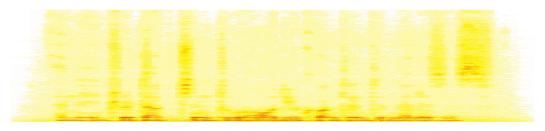
Introduction to Audio Content Analysis

Module 7.0: Dynamic Time Warping

alexander lerch



introduction

overview

Georgia Center for Music Tech College of Design

corresponding textbook section

Chapter 7: Alignment (pp. 139–146)

- lecture content
 - Dynamic Time Warping (DTW): synchronization of two sequences with similar content
- learning objectives
 - explain the standard DTW algorithm
 - discuss disadvantages of and modifications to the standard DTW algorithm
 - implement DTW



introduction

overview

Georgia Center for Music Tech Technology

corresponding textbook section

Chapter 7: Alignment (pp. 139–146)

lecture content

 Dynamic Time Warping (DTW): synchronization of two sequences with similar content

learning objectives

- explain the standard DTW algorithm
- discuss disadvantages of and modifications to the standard DTW algorithm
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dynamic time warping problem statement

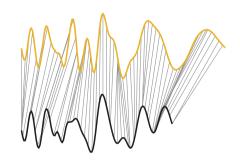
intro

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- align/synchronize two sequences
 - similar musical content
 - different tempo and timing

$$A(n_{\rm A})$$
 $n_{\rm A} \in [0; \mathcal{N}_{\rm A} - 1]$
 $B(n_{\rm B})$ $n_{\rm B} \in [0; \mathcal{N}_{\rm B} - 1]$

- ⇒ find the alignment path that
 - minimizes the pairwise distance between sequences
 - covers the whole sequence
 - does only move forward in time



dynamic time warping overview

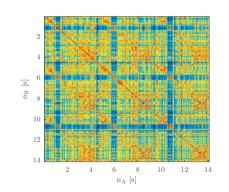
- dynamic programming technique
- time is warped non-linearly to match sequences
- finds optimal match between two sequences given a cost function
- the overall cost indicates the overall distance between the sequences

matlab source: matlab/displaySimMatrix.m

dynamic time warping processing steps

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- extract suitable features ⇒ two series of feature vectors
- compute distance matrix $\boldsymbol{D}_{\mathrm{AB}}(n_{\mathrm{A}},n_{\mathrm{B}})$
- compute alignment path $\boldsymbol{p}(n_{\rm P})$ with $n_{\rm P} \in [0; \mathcal{N}_{\rm P} - 1]$
 - minimal overall distance
- (align sequences using dynamic time stretching)

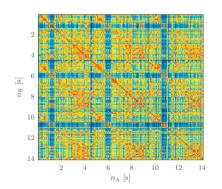


dynamic time warping distance matrix computation



given 2 sequences of vectors, compute the distance between all pairs of observations

- given 2 sequences of vectors, compute the distance between all pairs of observations
 - compute distance matrix $oldsymbol{D}_{ ext{AB}}(n_{ ext{A}}, n_{ ext{B}})$
 - example $D_{AB}(1, n_B)$ is the distance of the first vector in Seq. A to all vectors in Seq. B



dynamic time warping path properties 1/2



• boundaries: covers both A, B from beginning to end

$$m{
ho}(0) = [0,0] \ m{
ho}(\mathcal{N}_{
m P}-1) = [\mathcal{N}_{
m A}-1,\mathcal{N}_{
m B}-1]$$

causality: only forward movement

$$n_{\mathrm{A}}\big|_{\boldsymbol{p}(n_{\mathrm{P}})} \le n_{\mathrm{A}}\big|_{\boldsymbol{p}(n_{\mathrm{P}}+1)}$$
 $n_{\mathrm{B}}\big|_{\boldsymbol{p}(n_{\mathrm{P}})} \le n_{\mathrm{B}}\big|_{\boldsymbol{p}(n_{\mathrm{P}}+1)}$

continuity: no jumps

$$n_{\rm A}\big|_{p(n_{\rm P}+1)} \le (n_{\rm A}+1)\big|_{p(n_{\rm F}+1)}$$
 $n_{\rm B}\big|_{p(n_{\rm P}+1)} \le (n_{\rm B}+1)\big|_{p(n_{\rm F}+1)}$

dynamic time warping path properties 1/2



• boundaries: covers both A, B from beginning to end

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dynamic time warping path properties 1/2

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• causality: only forward movement

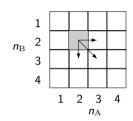
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continuity: no jumps

$$egin{aligned} n_{\mathrm{A}}ig|_{oldsymbol{
ho}(n_{\mathrm{P}}+1)} &\leq (n_{\mathrm{A}}+1)ig|_{oldsymbol{
ho}(n_{\mathrm{P}})} \ n_{\mathrm{B}}ig|_{oldsymbol{
ho}(n_{\mathrm{P}}+1)} &\leq (n_{\mathrm{B}}+1)ig|_{oldsymbol{
ho}(n_{\mathrm{P}})} \end{aligned}$$

alignment path properties 2/2



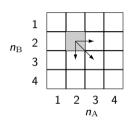


what is the minimum/maximum path length



alignment path properties 2/2





what is the minimum/maximum path length



$$\begin{split} \mathcal{N}_{\mathrm{P,min}} &= \mathsf{max}(\mathcal{N}_{\mathrm{A}}, \mathcal{N}_{\mathrm{B}}) \\ \mathcal{N}_{\mathrm{P,max}} &= \mathcal{N}_{\mathrm{A}} + \mathcal{N}_{\mathrm{B}} - 2 \end{split}$$

DTW: overall cost

alignment

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every path has an overall cost

$$\mathfrak{C}_{\mathrm{AB}}(j) = \sum_{n_{\mathrm{P}}=0}^{\mathcal{N}_{\mathrm{P}}-1} \boldsymbol{D}(\boldsymbol{p}_{j}(n_{\mathrm{P}}))$$

$$\mathfrak{C}_{\mathrm{AB},min} = \min_{egin{subarray}{c} \langle \mathfrak{C}_{\mathrm{AB}}(j) \rangle \\ j_{\mathrm{opt}} = \operatorname*{argmin}_{egin{subarray}{c} \langle \mathfrak{C}_{\mathrm{AB}}(j) \rangle \\ j_{ij} \end{array}}$$

how to determine the optimal path



DTW: overall cost

alignment

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every path has an overall cost

$$\mathfrak{C}_{\mathrm{AB}}(j) = \sum_{n_{\mathrm{P}}=0}^{\mathcal{N}_{\mathrm{P}}-1} oldsymbol{D}ig(oldsymbol{p}_{j}(n_{\mathrm{P}})ig)$$

optimal path minimizes the overall cost

$$egin{array}{lll} \mathfrak{C}_{{
m AB},min} &=& \min\limits_{orall j} \left(\mathfrak{C}_{{
m AB}}(j)
ight) \ j_{
m opt} &=& rgmin \left(\mathfrak{C}_{{
m AB}}(j)
ight) \end{array}$$

⇒ stay in the 'valleys' of distance matrix

how to determine the optimal path



alignment DTW: accumulated cost 1/2

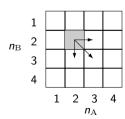
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accumulated cost: cost matrix

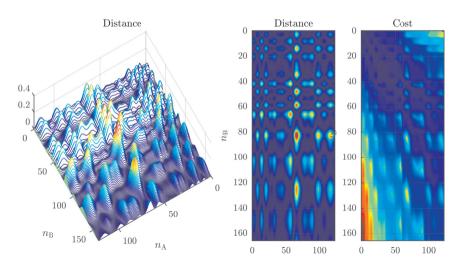
$$oldsymbol{C}_{ ext{AB}}(n_{ ext{A}},n_{ ext{B}}) = oldsymbol{D}_{ ext{AB}}(n_{ ext{A}},n_{ ext{B}}) + \min \left\{ egin{array}{l} oldsymbol{C}_{ ext{AB}}(n_{ ext{A}}-1,n_{ ext{B}}-1) \ oldsymbol{C}_{ ext{AB}}(n_{ ext{A}},n_{ ext{B}}-1) \end{array}
ight.$$

initialization

$$\boldsymbol{C}_{\mathrm{AB}}(0,0) = \boldsymbol{D}_{\mathrm{AB}}(0,0)$$





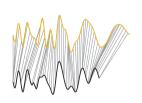


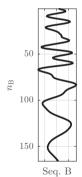
 $n_{\rm A}$

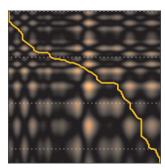
 $n_{\rm A}$

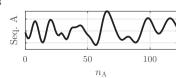
matlab source: matlab/displayDtwCost.m

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alignment DTW: algorithm description 1/2

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initialization:

$$C_{AB}(0,0) = D_{AB}(0,0), C_{AB}(n_A,-1) = \infty, C_{AB}(-1,n_B) = \infty$$

recursion

$$C_{AB}(n_{A}, n_{B}) = D_{AB}(n_{A}, n_{B}) + \min \begin{cases} C_{AB}(n_{A} - 1, n_{B} - 1) \\ C_{AB}(n_{A} - 1, n_{B}) \\ C_{AB}(n_{A}, n_{B} - 1) \end{cases}$$

$$j = \operatorname{argmin} \begin{cases} C_{AB}(n_{A} - 1, n_{B} - 1) \\ C_{AB}(n_{A} - 1, n_{B}) \\ C_{AB}(n_{A} - 1, n_{B}) \\ C_{AB}(n_{A}, n_{B} - 1) \end{cases}$$

$$\Delta p(n_{A}, n_{B}) = \begin{cases} [-1, -1] & \text{if } j = 0 \\ [-1, 0] & \text{if } j = 1 \\ [0, -1] & \text{if } j = 2 \end{cases}$$

alignment DTW: algorithm description 1/2

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initialization:

$$C_{AB}(0,0) = D_{AB}(0,0), C_{AB}(n_A,-1) = \infty, C_{AB}(-1,n_B) = \infty$$

recursion:

$$egin{array}{lcl} m{C}_{
m AB}(n_{
m A}, n_{
m B}) &=& m{D}_{
m AB}(n_{
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m AB}(n_{
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ight. \ \ m{j} &=& rgmin \left\{ egin{array}{lcl} m{C}_{
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ight. \end{array}
ight.$$

alignment DTW: algorithm description 2/2

termination:

$$n_{\Lambda} = \mathcal{N}_{\Lambda} - 1 \wedge n_{P} = \mathcal{N}_{P} - 1$$

path backtracking

$$p(n_{\rm P}) = p(n_{\rm P}+1) + \Delta p(p(n_{\rm P}+1)), \ n_{\rm P} = \mathcal{N}_{\rm P} - 2, \mathcal{N}_{\rm P} - 3, \dots, 0$$

alignment DTW: algorithm description 2/2

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termination:

$$n_{\rm A} = \mathcal{N}_{\rm A} - 1 \wedge n_{\rm B} = \mathcal{N}_{\rm B} - 1$$

path backtracking:

$$p(n_{\rm P}) = p(n_{\rm P}+1) + \Delta p(p(n_{\rm P}+1)), \ n_{\rm P} = \mathcal{N}_{\rm P} - 2, \mathcal{N}_{\rm P} - 3, \dots, 0$$

dynamic time warping example



$$A = [1, 2, 3, 0],$$

 $B = [1, 0, 2, 3, 1],$



dynamic time warping example

$$A = [1, 2, 3, 0],$$

 $B = [1, 0, 2, 3, 1],$



$$m{D}_{\mathrm{AB}} = \left[egin{array}{cccc} 0 & 1 & 2 & 1 \ 1 & 2 & 3 & 0 \ 1 & 0 & 1 & 2 \ 2 & 1 & 0 & 3 \ 0 & 1 & 2 & 1 \ \end{array}
ight]$$

dynamic time warping example

$$A = [1, 2, 3, 0],$$

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$$m{D}_{\mathrm{AB}} = \left[egin{array}{cccc} 0 & 1 & 2 & 1 \ 1 & 2 & 3 & 0 \ 1 & 0 & 1 & 2 \ 2 & 1 & 0 & 3 \ 0 & 1 & 2 & 1 \end{array}
ight]$$

$$\boldsymbol{D}_{\mathrm{AB}} = \begin{bmatrix} 0 & 1 & 2 & 1 \\ 1 & 2 & 3 & 0 \\ 1 & 0 & 1 & 2 \\ 2 & 1 & 0 & 3 \\ 0 & 1 & 2 & 1 \end{bmatrix} \qquad \boldsymbol{C}_{\mathrm{AB}} = \begin{bmatrix} 0 & \leftarrow 1 & \leftarrow 3 & \leftarrow 4 \\ \uparrow 1 & \nwarrow 2 & \nwarrow 4 & \nwarrow 3 \\ \uparrow 2 & \nwarrow 1 & \leftarrow 2 & \leftarrow 4 \\ \uparrow 4 & \uparrow 2 & \nwarrow 1 & \leftarrow 4 \\ \uparrow 4 & \uparrow 3 & \uparrow 3 & \nwarrow 2 \end{bmatrix}$$

dynamic time warping example

$$A = [1, 2, 3, 0],$$

 $B = [1, 0, 2, 3, 1],$



$$\mathbf{D}_{AB} = \begin{bmatrix} 0 & 1 & 2 & 1 \\ 1 & 2 & 3 & 0 \\ 1 & 0 & 1 & 2 \\ 2 & 1 & 0 & 3 \\ 0 & 1 & 2 & 1 \end{bmatrix}$$

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dynamic time warping variants

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transition weights: favor specific path directions

$$oldsymbol{C}_{\mathrm{AB}}(n_{\mathrm{A}},n_{\mathrm{B}}) = \min \left\{ egin{array}{lll} oldsymbol{C}_{\mathrm{AB}}(n_{\mathrm{A}}-1,n_{\mathrm{B}}-1) & + & \lambda_{\mathrm{d}} \cdot oldsymbol{D}_{\mathrm{AB}}(n_{\mathrm{A}},n_{\mathrm{B}}) \\ oldsymbol{C}_{\mathrm{AB}}(n_{\mathrm{A}}-1,n_{\mathrm{B}}) & + & \lambda_{\mathrm{v}} \cdot oldsymbol{D}_{\mathrm{AB}}(n_{\mathrm{A}},n_{\mathrm{B}}) \\ oldsymbol{C}_{\mathrm{AB}}(n_{\mathrm{A}},n_{\mathrm{B}}-1) & + & \lambda_{\mathrm{h}} \cdot oldsymbol{D}_{\mathrm{AB}}(n_{\mathrm{A}},n_{\mathrm{B}}) \end{array}
ight.$$

step types





dynamic time warping variants

transition weights: favor specific path directions

$$oldsymbol{C}_{\mathrm{AB}}(n_{\mathrm{A}},n_{\mathrm{B}}) = \min \left\{ egin{array}{lll} oldsymbol{C}_{\mathrm{AB}}(n_{\mathrm{A}}-1,n_{\mathrm{B}}-1) & + & \lambda_{\mathrm{d}} \cdot oldsymbol{D}_{\mathrm{AB}}(n_{\mathrm{A}},n_{\mathrm{B}}) \\ oldsymbol{C}_{\mathrm{AB}}(n_{\mathrm{A}}-1,n_{\mathrm{B}}) & + & \lambda_{\mathrm{v}} \cdot oldsymbol{D}_{\mathrm{AB}}(n_{\mathrm{A}},n_{\mathrm{B}}) \\ oldsymbol{C}_{\mathrm{AB}}(n_{\mathrm{A}},n_{\mathrm{B}}-1) & + & \lambda_{\mathrm{h}} \cdot oldsymbol{D}_{\mathrm{AB}}(n_{\mathrm{A}},n_{\mathrm{B}}) \end{array}
ight.$$

step types





dynamic time warping optimization



- ullet challenge: distance matrix dimensions $\mathcal{N}_A\cdot\mathcal{N}_B$
- ⇒ DTW *inefficient* for long sequences
 - high memory requirements
 - large number of operations

optimizations

- maximum time and tempo deviation
- sliding window
- multi-scale DTW (several downsampled iterations)

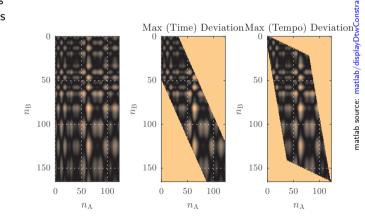
dynamic time warping optimization

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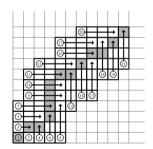
variants

- optimization
 - challenge: distance matrix dimensions $\mathcal{N}_A \cdot \mathcal{N}_B$
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 - large number of operations

optimizations:

- maximum time and tempo deviation
- sliding window
- multi-scale DTW (severa downsampled iterations)

Information Retrieval (ISMIR), London, Sep. 2005.



example

¹S. Dixon and G. Widmer, "MATCH: A Music Alignment Tool Chest," in *Proceedings of the 6th International Conference on Music*

variants

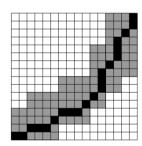
optimization

overview

- ullet challenge: distance matrix dimensions $\mathcal{N}_{A}\cdot\mathcal{N}_{B}$
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optimizations:

- maximum time and tempo deviation
- sliding window
- multi-scale DTW (several downsampled iterations)



example

¹M. Mller, H. Mattes, and F. Kurth, "An Efficient Multiscale Approach to Audio Synchronization," in *Proceedings of the International Society*

summary

lecture content



- dynamic time warping
 - find globally optimal alignment path between two sequences
- processing steps
 - compute distance matrix
 - compute cost matrix
 - back-track path

