

CS5314 RANDOMIZED ALGORITHMS

Homework 2 (Suggested Solution)

1. **Ans:**

(a) Firstly, by linearity of expectation, we have

$$E[X^2] = E[X \sum_i X_i] = \sum_i E[XX_i].$$

Then,

$$\begin{aligned} E[XX_i] &= E[E[XX_i | X_i]] \\ &= \Pr(X_i = 1)E[XX_i | X_i = 1] + \Pr(X_i = 0)E[XX_i | X_i = 0] \\ &= \Pr(X_i = 1)E[X | X_i = 1] + \Pr(X_i = 0) \cdot 0 = \Pr(X_i)E[X | X_i = 1]. \end{aligned}$$

Combining the above gives the desired result.

(b) $E[X | X_i = 1] = (n - 1)p + 1.$

(c) $\text{Var}(X) = E[X^2] - (E[X])^2 = np((n - 1)p + 1) - (np)^2 = np(1 - p).$

2. **Ans:**

(a) $\Pr(Y_i = 0) = \Pr(HH) + \Pr(TT) = 1/2$, and $\Pr(Y_i = 1) = 1 - \Pr(Y_i = 0) = 1/2$.

(b) If i th pair and j th pair do not share any bit,

$$E[Y_i Y_j] = \Pr(Y_i = 1 \cap Y_j = 1) = \Pr(Y_i = 1)\Pr(Y_j = 1) = E[Y_i]E[Y_j].$$

Otherwise, they share exactly one bit, so that

$$\begin{aligned} E[Y_i Y_j] &= \Pr(Y_i = 1 \cap Y_j = 1) \\ &= \Pr(\text{unshared bits have same value, and opposite to shared bit}) \\ &= 1/4 = E[Y_i]E[Y_j]. \end{aligned}$$

(c) $\text{Cov}(X, Y) = E[(Y_i - E[Y_i])(Y_j - E[Y_j])] = E[Y_i Y_j] - E[Y_i]E[Y_j] = 0.$

(d)

$$\begin{aligned} \Pr(|Y - E[Y]| \geq n) &\leq \frac{\text{Var}(Y)}{n^2} \\ &= \frac{\sum_i \text{Var}(Y_i) + \sum_{i,j} \text{Cov}(Y_i, Y_j)}{n^2} \\ &= \frac{\binom{n}{2} \cdot 1/4 + 0}{n^2} \leq 1/8. \end{aligned}$$

3. **Ans.** For each i , we have

$$E[X_i] = \Pr(X_i = 1) = 1/n \quad \text{and} \quad E[X_i^2] = \Pr(X_i^2 = 1) = 1/n.$$

For any i, j with $i \neq j$, we have

$$\begin{aligned} E[X_i X_j] &= \Pr(X_i X_j = 1) = \Pr(X_i = 1 \cap X_j = 1) \\ &= \Pr(X_i = 1) \Pr(X_j = 1 \mid X_i = 1) = \frac{1}{n(n-1)}. \end{aligned}$$

Let X be the number of fixed points. So, $X = \sum_{i=1}^n X_i$, and $E[X] = \sum_{i=1}^n E[X_i] = 1$. Then, we have

$$\begin{aligned} \text{Var}[X] &= E[X^2] - (E[X])^2 = E \left[\left(\sum_{i=1}^n X_i \right)^2 \right] - 1 \\ &= \sum_{i=1}^n E[X_i^2] + \sum_{i \neq j} E[X_i X_j] - 1 = n \cdot \frac{1}{n} + n(n-1) \cdot \frac{1}{n(n-1)} - 1 = 1. \end{aligned}$$

4. (a) **Ans.** Let $X = n\tilde{p}$ denote the number of heads that came up. So, $X = \text{Bin}(n, p)$ and $E[X] = np$. Then we have:

$$\begin{aligned} \Pr(|p - \tilde{p}| > \varepsilon p) &= \Pr(|np - n\tilde{p}| > n\varepsilon p) \\ &= \Pr(np - n\tilde{p} > n\varepsilon p) + \Pr(n\tilde{p} - np > n\varepsilon p) \\ &= \Pr(X < (1 - \varepsilon)E[X]) + \Pr(X > (1 + \varepsilon)E[X]) \\ &\leq \exp\left(\frac{-na\varepsilon^2}{2}\right) + \exp\left(\frac{-na\varepsilon^2}{3}\right). \end{aligned}$$

- (b) **Ans.** When

$$n > \frac{3 \ln(2/\delta)}{a\varepsilon^2},$$

we have:

$$\frac{na\varepsilon^2}{3} > \ln(2/\delta) \quad \text{so that} \quad \delta > 2 \exp\left(\frac{-na\varepsilon^2}{3}\right).$$

Combining this with the result of part (a), we have:

$$\Pr(|p - \tilde{p}| > \varepsilon p) \leq \exp\left(\frac{-na\varepsilon^2}{2}\right) + \exp\left(\frac{-na\varepsilon^2}{3}\right) < 2 \exp\left(\frac{-na\varepsilon^2}{3}\right) < \delta.$$

5. (a) **Ans.** We shall make use of the following claim:

Claim 1. For any $r \in [0, 1]$, $e^{tr} - 1 \leq r(e^t - 1)$.

Proof. Let $f(r) = r(e^t - 1) - e^{tr} + 1$. Then we have $f'(r) = (e^t - 1) - te^{tr}$, and $f''(r) = -t^2 e^{tr} \leq 0$. This implies that f is a concave function.

In other words, for $r \in [0, 1]$, f achieves minimum value either at the boundaries $f(0)$ or $f(1)$. Thus, $f(r) \geq \min\{f(0), f(1)\} = 0$ for all $r \in [0, 1]$, and the claim follows. \square

Back to the answer. Since $W = \sum_{i=1}^n a_i X_i$, we have

$$\nu = E[W] = \sum_{i=1}^n a_i E[X_i] = \sum_{i=1}^n a_i p_i.$$

For any i ,

$$E[e^{ta_i X_i}] = p_i e^{ta_i} + (1 - p_i) = 1 + p_i(e^{ta_i} - 1) \leq 1 + p_i a_i (e^t - 1),$$

where the last inequality is from Claim 1.

Hence,

$$E[e^{ta_i X_i}] \leq e^{p_i a_i (e^t - 1)},$$

and by the independence of X_i 's and property of MGF,

$$E[e^{tW}] = \prod_{i=1}^n E[e^{ta_i X_i}] \leq \prod_{i=1}^n e^{a_i p_i (e^t - 1)} = e^{\nu(e^t - 1)}.$$

For any $t > 0$, we have

$$\Pr(W \geq (1 + \delta)\nu) = \Pr(e^{tW} \geq e^{t(1+\delta)\nu}) \leq \frac{E[e^{tW}]}{e^{t(1+\delta)\nu}} \leq \frac{e^{\nu(e^t - 1)}}{e^{t(1+\delta)\nu}}.$$

Then, for any $\delta > 0$, we can set $t = \ln(1 + \delta) > 0$ and obtain:

$$\Pr(W \geq (1 + \delta)\nu) < \left(\frac{e^\delta}{(1 + \delta)^{(1+\delta)}} \right)^\nu.$$

(b) **Ans.** For any $t < 0$, we have

$$\Pr(W \leq (1 - \delta)\nu) = \Pr(e^{tW} \geq e^{t(1-\delta)\nu}) \leq \frac{E[e^{tW}]}{e^{t(1-\delta)\nu}} \leq \frac{e^{\nu(e^t - 1)}}{e^{t(1-\delta)\nu}}.$$

Then, for any $0 < \delta < 1$, we can set $t = \ln(1 - \delta) < 0$ and obtain:

$$\Pr(W \leq (1 - \delta)\nu) < \left(\frac{e^{-\delta}}{(1 - \delta)^{(1-\delta)}} \right)^\nu.$$