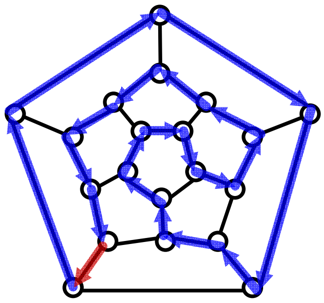
The Hamiltonian cycle problem

Input: G = (V, E)



a Hamiltonian cycle H

A nondeterministic algorithm

Step 1: Guess a cycle H

Step 2: Verify whether H is a Hamiltonian cycle

(i) all edges exist?

(ii) visit each vertex exactly once?

What dose "solve" mean?

⇨ an algorithm that runs in polynomial time (P)

What is the target?

Set of all problems

- ~~un-computable~~
- computable
  - ~~problems whose answers can not be verified in $O(n^k)$ time~~
  - problems whose answers can be verified in $O(n^k)$ time (NP)

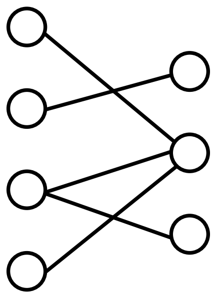
Reduction:

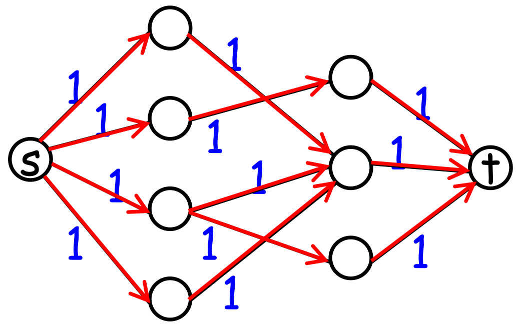
Problem A $\xrightarrow[\text{reduction}]{O(T)}$ Problem B

Example:

max bipartite matching $\xrightarrow[\text{reduction}]{O(V + E)}$ maximum flow

G = (V, E)





Partition Problem:

$S = (a_1, a_2, \ldots, a_n)$ ⇨ S_1 and S_2 such that $Sum(S_1) = Sum(S_2)$

3-Partition Problem:

$S = (a_1, a_2, \ldots, a_n)$ ⇨ S_1, S_2, S_3 such that $Sum(S_1) = Sum(S_2) = Sum(S_3)$

Partition problem $\xrightarrow[\text{reduction}]{O(n)}$ 3-partition problem

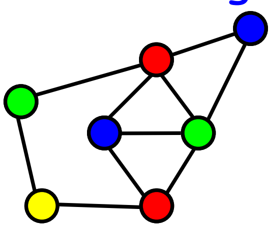
$S = (a_1, a_2, \ldots, a_n)$ $\qquad$ $S' = (a_1, a_2, \ldots, a_n, Sum(S)/2)$

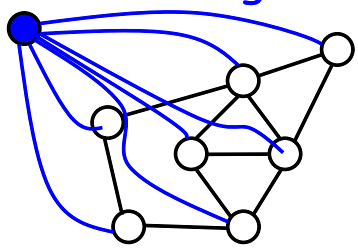
$S = (3, 5, 7, 9)$ $\qquad$ $S' = (3, 5, 7, 9, 24/2)$

$= (3, 5, 7, 9, 12)$

Coloring: Given G, assign color to each node such that adjacent nodes have different color.

3-coloring $\xrightarrow[\text{reduction}]{O(n)}$ 4-coloring





3-coloring: Determine whether a given G can be colored by using {0, 1, 2}.

4-coloring: Determine whether a given G can be colored by using {0, 1, 2, 3}.

Reduction:

Problem A $\xrightarrow[\text{reduction}]{O(T)}$ Problem B

* usually: easy $\Rightarrow$ hard (or, as hard as)
  A can be solved by solving B
  B is more difficult

* may: hard $\Rightarrow$ easy

Example:

embed the difficulty

Find(A, x) $\xrightarrow[\text{sort } A]{O(n \lg n)}$ Binary Search(S, x)

O(n) O(lg n)

(i) If $A \xrightarrow{O(n^a)} B$ then $A: O(n^a + n^b)$

? $O(n^b)$

imply a solution of A

(ii) If $A \xrightarrow{O(n^b)} B$ then B: ?

$O(n^a)$?

imply nothing for B

(iii) If $A \xrightarrow{O(n^b)} B$ then

$\Omega(n^a)$?

may imply difficulty of B

case 1. $\Omega(n^{10})$ $O(n^5)$

case 2. $\Omega(n^{10})$ $O(n^{11})$

- case 1: $b < a \Rightarrow B: \Omega(n^a)$

  /* e.g. B: $O(n^9)$ then A: $O(n^5) + O(n^9) = \cancel{O}(n^9)$

- case 2: $b \geq a \Rightarrow B$: ?

  /* may: hard $\Rightarrow$ easy, or easy $\Rightarrow$ hard

Problem $A \xrightarrow{O(T_1)}$ Problem B

$O(T_2)$

imply a solution of A

* $A: O(T_1 + T_2)$

* Assume that T_1 is polynomial

數學家心中的 reduction 都是 polynomial

$\Rightarrow$ $A \in P$ if T_2 is polynomial

$\Rightarrow$ $A \in P$ if $B \in P$

解掉 B 就解掉 A

* If $A \xrightarrow{O(n^k)} B$, then $A \in P$ if $B \in P$

(i.e., A can be solved by solving B.)

解掉 B 就解掉全世界 (NP)

* If all NP $\xrightarrow{O(n^k)} B$, then all NP $\in P$ if $B \in P$

(i.e., all NP can be solved by only solving B.)

(i.e., NP = P if $B \in P$.)

假設真的存在

What is the goal?

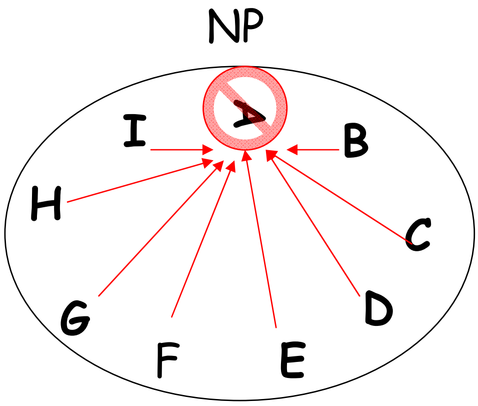
⇨ solve all problems in NP (in polynomial time) at a time
(i.e., prove NP = P)

$A \in$ NP 表示 A 是一個對手
$A \in$ P 表示 A 被解決了

How can "all problems" be solved at a time?

⇨ Idea: reduction

(1) find a problem A in NP such that all problems in NP can be reduced to A in polynomial time
(2) solve A in polynomial time

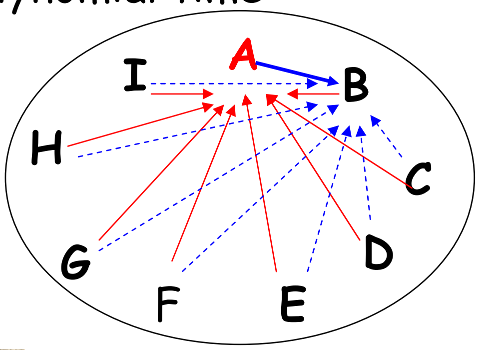


such a problem A is NP-C
(if exists)

A is NP-C:
(1) A is in NP
(2) all NP problems can be reduced to A in polynomial time

Assume that there exists an NP-C A.

If A can be reduced to B in polynomial time
(1) all NP problems can also be reduced to B
(2) B is NP-C
(3) B is as hard as A

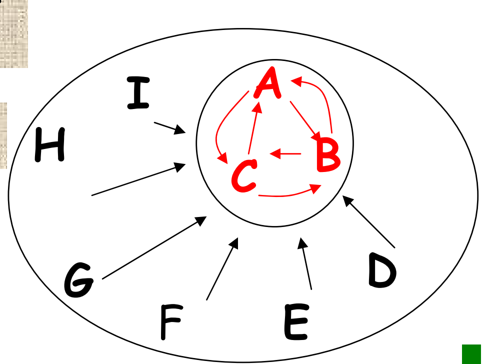


All NP-C problems are of the same difficulty.
(They can be reduced to each other.)

all NP $\Rightarrow_p$ B $\cong$ an NP-C $\Rightarrow_p$ B $\cong$ all NP-C $\Rightarrow_p$ B

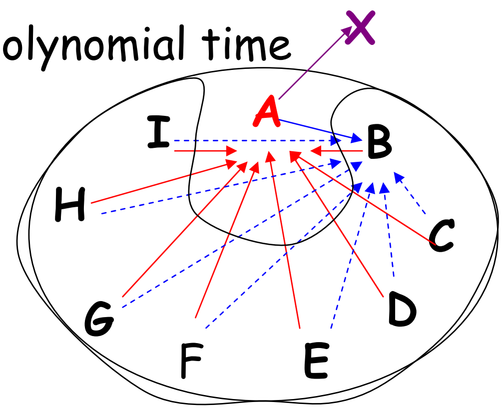
If an NP-C is solved

⇨ all NP (including all NP-C) are solved,



A is NP-C:
(1) A is in NP
(2) all NP problems can be reduced to A in polynomial time

A is NP-H: only (2)

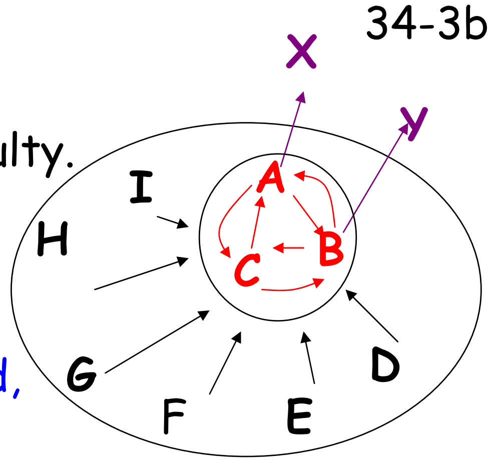


Assume that there exists an NP-C A.

If A can be reduced to B in polynomial time
(1) all NP problems can also be reduced to B
(2) B is NP-C
(3) B is as hard as A

If A can be reduced to an X $\notin$ NP in polynomial time
(1) all NP problems can be reduced to X
(2) X is NP-H, but not NP-C
(3) X is harder than A

NP-H: A, B, C, X, Y
NP-C: A, B, C



All NP-C problems are of the same difficulty.
But, all NP-H problems are not.

If an NP-H or NP-C is solved

⇨ all NP (including all NP-C) are solved,
but not all NP-H

Goal: Solve all problems in NP at a time.

How: Solve a problem in NP-C?

QUESTION: Dose NP-C exist?

Note that it is impossible to solve a
problem in NP-H, but not in NP-C. Why?

How to prove a problem A is in NP-C (NP-H)?

① Show that A is in NP.

give an $O(n^k)$-time nondeterministic algo for A

② Show that all in NP $\xrightarrow{O(n^k)}$ A.

(a) Find a problem $Y \in$ NP-C

(b) Show $Y \xrightarrow{O(n^k)} A$

all NP $\xrightarrow{O(n^k)} Y \xrightarrow{O(n^k)} A$

$O(n^k)$

(omit ① for NP-H)

Example: 34-4a

Prove 3-partition $\in$ NP-C.

① Show that A is in NP.

(i) guess S_1, S_2, S_3

(ii) check $S_1 \cup S_2 \cup S_3 = S$ and $\mathrm{Sum}(S_1) = \mathrm{Sum}(S_2) = \mathrm{Sum}(S_3)$

$O(n \lg n)$ time

② Show that all in NP $\xrightarrow{O(n^k)}$ A.

(a) It is known 2-partition $\in$ NP-C

(b) 2-partition $\xrightarrow{O(n)}$ 3-partition

(a) NP-C: No one knows how to solve these problems. (Y/N) 34-4b

(No algorithms exist for these problems.)

(b) If we consider NP as an army, then

NP-C: ? NP-H but not NP-C: ?

If an NP-C surrenders, then

all NP too? all NP-C too? all NP-H too?

If an NP-H surrenders, then

all NP too? all NP-C too? all NP-H too?

(1) How to prove a problem is P?

(2) How to prove a problem is NP?

(3) How to prove a problem is NP-C?

(4) How to prove a problem is NP-H?

(5) How to prove NP = P?

(6) How to prove NP $\neq$ P?

(7) NP = P or NP $\neq$ P ?

(8) What do we learn ?

Optimization Problems

P
* trivial
* greedy
* DP

NP-hard
* brute-force ($n \leq 20$) ⇒ exact solution
* B & B ($n \leq 30 \sim 100$)
* n > 100 ???

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* approximation ⇒ near-optimal solution
* heuristic