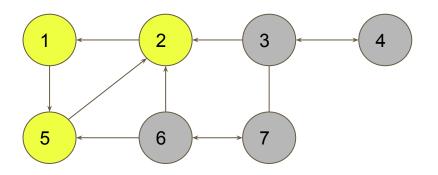
Stack



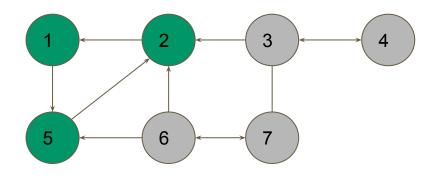
Stack



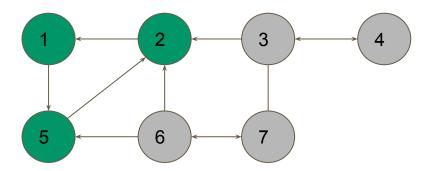
Stack



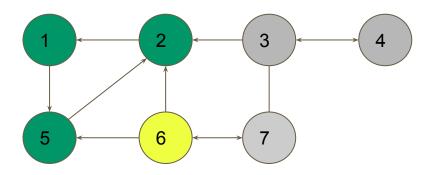
Stack



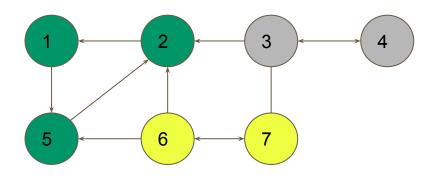
Stack



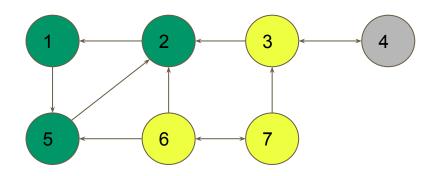
Stack



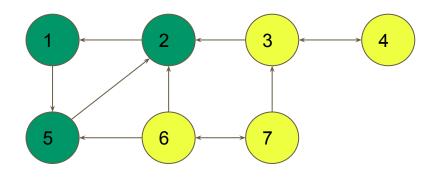
Stack



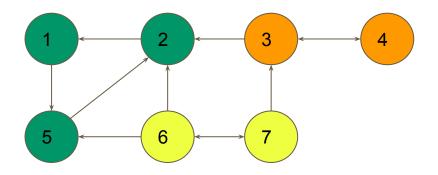
Stack



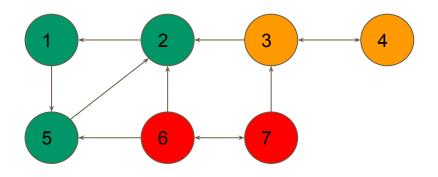
Stack



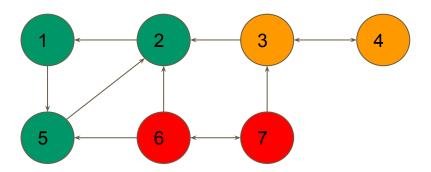
Stack



Stack



Stack



Stack

```
stack s
id = 0
tarjan(graph G)
      s = empty stack
      id = 0
      for v in G
            if v.id = -1
                  dfs(G,v)
dfs(graph G, node v)
      v.id = id
      v.lowlink = id
      id++
      s.push(v)
      v.onStack = true
```

```
for u in G.adjacent[v]
      if v.id = -1
            dfs(G, v)
            v.lowlink =min(v.lowlink, u.lowlink)
      else if v.onStack
            v.lowlink =min(v.lowlink, u.lowlink)
if v.lowlink = v.id
      do
            u = s.pop()
            u.onStack = false
            add u to current SCC
      while u≠v
      output current SCC
```