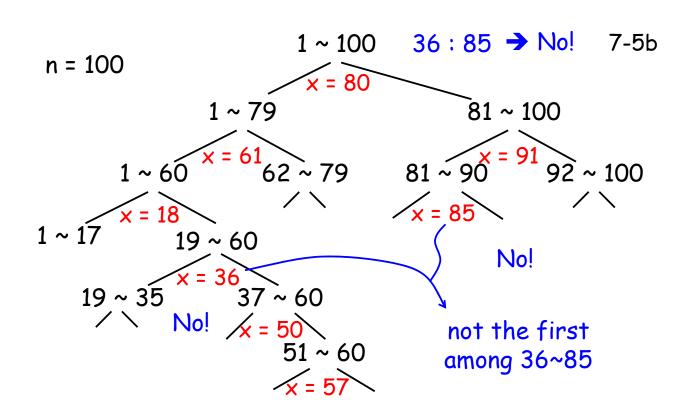
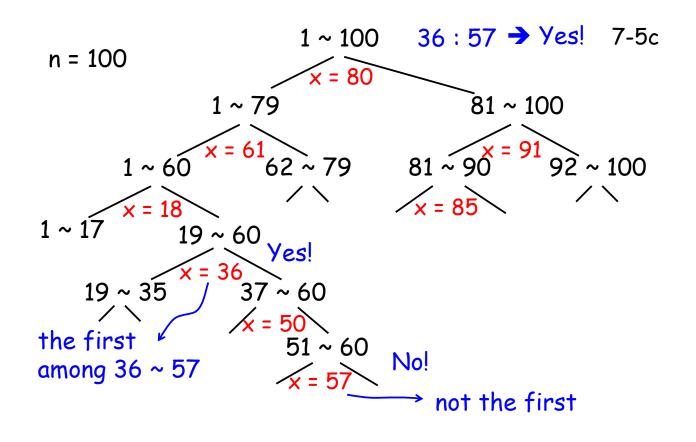
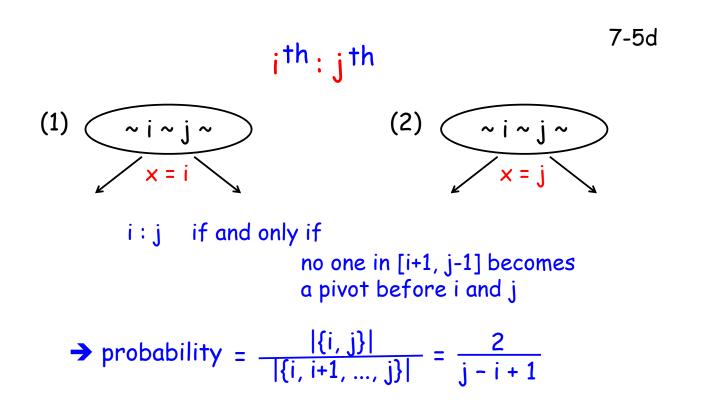


- -*Every number serves as a pivot at most once!
- *Each call is on "consecutive" numbers!

tat most n internal nodes at most 2n+1 calls to Quicksort at most n+1 leaves







Quick Sort: Average Case

$$E(n) = (n - 1) +$$

rage Case
$$E(n) = (n-1) + \begin{cases}
1/n * (E(1) + E(n-1)) \\
1/n * (E(2) + E(n-3)) \\
\vdots & \vdots \\
1/n * (E(n-1) + E(n-3))
\end{cases}$$

=
$$(n-1) + \frac{1}{n} \sum_{q=1}^{n} (E(q-1) + E(n-q))$$

=
$$(n - 1) + \frac{2}{n} \sum_{k=1}^{n-1} E(k)$$

n, bn, $\Theta(n)$

Classic approach:

$$E(n) = \begin{cases} \Theta(n) \\ bn \\ n \\ n-1 \end{cases} + \frac{2}{n} \sum_{k=1}^{n-1} E(k) \quad \begin{array}{l} \text{guess } E(n) = O(n \mid g \mid n) \\ \text{prove by substitution method} \\ \text{(very hard to understand!!!)} \end{cases}$$

Knuth's approach:
$$E(n) = n+1 + \frac{2}{n} \sum_{k=1}^{n-1} E(k)$$

$$E(n) = n+1 + \frac{2}{n} \sum_{k=1}^{n-1} E(k)$$

$$E(n) = \frac{n+1}{n} E(n-1) + 2$$

$$E(n) = \frac{n+1}{n} E(n-1$$

Randomized-variable approach:

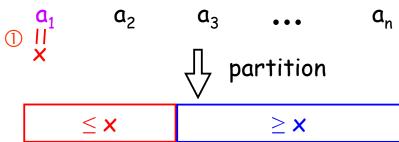
| lomized-variable approach:

$$E(n) = O(n + X) = O(n + \sum_{i,j} Pro(i : j))$$
| assume independent

7-8b

7-8a

Pro. 7-1 7-10a



 $^{\circ}$ *a₁ may be either $^{\square}$ or $^{\square}$

*guarantee
$$|\Box| \ge 1$$
 $|\Box| \ge 1$