

# Voronoi diagrams in information geometry

— Statistical Voronoi diagrams and their applications —

Frank NIELSEN

École Polytechnique  
Sony CSL  
e-mail: [Frank.Nielsen@acm.org](mailto:Frank.Nielsen@acm.org)

MaxEnt 2014

# Outline of the tutorial

1. **Euclidean Voronoi diagrams**
2. Discriminant analysis, **Mahalanobis Voronoi diagrams**, and **anisotropic Voronoi diagrams**
3. **Fisher-Hotelling-Rao Voronoi diagrams** (Riemannian curved geometries)
4. **Kullback-Leibler Voronoi diagrams** (Bregman Voronoi diagrams, dually flat geometries)
5. Bayes' error, Chernoff information and **statistical Voronoi diagrams**

## Euclidean (ordinary) Voronoi diagrams

$\mathcal{P} = \{P_1, \dots, P_n\} : n \text{ distinct point generators in Euclidean space } \mathbb{E}^d$



$$V(P_i) = \{X : D_E(P_i, X) \leq D_E(P_j, X), \forall j \neq i\} = \cap_{i=1}^n \text{Bi}^+(P_i, P_j)$$

$$D_E(P, Q) = \|\theta(P) - \theta(Q)\|_2 = \sqrt{\sum_{i=1}^d (\theta_i(P) - \theta_i(Q))^2}$$

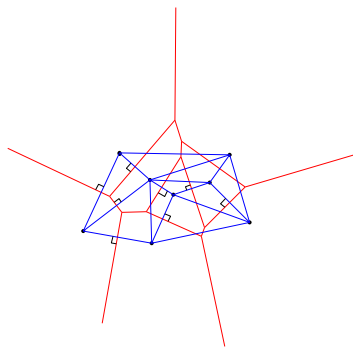
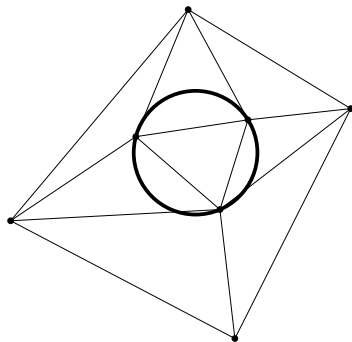
$\theta(P) = p$  : Cartesian coordinate system with  $\theta_j(P_i) = p_i^{(j)}$ .

**Bisectors**  $\text{Bi}(P, Q) = \{X : D_E(P, X) \leq D_E(Q, X)\}$  : **hyperplanes**

**Voronoi diagram** = cell complex  $V(P_i)$ 's with their faces

$\Rightarrow$  Many applications : crystal growth, codebook/quantization, molecule interfaces/docking, motion planning, etc.

# Voronoi diagrams and dual Delaunay simplicial complex

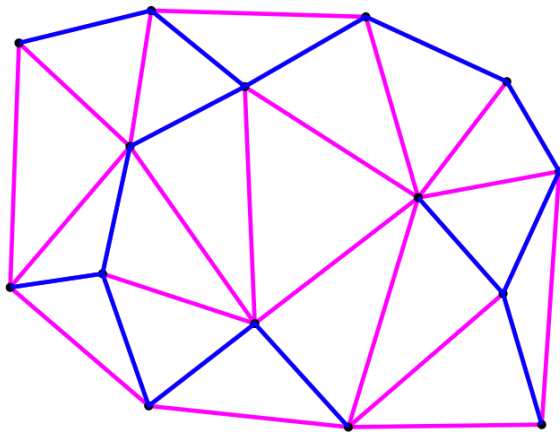


- ⇒ **Empty sphere** property, **max min angle** triangulation, etc
- ⇒ General position : no  $(d + 2)$  points cospherical
- ⇒ Voronoi & dual **Delaunay triangulation**
- ⇒ Bisector  $\text{Bi}(P, Q)$  **perpendicular**  $\perp$  to segment  $[PQ]$

## Minimum spanning tree $\subset$ Delaunay triangulation

All edges of the Euclidean MST are Delaunay edges

→ Prim's greedy algorithm in Delaunay graph in  $O(n \log n)$

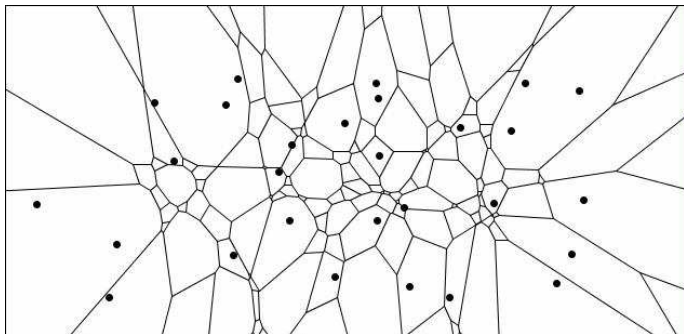


**Convex hull** : boundary  $\partial$  of the Delaunay triangulation

## Order $k$ Voronoi diagrams

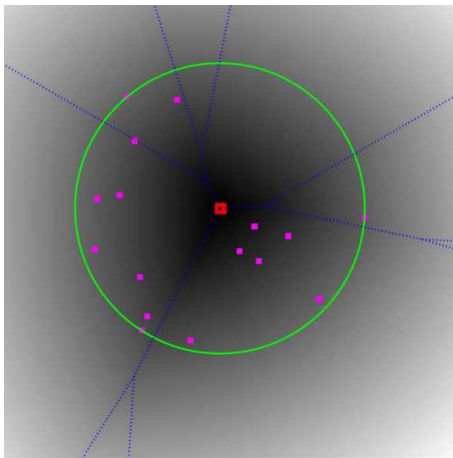
All subsets of size  $k$ ,  $\mathcal{P}_k = \binom{\mathcal{P}}{k} = \{\mathcal{K}_1, \dots, \mathcal{K}_N\}$  with  $N = \binom{n}{k}$ .  
partition the space into **non-empty  $k$ -order Voronoi cells** :

$$\text{Vor}_k(\mathcal{K}_i) = \{x : \forall q \in \mathcal{K}_i, \forall r \in \mathcal{P} \setminus \mathcal{K}_i, D(x, q) \leq D(x, r)\}$$



Combinatorial complexity not yet settled!!! ( $k$ -sets/levels)  
 $\equiv$  projection of  $k$ -levels of arrangements of hyperplanes in  $\mathbb{R}^{d+1}$ .

# Farthest order Voronoi and Minimum Enclosing Ball



Circumcenter/minmax center on the farthest Voronoi diagram.  
*Non-differentiable* optim. at the farthest Voronoi [16]

# Voronoi & Delaunay : Complexity and algorithms

- ▶ Complexity :  $\Theta(n^{\lceil \frac{d}{2} \rceil})$  (quadratic in 3D)  
Moment curve :  $t \mapsto (t, t^2, \dots, t^d)$ , etc.
- ▶ Construction :  $\Theta(n \log n + n^{\lceil \frac{d}{2} \rceil})$ , output-sensitive algorithms  
 $\Omega(n \log n + f)$ , **not yet optimal output-sensitive algorithms.**
- ▶ Voronoi diagram :  $\boxed{\text{Vor}_D(\mathcal{P}) = \text{Vor}_{f(D(\cdot, \cdot))}(\mathcal{P})}$  for a strictly monotonically increasing function  $f(\cdot)$   
Vor. Euclidean  $\equiv$  Vor. squared Euclidean  $\subset$  Vor. Bregman
- ▶ Geometric toolbox : **space of spheres** (polarity, orthogonality between spheres, etc.)



# Multiple Hypothesis Testing [9, 10]

Given a random variable  $X$  with  $n$  hypothesis

$$H_1 : X \sim P_1$$

$$: \dots$$

$$H_n : X \sim P_n$$

decide for a IID sample  $x_1, \dots, x_m \sim X$  which hypothesis holds true?

$$P_{\text{correct}}^m = 1 - P_{\text{error}}^m$$

Asymptotic regime : **Error exponent**  $\alpha$

$$\lim_{m \rightarrow \infty} -\frac{1}{m} \log P_e^m = \alpha$$

# Bayesian hypothesis testing

- ▶ **Prior probabilities** :  $w_i = \mathbb{P}(X \sim P_i) > 0$  (with  $\sum_{i=1}^n w_i = 1$ )
- ▶ **Conditional probabilities** :  $\mathbb{P}(X = x | X \sim P_i)$ .

$$\mathbb{P}(X = x) = \sum_{i=1}^n \mathbb{P}(X \sim P_i) \mathbb{P}(X = x | X \sim P_i) = \sum_{i=1}^n w_i \mathbb{P}(X | P_i)$$

- ▶ **Cost design matrix**  $[c_{ij}]$  = with  $c_{i,j}$  = **cost** of deciding  $H_i$  when in fact  $H_j$  is true
- ▶  $p_{i,j}(u(x))$  = probability of making this decision using **rule**  $u$ .

## Bayesian detector : MAP rule

- ▶ Minimize the *expected cost* :

$$E_X[c(r(x))], \quad c(r(x)) = \sum_i \left( w_i \sum_{j \neq i} c_{i,j} p_{i,j}(r(x)) \right)$$

- ▶ Special case : **Probability of error**  $P_e$  :  $c_{i,i} = 0$  and  $c_{i,j} = 1$  for  $i \neq j$  :

$$P_e = E_X \left[ \sum_i \left( w_i \sum_{j \neq i} p_{i,j}(r(x)) \right) \right]$$

- ▶ **Maximum a posteriori probability** (MAP) rule :

$$\text{map}(x) = \operatorname{argmax}_{i \in \{1, \dots, n\}} w_i p_i(x)$$

where  $p_i(x) = \mathbb{P}(X = x | X \sim P_i)$  are the conditional probabilities.

→ **MAP Bayesian detector minimizes  $P_e$  over all rules [7]**

# MVNs, MAP rule, and Mahalanobis distance (1930)

## MultiVariate Normals

(MVNs)

- ▶ Say,  $X_1 \sim N(\mu_1, \Sigma)$  and  $X_2 \sim N(\mu_2, \Sigma)$
- ▶ Probability of error (misclassification) for  $w_1 = w_2 = \frac{1}{2}$  :

$$P_e = \frac{1}{2}(1 - \text{TV}(P_1, P_2)) = \Phi\left(-\frac{1}{2}\Delta(P_1, P_2)\right)$$

$\Phi(\cdot)$  : standard normal distribution function

- ▶ Mahalanobis metric distance  $\Delta(\cdot, \cdot)$  :

$$\Delta^2(X_1, X_2) = (\mu_1 - \mu_2)^\top \Sigma^{-1}(\mu_1 - \mu_2) = \Delta\mu^\top \Sigma^{-1} \Delta\mu$$

- ▶ Measure the density overlapping (the greater, the lesser)
- ▶ generalize Euclidean distance ( $\Sigma = I$ )
- ▶  $\Delta^2$  : only symmetric Bregman divergence.

$\Rightarrow$  easy to approximate  $\tilde{P}_e$  experimentally

# Mahalanobis Voronoi diagrams

MHT for isotropic MVNs with uniform weight  $w_i = \frac{1}{n}$

Cholesky decomposition  $\Sigma = LL^\top$

$$\begin{aligned}\Delta^2(X_1, X_2) &= (\mu_1 - \mu_2)^\top \Sigma^{-1} (\mu_1 - \mu_2), \\ &= (\mu_1 - \mu_2)^\top (L^{-1})^\top L^{-1} (\mu_1 - \mu_1), \\ &= (L^{-1}\mu_1 - L^{-1}\mu_2)^\top (L^{-1}\mu_1 - L^{-1}\mu_2)\end{aligned}$$

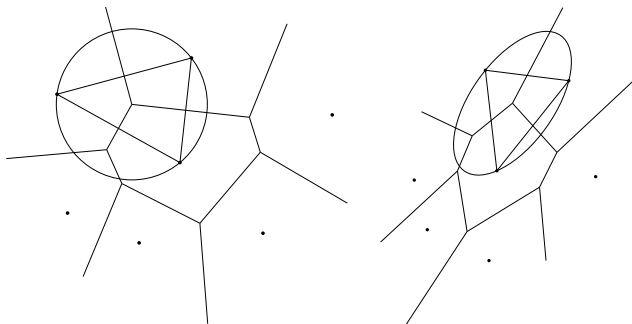
$$\Delta^2(X_1, X_2) = D_E(L^{-1}\mu_1, L^{-1}\mu_2)$$

Mahalanobis Voronoi diagram :

- ▶ Cholesky decomposition :  $\Sigma = LL^\top$
- ▶ Map  $\mathcal{P} = \{p_1, \dots, p_n\}$  to  $\mathcal{P}' = \{p'_1, \dots, p'_n\}$  with  $p'_i = L^{-1}x_i$ .
- ▶ Ordinary Voronoi diagram  $\text{Vor}_E(\mathcal{P}')$
- ▶ Map  $\text{Vor}_E(\mathcal{P}')$  to  $\text{Vor}_\Sigma(\mathcal{P}) = L\text{Vor}_E(\mathcal{P}')$ .

# Mahalanobis Voronoi diagrams

$\Sigma$  account for both **correlation** and dimension (feature) **scaling**



Dual structure  $\equiv$  anisotropic Delaunay triangulation

$\Rightarrow$  "empty circumellipse" property

Recall : MAP rule from Voronoi diagram

# Anisotropic Voronoi diagrams [6] (MHT for MVNs)

Classification : Generator  $X_i \sim N(\mu_i, \Sigma_i)$ ,  $w_i = \frac{1}{n}$ .



Discriminant functions are **quadratic bisectors**

Orphans, islands, dual not a triangulation

(need wedged/visibility conditions)

Applications : Crystal growth in fields (not Riemannian smooth metric)

# Statistics : Estimators

Given  $x_1, \dots, x_n$  IID observations (from a population), estimate the underlying distribution  $X \sim p$ .

- ▶ **empirical CDF** :  $\hat{F}_n(x) = \frac{1}{n} \sum_i I_{(-\infty, x)}(x_i)$ .  
Glivenko-Cantelli theorem :  $\sup_x |\hat{F}_n(x) - F(x)| \rightarrow 0$  a.s.
- ▶ Fisher approach : Density  $p$  belongs to a **parametric family**  $p(x|\theta)$ ,  $D = \# \text{parameters (order)}$ 
  - ▶ **Method of moments** : Match distribution moments with sample moments :

$$\mathbb{E}_X[X^l] = \frac{1}{n} \sum_{i=1}^n x_i^l,$$

$\Rightarrow$  any  $D$  independent equations yields an estimator

- ▶ Fisher Maximum (log-)Likelihood Estimator **MLE** :

$$\max_{\theta} l(\theta; x_1, \dots, x_n) = \max_{\theta} \prod_{i=1}^n p(x_i; \theta)$$



## Estimation : Variance Lower Bound

Which estimator  $\hat{\theta}_n$  shall we choose ?

**Estimator is also a random variable on a random vector**

- ▶ Consistent :  $\lim_{n \rightarrow \infty} \hat{\theta}_n \rightarrow \theta$
- ▶ Unbiased :  $\mathbb{E}_{\theta}[\hat{\theta}_n] - \theta = 0$
- ▶ Minimum square error (MSE) :  $\mathbb{E}[(\hat{\theta} - \theta)^2] = \mathbb{B}[\hat{\theta}]^2 + \mathbb{V}[\hat{\theta}]$
- ▶ **Fisher information matrix :**

$$I(\theta) = \left[ I_{i,j}(\theta) = \mathbb{E}_{\theta} \left[ \frac{\partial}{\partial \theta_i} \log p(x|\theta) \frac{\partial}{\partial \theta_j} \log p(x|\theta) \right] \right]$$

- ▶ **Fréchet Darmois Cramér-Rao lower bound** for unbiased estimators :

$$\mathbb{V}[\hat{\theta}_n] \succeq \frac{1}{n} I^{-1}(\theta)$$

- ▶ **Efficient** : estimator reaches the Cramér-Rao lower bound

# Sufficient statistics and exponential families

- **Sufficient statistic**  $t(x)$  :

$$\mathbb{P}(x|t(x) = t, \theta) = \mathbb{P}(x|t(x) = t)$$

All information for inference contained in  $t(x)$ . ( $\neq$  ancillary)

- For univariate Gaussians,  $D = 2$  :  $t_1(x) = x$  and  $t_2(x) = x^2$ .

$$\hat{\mu} = \bar{x} = \frac{1}{n} \sum_i x_i, \quad \hat{\sigma}^2 = \frac{1}{n} \sum_i (x_i - \bar{x})^2 = \left( \frac{1}{n} \sum_i x_i^2 \right) - (\bar{x})^2$$

- **Exponential families** have finite dimensional sufficient statistics (Koopman) : Reduce  $n$  data to  $D$  statistics.

$$\forall x \in \mathcal{X}, \mathbb{P}(x|\theta) = \exp(\theta^\top t(x) - F(\theta) + k(x))$$

$F(\cdot)$  : log-normalizer/cumulant/partition function

$k(x)$  : auxiliary term for carrier measure

# Exponential families : Fisher information and MLE

Common distributions are exponential families :  
Poisson, Gaussians, Gamma, Beta, Dirichlet, etc.

- ▶ Fisher information :

$$I(\theta) = \nabla^2 F(\theta)$$

- ▶ MLE for an exponential family :

$$\nabla F(\theta) = \frac{1}{n} \sum_i t(X_i) = \eta$$

MLE exists iff.

$$\boxed{\eta = \nabla F(\theta) \in \text{int}(\text{CH}(\mathcal{X}))}$$

# Population space : Hotelling (1930) [5] & Rao (1945) [25]

## Birth of **differential-geometric methods in statistics**.

- ▶ Fisher information matrix (positive definite) can be used as a (smooth) Riemannian metric tensor.
- ▶ Distance between two populations indexed by  $\theta_1$  and  $\theta_2$  : Riemannian distance (metric length)
- ▶ Fisher-Hotelling-Rao (FHR) geodesic distance used in classification : Find the closest population to a given population
- ▶ Used in tests of significance (null versus alternative hypothesis), power of a test :  $\mathbb{P}(\text{reject } H_0 | H_0 \text{ is false})$  (surfaces in population spaces)

## Rao's distance (1945, introduced by Hotelling 1930 [5])

- ▶ Infinitesimal length element :

$$ds^2 = \sum_{ij} g_{ij}(\theta) d\theta_i d\theta_j = d\theta^T I(\theta) d\theta$$

- ▶ Geodesic and distance are hard to explicitly calculate :

$$\rho(p(x; \theta_1), p(x; \theta_2)) = \min_{\substack{\theta(s) \\ \theta(0)=\theta_1 \\ \theta(1)=\theta_2}} \int_0^1 \sqrt{\left(\frac{d\theta}{ds}\right)^T I(\theta) \frac{d\theta}{ds}} ds$$

- ▶ Metric property of  $\rho$ , many tools [1] : Riemannian Log/Exp tangent/manifold mapping

# Fisher Rao Hotelling Voronoi : Riemannian Voronoi diagrams

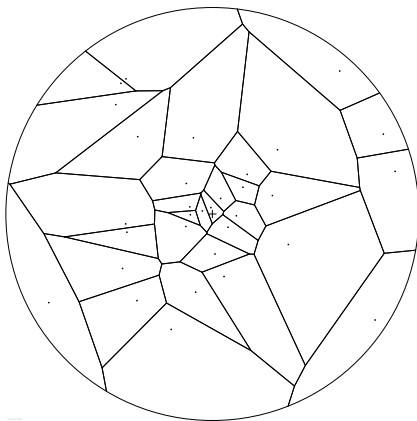
- ▶ Location-scale  $2D$  families have *non-positive curvature* (Hotelling, 1930) : **FHR Voronoi diagrams** amount to **hyperbolic Voronoi diagrams** or Euclidean diagrams (location families only like isotropic Gaussians)
- ▶ Multinomial family has **spherical** geometry on the positive orthant : **Spherical Voronoi diagram** (stereographic projection  $\propto$  Euclidean Voronoi diagrams)
- ▶ Arbitrary families  $p(x|\theta)$  : Geodesics not in closed forms  $\rightarrow$  limited computational framework in practice...

## Hyperbolic Voronoi diagrams [19, 22]

- ▶ In Klein disk, the hyperbolic Voronoi diagram amounts to a **clipped affine Voronoi diagram**, or a **clipped power diagram**. Efficient clipping algorithm [2].
- ▶ Convert to other models of hyperbolic geometry : Poincaré disk, upper half space, hyperboloid, **Beltrami** hemisphere, etc.
- ▶ **Conformal** (good for vizualizing) versus **non-conformal** (good for computing) models. (conformal metric  $G(x) = \lambda(x)I$  is scaled identity metric)

# Hyperbolic Voronoi diagrams [19, 22]

Hyperbolic Voronoi diagram in Klein disk :

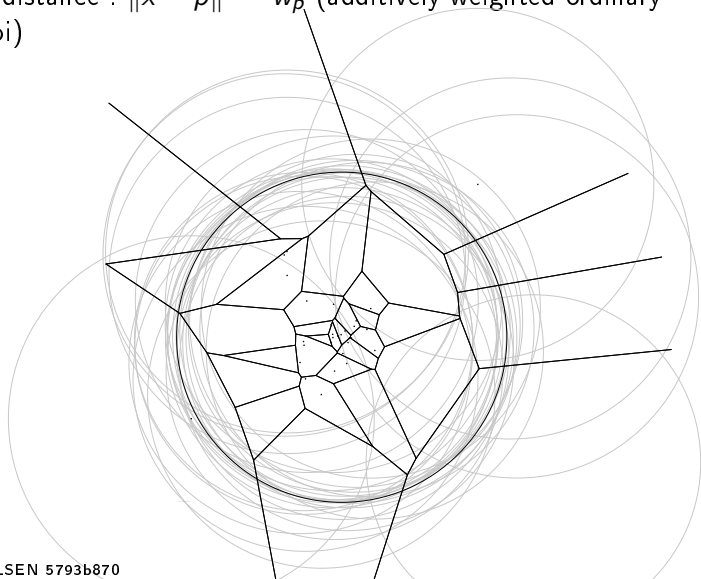




# Hyperbolic Voronoi diagrams [19, 22]

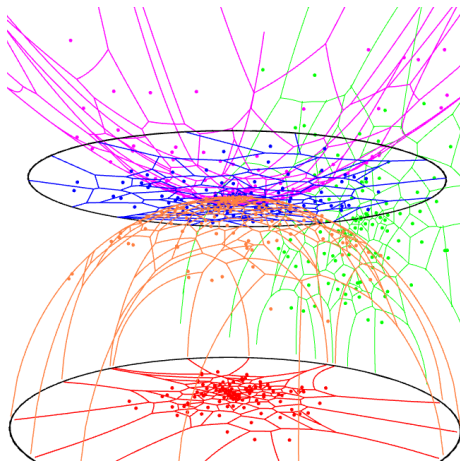
Affine Voronoi diagram equivalent to power diagram.

Power distance :  $\|x - p\|^2 - w_p$  (additively weighted ordinary Voronoi)



# Hyperbolic Voronoi diagrams [19, 22]

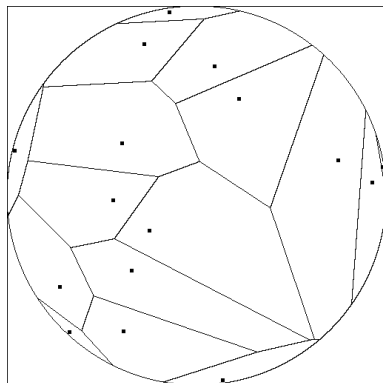
5 common models of the abstract hyperbolic geometry



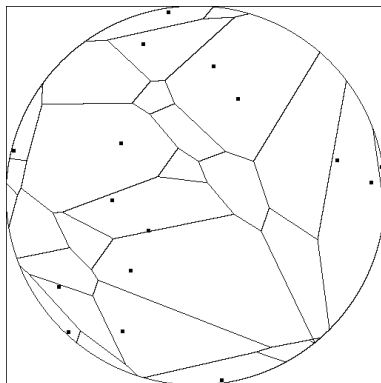
<https://www.youtube.com/watch?v=i9IUzNxeH4o>

ACM Symposium on Computational Geometry (SoCG'14)

## $k$ -order hyperbolic Voronoi diagram



$k = 1$



$k = 2$

Affine diagram  $\Rightarrow$  equivalent to a power diagram [21]

Watch  $k$ -order Voronoi :

<https://www.youtube.com/watch?v=i9IUzNxeH4o>

## Constrast function and dually flat space

- ▶ Convex and strictly differentiable function  $F(\theta)$  admits a Legendre-Fenchel convex conjugate  $F^*(\eta)$  :

$$F^*(\eta) = \sup_{\theta} (\theta^\top \eta - F(\theta)), \quad \nabla F(\theta) = \eta = (\nabla F^*)^{-1}(\eta)$$

- ▶ Young's inequality gives rise to **canonical divergence** [8] :

$$F(\theta) + F^*(\eta') \geq \theta^\top \eta' \Rightarrow A_{F,F^*}(\theta, \eta') = F(\theta) + F^*(\eta') - \theta^\top \eta'$$

- ▶ Writing using single coordinate system, get dual **Bregman divergences** :

$$\begin{aligned} B_F(\theta_p : \theta_q) &= F(\theta_p) - F(\theta_q) - (\theta_p - \theta_q)^\top \nabla F(\theta_q) \\ &= B_{F^*}(\eta_q : \eta_p) \\ &= A_{F,F^*}(\theta_p, \eta_q) = A_{F^*,F}(\eta_q : \theta_p) \end{aligned}$$

# Discriminant analysis exponential families

- ▶ Consider  $X_1 \sim E_F(\theta_1)$ ,  $X_2 \sim E_F(\theta_2)$  (or moment parameterization  $\eta_i = \nabla F(\theta_i)$ )
- ▶ Bayes' rule : Classify  $x$  from  $X_1$  iff.  $\mathbb{P}(x|\theta_1) > \mathbb{P}(x|\theta_2)$
- ▶ Use bijection between exponential families and Bregman divergences :

$$\log p(x|\theta) = -B_{F^*}(t(x) : \eta) + F^*(t(x)) + k(x)$$

- ▶ MAP rule ( $w_1 = w_2 = \frac{1}{2}$ ) :

$$B_{F^*}(t(x) : \eta_1) < B_{F^*}(t(x) : \eta_2)$$

- ▶ Bregman bisector :

$$\text{Bi}_{F^*}(\eta_1, \eta_2) = \{\eta \mid B_{F^*}(\eta : \eta_1) = B_{F^*}(\eta : \eta_2)\}$$

## Bregman dual bisectors [3, 17, 20]

Right-sided bisector :  $\rightarrow$  Hyperplane ( $\theta$ -hyperplane)

$$H_F(p, q) = \{x \in \mathcal{X} \mid B_F(x : p) = B_F(x : q)\}.$$

$H_F :$

$$\langle \nabla F(p) - \nabla F(q), x \rangle + (F(p) - F(q) + \langle q, \nabla F(q) \rangle - \langle p, \nabla F(p) \rangle) = 0$$

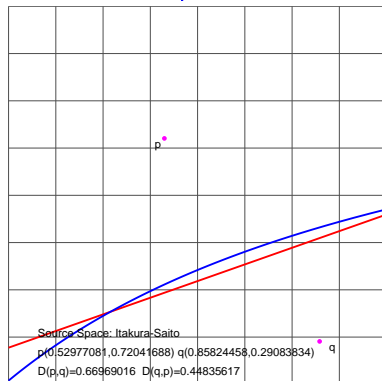
Left-sided bisector :  $\rightarrow$   $\theta$ -Hypersurface ( $\eta$ -hyperplane)

$$H'_F(p, q) = \{x \in \mathcal{X} \mid B_F(p : x) = B_F(q : x)\}$$

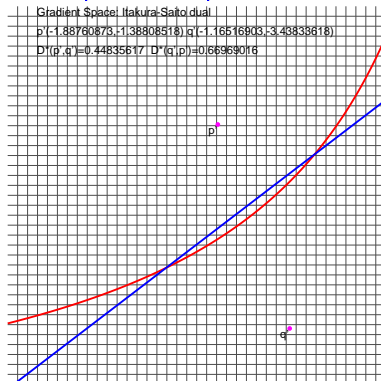
$$H'_F : \langle \nabla F(x), q - p \rangle + F(p) - F(q) = 0$$

# Visualizing Bregman bisectors

Primal coordinates  $\theta$   
natural parameters



Dual coordinates  $\eta$   
expectation parameters



# Space of Bregman spheres and Bregman balls [3]

Dual Bregman balls (bounding Bregman spheres) :

$$\begin{aligned}\text{Ball}_F^r(c, r) &= \{x \in \mathcal{X} \mid B_F(x : c) \leq r\} \\ \text{and } \text{Ball}_F^l(c, r) &= \{x \in \mathcal{X} \mid B_F(c : x) \leq r\}\end{aligned}$$

Legendre duality :

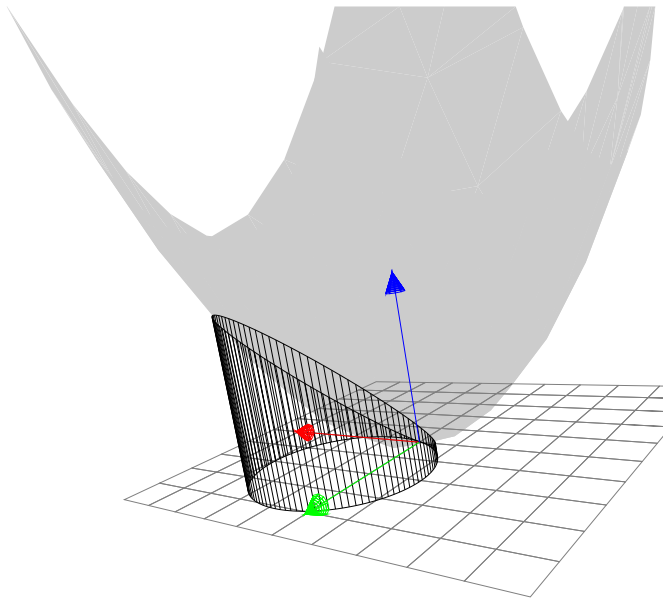
$$\text{Ball}_F^l(c, r) = (\nabla F)^{-1}(\text{Ball}_{F^*}^r(\nabla F(c), r))$$



Illustration for Itakura-Saito divergence,  $F(x) = -\log x$



## Lifting/Polarity : Potential function graph $\mathcal{F}$



## Space of Bregman spheres : Lifting map [3]

$\mathcal{F} : x \mapsto \hat{x} = (x, F(x))$ , hypersurface in  $\mathbb{R}^{d+1}$ .

$H_p$  : Tangent hyperplane at  $\hat{p}$ ,  $z = H_p(x) = \langle x - p, \nabla F(p) \rangle + F(p)$

- ▶ Bregman sphere  $\sigma \longrightarrow \hat{\sigma}$  with supporting hyperplane

$$H_\sigma : z = \langle x - c, \nabla F(c) \rangle + F(c) + r.$$

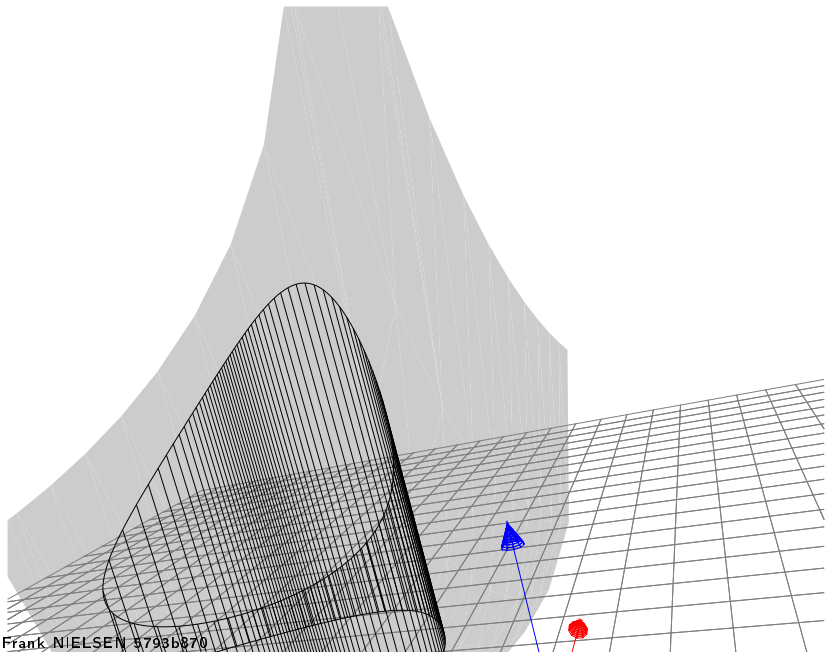
(// to  $H_c$  and shifted vertically by  $r$ )

$$\hat{\sigma} = \mathcal{F} \cap H_\sigma.$$

- ▶ intersection of any hyperplane  $H$  with  $\mathcal{F}$  projects onto  $\mathcal{X}$  as a Bregman sphere :

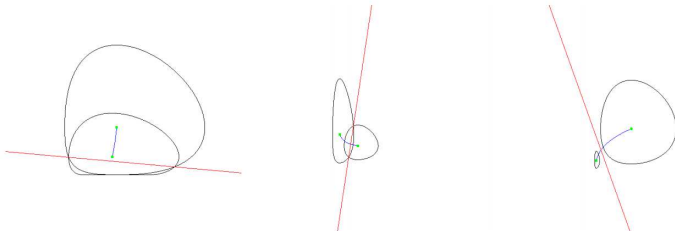
$$H : z = \langle x, a \rangle + b \rightarrow \sigma : \text{Ball}_F(c = (\nabla F)^{-1}(a), r = \langle a, c \rangle - F(c) + b)$$

# Lifting with Itakura-Saito potential function



## Space of Bregman spheres [3]

- ▶ Union/intersection of Bregman spheres from representational polytope [3]
- ▶ Radical axis of two Bregman balls is an hyperplane :  
Applications to Nearest Neighbor search trees like Bregman ball trees or Bregman vantage point trees [23].



## Space of spheres : **Minimum enclosing ball** [14, 24]

To a hyperplane  $H_\sigma = H(a, b) : z = \langle a, x \rangle + b$  in  $\mathbb{R}^{d+1}$ , corresponds a ball  $\sigma = \text{Ball}(c, r)$  in  $\mathbb{R}^d$  with center  $c = \nabla F^*(a)$  and radius :

$$r = \langle a, c \rangle - F(c) + b = \langle a, \nabla F^*(a) \rangle - F(\nabla F^*(a)) + b = F^*(a) + b$$

since  $F(\nabla F^*(a)) = \langle \nabla F^*(a), a \rangle - F^*(a)$  (Young equality)

SEB : Find halfspace  $H(a, b)^- : z \leq \langle a, x \rangle + b$  that contains all lifted points :

$$\min_{a,b} r = F^*(a) + b,$$

$$\forall i \in \{1, \dots, n\}, \quad \langle a, x_i \rangle + b - F(x_i) \geq 0.$$

$\Rightarrow$  **Convex Program (CP) with linear inequality constraints**

$\Rightarrow F(\theta) = F^*(\eta) = \frac{1}{2}x^\top x : \text{CP} \rightarrow \text{Quadratic Programming (QP)} [4].$

## InSphere predicates wrt Bregman divergences [3]

Implicit representation of Bregman spheres/balls :

- Is  $x$  inside the Bregman ball defined by  $d + 1$  support points?

$$\text{InSphere}(x; p_0, \dots, p_d) = \begin{vmatrix} 1 & \dots & 1 & 1 \\ p_0 & \dots & p_d & x \\ F(p_0) & \dots & F(p_d) & F(x) \end{vmatrix}$$

- $\text{InSphere}(x; p_0, \dots, p_d)$  is negative, null or positive depending on whether  $x$  lies inside, on, or outside  $\sigma$ .

# Bregman Voronoi diagrams as minimization diagrams [3]

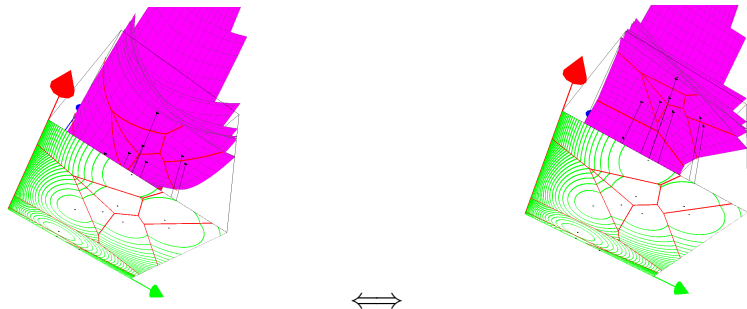
A subclass of affine diagrams which have all non-empty cells .

Minimization diagram of the  $n$  functions

$$D_i(x) = B_F(x : p_i) = F(x) - F(p_i) - \langle x - p_i, \nabla F(p_i) \rangle.$$

$\equiv$  minimization of  $n$  linear functions :

$$H_i(x) = (p_i - x)^T \nabla F(p_i) - F(p_i)$$



Watch [https://www.youtube.com/watch?v=L7v1wuN\\_9Wg](https://www.youtube.com/watch?v=L7v1wuN_9Wg)

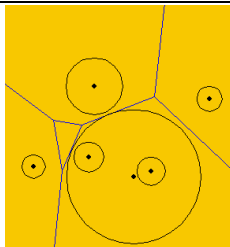
## Bregman Voronoi from Power diagrams [15, 20]

Any affine diagram can be built from a **power diagram**.  
(power diagrams defined in full space  $\mathbb{R}^d$ )

- ▶ Power distance of  $x$  to  $\text{Ball}(p, r)$  :  $\|p - x\|^2 - r^2$ .
- ▶ Laguerre diagram : minimization diagram of  $D_i(x) = \|p_i - x\|^2 - r_i^2$
- ▶ Power bisector of  $\text{Ball}(p_i, r_i)$  and  $\text{Ball}(p_j, r_j)$  = radical hyperplane :

$$2\langle x, p_j - p_i \rangle + \|p_i\|^2 - \|p_j\|^2 + r_j^2 - r_i^2 = 0.$$

**Universality :** Affine Bregman Voronoi diagram  $\Leftrightarrow$  Power diagram



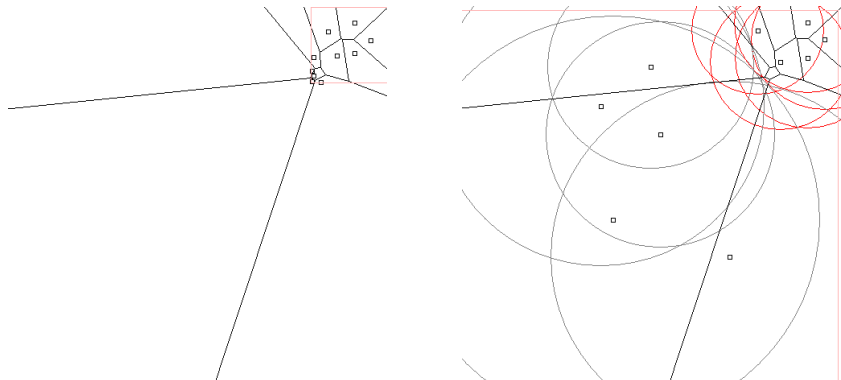


# Affine Bregman Voronoi diagrams as power diagrams

Equivalence :  $B(\nabla F(p_i), r_i)$  with

$$r_i^2 = \langle \nabla F(p_i), \nabla F(p_i) \rangle + 2(F(p_i) - \langle p_i, \nabla F(p_i) \rangle)$$

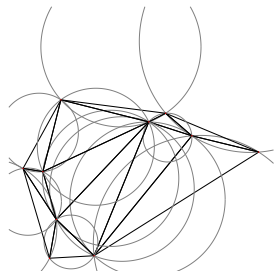
(imaginary radii shown in red, WLOG.  $r_i \geq 0$  by shifting)



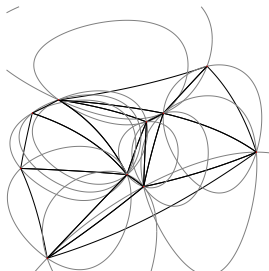
Some cells may be empty in the Laguerre diagram (power diagram) but not in the Bregman diagram

<http://www.csl.sony.co.jp/person/nielsen/BVDapplet/>

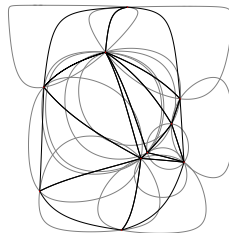
# Bregman dual Delaunay triangulations



Delaunay



Exponential



Hellinger-like

- ▶ empty Bregman sphere property,
- ▶ geodesic triangles.

BVDs extends Euclidean Voronoi diagrams with similar complexity/algorithms.

# Riemannian tensor and orthogonality

Dually flat space.

- ▶  $F(\theta)$  convex on  $\Theta$  yields convex conjugate  $F^*(\eta)$  on  $H$
- ▶ **Same** Riemannian tensor can be expressed in both coordinate systems :

$$g(\theta) = \nabla^2 F(\theta), \quad g^*(\eta) = \nabla^2 F(\eta), \quad g(\theta(X))g^*(\eta(X)) = I, \forall X$$

Same infinitesimal length element  $ds^2 = (ds^*)^2$ .

- ▶ Two curves  $\gamma_1(t)$  and  $\gamma_2(t)$  are orthogonal at  $Q = \gamma_1(t_0) = \gamma_2(t_0)$  iff.

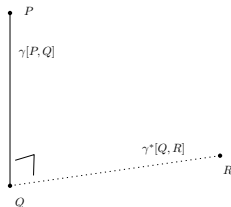
$$\langle \dot{\gamma}_1(t_0), \dot{\gamma}_2(t_0) \rangle_{g(\theta(Q))} = \langle \dot{\gamma}_1(t_0), \dot{\gamma}_2(t_0) \rangle_{g^*(\eta(Q))}^* = 0$$

- ▶ Geodesic  $\gamma(P, Q)$  and dual geodesic  $\gamma^*(P, Q)$  (and geodesic segments  $\gamma[P, Q]$  and  $\gamma^*[P, Q]$ )

# Orthogonality

3-point property (generalized law of cosines) :

$$B_F(p : r) = B_F(p : q) + B_F(q : r) - (p - q)^T (\nabla F(r) - \nabla F(q))$$



$\gamma(P, Q)$  orthogonal to  $\gamma^*(Q, R)$  iff.

$$B_F(p : r) = B_F(p : q) + B_F(q : r)$$

Equivalent to  $\langle \theta_p - \theta_q, \eta_r - \eta_q \rangle = 0$

Extend Pythagoras' theorem

$$\gamma(P, Q) \perp \gamma^*(Q, R)$$

# Dually orthogonal Bregman Voronoi & Triangulations

Ordinary Voronoi diagram is perpendicular to Delaunay triangulation. (Vor  $k$ -face  $\perp$  Del  $d - k$ -face)

Dual line segment geodesics :

$$\begin{aligned}\gamma(P, Q) &= \{\theta = \theta_p + (1 - \lambda)\theta_q \mid \lambda \in [0, 1]\} \\ \gamma^*(P, Q) &= \{\eta = \eta_p + (1 - \lambda)\eta_q \mid \lambda \in [0, 1]\}\end{aligned}$$

Bisectors :

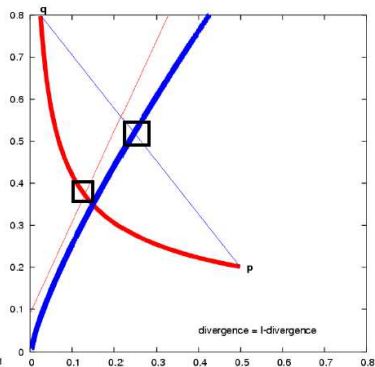
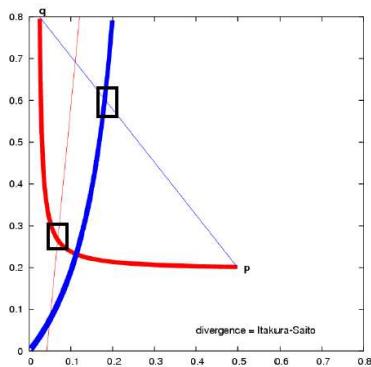
$$\begin{aligned}\text{Bi}_\theta(p, q) &: \langle x, \theta_q - \theta_p \rangle + F(\theta_p) - F(\theta_q) = 0 \\ \text{Bi}_\eta(p, q) &: \langle x, \eta_q - \eta_p \rangle + F^*(\eta_p) - F^*(\eta_q) = 0\end{aligned}$$

Dual orthogonality :

$$\begin{aligned}\text{Bi}_\eta(p, q) &\perp \gamma^*(P, Q) \\ \gamma(P, Q) &\perp \text{Bi}_\theta(p, q)\end{aligned}$$

# Dually orthogonal Bregman Voronoi & Triangulations

$$\begin{aligned} \text{Bi}_\eta(p, q) &\perp \gamma^*(P, Q) \\ \gamma(P, Q) &\perp \text{Bi}_\theta(p, q) \end{aligned}$$



# MAP rule and additive Bregman Voronoi diagrams

Optimal **MAP decision rule** :

$$\begin{aligned}\text{map}(x) &= \operatorname{argmax}_{i \in \{1, \dots, n\}} w_i p_i(x) \\ &= \operatorname{argmax}_{i \in \{1, \dots, n\}} -B^*(t(x) : \eta_i) + \log w_i, \\ &= \operatorname{argmin}_{i \in \{1, \dots, n\}} B^*(t(x) : \eta_i) - \log w_i\end{aligned}$$

$\Rightarrow$  **dual Bregman Voronoi with additive weights** (affine diagram).

Vertical projection of the intersections of the lifted half-spaces on the potential  $\mathcal{F}^*$  shifted by the weights  $-\log w_i$ .

## Relative entropy for exponential families [18]

- Kullback-Leibler divergence (cross-entropy minus entropy) :

$$\begin{aligned}\text{KL}(P : Q) &= \int p(x) \log \frac{p(x)}{q(x)} dx \geq 0 \\&= \underbrace{\int p(x) \log \frac{1}{q(x)} dx}_{H^\times(P:Q)+c} - \underbrace{\int p(x) \log \frac{1}{p(x)} dx}_{H(p)=H^\times(P:P)-c} \\&= F(\theta_Q) - F(\theta_P) - \langle \theta_Q - \theta_P, \nabla F(\theta_P) \rangle \\&= B_F(\theta_Q : \theta_P) = B_{F^*}(\eta_P : \eta_Q)\end{aligned}$$

Bregman divergence  $B_F$  defined for a strictly convex and differentiable function up to some affine terms.

- Proof  $\text{KL}(P : Q) = B_F(\theta_Q : \theta_P)$  follows from

$$X \sim E_F(\theta) \implies \boxed{E[t(X)] = \nabla F(\theta)} = \eta$$



## Upper bounds $P_e$ using Chernoff Information [11]

- ▶ Trick :  $\min(a, b) \leq a^\alpha b^{1-\alpha}, \forall \alpha \in (0, 1)$  (for  $a, b > 0$ )

$$E^* = \int \min(\mathbb{P}(C_1|x), \mathbb{P}(C_2|x))p(x)dx \leq w_1^\alpha w_2^{1-\alpha} \int p_1^\alpha(x)p_2^{1-\alpha}(x)dx$$

- ▶ Upper bound the minimum error  $E^*$

$$E^* \leq w_1^\alpha w_2^{1-\alpha} c_\alpha(p_1 : p_2),$$

$$c_\alpha(p_1 : p_2) = \int p_1^\alpha(x)p_2^{1-\alpha}(x)dx : \text{Chernoff } \alpha\text{-coefficient.}$$

$$c^*(p_1 : p_2) = c_{\alpha^*}(p_1 : p_2) = \min_{\alpha \in (0,1)} \int p_1^\alpha(x)p_2^{1-\alpha}(x)dx.$$

- ▶ *Chernoff information* (or Chernoff divergence) :

$$C^*(p_1 : p_2) = C_{\alpha^*}(p_1 : p_2) = \max_{\alpha \in (0,1)} -\log \int p_1^\alpha(x)p_2^{1-\alpha}(x)dx$$

extend methodology with **quasi-arithmetic means** [11]

## Chernoff coefficient/information : exponential families

$$\begin{aligned}C_{\alpha}(p : q) &= -\log c_{\alpha}(p, q) = J_F^{(\alpha)}(\theta_p : \theta_q), \\c_{\alpha}(p : q) &= e^{-C_{\alpha}(p:q)} = e^{-J_F^{(\alpha)}(\theta_p:\theta_q)}.\end{aligned}$$

Skewed Jensen divergence (on natural parameters) :

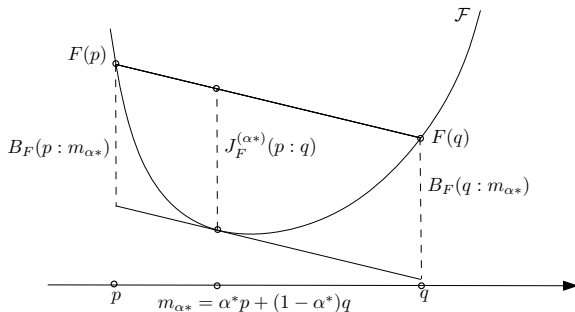
$$J_F^{(\alpha)}(\theta_p : \theta_q) = (\alpha F(\theta_p) + (1 - \alpha)F(\theta_q)) - F(\alpha\theta_p + (1 - \alpha)\theta_q)$$

# Maximizing skew Jensen divergence

$$\alpha^* = \arg \max_{0 < \alpha < 1} J_F^{(\alpha)}(p : q)$$

$$J_F^{(\alpha^*)}(p : q) = B_F(p : m_{\alpha^*}) = B_F(q : m_{\alpha^*})$$

$m_\alpha = \alpha p + (1 - \alpha)q$  :  $\alpha$ -mixing of  $p$  and  $q$ .



Maximum skew Jensen divergence amounts to Bregman divergences.

# Geometry of the best error exponent : binary hypothesis [11]

**Chernoff distribution  $P^*$  :**

$$P^* = P_{\theta_{12}^*} = G_e(P_1, P_2) \cap \text{Bi}_m(P_1, P_2)$$

**e-geodesic :**

$$G_e(P_1, P_2) = \{E_{12}^{(\lambda)} \mid \theta(E_{12}^{(\lambda)}) = (1 - \lambda)\theta_1 + \lambda\theta_2, \lambda \in [0, 1]\},$$

**m-bisector :**

$$\text{Bi}_m(P_1, P_2) : \{P \mid F(\theta_1) - F(\theta_2) + \eta(P)^\top \Delta\theta = 0\},$$

**Optimal natural parameter of  $P^*$  :**

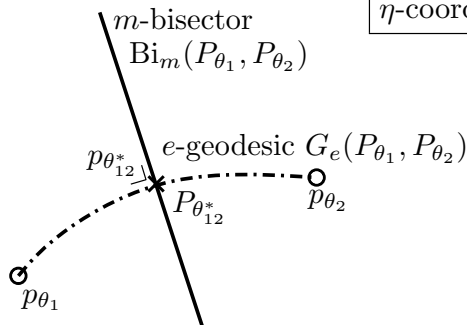
$$\theta^* = \theta_{12}^{(\alpha^*)} = \operatorname{argmin}_{\theta \in \Theta} B(\theta_1 : \theta) = \operatorname{argmin}_{\theta \in \Theta} B(\theta_2 : \theta).$$

→ **closed-form** for order-1 family, or efficient **bisection search**.

# Geometry of the best error exponent : binary hypothesis

$$P^* = P_{\theta_{12}^*} = G_e(P_1, P_2) \cap \text{Bi}_m(P_1, P_2)$$

$\eta$ -coordinate system



$$C(\theta_1 : \theta_2) = B(\theta_1 : \theta_{12}^*)$$

BHT :  $P_e$  bounded using Bregman divergence between Chernoff distribution and class-conditional distributions.

# Geometry of the best error exponent : multiple hypothesis

$n$ -ary MHT [7] from *minimum pairwise Chernoff distance* :

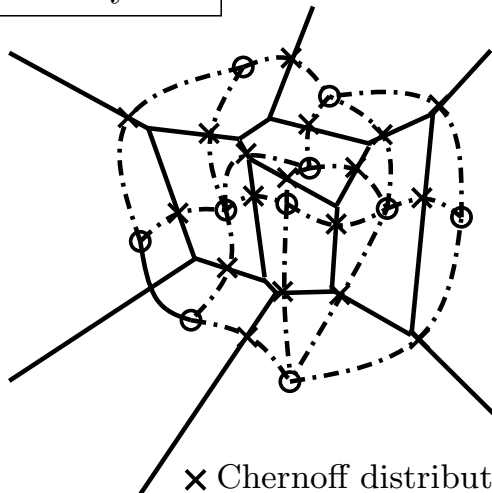
$$C(P_1, \dots, P_n) = \min_{i,j \neq i} C(P_i, P_j)$$

$$P_e^m \leq e^{-mC(P_{i^*}, P_{j^*})}, \quad (i^*, j^*) = \operatorname{argmin}_{i,j \neq i} C(P_i, P_j)$$

Compute for each pair of **natural neighbors**  $P_{\theta_i}$  and  $P_{\theta_j}$ , the Chernoff distance  $C(P_{\theta_i}, P_{\theta_j})$ , and choose the pair with minimal distance. (Proof by contradiction using **Bregman Pythagoras** theorem.)

# Multiple Hypothesis testing & Chernoff information

$\eta$ -coordinate system



x Chernoff distribution between  
natural neighbours

## Summary : Statistical Voronoi diagrams

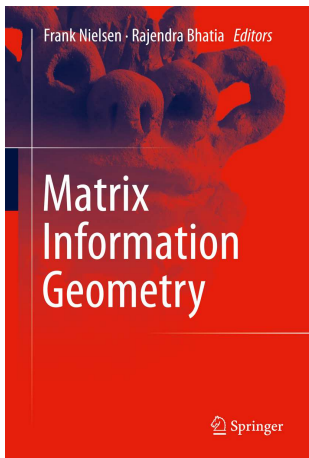
- ▶ Mahalanobis Voronoi diagrams, anisotropic Voronoi diagrams (additively-weighted)
- ▶ Fisher-Hotelling-Rao Riemannian Voronoi diagrams
- ▶ Kullback-Leibler Voronoi diagrams for exponential families = Bregman Voronoi diagrams
  - ▶ Extend ordinary Voronoi diagrams ( $\equiv$  isotropic Gaussians)
  - ▶ Space of spheres : Lifting/polarity
  - ▶ bisectors  $\perp$  geodesics
  - ▶ dual regular triangulation with empty spheres

Information geometry : affine differential geometry of “parameter spaces”, invariance principles

Many other kinds of Voronoi : Jensen-Bregman,  $(u, v)$ -structures , conformal divergences, total Bregman & total Jensen divergences, etc.



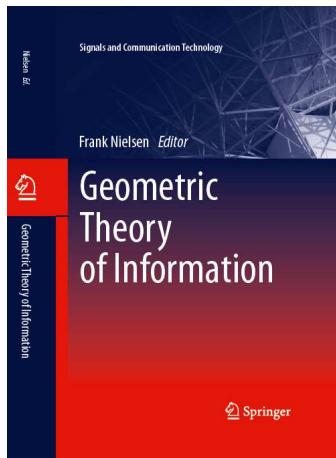
# Computational Information Geometry : Edited books



[13]

<http://www.springer.com/engineering/signals/book/978-3-642-30231-2>

<http://www.sonycs1.co.jp/person/nielsen/infogeo/MIG/MIGBOOKWEB/>



[12]

<http://www.springer.com/engineering/signals/book/978-3-319-05316-5>

<http://www.sonycs1.co.jp/person/nielsen/infogeo/GTI/GeometricTheoryOfInformation.html>



**Marc Arnaudon and Frank Nielsen.**

On approximating the Riemannian 1-center.

*Comput. Geom. Theory Appl.*, 46(1) :93–104, January 2013.



**Jean-Daniel Boissonnat and Christophe Delage.**

Convex hull and Voronoi diagram of additively weighted points.

In Gerth Ståhlting Brodal and Stefano Leonardi, editors, *ESA*, volume 3669 of *Lecture Notes in Computer Science*, pages 367–378. Springer, 2005.



**Jean-Daniel Boissonnat, Frank Nielsen, and Richard Nock.**

Bregman Voronoi diagrams.

*Discrete and Computational Geometry*, 44(2) :281–307, April 2010.



**Bernd Gärtner and Sven Schönherr.**

An efficient, exact, and generic quadratic programming solver for geometric optimization.

In *Proceedings of the sixteenth annual symposium on Computational geometry*, pages 110–118. ACM, 2000.



**Harold Hotelling.**



**Francois Labelle and Jonathan Richard Shewchuk.**

Anisotropic Voronoi diagrams and guaranteed-quality anisotropic mesh generation.

In *Proceedings of the nineteenth annual symposium on Computational geometry*, pages 191–200. ACM, 2003.



**C. C. Leang and D. H. Johnson.**

On the asymptotics of  $M$ -hypothesis Bayesian detection.

*IEEE Transactions on Information Theory*, 43(1) :280–282, January 1997.



**Frank Nielsen.**

Legendre transformation and information geometry.

Technical Report CIG-MEMO2, September 2010.



**Frank Nielsen.**

Hypothesis testing, information divergence and computational geometry.

In Frank Nielsen and Fr    ric Barbaresco, editors, *Geometric Science of Information*, volume 8085 of *Lecture Notes in Computer Science*, pages 241–248. Springer Berlin Heidelberg, 2013.



Frank Nielsen.

Generalized bhattacharyya and chernoff upper bounds on bayes error using quasi-arithmetic means.

*Pattern Recognition Letters*, 42 :25–34, 2014.



Frank Nielsen.

Generalized bhattacharyya and chernoff upper bounds on bayes error using quasi-arithmetic means.

*Pattern Recognition Letters*, 42 :25–34, 2014.



Frank Nielsen.

*Geometric Theory of Information*.

Springer, 2014.



Frank Nielsen and Rajendra Bhatia, editors.

*Matrix Information Geometry (Revised Invited Papers)*. Springer, 2012.



Frank Nielsen and Richard Nock.

On the smallest enclosing information disk.

*Information Processing Letters (IPL)*, 105(3) :93–97, 2008.



Frank Nielsen and Richard Nock.

Quantum Voronoi diagrams and Holevo channel capacity for 1-qubit quantum states.

In *IEEE International Symposium on Information Theory (ISIT)*, pages 96–100, Toronto, Canada, July 2008. IEEE.



Frank Nielsen and Richard Nock.

Approximating smallest enclosing balls with applications to machine learning.

*Int. J. Comput. Geometry Appl.*, 19(5) :389–414, 2009.



Frank Nielsen and Richard Nock.

The dual Voronoi diagrams with respect to representational Bregman divergences.

In *International Symposium on Voronoi Diagrams (ISVD)*, pages 71–78, 2009.



Frank Nielsen and Richard Nock.

Entropies and cross-entropies of exponential families.

In *Proceedings of the International Conference on Image Processing, ICIP 2010, September 26-29, Hong Kong, China*, pages 3621–3624, 2010.



Frank Nielsen and Richard Nock.

Hyperbolic Voronoi diagrams made easy.

In *2013 13th International Conference on Computational Science and Its Applications*, pages 74–80. IEEE, 2010.



Frank Nielsen and Richard Nock.

Hyperbolic Voronoi diagrams made easy.

In *International Conference on Computational Science and its Applications (ICCSA)*, volume 1, pages 74–80, Los Alamitos, CA, USA, march 2010. IEEE Computer Society.



Frank Nielsen and Richard Nock.

Further results on the hyperbolic Voronoi diagrams.

*ISVD*, 2014.



Frank Nielsen and Richard Nock.

Visualizing hyperbolic Voronoi diagrams.

In *Symposium on Computational Geometry*, page 90, 2014.



Frank Nielsen, Paolo Piro, and Michel Barlaud.

Bregman vantage point trees for efficient nearest neighbor queries.

In *Proceedings of the 2009 IEEE International Conference on Multimedia and Expo (ICME)*, pages 878–881, 2009.



Richard Nock and Frank Nielsen.

Fitting the smallest enclosing bregman balls.

In *16th European Conference on Machine Learning (ECML)*, pages 649–656, October 2005.



Calyampudi Radhakrishna Rao.

Information and the accuracy attainable in the estimation of statistical parameters.

*Bulletin of the Calcutta Mathematical Society*, 37 :81–89, 1945.