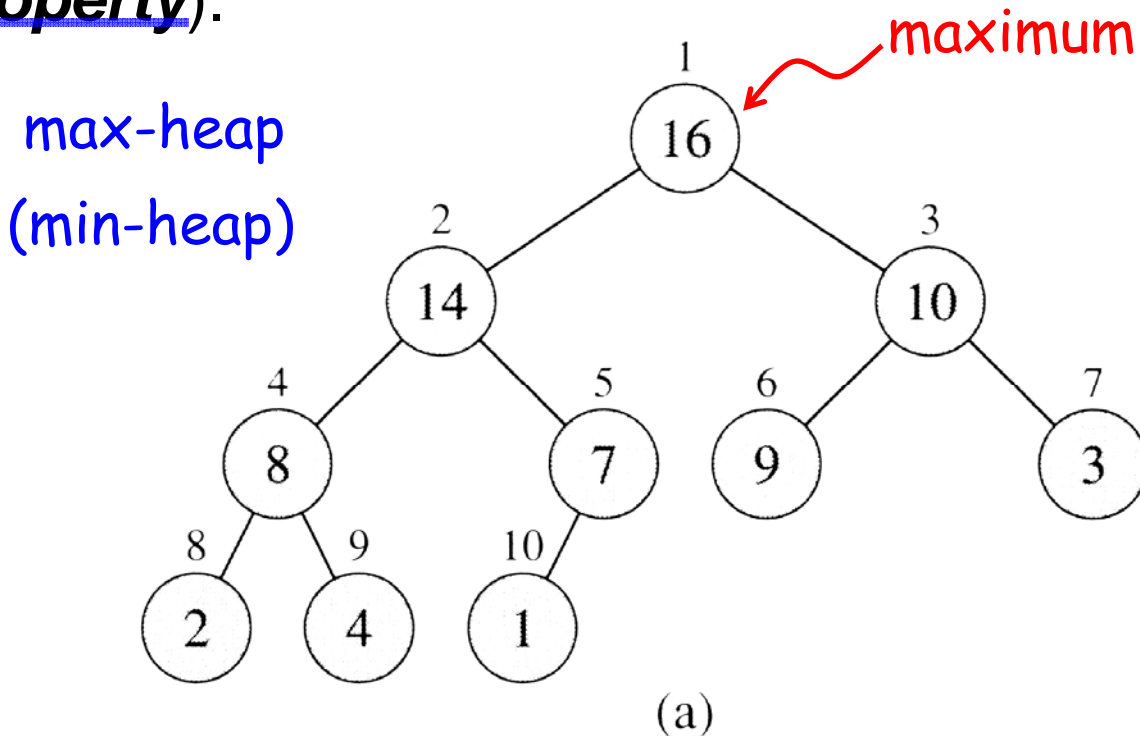


Heapsort

(binary) ①

6.1 **Heap**: an array that can be viewed as a complete binary tree in which each node has a key not smaller than those of its children ② (heap property).



1	2	3	4	5	6	7	8	9	10
16	14	10	8	7	9	3	2	4	1

(b)

- $A[1]$ is the root
- $Parent(i) = \lfloor i/2 \rfloor$ (shifting i right one bit)
- $Left(i) = 2i$ (shifting i left one bit)
- $Right(i) = 2i+1$
- A heap of n nodes has height $\Theta(\lg n)$ ≡ $\lfloor \lg n \rfloor$
max # of edges from root to a leaf

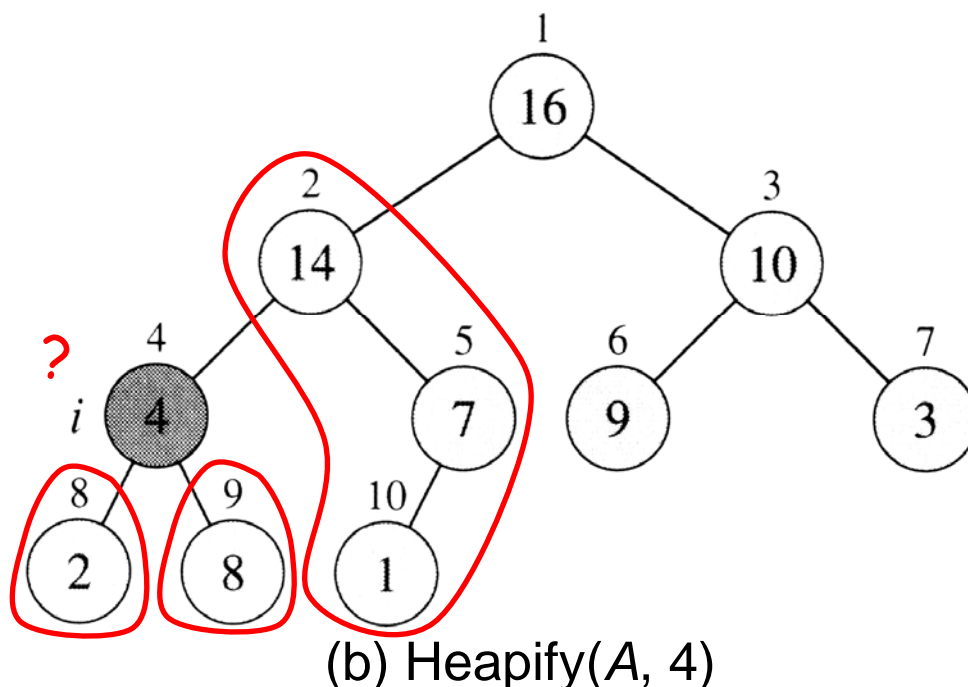
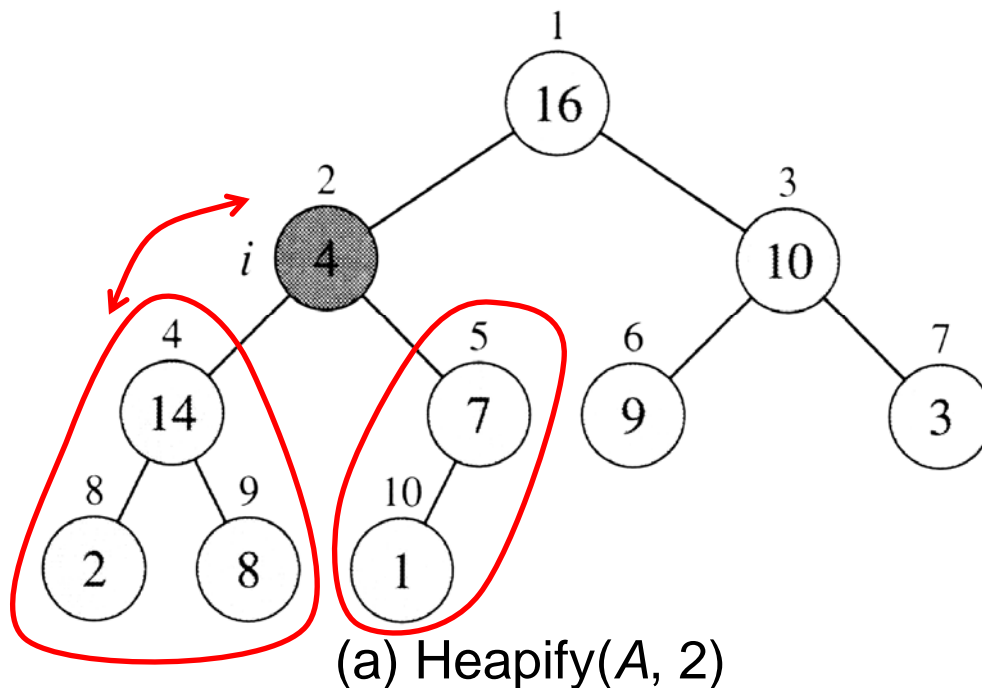
6.2 Maintaining the heap property

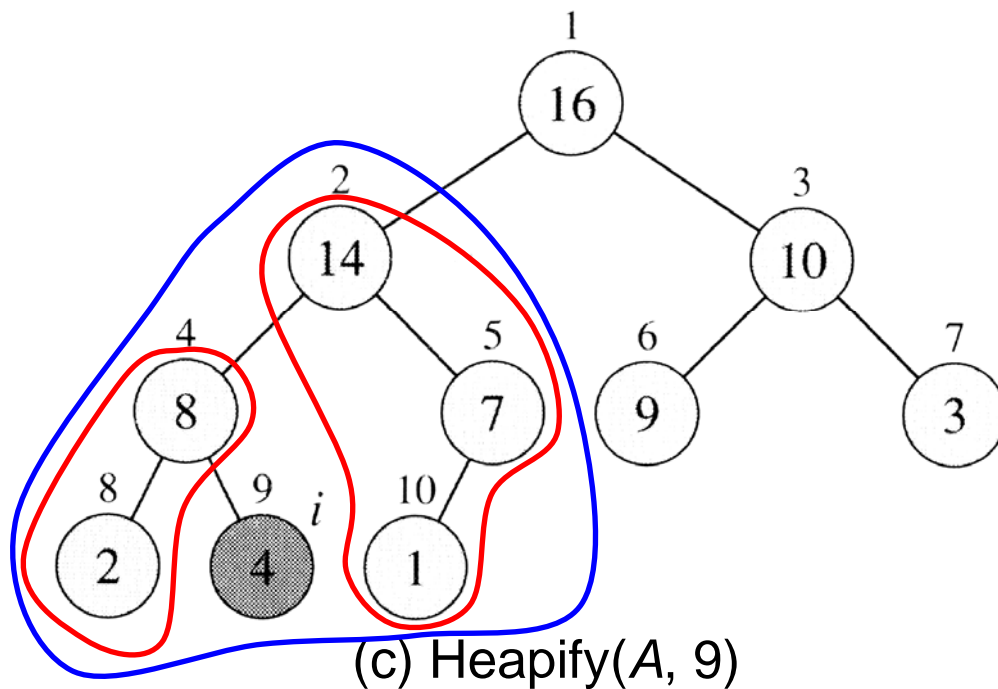
Heapify(A, i): Binary trees rooted at Left(i) and Right(i) are heaps, but A[i] may be smaller than its children.

6-2a

* merge two smaller heap into one

Example:





6-3x

① 用 recurrence \rightarrow subtree i 中 node # $\rightarrow i$ 的高度

- $T(n) \leq T(2n/3) + \Theta(1) = O(\lg n) = O(h)$,
where n is the number of nodes in the subtree rooted at $A[i]$ and h is the height of $A[i]$.

② 用 高度 說明

6.3 Building a heap

Build-Heap(A)

```

1  heap-size[ $A$ ]  $\leftarrow$  length[ $A$ ]
2  for  $i \leftarrow \lfloor \text{length}[A]/2 \rfloor$  downto 1 do
3    Heapify( $A, i$ )
  
```

$$T(n) \leq \sum_{h=0}^{\lfloor \lg n \rfloor} \left[\frac{n}{2^{h+1}} \right] O(h) = O(n \times \sum_{h=0}^{\lfloor \lg n \rfloor} \frac{h}{2^h}) = O(n \times 2)$$

6-3a

$= O(n)$. (Ex. 6.3-3)

(at most # nodes at h)

$$\sum_{k=1}^{\infty} kx^k = \frac{x}{(1-x)^2} \quad (x < 1)$$

OR: $T(n) = 2T(\lfloor n/2 \rfloor) + \lg n$ (Assume $n = 2^{h+1} - 1$.)

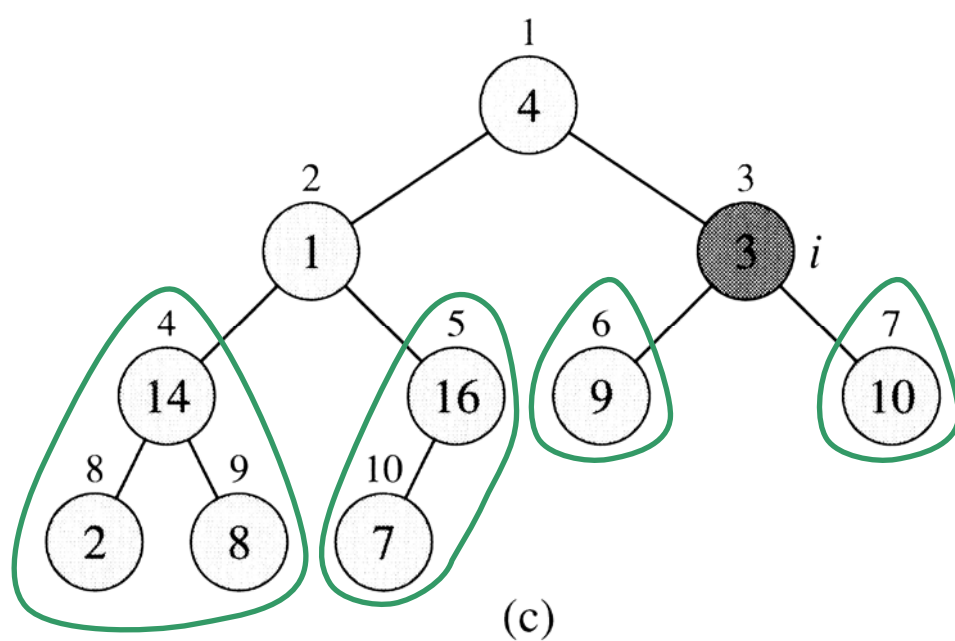
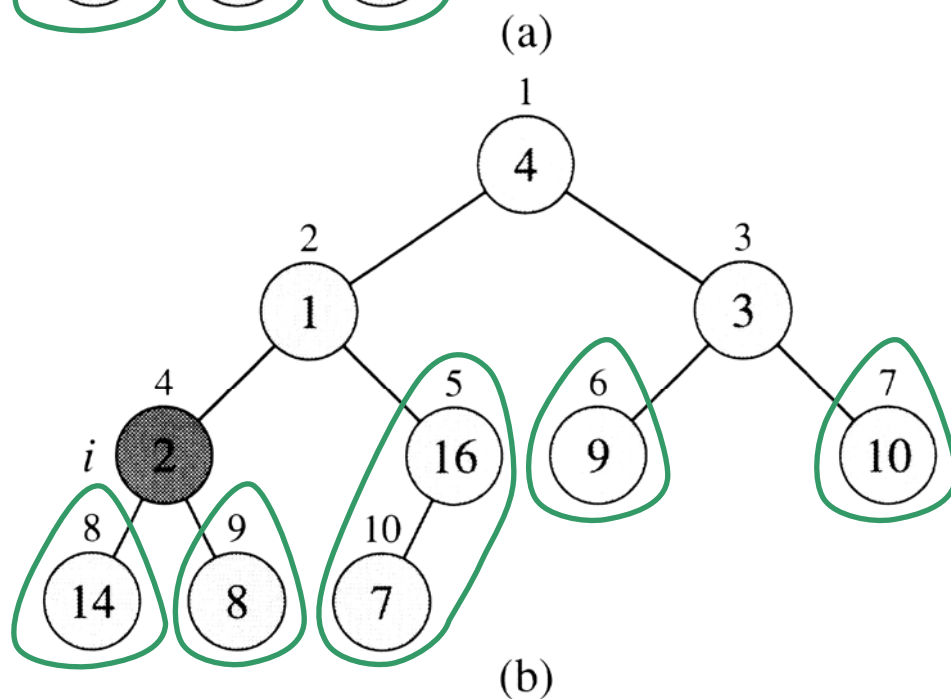
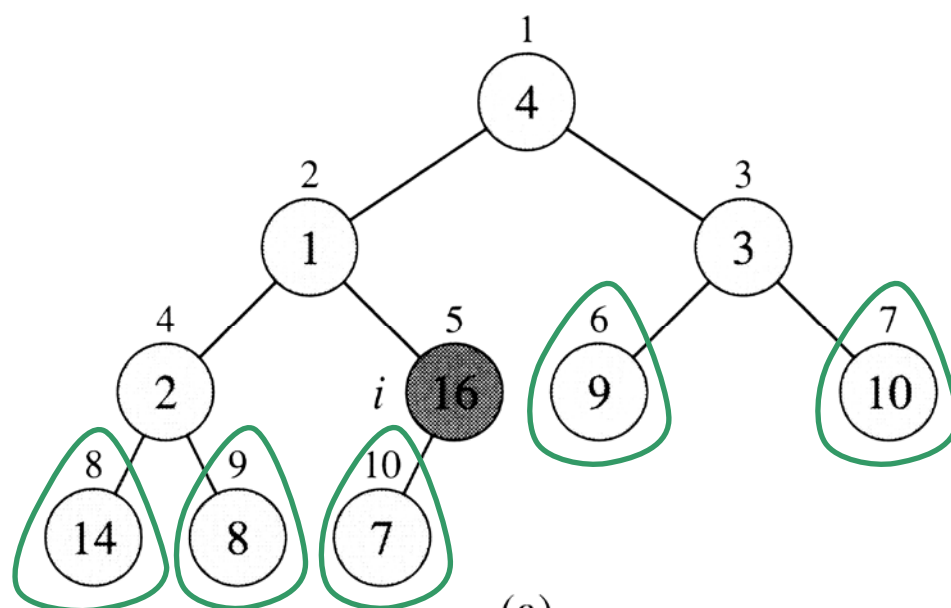
6-3y

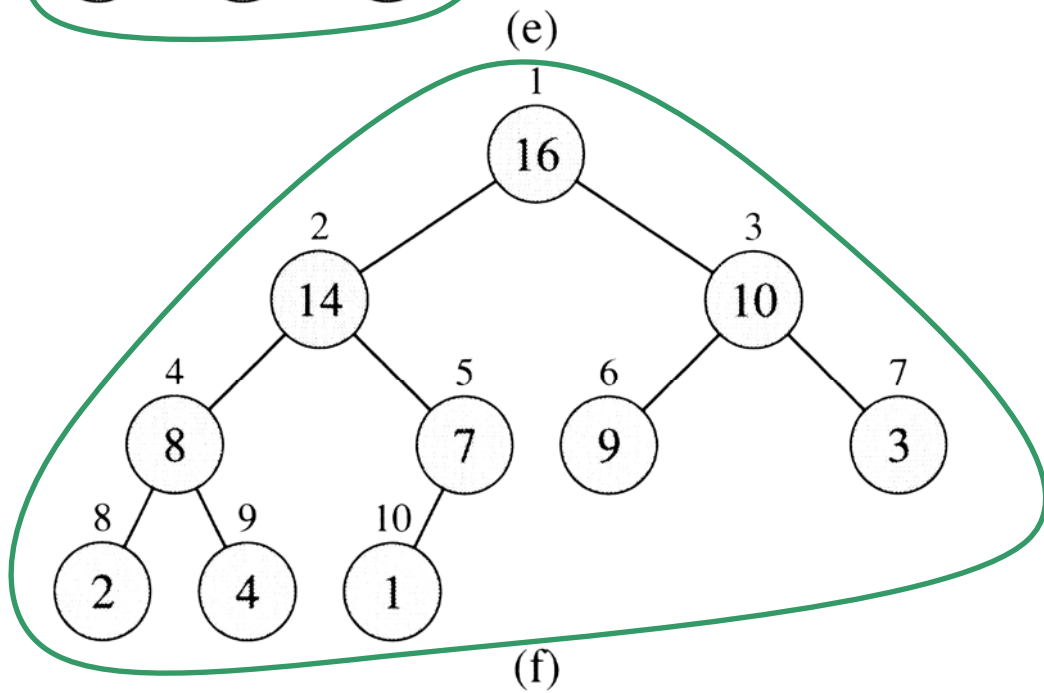
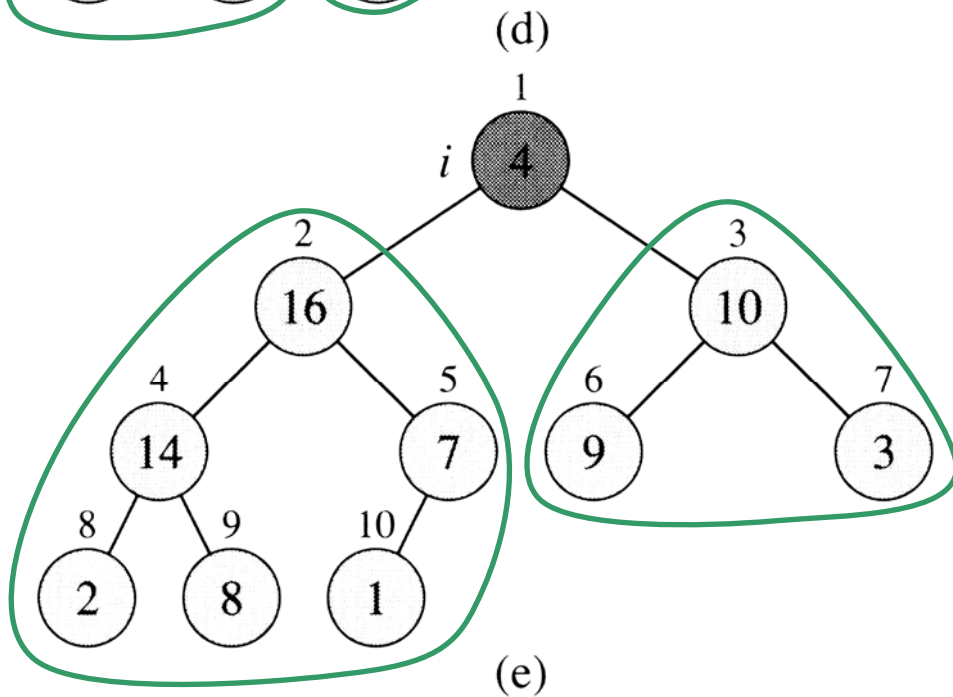
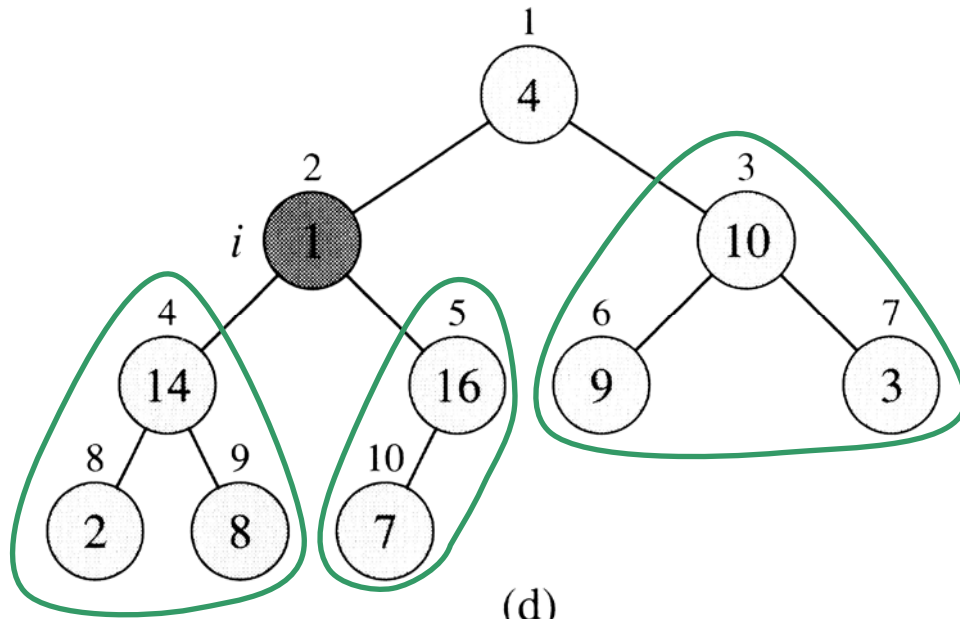
Append dummy nodes

Example:

A	4	1	3	2	16	9	10	14	8	7
---	---	---	---	---	----	---	----	----	---	---

6-4





6.4 Heapsort

Stage 1: Build a heap

Stage 2: Repeatedly delete the root of the heap. 6-6x

HeapSort(*A*)

```

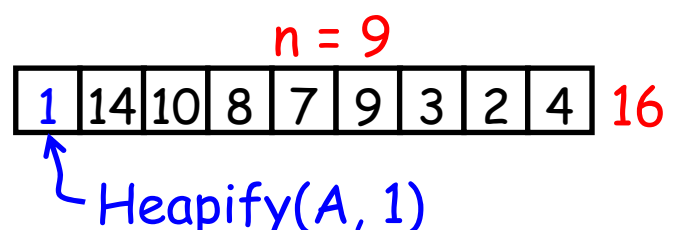
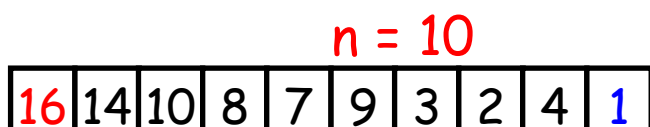
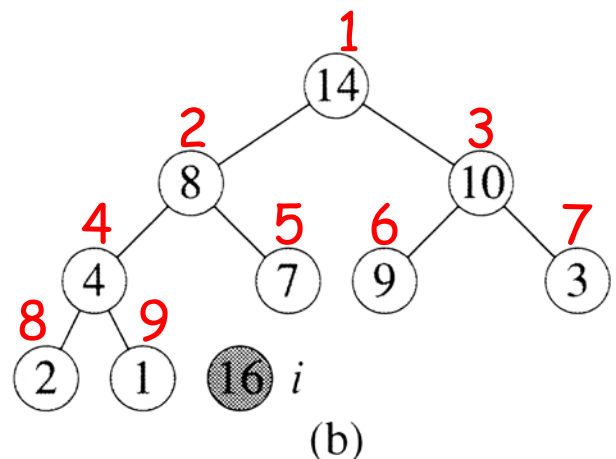
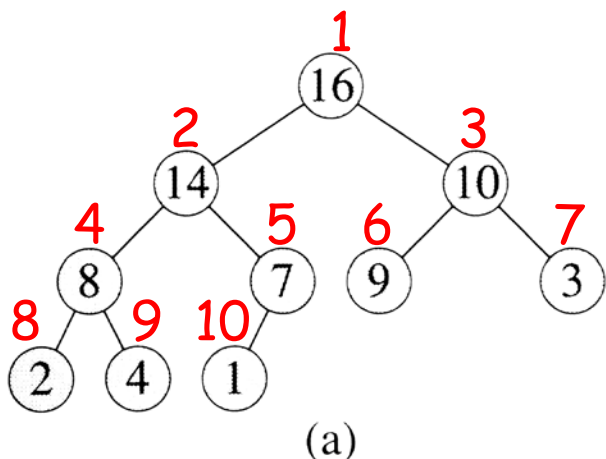
1  Build-Heap(A)
2  for i ← length[A] downto 2
3  { do exchange A[1] ↔ A[i]
4    heap-size[A] ← heap-size[A] − 1
5    Heapify(A, 1)

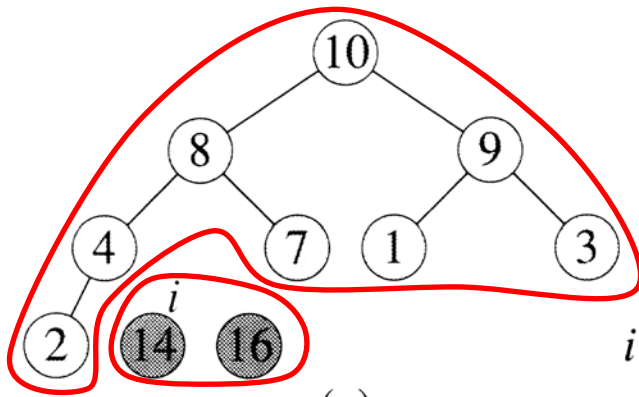
```

- $T(n) = O(n) + O((n-1) \times \lg n)$
 $= \underline{O(n \lg n)}$

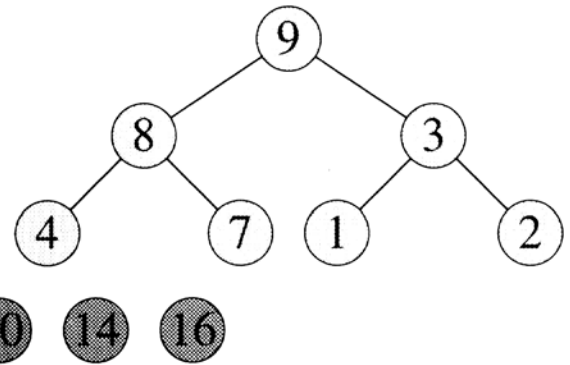
6-6y

Example: sort $A = \{4, 1, 3, 2, 16, 9, 10, 14, 8, 7\}$.

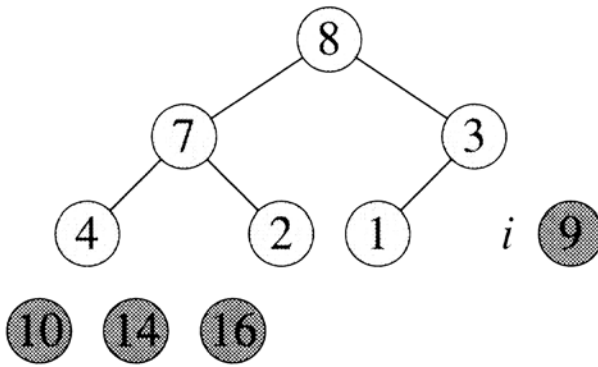




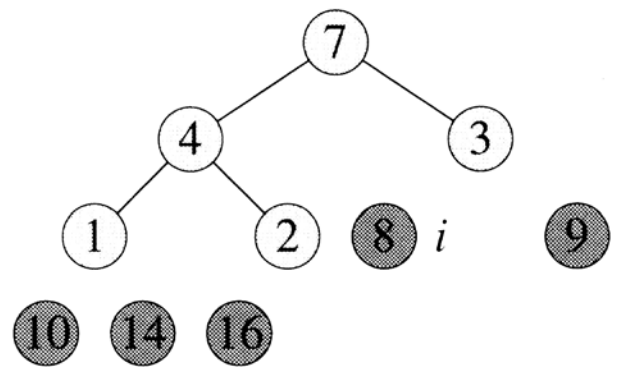
(c)



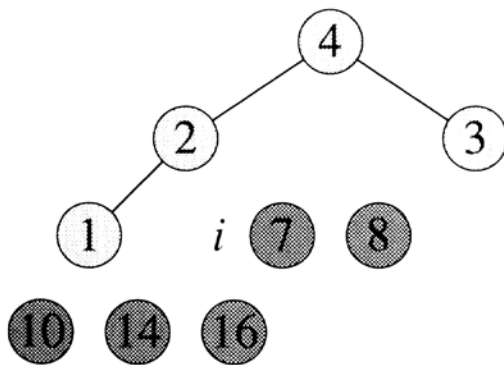
(d)



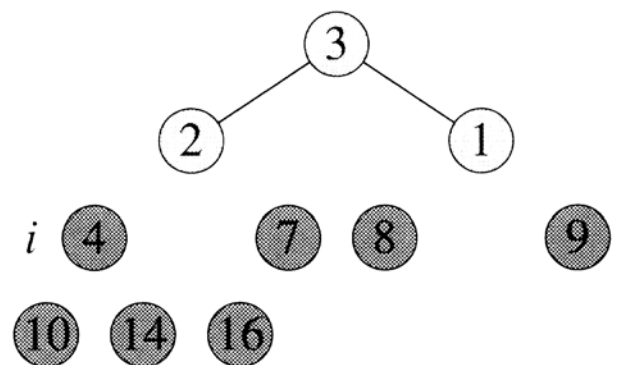
(e)



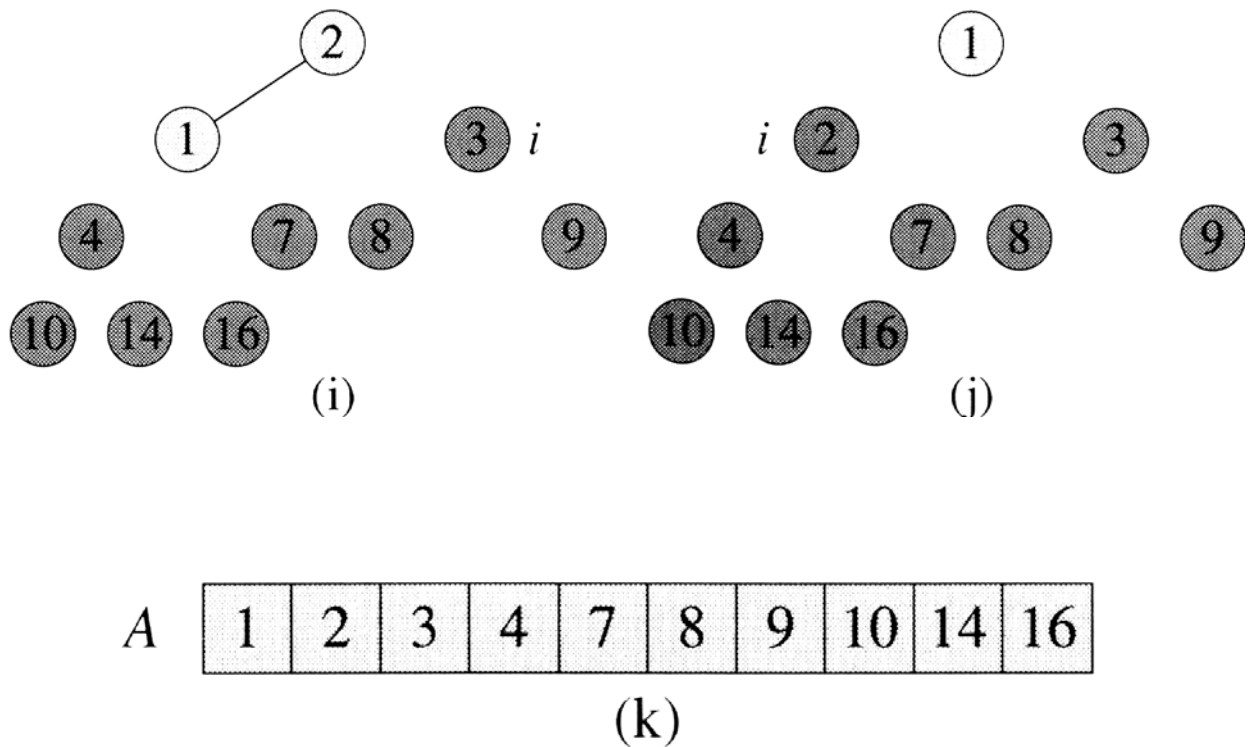
(f)



(g)



(h)



6.5 Priority queues

Priority queue: a data structure for maintaining a set A of elements, each has a value called a **key**. It should support the following operations.

*for convenience, assume "element = key"

Insert(A, x): * x has a key

Maximum(A): (return)

Extract-Max(A): (return and remove)

Increase-Key(A, a, k): increase a 's key to larger key k , where a is an element of A

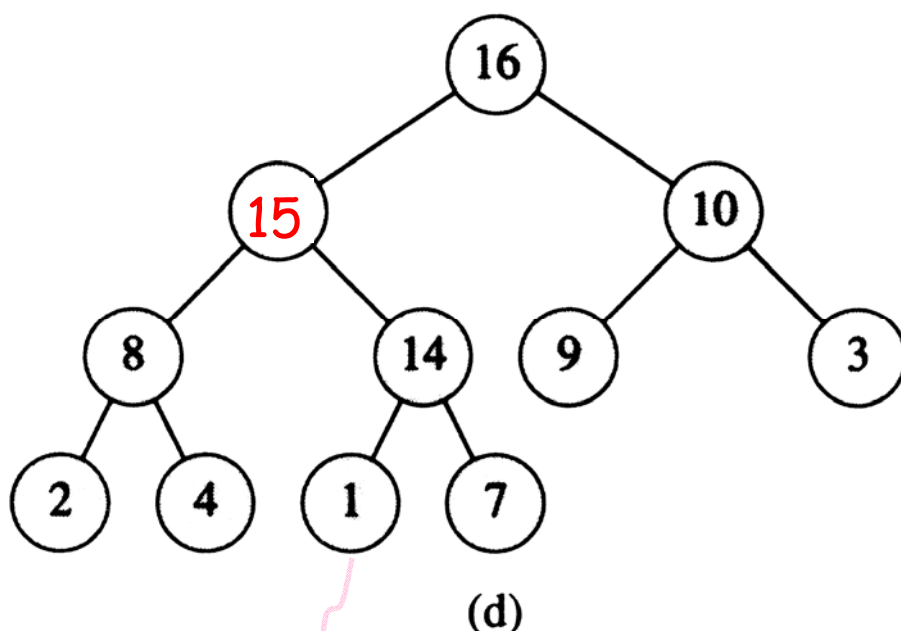
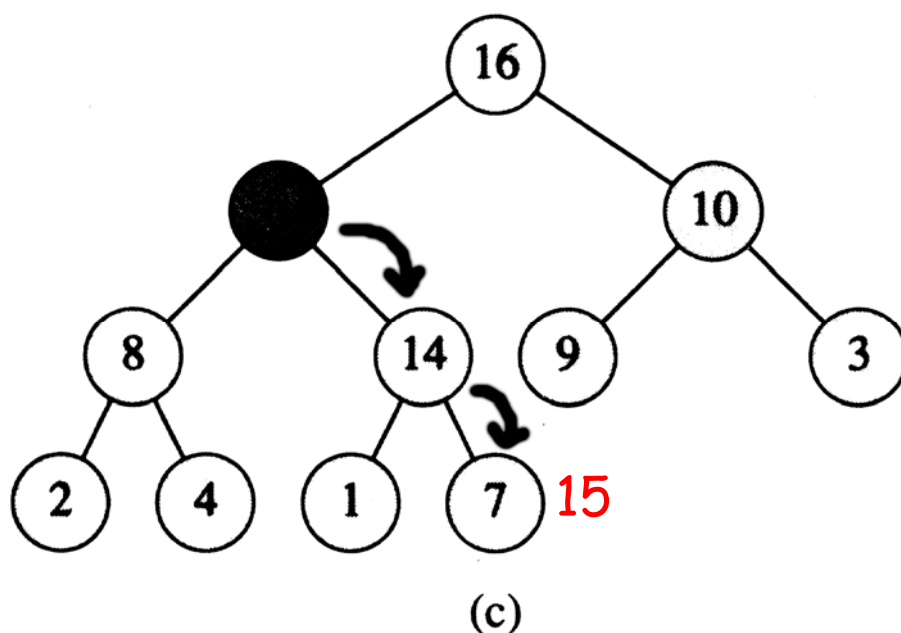
Applications: Job scheduling on a shared computer based upon "priorities."


```
graph TD; 16((16)) --- 14((14)); 16 --- 10((10)); 14 --- 8((8)); 14 --- 7((7)); 8 --- 2((2)); 8 --- 4((4)); 7 --- 1((1)); 10 --- 9((9)); 10 --- 3((3));
```

```
graph TD; 16((16)) --- 14((14)); 16 --- 10((10)); 14 --- 8((8)); 14 --- 7((7)); 8 --- 2((2)); 8 --- 4((4)); 7 --- 1((1)); 7 --- black(( )); 10 --- 9((9)); 10 --- 3((3));
```

15 (do nothing)

Diagram illustrating a heap data structure. The array contains the values 16, 14, 10, 8, 7, 9, 3, 2, 4, 1. A blue bracket underneath the first nine elements is labeled "heap". A vertical dashed line is positioned between the last element (1) and the next element, which is a shaded box. The array continues with an empty box.



increase to 15.5

Insert(A, x)

```

1  heap-size(A) ← heap-size(A)+1
2  i ← heap-size(A)
3  { while i > 1 and A[Parent(i)] < x
4    do A(i) ← A[Parent(i)]  往下拉
5    i ← Parent(i)  向上 check
6  A[i] ← x
```

Increase-Key(A, i, k): $O(\lg n)$ time
(similar to *Insert*)

Maximum(A): $O(1)$ time

Extract-Max(A): $O(\lg n)$ time

Step 1: Exchange $A[1]$ and $A[\text{heap-size}]$ $n = 10$

Step 2: $\text{heap-size} \leftarrow \text{heap-size} - 1$ $n = 9$

Step 3: *Heapify*($A, 1$)

Step 4: return $A[\text{heap-size} + 1]$ $n+1 = 10$

Homework: Ex. 6.2-5, 6.5-9, Prob. 6-2, **6-3**

* array implementation

* handles: pointers to the objects in a data structure